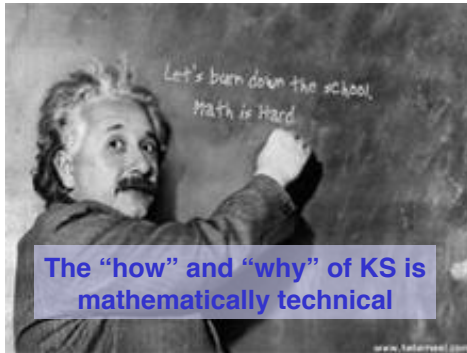


Sparse Linear Algebra: Introduction to Krylov subspace methods

- ▶ What is a Krylov subspace?
- ▶ Iterative versus direct methods
- ▶ Measuring the error
- ▶ Conjugate Gradient
 - Mathematical overview
 - Toy example
- ▶ Other KS methods
 - The GMRES and BiCGSTAB methods



Mathematical proofs and derivations are beyond the scope of this course

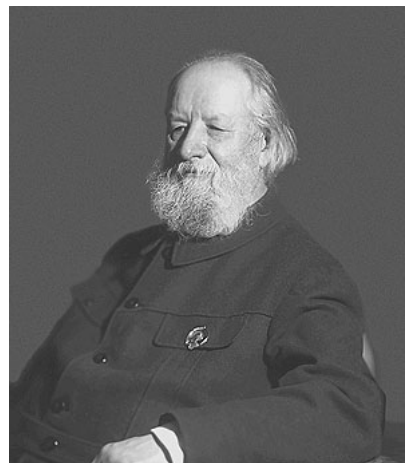
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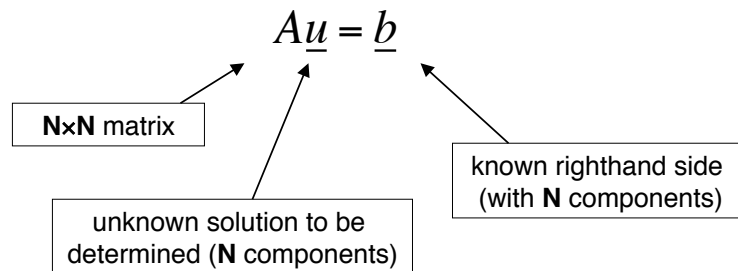
Implementing KS is easy

► **Alexei Krylov**

- Russian Naval engineer and applied mathematician
- Photo taken in 1930s Krylov in his 60s
- One of first people to classify the amount of work required for a given computation
- Krylov subspaces are constructed from linear systems



- ▶ Recall that a linear system (of size **N**) can be represented by a matrix equation, of the form:



- ▶ This is simply a representation of **N** equations linking **N** unknown quantities: $\underline{u}_1, \underline{u}_2, \dots, \underline{u}_N$.

- ▶ Definition: The Krylov subspace is said to be **spanned** by the vectors:

$$\left\{ \underline{v}, A\underline{v}, \dots, A^{N-1}\underline{v} \right\}$$

- ▶ Krylov sub-space is a property of matrix and vector
- ▶ Repeated application of Matrix **A** to \underline{v}

- ▶ Note that we never perform “matrix-matrix” operations
- ▶ Only ever need “matrix-vector” multiplications
- ▶ Multiplying a matrix by a vector always produces another vector
- ▶ If label the KS vectors as

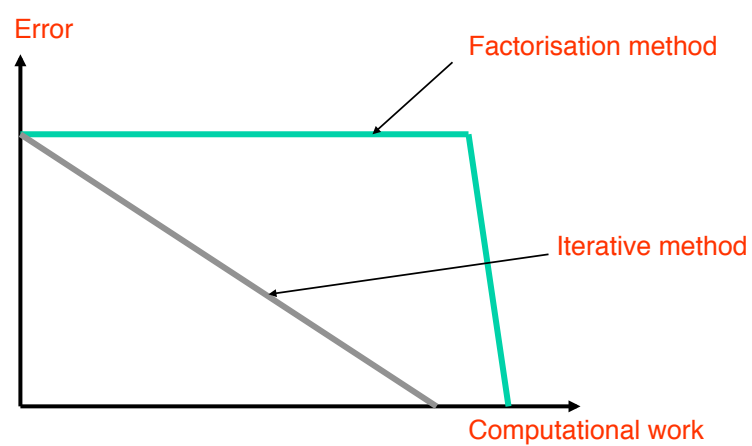
$$\{v_1, v_2 = Av_1, v_3 = A^2v_1, \dots, v_N = A^{N-1}v_1\}$$

- ▶ considering the third vector for example, we note that
$$v_3 = A^2v_1 = A(Av_1) = Av_2$$
- ▶ So even though we have an A^2 term in there, we never calculate it explicitly.

- ▶ Iterative methods
 - do not modify source matrix, involve matrix only through matrix-vector multiplication (possibly with transpose).
 - preserve sparsity and structure
 - Memory: Good for storage
 - progressively refine solution allowing user to impose accuracy constraints interactively
 - operate on individual righthand sides

► Direct (factorisation) methods

- act on source matrix, destroying structure such as sparsity
- involve redundant calculations on zero elements
 - Memory: storage of zero elements
- produce fixed accuracy solutions in a prescribed number of steps
- can be used efficiently for multiple righthand sides



Given a vector \underline{v} , how close is it to solution \underline{u} ?

$$\begin{aligned}\underline{\text{error}} &= \underline{u} - \underline{v} \\ &= A^{-1}(\underline{b} - A\underline{v})\end{aligned}$$

residual or backward error

"Take a candidate solution, apply matrix \mathbf{A} to it, and see how close it is to \underline{b} ."

$$\underline{\text{residual}} = \underline{b} - A\underline{v}$$

We need a way to quantify size of a vector -- a *norm*.

A number of different versions exist. We will consider the *Euclidean* (or L_2) norm:

$$\|\underline{v}\|_2 = \sqrt{\sum_{i=1}^N |v_i|^2}$$

Pythagoras' Theorem!

Thus, one can quantify the residual:

$$r = \frac{\|b - A\underline{v}\|_2}{\|b\|_2}$$

normalising factor

Facts:

- ▶ If $\mathbf{r}=\mathbf{0}$, then $\underline{\mathbf{v}}$ is exact solution.
- ▶ If \mathbf{r} is "small", then $\underline{\mathbf{v}}$ is likely to be close to solution.

- ▶ KS methods are class of iterative *search algorithms*.
- ▶ At iteration \mathbf{k} , take existing candidate solution $\underline{\mathbf{v}}_k$ and improve it by "minimising error" in prescribed direction $\underline{\mathbf{p}}_k$:

$$\underline{\mathbf{v}}_{k+1} = \underline{\mathbf{v}}_k + \alpha \underline{\mathbf{p}}_k$$

new candidate solution

old candidate solution

search direction

size of correction

Search algorithm = minimisation problem

For example, aim to minimise:

$$\phi(\underline{v}) = \frac{1}{2} \underline{v}^T \underline{A} \underline{v} - \underline{v}^T \underline{b}$$

[Quadratic form]

This is equivalent to minimising the error:

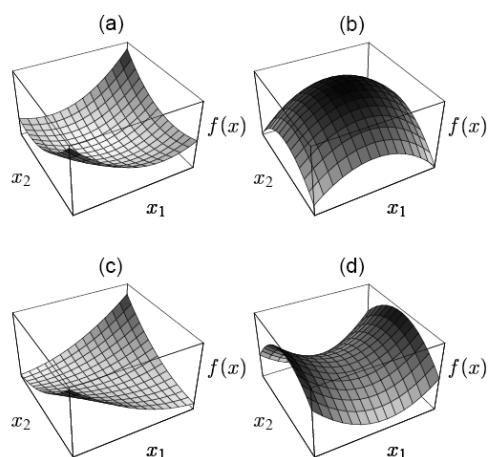
$$\|\underline{u} - \underline{v}\|$$

[only true if **A** is symmetric!]

The Quadratic Form

- (a) Positive definite
- (b) Negative definite
- (c) Singular
- (d) Indefinite

Choose search
directions to
locate the
minimum of ϕ



Need to solve minimisation problem (for new search direction). Choice of

$$\alpha = \frac{\overbrace{\underline{r}_k \cdot \underline{r}_k}^{\text{scalar product}}}{\underbrace{\underline{p}_k \cdot \underline{A} \underline{p}_k}_{\text{matrix multiplication}}}$$

Diagram annotations: "residual" points to \underline{r}_k ; "scalar product" points to $\underline{r}_k \cdot \underline{r}_k$; "matrix multiplication" points to $\underline{p}_k \cdot \underline{A} \underline{p}_k$.

will minimise:

$$\phi(\underline{v}) \quad \text{where} \quad \underline{v} = \underline{v}_k + \alpha \underline{p}_k$$

Diagram annotation: "New search direction" points to \underline{p}_k .

The scalar (dot) product of two vectors is:

$$\underline{u} \cdot \underline{v} = u_1 \times v_1 + \dots + u_N \times v_N$$

Notice that:

$$\|\underline{u}\|_2 = \sqrt{\underline{u} \cdot \underline{u}}$$

We say two vectors \underline{u} and \underline{v} are orthogonal if

$$\underline{u} \cdot \underline{v} = 0$$

and conjugate with respect to the matrix \mathbf{A} if

$$\underline{u} \cdot \underline{A} \underline{v} = 0$$

Need to generate "good spread" of search directions. Obvious choice is:

$$\underline{p}_{k+1} = \underline{r}_{k+1}$$

Better choice is:

$$\underline{p}_{k+1} = \underline{r}_{k+1} + \beta \underline{p}_k$$

Scalar β chosen to give a "good spread" of conjugate search directions.

[search directions are mutually conjugate]

- ▶ **Effective:** Method should minimise error (or something associated to error).
- ▶ **Bounded convergent:** Method should search solution space effectively, converging within pre-determined number of steps.
- ▶ **Efficient:** Cost of individual iteration should be small (and consistent).
- ▶ **Progressive:** Each iteration should improve the solution.

The Conjugate Gradient (CG) method:

- ▶ converges to solution of equation; effective ✓
- ▶ converges in $\leq N$ iterations; bounded convergent ✓
- ▶ requires 1 matrix-vector multiplication and 2 scalar products per iteration; efficient ✓
- ▶ improves the solution at each iteration. progressive ✓

But only for symmetric, positive definite matrices!

Set $k=0$ and choose \underline{v}_0 .
 Compute $\underline{r}_0 = \underline{b} - \underline{A}\underline{v}_0$, set $\underline{p}_0 = \underline{r}_0$. initial setup

While ($k < \text{maxiter}$)

$\alpha = \underline{r}_k \cdot \underline{r}_k / \underline{p}_k \cdot \underline{A}\underline{p}_k$ minimisation: compute correction and apply
 $\underline{v}_{k+1} = \underline{v}_k + \alpha \underline{p}_k$
 $\underline{r}_{k+1} = \underline{r}_k - \alpha \underline{A}\underline{p}_k$
 if ($\|\underline{r}_{k+1}\|_2 / \|\underline{b}\|_2 < \text{tol}$) break test new solution -- are we close enough?
 $\beta = \underline{r}_{k+1} \cdot \underline{r}_{k+1} / \underline{r}_k \cdot \underline{r}_k$ compute new search direction
 $\underline{p}_{k+1} = \underline{r}_{k+1} + \beta \underline{p}_k$
 $k = k + 1$
 end while

Recall the 'apples and pears' example:

- ▶ 2 apples and 3 pears costs 40p
- ▶ 3 apples and 5 pears costs 65p.

Solution is apples cost 5p, pears cost 10p.

Associated linear system is:

$$\begin{pmatrix} 2 & 3 \\ 3 & 5 \end{pmatrix} \begin{pmatrix} a \\ p \end{pmatrix} = \begin{pmatrix} 40 \\ 65 \end{pmatrix}$$

A symmetric matrix!

Set $k=0$ and choose \underline{v}_0 .

Compute $\underline{r}_0 = \underline{b} - \underline{A}\underline{v}_0$, set $\underline{p}_0 = \underline{r}_0$.

initial setup

While ($k < \text{maxiter}$)

$$\alpha = \underline{r}_k \cdot \underline{r}_k / \underline{p}_k \cdot \underline{A} \underline{p}_k$$

$$\underline{v}_{k+1} = \underline{v}_k + \alpha \underline{p}_k$$

$$\underline{r}_{k+1} = \underline{r}_k - \alpha \underline{A} \underline{p}_k$$

if ($\|\underline{r}_{k+1}\|_2 / \|\underline{b}\|_2 < \text{tol}$) break

$$\beta = \underline{r}_{k+1} \cdot \underline{r}_{k+1} / \underline{r}_k \cdot \underline{r}_k$$

$$\underline{p}_{k+1} = \underline{r}_{k+1} + \beta \underline{p}_k$$

$k=k+1$

end while

$$\begin{pmatrix} 2 & 3 \\ 3 & 5 \end{pmatrix} \begin{pmatrix} a \\ p \end{pmatrix} = \begin{pmatrix} 40 \\ 65 \end{pmatrix}$$

Initial setup:

Guess $a=0$ and $p=0$;

Compute residual: $\underline{r}_0 = (40, 65) = \underline{p}_0$.

Set $k=0$ and choose \underline{v}_0 .

Compute $\underline{r}_0 = \underline{b} - \underline{A}\underline{v}_0$, set $\underline{p}_0 = \underline{r}_0$.

While ($k < \text{maxiter}$)

$$\alpha = \underline{r}_k \cdot \underline{r}_k / \underline{p}_k \cdot \underline{A}\underline{p}_k$$

$$\underline{v}_{k+1} = \underline{v}_k + \alpha \underline{p}_k$$

$$\underline{r}_{k+1} = \underline{r}_k - \alpha \underline{A}\underline{p}_k$$

if $(\|\underline{r}_{k+1}\|_2 / \|\underline{b}\|_2 < \text{tol})$ break

$$\beta = \underline{r}_{k+1} \cdot \underline{r}_{k+1} / \underline{r}_k \cdot \underline{r}_k$$

$$\underline{p}_{k+1} = \underline{r}_{k+1} + \beta \underline{p}_k$$

$k=k+1$

end while

minimisation: compute
correction and apply

► $\underline{r}_0 \cdot \underline{r}_0 = 40^2 + 65^2 = 5,825$

Matrix-vector multiply

$$\mathbf{A}\underline{p}_0 = \begin{pmatrix} 2 & 3 \\ 3 & 5 \end{pmatrix} \begin{pmatrix} 40 \\ 65 \end{pmatrix} = \begin{pmatrix} 2 \times 40 + 3 \times 65 \\ 3 \times 40 + 5 \times 65 \end{pmatrix} = \begin{pmatrix} 275 \\ 445 \end{pmatrix}$$

$$\underline{p}_0 \cdot \mathbf{A}\underline{p}_0 = 40 \times 275 + 65 \times 445 = 39,925$$

$$\alpha = 5,825 \div 39,925 = 0.145899 \text{ (6 s.f.)}$$

$$\underline{v}_1 = \underline{v}_0 + 0.145899 \times \underline{p}_0 = (5.83594, 9.48341)$$

Nearly there in one step!

Set $k=0$ and choose \underline{v}_0 .

Compute $\underline{r}_0 = \underline{b} - \mathbf{A}\underline{v}_0$, set $\underline{p}_0 = \underline{r}_0$.

While ($k < \text{maxiter}$)

$$\alpha = \underline{r}_k \cdot \underline{r}_k / \underline{p}_k \cdot \mathbf{A}\underline{p}_k$$

$$\underline{v}_{k+1} = \underline{v}_k + \alpha \underline{p}_k$$

$$\underline{r}_{k+1} = \underline{r}_k - \alpha \mathbf{A}\underline{p}_k$$

if ($\|\underline{r}_{k+1}\|_2 / \|\underline{b}\|_2 < \text{tol}$) break

test new solution -- are we close enough?

$$\beta = \underline{r}_{k+1} \cdot \underline{r}_{k+1} / \underline{r}_k \cdot \underline{r}_k$$

$$\underline{p}_{k+1} = \underline{r}_{k+1} + \beta \underline{p}_k$$

$$k = k + 1$$

end while

$$\underline{\mathbf{r}}_1 = \underline{\mathbf{r}}_0 - \alpha \mathbf{A} \underline{\mathbf{p}}_0 = \begin{pmatrix} 40 \\ 65 \end{pmatrix} - 0.145899 \begin{pmatrix} 275 \\ 445 \end{pmatrix} = \begin{pmatrix} -0.122104 \\ 0.0751409 \end{pmatrix}$$

$$\|\underline{\mathbf{r}}_1\|_2 = 0.00187852$$

$$\|\underline{\mathbf{b}}\|_2 = 5825$$

$$\|\underline{\mathbf{r}}_1\|_2 / \|\underline{\mathbf{b}}\|_2 = 3.52885\text{e-}06 \text{ (very close to stopping!)}$$

Set $\mathbf{k}=0$ and choose $\underline{\mathbf{v}}_0$.

Compute $\underline{\mathbf{r}}_0 = \underline{\mathbf{b}} - \mathbf{A}\underline{\mathbf{v}}_0$, set $\underline{\mathbf{p}}_0 = \underline{\mathbf{r}}_0$.

While ($\mathbf{k} < \text{maxiter}$)

$$\alpha = \underline{\mathbf{r}}_k \cdot \underline{\mathbf{r}}_k / \underline{\mathbf{p}}_k \cdot \mathbf{A} \underline{\mathbf{p}}_k$$

$$\underline{\mathbf{v}}_{k+1} = \underline{\mathbf{v}}_k + \alpha \underline{\mathbf{p}}_k$$

$$\underline{\mathbf{r}}_{k+1} = \underline{\mathbf{r}}_k - \alpha \mathbf{A} \underline{\mathbf{p}}_k$$

if ($\|\underline{\mathbf{r}}_{k+1}\|_2 / \|\underline{\mathbf{b}}\|_2 < \text{tol}$) break

$$\beta = \underline{\mathbf{r}}_{k+1} \cdot \underline{\mathbf{r}}_{k+1} / \underline{\mathbf{r}}_k \cdot \underline{\mathbf{r}}_k$$

$$\underline{\mathbf{p}}_{k+1} = \underline{\mathbf{r}}_{k+1} + \beta \underline{\mathbf{p}}_k$$

$$\mathbf{k} = \mathbf{k} + 1$$

end while

compute new search
direction

$$\beta = \underline{r}_1 \cdot \underline{r}_1 / \underline{r}_0 \cdot \underline{r}_0 = 0.00187852/5825$$

$$= 3.52885e-06$$

$$\underline{p}_1 = \underline{r}_1 + \beta \underline{p}_0$$

$$\underline{p}_1 = \begin{pmatrix} -0.122104 \\ 0.0751409 \end{pmatrix} + 3.52885e-06 \begin{pmatrix} 40 \\ 65 \end{pmatrix} = \begin{pmatrix} -0.121963 \\ 0.0753703 \end{pmatrix}$$

Start next iteration

$$\begin{pmatrix} 2 & 3 \\ 3 & 5 \end{pmatrix} \begin{pmatrix} a \\ p \end{pmatrix} = \begin{pmatrix} 40 \\ 65 \end{pmatrix}$$

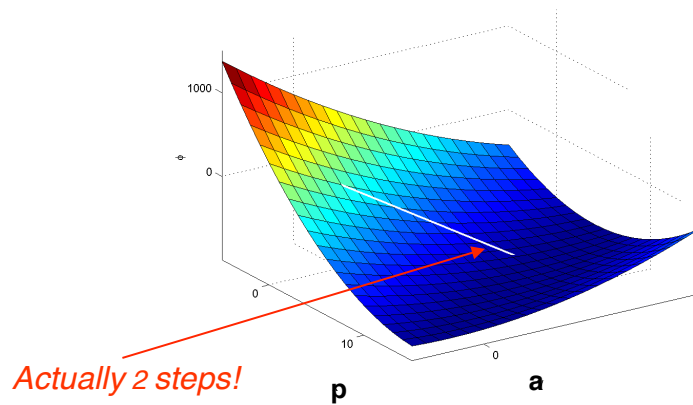
$$\underline{p}_1 \cdot \underline{A} \underline{p}_1 = 0.00299902$$

$$\alpha = 0.0205555 \div 0.00299902 = 6.85408$$

$$\underline{v}_2 = \underline{v}_1 + 6.85408 \times \underline{p}_1 = (5.000, 10.000)$$

$$\|\underline{r}_2\|_2 / \|\underline{b}\|_2 = 5.9692e-20$$

EXACT SOLUTION!



► Convergence

- $N=2 \rightarrow k \leq 2$ Exact solution after 2 iterations
- $\|r^2\|$ is very small

► Precision

- Rounding errors will produce only approximate solution
- Single versus double precision

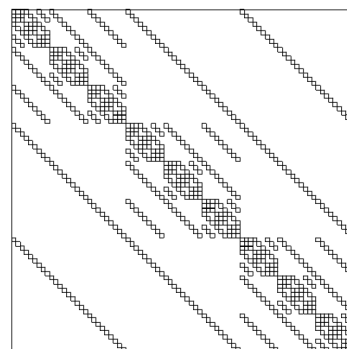
► Real applications

- Use finite precision
- Stop when residual $< \epsilon$
- Approximate solutions
- $N_{\text{iter}} \ll N$

- ▶ Cost
 - involves one matrix-vector multiplication and two scalar products per iteration.
- ▶ does not modify the source matrix **A**
- ▶ requires user to provide routine for matrix-vector product

But only for symmetric, positive definite matrices!

- ▶ Particle physics example
 - Quantum Chromodynamics (QCD)
- ▶ Matrix is highly structured
 - $N \sim O(10^3) \rightarrow$
 - SM Pickles PhD thesis UoE 1998
- ▶ Real calculation $N > O(10^7)$
- ▶ CG converge $\leq N$ iters
 - 4096 cores of HECToR
 - $O(10^5)$ iters takes 30 mins
 - Calculation done in double-prec appropriate finite residual 10^{-8}



- ▶ For more general class of matrix, cannot achieve all *desirable properties*, though can fulfil most.
- ▶ Two popular methods considered:
 - Generalised Minimum RESidual method, **GMRES**
 - Bi-Conjugate Gradient method with STABilisation, **BiCGSTAB**

The GMRES method:

- ▶ converges to solution of equation; effective ✓
- ▶ converges in $\leq N$ iterations; bounded convergent ✓
- ▶ requires expensive orthogonalisation of search directions at each iteration, -- depends on all previous iterations/search directions -- computationally and memory intensive; efficient ✗
- ▶ improves the solution at each iteration. progressive ✓

The BiCGSTAB method:

- ▶ may not converge to solution; effective ×
- ▶ convergence is unbounded; bounded convergent ×
- ▶ requires 2 matrix multiplications and 4 scalar products per iteration; efficient ✓
- ▶ not guaranteed to improve solution at each iteration. progressive ×

However, very often it works!

Set $k=0$ and choose \underline{v}_0 .

Compute $\underline{r}_0 = \underline{b} - A\underline{v}_0$, $\underline{l}_0 = \underline{r}_0$, $\underline{q}_0 = \underline{p}_0 = \underline{0}$, $\rho_0 = \alpha = \omega_0 = 1$.

While ($k < \text{maxiter}$)

$$\rho_{k+1} = \underline{l}_0 \cdot \underline{r}_k, \quad \beta = (\rho_{k+1} / \rho_k) \times (\alpha / \omega_k)$$

$$\underline{p}_{k+1} = \underline{r}_k + \beta(\underline{p}_k - \omega_k \underline{q}_k), \quad \underline{q}_{k+1} = A\underline{p}_{k+1}$$

$$\alpha = \rho_{k+1} / (\underline{l}_0 \cdot \underline{q}_{k+1}), \quad \underline{s} = \underline{r}_k - \alpha \underline{q}_{k+1}, \quad \underline{t} = A\underline{s}$$

$$\omega_{k+1} = (\underline{t} \cdot \underline{s}) / (\underline{t} \cdot \underline{t})$$

$$\underline{v}_{k+1} = \underline{v}_k + \alpha \underline{p}_{k+1} + \omega_{k+1} \underline{s}$$

$$\underline{r}_{k+1} = \underline{s} - \omega_{k+1} \underline{t}$$

if ($\|\underline{r}_{k+1}\|_2 / \|\underline{b}\|_2 < \text{tol}$) break

$k = k + 1$

end while

\underline{l} is not updated (i.e. \underline{l}_0 not \underline{l}_k)

Three scalars
initialised to 1

4 update vectors
2 temporary vectors

Set $k=0$ and choose \underline{v}_0 .

Compute $\underline{r}_0 = \underline{b} - A\underline{v}_0$, $\underline{l}_0 = \underline{r}_0$, $\underline{q}_0 = \underline{p}_0 = \underline{0}$, $\rho_0 = \alpha = \omega_0 = 1$.

While ($k < \text{maxiter}$)

$$\rho_{k+1} = \underline{l}_0 \cdot \underline{r}_k, \quad \beta = (\rho_{k+1} / \rho_k) \times (\alpha / \omega_k)$$

$$\underline{p}_{k+1} = \underline{r}_k + \beta(\underline{p}_k - \omega_k \underline{q}_k), \quad \underline{q}_{k+1} = A\underline{p}_{k+1}$$

$$\alpha = \rho_{k+1} / (\underline{l}_0 \cdot \underline{q}_{k+1}), \quad \underline{s} = \underline{r}_k - \alpha \underline{q}_{k+1}, \quad \underline{t} = A\underline{s}$$

$$\omega_{k+1} = (\underline{t} \cdot \underline{s}) / (\underline{t} \cdot \underline{t})$$

$$\underline{v}_{k+1} = \underline{v}_k + \alpha \underline{p}_{k+1} + \omega_{k+1} \underline{s}$$

$$\underline{r}_{k+1} = \underline{s} - \omega_{k+1} \underline{t}$$

if ($\|\underline{r}_{k+1}\|_2 / \|\underline{b}\|_2 < \text{tol}$) break

$k = k + 1$

end while

2 matrix-vector
operations

PNA L14 Advanced KS methods

Set $k=0$ and choose \underline{v}_0 .

Compute $\underline{r}_0 = \underline{b} - A\underline{v}_0$, $\underline{l}_0 = \underline{r}_0$, $\underline{q}_0 = \underline{p}_0 = \underline{0}$, $\rho_0 = \alpha = \omega_0 = 1$.

While ($k < \text{maxiter}$)

$$\rho_{k+1} = \underline{l}_0 \cdot \underline{r}_k, \quad \beta = (\rho_{k+1} / \rho_k) \times (\alpha / \omega_k)$$

$$\underline{p}_{k+1} = \underline{r}_k + \beta(\underline{p}_k - \omega_k \underline{q}_k), \quad \underline{q}_{k+1} = A\underline{p}_{k+1}$$

$$\alpha = \rho_{k+1} / (\underline{l}_0 \cdot \underline{q}_{k+1}), \quad \underline{s} = \underline{r}_k - \alpha \underline{q}_{k+1}, \quad \underline{t} = A\underline{s}$$

$$\omega_{k+1} = (\underline{t} \cdot \underline{s}) / (\underline{t} \cdot \underline{t})$$

$$\underline{v}_{k+1} = \underline{v}_k + \alpha \underline{p}_{k+1} + \omega_{k+1} \underline{s}$$

$$\underline{r}_{k+1} = \underline{s} - \omega_{k+1} \underline{t}$$

if ($\|\underline{r}_{k+1}\|_2 / \|\underline{b}\|_2 < \text{tol}$) break

$k = k + 1$

end while

4 scalar products

1 vector norm

PNA L14 Advanced KS methods

- ▶ Other KS methods exist and are used: Steepest Descent; MINRES; GMRES with restart; Conjugate Residual; BiCG.
- ▶ Generally, follow same principles.
- ▶ In practice, call numerical library (NAG, PetSC, ARPACK)
- ▶ Still need to provide the matrix-vector multiplier!

- ▶ Krylov subspace method is a form of iterative improvement. **You decide when to stop!**
- ▶ Replace linear solver with minimisation method
- ▶ For SPD matrices, standard implementation is Conjugate Gradient method.
- ▶ Otherwise, have a choice between **GMRES** and **BiCGSTAB**.
- ▶ Other methods do exist.

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