



Objectives of the coming three lectures:

- understand the basic principles of UPC
- motivation behind PGAS
- learn about data distribution, synchronisation
- advanced features (dynamic memory allocation, collectives)

→ Practicals will try and emphasise the most important aspects of UPC



Unified Parallel C

Parallel extension to ISO C 99, adding

- o explicit parallelism
- o global shared address space
- synchronisation

Both commercial and open source compilers available

- o Cray, IBM, SGI, HP
- o GWU, LBNL, GCCUPC

UPC and the World of PGAS



UPC != PGAS

- PGAS is a programming model
- UPC is one implementation of this model

Many other implementations

- Language extension: Coarray Fortran
- New languages: Chapel, X10, Fortress, Titanium
- PGAS-like libraries: OpenSHMEM, Global Arrays

All implementations are different, but follow the same model!



UPC uses threads that operate independently in a SPMD fashion

→ threads execute the same UPC program

Identifiers that return information about the program environment:

THREADS: holds total number of threads

MYTHREAD: stores thread index

→ index runs from 0 to THREADS-1



```
#include <upc.h>
#include <stdio.h>

void main() {
   printf("Thread %d of %d says: Hello!", MYTHREAD, THREADS);
}
```

Private vs. shared memory space



Concept of two memory spaces: private and shared

objects declared in **private** memory space are only accessible by a single thread

objects declared in **shared** memory space are accessible by all threads

→ shared memory space is used to communicate information between threads



private variables declared as normal C variable

o multiple instances of variable will exist

```
int x; // private variable
```

shared variables declared with **shared** qualifier

- o only allocated once, in shared memory space
- o accessible by all threads

```
shared int y; // shared variable
```



If shared variable is scalar, space only allocated on thread 0

```
int x;
shared int y;
```

thread 0	thread 1	thread 2	thread 3
x	x	x	x
У			

Private Memory Space

Shared Memory Space



all threads can directly access shared data, even if it resides in a remote location

UPC creates logical partitioning of the shared memory space

- → objects have *affinity* to one thread
- → shared scalars always have affinity to thread 0

better performance if a thread access data to which it has affinity

→ always keep data locality and affinity in mind



If a shared variable is an array, space allocated across shared memory space in a *cyclic* fashion by default

```
int x;
shared int y[16];
```

thread 0	thread 1	thread 2	thread 3
X	x	x	x
y[0]	y[1]	y[2]	y[3]
y[4]	y[5]	y[6]	y[7]
y[8]	y[9]	y[10	y[11
y[12]	y[13]	y[14]	y[15]

Private Memory Space

Shared Memory Space

Shared array distribution (2)

int x;



Change data layout by adding a "blocking factor" to shared arrays shared[blocksize] type array[n]

```
shared[2] int y[16];
thread 0
            thread 1
                        thread 2
                                     thread 3
                                         X
    X
                X
                             X
   y[0]
               y[2]
                            y[4]
                                        y[6]
   y[1]
               y[3]
                            y[5]
                                        y[7]
   y[8]
               y[10]
                                       y[14]
                           y[12]
   y[9]
               y[11]
                           y[13]
                                       y[15]
```

Private Memory Space

Shared Memory Space

Work sharing



Shared data means shared workload!

If shared data is distributed between threads, threads can distribute work on this data between them

UPC has built-in mechanism for explicitly distributing and sharing work



Statement for work distribution

- allows loop assignment of tasks to threads
- o parallel for loop

4th parameter defines affinity to thread

 if "affinity % THREADS" matches MYTHREAD, execute iteration for that THREAD

```
upc_forall(expression; expression; affinity)
```

Condition: iterations of upc_forall must be independent!



```
#define N 10 * THREADS
shared int vector1[N];
shared int vector2[N];
shared int sum[N];
void main() {
   int i;
  for(i=0; i<N; i++){
      sum[i] = vector1[i] + vector2[i];
```



```
#define N 10 * THREADS
shared int vector1[N];
shared int vector2[N];
shared int sum[N];
                                        evaluated as i%THREADS
void main() {
   int i;
   upc_forall(i=0; i<N; i++; i){</pre>
      sum[i] = vector1[i] + vector2[i];
```

Side effects of shared data



Holding data in a shared memory space has implications

- 1) the lifetime of shared data needs to extend beyond the scope it was defined in (unless this is program scope)
 - → storage duration

- 2) the shared data needs to be keep up-to-date
 - → synchronisation

Storage duration of shared objects



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Shared objects cannot have automatic storage duration

any variable defined inside a function!

Why?

SPMD model means a shared variable may be accessed outside lifetime of the function!

Conclusion

shared variables must either

- have file scope;
- o or be declared as **static** if defined inside a function.



ensures shared objects are accessible throughout program execution

→ objects are not linked to the scope of a thread

→ objects will not simply "disappear" after a thread exists the scope in which the object was defined



```
#define max(a,b) (((a)>(b)) ? (a) : (b))
shared int maximum[THREADS];
                                                  Here: shared variables have file scope!
shared int globalMax = 0;
shared int a[THREADS*10];
void main(int argc, char **argv) {
        ... // initialise array a
        upc forall(int i=0; i<THREADS*10; i++; i){</pre>
                   maximum[MYTHREAD] = max(maximum[MYTHREAD], a[i]);
        }
        if (MYTHREAD == 0){
           for (int thread=0; thread<THREADS; thread++){</pre>
                     globalMax = max(globalMax,maximum[thread]);
                              This code will not work!!
```

Synchronisation



Ensure all threads reach same point in execution

necessary for memory and data consistency

Barriers used for synchronisation

- blocking
- split-phase (non-blocking)

```
upc_barrier  → blocking

upc_notify, upc_wait → non-blocking
```

Example: maximum of an array



```
#define max(a,b) (((a)>(b)) ? (a) : (b))
shared int maximum[THREADS];
                                                   Here: shared variables have file scope!
shared int globalMax = 0;
shared int a[THREADS*10];
void main(int argc, char **argv) {
        ... // initialise array a
        upc barrier;
        upc forall(int i=0; i<THREADS*10; i++; i){</pre>
                   maximum[MYTHREAD] = max(maximum[MYTHREAD], a[i]);
                                                   Ensure all threads found local maximum
        upc barrier;
        if (MYTHREAD == 0){
           for (int thread=0; thread<THREADS; thread++){</pre>
                     globalMax = max(globalMax,maximum[thread]);
                                                   Makes sure globalMax is found before
        upc_barrier;
                                                   being used
```

References



UPC Language Specification (Version 1.2):
 http://upc.gwu.edu/docs/upc_specs_1.2.pdf

UPC homepage:

http://upc.gwu.edu/

- GCCUPC compiler:
- http://www.gccupc.org