



# PA25 | 6: COVERAGE AND INTERVENTION 3: FAULT-BASED (MUTATION) TESTING

**Kai Petersen,**  
Blekinge Institute of Technology

Based on Gordan Fraser's slides, Mutation testing  
as well as Ammann & Offutt - Introduction to Software Testing

# KEY QUESTION

- Tests show only the presence, not the absence of faults
- We have seen how we find tests based on coverage (considering requirements and source code structure)
- How can we know **how good our tests are?**
  - Calculation:  $\frac{\text{\#found faults}}{\text{\#total faults}}$ , but: **we do not know the total number of faults**
- How to solve this?
  - Use known or inserted faults to evaluate the goodness of test cases, also known as mutation testing

# FAULT-BASED TESTING

- Mutation testing utilizes faults in the software to:
  - generate test data
  - evaluation of testing effectiveness



# DEFINITIONS



- **Mutant:** Single change in a program leads to a new mutant that is **syntactically** legal
- We would like to introduce mutants that lead to a different behaviour of the program, but the program should still compile (syntactical correctness)
- Goal: **Kill mutants**
  - If a test case causing the **mutant to fail means** that the **mutant is dead**
  - A mutant is called **stubborn** if the existing set of test cases are insufficient to kill it
  - Test cases that identified mutants are useful
  - Consensus that mutants increase the strength of test cases

# MUTATION TESTING EXAMPLE

**Test  
verdict**

```
int do_something(int x, int y)
{
    if(x < y)
        return x+y;
    else
        return x*y;
}
```

```
int a = do_something(5, 10);
assertEquals(a, 15);
```

**Pass**

```
int do_something(int x, int y)
{
    if(x < y)
        return x-y;
    else
        return x*y;
}
```

**Mutant with  
an operator  
change**

```
int a = do_something(5, 10);
assertEquals(a, 15);
```

**Fail**

# GENERATING MUTANTS

Mutation operators



# ABSOLUTE VALUE INSERTION

**Def:** Each arithmetic expression (and subexpression) is modified by the functions `abs()`, `negAbs()`, and `failOnZero()`

```
int gcd(int x, int y) {  
    int tmp;  
    while(y != 0) {  
        tmp = x % y;  
        x = y;  
        y = tmp;  
    }  
    return x;  
}
```

**Abs()**

```
int gcd(int x, int y) {  
    int tmp;  
    while(y != 0) {  
        tmp = x % y;  
        x = abs(y);  
        y = tmp;  
    }  
    return x;  
}
```

**negAbs()**

```
int gcd(int x, int y) {  
    int tmp;  
    while(y != 0) {  
        tmp = x % y;  
        x = y;  
        y = tmp;  
    }  
    return -abs(x);  
}
```

**failOnZero()**

```
int gcd(int x, int y) {  
    int tmp;  
    while(y != 0) {  
        tmp = x % y;  
        x = y;  
        y = tmp;  
    }  
    return 0;  
}
```



# ARITHMETIC OPERATOR REPLACEMENT

**Def:** Each occurrence of one of the arithmetic operators  $+$ ,  $-$ ,  $*$ ,  $/$ ,  $**$ , and  $\%$  is replaced by each of the other operators. In addition, each is replaced by the special mutation operators `leftOp`, `rightOp`, and `mod`.

```
int gcd(int x, int y) {  
    int tmp;  
    while(y != 0) {  
        tmp = x % y;  
        x = y;  
        y = tmp;  
    }  
    return x;  
}
```

```
int gcd(int x, int y) {  
    int tmp;  
    while(y != 0) {  
        tmp = x * y;  
        x = y;  
        y = tmp;  
    }  
    return x;  
}
```

Also, left and right operators are dropped, i.e. `tmp = x;` and `tmp = y;`



# RELATIONAL OPERATOR REPLACEMENT

**Def:** Each occurrence of one of the relational operators ( $<$ ,  $\leq$ ,  $>$ ,  $\geq$ ,  $=$ ,  $\neq$ ) is replaced by each of the other operators and by falseOp and trueOp.

```
int gcd(int x, int y) {  
    int tmp;  
    while(y != 0) {  
        tmp = x % y;  
        x = y;  
        y = tmp;  
    }  
    return x;  
}
```

```
int gcd(int x, int y) {  
    int tmp;  
    while(y > 0) {  
        tmp = x % y;  
        x = y;  
        y = tmp;  
    }  
    return x;  
}
```

Also, true and false are put directly into the expression, i.e.  
while(true) as well as while(false)

# CONDITIONAL OPERATOR REPLACEMENT

**Def:** Each occurrence of each logical operator (and–&&, or–, and with no conditional evaluation–&, or with no conditional evaluation–|, not equivalent – ^) is replaced by each of the other operators; in addition, each is replaced by falseOp, trueOp, leftOp, and rightOp.

if (a && b)

if (a || b)

if (a & b)

if (a | b)

if (a ^ b)

if (false)

if (true)

if (a)

if (b)

# SHIFT OPERATOR REPLACEMENT

**Def:** Each occurrence of one of the shift operators <<, >>, and >>> is replaced by each of the other operators. In addition, each is replaced by the special mutation operator leftOp.

**m<<a**

**x = m >> a;**

**x = m >>> a;**

**x = m;**

## Excursion: The bitwise-shift operator

Syntax: [variable]<<[number of places]

```
int mult_power(int number, int power)
{
    return number<<power;
}
```

Example: 00001111

Shift by 2: 00111100

Distinction between signed and unsigned shift operator to be made.

<< and >> are signed shift operators

<<< and >>> are unsigned shift operators



# LOGICAL OPERATOR REPLACEMENT

**Def:** Each occurrence of each bitwise logical operator (bitwise and (&), bitwise or (|), and exclusive or (^)) is replaced by each of the other operators; in addition, each is replaced by leftOp and rightOp.

$x = m \& n;$

$x = m | n;$

$x = m ^ n;$

$x = m;$

$x = n;$



# ASSIGNMENT OPERATOR REPLACEMENT

**Def:** Each occurrence of one of the assignment operators ( $+=$ ,  $-=$ ,  $*=$ ,  $/=$ ,  $\%=$ ,  $\&=$ ,  $\|=$ ,  $\hat{=}$ ,  $\ll=$ ,  $\gg=$ ,  $\ggg=$ ) is replaced by each of the other operators.

$x += 3;$

$x -= 3;$

$x *= 3;$

$x /= 3;$

$x \%= 3;$

$x \&= 3;$

$x \|= 3;$

$x \hat{=} 3;$

$x \ll= 3;$

$x \gg= 3;$

$x \ggg= 3;$

# UNARY OPERATOR INSERTION

**Def:** Each unary operator (arithmetic  $+$ , arithmetic  $-$ , conditional  $!$ , logical  $\sim$ ) is inserted before each expression of the correct type.

$x = 3 * a;$

$x = 3 * +a;$

$x = 3 * -a;$

$x = +3 * a;$

$x = -3 * a;$



# UNARY OPERATOR DELETION

**Def:** Each unary operator (arithmetic  $+$ , arithmetic  $-$ , conditional  $!$ , logical  $\sim$ ) is deleted.

if  $!(a > -b)$

if  $(a > -b)$

if  $!(a > b)$

# SCALAR VARIABLE REPLACEMENT

**Def:** Each variable reference is replaced by every other variable of the appropriate type that is declared in the current scope.

```
x = a * b;
```

```
x = a * a;
```

```
a = a * b;
```

```
x = x * b;
```

```
x = a * x;
```

```
x = b * b;
```

```
b = a * b;
```



# GENERATING MUTANTS

Integration/Interface Mutation



# INTEGRATION MUTATION

- Integration mutation: Often misunderstandings on
  - **delivery side:** deliver what is not desired by the caller (e.g caller expects records in ascending order, but delivered in random order)
  - **receiver side:** recipient makes assumptions on what is delivered (e.g. expects to receive data in a specific format, which was not the case)
  - **Integration (interface) mutation:** Mutate connections between components



# TYPES OF MUTATION OPERATORS

- T1: Change a calling method by modifying the values sent to a called method
- T2: Change a calling method by modifying the call
- T3: Change a called method by modifying the values that enter and leave a method
- T4: Change a called method by modifying the return statements



## Integration parameter variable replacement:

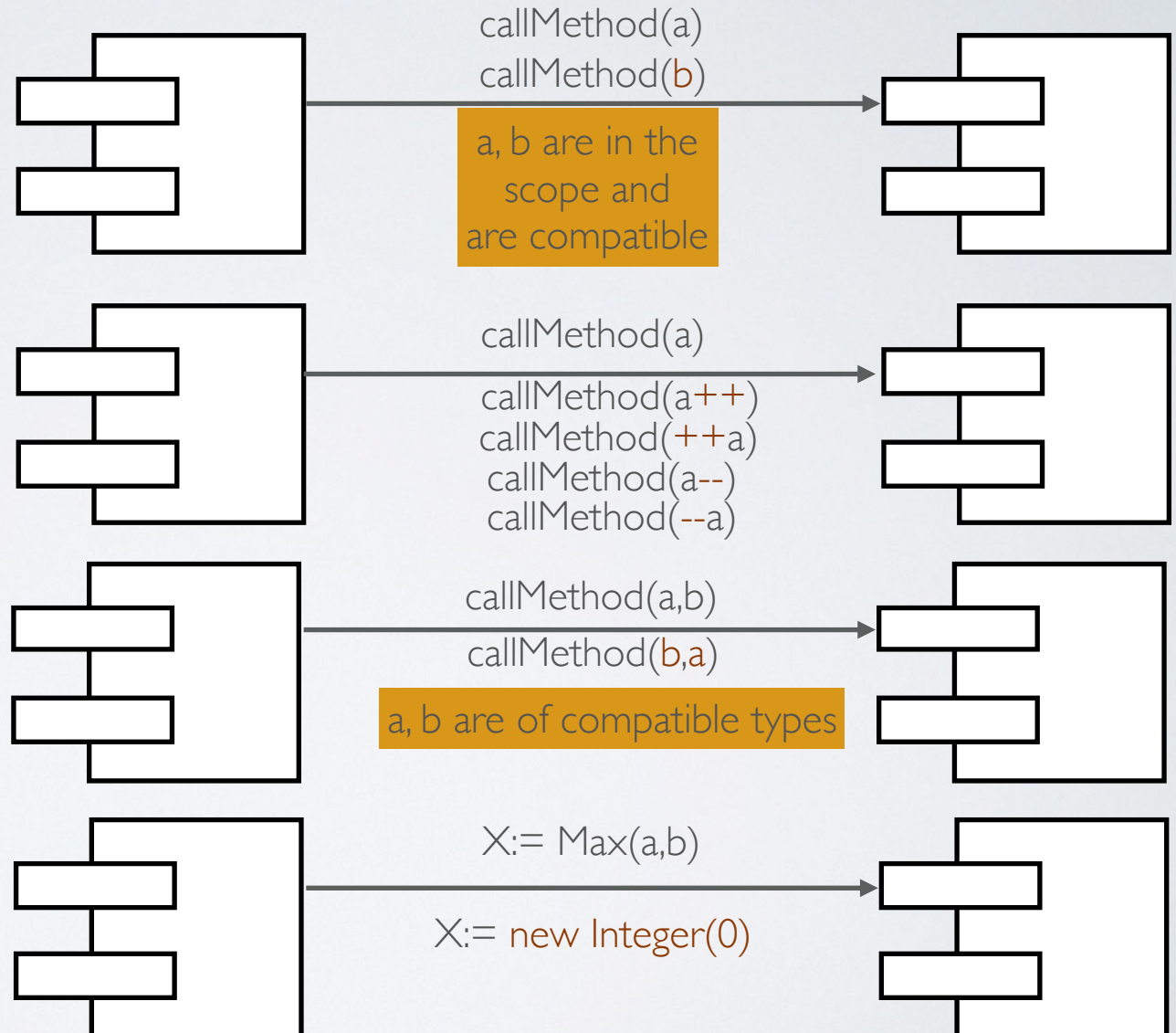
Each parameter in the method call is replaced by another variable in the scope of the method call

## Integration unary insertion:

Each expression in the method call is modified by inserting all possible unary operators in front and behind it

**Integration Parameter Exchange:** Each parameter in a method call is exchanged with each parameter of compatible types in that method call.

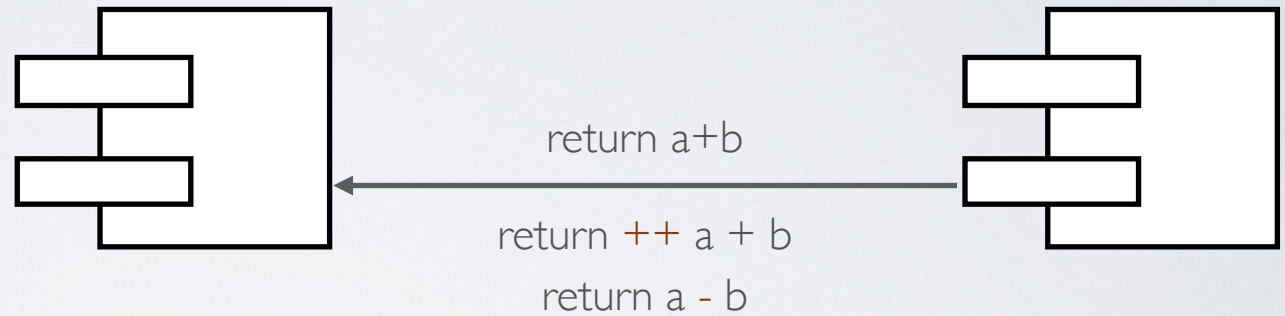
**Integration method call deletion:** Each method call is deleted. If the method returns a value and it is used in an expression, the method call is replaced with an appropriate constant value.





## Integration Return Expression

**Modification:** Each expression in each return statement in a method is modified by applying the unary operator insertion (UOI) and arithmetic operator replacement (AOR) operators.



# GENERATING MUTANTS

## OO Mutation



Earlier we only considered mutation within methods or functions,  
we now consider object oriented concepts related to  
information hiding, inheritance, polymorphism, dynamic binding, and  
method overloading



# EXCURSION: OO CONCEPTS

- **Encapsulation:** Information hiding = restriction to member variables (modification of variables through methods)
- **Method overriding:** Method of subclass can be redefined (different implementation), but has the same name, arguments, and result type
- **Variable hiding:** Variable defined in the child class with same name hides variable in the inheriting class
- **Polymorphism:** Interface that can take different types (templates/generics), e.g. in sorting one could hand over different object types (strings, ints, etc.)
- **Method overloading:** In the **same class** we use the same name for constructors and methods, but with different parameters
- **Method overriding:** Child class declares a method with **the same name** as in the mother class and hence overrides the method of the parent class



# OO MUTATIONS

## ENCAPSULATION

**Access Modifier Change:** The access level for each instance variable and method is changed to other access levels.

Original

myclass
private int myVar;

**Goal: Check whether the accessibility of variables is correct**

myclass
public int myVar;

myclass
protected int myVar;

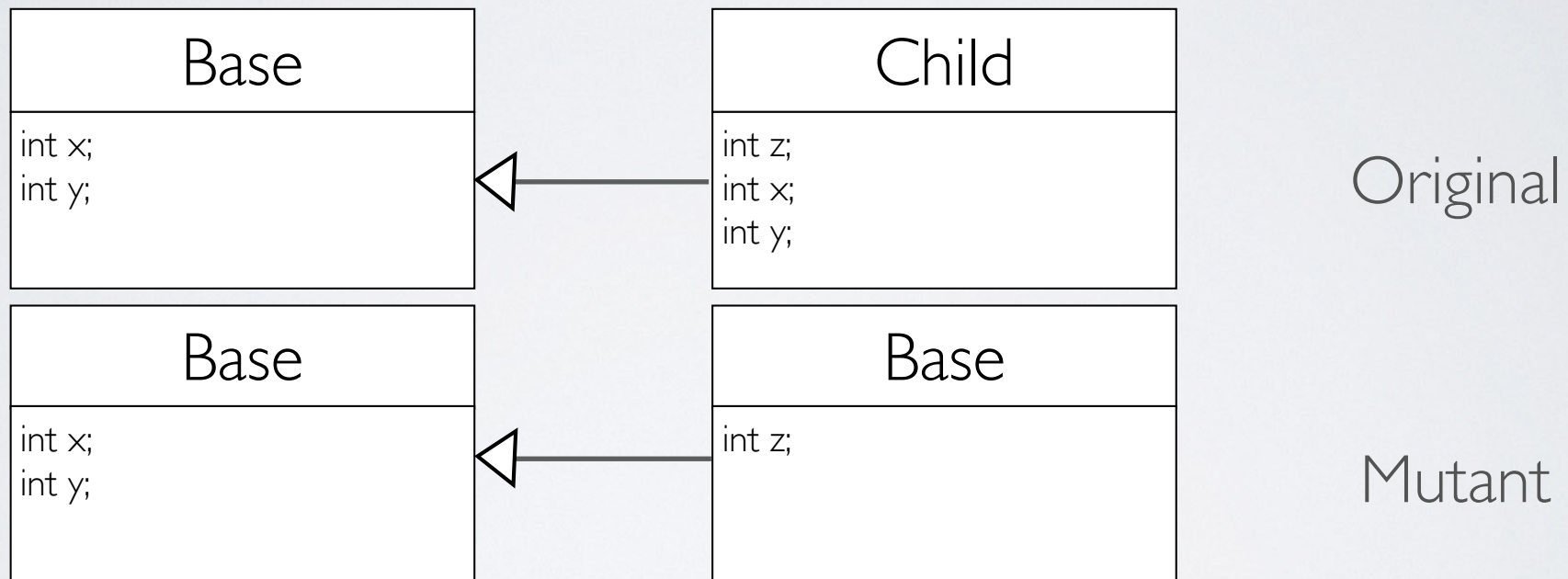
myclass
int myVar;

Mutants



# OO MUTATIONS INHERITANCE

**Hiding variable deletion:** Each declaration of an overriding or hiding variable is deleted.

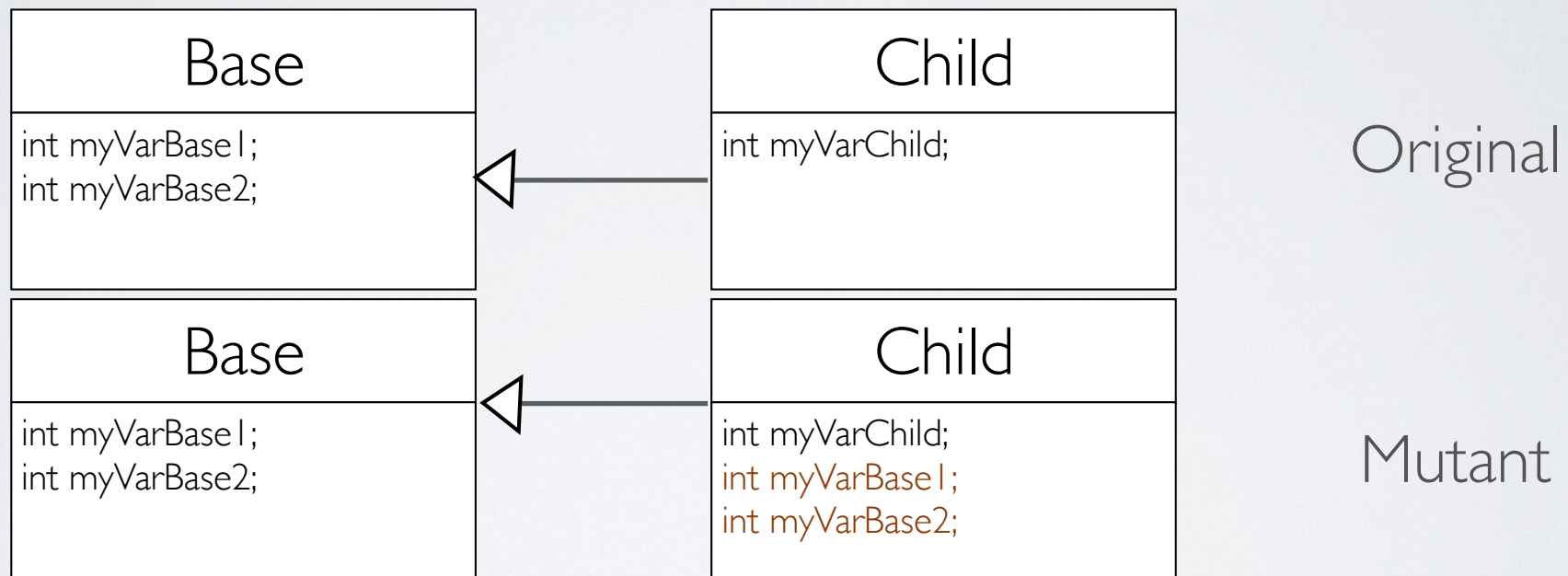


**Goal: Check whether tests can detect wrong references  
(common mistake made)**

# OO MUTATIONS

## INHERITANCE

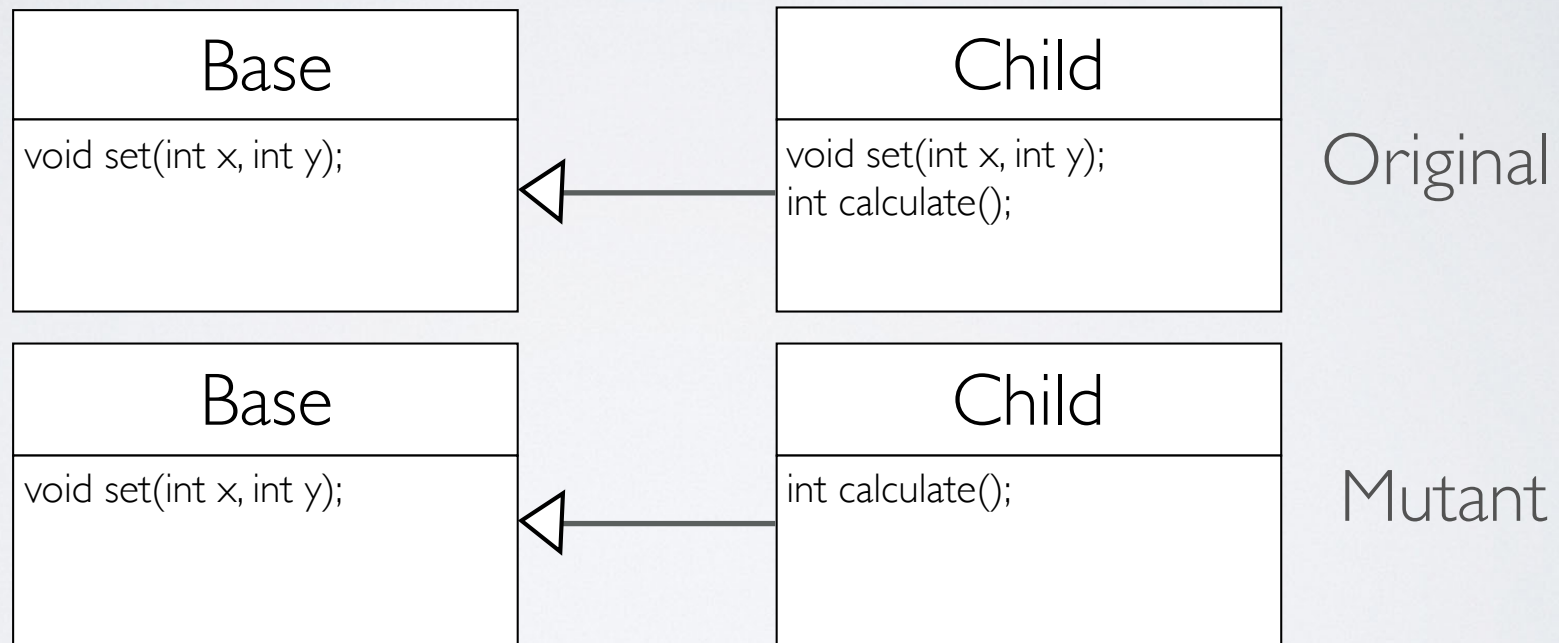
**Hiding variable insertion:** A declaration is added to hide the declaration of each variable declared in an ancestor.



**Goal:** Find test cases that show that reference to the overriding variable is incorrect

# OO MUTATIONS INHERITANCE

**Overriding method deletion:** Each entire declaration of an overriding method is deleted.



**Goal:** Allows to evaluate whether the invocation of the method is to the intended method

# OO MUTATIONS CONT.

- Overridden method moving
- Overridden method rename
- Super keyword deletion
- Parent constructor deletion
- Actual type change
- Declared type change
- Parameter type change
- Reference type change
- Overloading method deletion
- Argument order change
- Argument number change
- Keyword deletion
- Static modifier change
- Variable initialisation deletion
- Default constructor delete
- Overloading method change



# WEAK AND STRONG MUTATION

- Criteria of reachability, interaction, and propagation **RIP**
  - Reachability (R): **Location** of the fault must be **reachable**
  - Infection (I): Fault is executed and leads to an **incorrect state** in the program
  - Propagation (P): **Output** of the program is **incorrect**
- **Strong mutation**: An incorrect state should propagate to the output of the program (should fulfil R+I+P)
- **Weak mutation**: Test case reaches the mutant, but does not require propagation to the output (should fulfil R+I)

# EXAMPLE RIP

```
public int lastIndexOf(int[] x, int y) {
    for (int i = x.length-1; i > 0; i--){
        if( x[i]==y ) return i;
    }

    return -1;
}
```

State (program counter)	Line	Expected	Actual
S1	lastIndexOf()	x=[1]	x=[1]
S2	for (int i =...)	i=0	i=0
S3		Program counter = if (...)	Program counter = return -1

Error state in S3, but we have the expected output

Input: x=[1], y=2  
Expected output: -1

# BASED ON WEAK AND STRONG MUTATION WE DEFINE COVERAGE CRITERIA

- **Weak mutation coverage:** For each mutant there exist a test case to weakly kill the mutant.
- **Strong mutation coverage:** For each mutant there exist a test case to strongly kill the mutant.



# EQUIVALENT MUTANTS

- **Mutations are equivalent** if the behaviour stays the same after introducing a mutation to the code.

```
int max(int[] values) {  
    int r, i;  
    r = 0;  
    for(i = 1; i < values.length; i++) {  
        if (values[i] >= values[r])  
            r = i;  
    }  
    return values[r];  
}
```

The function will give the correct max value for > as well as >=

- **Notes on equivalent mutants**
  - hard and costly to detect (equivalence mutant problem) —> e.g. Adamopoulos et al. propose a solution
  - if we could find the equivalent mutants, we could reduce the overall testing effort
  - what is more expensive, finding the equivalent mutants and removing them, or including them in the test suite? -> did not find an answer to that question



# FIRST AND SECOND ORDER MUTATIONS

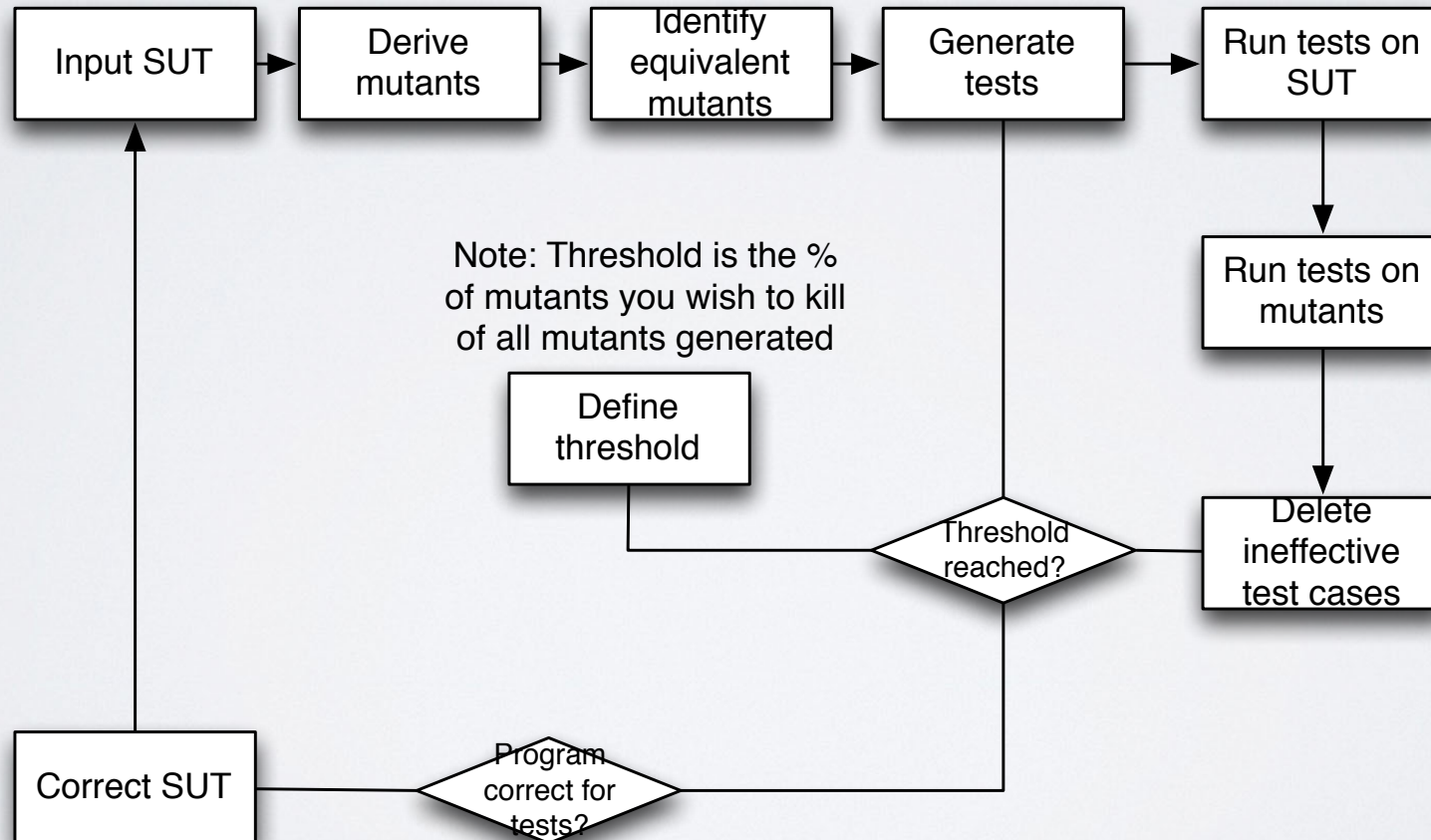
- Definitions:
  - **First order mutation:** insertion of a single mutant
  - **Higher order mutation:** Insertion of two more more mutants
- **Problem** of first order mutants: Trivial and easy to detect (does not consider fault combinations that make the testing interesting/subtle)
- First order mutants do not represent a realistic situation

# EVALUATING MUTANTS

- **Mutation score** =  $\frac{\text{\#dead mutants}}{\text{\# non-equivalent mutants}}$  (mutant coverage)

# MUTATION TESTING PROCESS

Note: If you have tests  
that are not able to kill mutants  
in the test run you would add new ones  
you believe will do that





Mutation testing is considered one of the strongest approaches to arrive at an effective test suite

–But: What is the challenge with it?



# EVALUATION OF HIGHER ORDER MUTATIONS

$a + b > c$

Original

$a * b > c$

$a / b > c$

$a \% b > c$

$a > c$

$b > c$

$\text{abs}(a) - b > c$

$a - \text{abs}(b) > c$

$a - b > \text{abs}(c)$

$\text{abs}(a - b) > c$

$0 - b > c$

$a - 0 > c$

$a - b \geq c$

$--a - b > c$

$a - b < c$

$a - b \leq c$

$a - b = c$

$a - b \neq c$

$b - b > c$

$a - a > c$

$c - b > c$

$a - c > c$

$a - b > a$

$a - b > b$

$a - b > c$

$++a - b > c$

$a - ++b > c$

$a - --b > c$

$a - b > --c$

$++(a - b) > c$

$--(a - b) > c$

$-a - b > c$

$a - -b > c$

$a - b > -c$

$-(a - b) > c$

$a - b > 0$

$-\text{abs}(a) - b > c$

$a - -\text{abs}(b) > c$

$a - b > -\text{abs}(c)$

$-\text{abs}(a - b) > c$

Mutants  
(selection)

# REASONS FOR LOW TIME EFFICIENCY

- High number of mutations for each mutation operator
- To be tested **dynamically**, mutants have to be **compiled** and tests have to be run **against each mutant**
- Consequence:
  - Need **strategies for selection**
  - Need for **high computing power** (one of the reasons because there was a large break in research on mutation, invented in 1976, continued in the 90s/2000s)



# STRATEGIES I

- Strategies
  - Coverage
  - Mutation sampling
  - Selective mutation
  - Parallelization
  - Weak coverage
  - Mutate bytecode





# STRATEGIES II

- Coverage
  - Determine statement coverage
  - Only execute those test cases that reach the statement including a mutation
- Strong vs. weak mutation
  - Compare internal states after mutation
  - Reported to save 50% execution time
- Mutant schema
  - Create a mutant “product line”, using a meta mutant
  - One compilation, mutant activation and deactivation (e.g. switch function)



# STRATEGIES I II

- Sampling
  - **Strategy 1**: Identify test cases that kill all mutants (very effective, but not efficient)
  - **Strategy 2**: Only focus on test cases focusing on those mutation operators that subsume all other mutation operators
- For **Strategy 2**, a subset of mutants will detect more than 99% of all mutants
  - ABS(), Arithmetic operator replacement, Conditional operator replacement, Relational operator replacement, Unary operator replacement

# TOOLS

- Evo-suite - <http://www.evosuite.org/>
  - Research shows that (based on a sample selection of programs) mutation testing is scalable for the application of novel techniques
  - Generated 1,380,302 mutants for 8,963 classes
- Jester - <http://jester.sourceforge.net/>

**Recommended reading:** Gordon Fraser and Andrea Arcuri. Achieving Scalable Mutation-based Generation of Whole Test Suites. Empirical Software Engineering



# OVERALL REFLECTIONS

- Research provides evidence that superior test suites can be generated (compared with other test design techniques)
  - structural code coverage
  - data-flow coverage
- Convincing work is done (e.g. through search-based algorithms) to make mutation testing beneficial and cost efficient
- Though: Need more industria application (often open source studied)