

Notes Made By

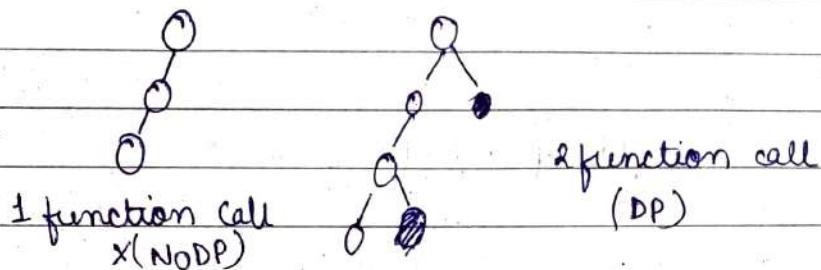
- RITI Kumari

Dynamic Programming (Aditya Verma)

$\text{DP} = \text{enhanced recursion}$

How to identify DP problem (2 cases)

- 1) where there is recursion, DP is used (for overlapping problem)
 - a) choice



- b) Optimal - Min, max, largest

How to write DP code?

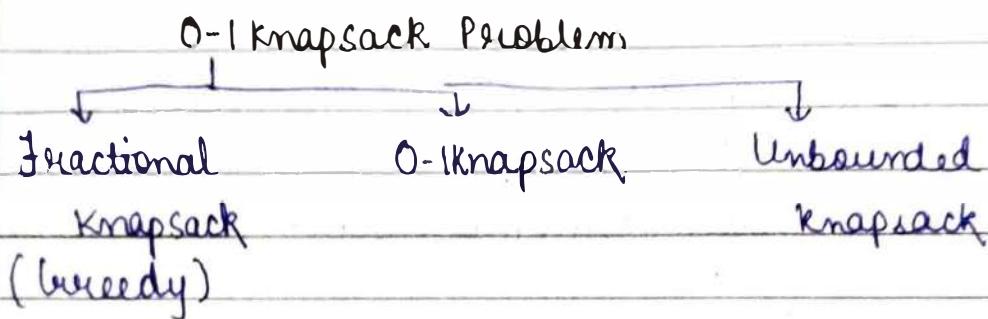
Recursive solution \rightarrow memoization \rightarrow Top down approach

Questions on DP.

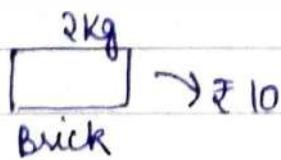
- 1) 0-1 knapsack (6)
- 2) Unbounded knapsack (5)
- 3) Fibonacci (7)
- 4) LCS (15) (longest common subsequence)
- 5) LIS (10) (longest increasing subsequence)
- 6) Kadane's Algorithm (6)
- 7) Matrix chain multiplication (7)
- 8) DP on tree (4)
- 9) DP on grid (14)
- 10) Others (5)

Types of knapsack

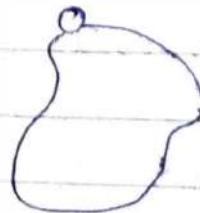
1. Subset sum
2. Equal sum partition
3. Count of subset sum
4. Minimum subset diff
5. Target sum
- 6.



0-1 Knapsack \rightarrow We are given some weight & some value. Then a max weight W : pick items so that the profit is maximum. And the weight has a given bound W .



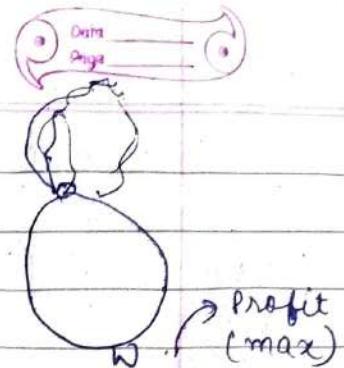
	I_1	I_2	I_3	I_4	
$wt[] =$	1	3	4	5	
$val[] =$	1	4	5	7	



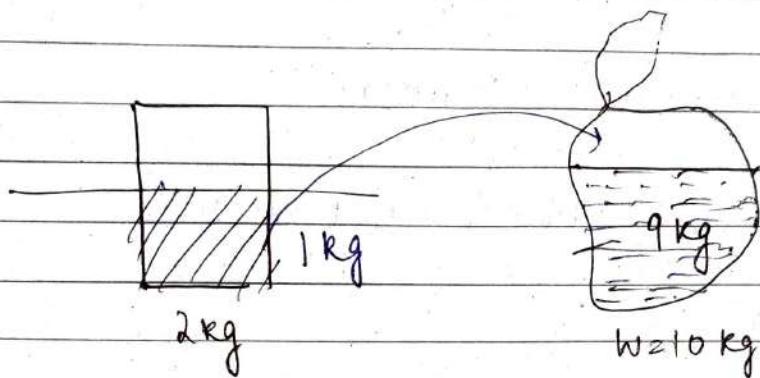
$W=7\text{kg}$

Max Profit = ?

$P_1 \quad P_2 \quad P_3 \quad P_4$
 $w_1 \quad w_2 \quad w_3 \quad w_4$



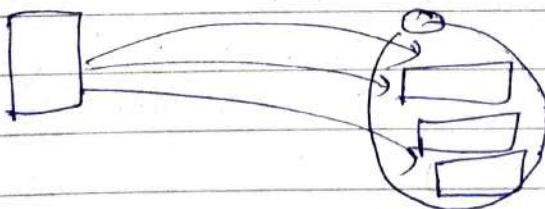
Fractional knapsack



Greedy approach

Unbounded knapsack

unlimited supply of every item

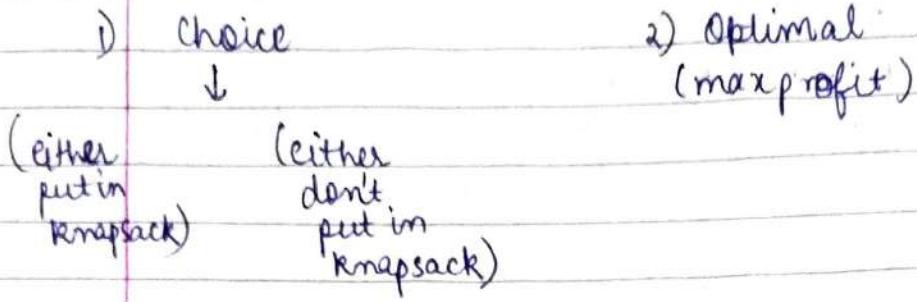


0-1 knapsack

i) How to identify

$wt[] : 1 \quad 3 \quad 4 \quad 5$ $W : 7 \text{ kg}$
 $val[] : 1 \quad 4 \quad 5 \quad 7$

$\% : \text{max profit}$



DP: Recursive \rightarrow Memoization \rightarrow Top down (DP)

DP \rightarrow recursion storage

0-1 knapsack Recursive

Identify

DP \rightarrow Recursive $\begin{cases} \xrightarrow{\text{DP (topdown)}} \\ \xrightarrow{\text{DP (memoization)}} \end{cases}$

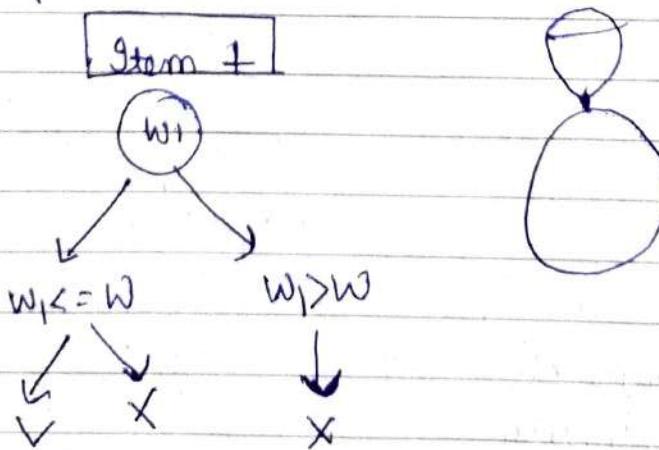
I/P $\text{wt}[] = [1|3|4|5]$

$\text{val}[] = [1|4|5|7]$

O/p \rightarrow Max Profit

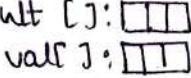
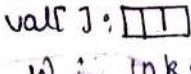
(Capacity of knapsack) $w = 7 \text{ kg}$

choice diagram



We have to return the max profit so return type would be int.

Base condⁿ → think of the smallest valid ip.

IP wt []:  val []: ] → n → 0.

w: 10 kg → 0 kg



max profit

```
int Knapsack( int wt[], int val[], int w, int n ) {  
    // base condition  
    if (n == 0 || w == 0)  
        return 0;
```

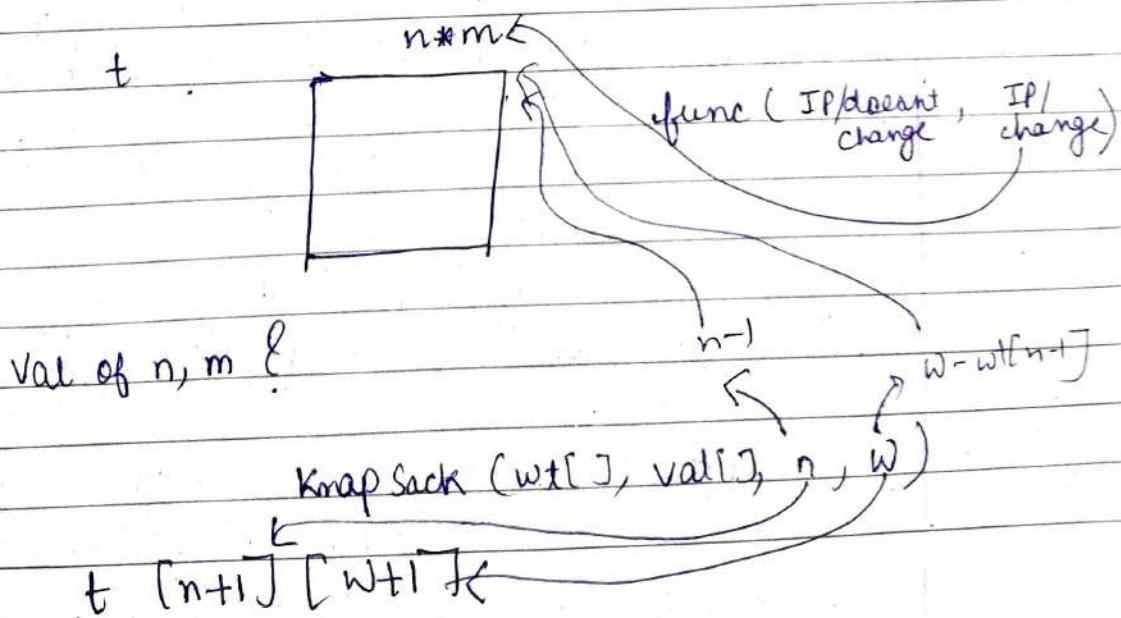
// choice diagram

if (wt[n-1] <= w) { → weight vs gain
 & including
 return max(val[n-1] + Knapsack(wt, val, w - wt[n-1]),

n-1)
 }
 else if (wt[n-1] > w) { weight less hai,
 but dont include
 return Knapsack(wt, val, w, n-1);
 }

Q1 Knapsack Memoization

Memoization = Recursive + 2 lines



\uparrow	-1	-1	-1	-1	-1	
	-1	-1	-1	-1	-1	
$n+1$	-1	-1	-2	-1	-1	
	-1	-1	-1	-1	-1	
\downarrow	-1	-1	-1	-1	-1	
	-1	-1	-1	-1	-1	
\leftarrow	$w+1$	\rightarrow				

if (-1) is not present then val exists so, return

initialise this matrix with -1.

```
int t[n+1][w+1]
memset(t, -1, sizeof(t))
```

Changes

Now declare the matrix globally.

```
int static t[102][002]; constraint  
n <= 100  
memset(t, -1, sizeof t) w <= 1000
```

```
int knapsack (int wt[], int val[], int w, int n)  
{
```

if ($n == 0$ || $w == 0$)

return 0;

if ($t[n][w] \neq -1$)

return t[n][w];

if ($\text{wt}[n-1] \leq w$)

return $t[n][w] = \max \{ val[n-1] + knapsack(wt, val, w-wt[n-1]), t[n][w] \}$

1 knapsack(wt, val, v, n-1))

else if ($wt[n-1] > w$)

return $t[n][w] = \text{Knapsack}(wt, val, w, n-1);$

}

The complexity of top down & memoization remains same but the problem with memoization is the stack gets full due to repeated funcn' calls.

(0-1 Top Down
Knapsack)

Real DP

Recursive \rightarrow Memoize \rightarrow Top down \rightarrow 6 Problems.

\downarrow

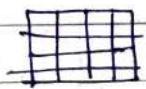
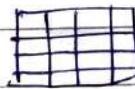
\downarrow

\downarrow

BC + recursive
calls

RC +
table

only
Table

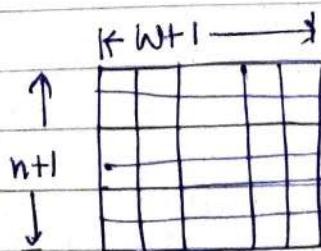


Recursive

Memoization

Initialising
matrix with -1

Top-down (totally omit the
recursive call &
use the table only)



We only make table for
the ip which is
changing.

2 steps to make table from top down

Step 1: Initialization

Step 2: Recursive code changes Iterative code

Step 1: Initialize

$$w = 7$$

$$n = 4$$

$$wt[] = [1 \ 3 \ 4 \ 5]$$

$$val[] = [1 \ 4 \ 5 \ 7]$$

w → (j)

		0	1	2	3	4	5	6	7
wt		0							
val		1							
		2							
		3							
		4							
n									
(i)									

t[n][w] → it will give ans

$$wt[] = [1 \ 3 \ 4 \ 5] \quad wt[] = [1 \ 3]$$

$$val[] = [1 \ 4 \ 5 \ 7] \quad val[] = [1 \ 4]$$

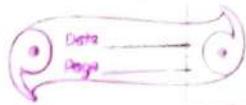
w=3

w = 3

$$wt[] = [1 \ 3 \ 4]$$

$$val[] = [1 \ 4 \ 5]$$

w=6



RC + table \longrightarrow table

Base condⁿ \longrightarrow Initialization

Base condⁿ RC
 $\text{if } (n == 0 \text{ || } w == 0)$
 $\quad \downarrow \quad \text{between 0;}$

table

$\text{for (int } i=0; i < n+1; i++) \{$

$\quad \text{for (int } j=0; j < w+1; j++)$

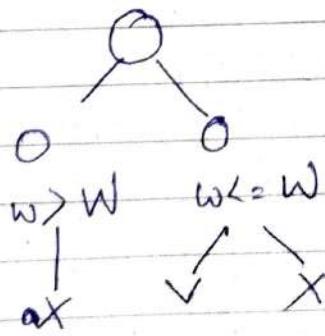
$\quad \text{if } (i==0 \text{ || } j==0)$

$\quad \quad \quad t[i][j] = 0;$

$\}$

w\o	0	1	2	...
0	0	0	0	
1	0			
2	0			
3	0			
4	0			

Choice
diagram



$n, w \rightarrow i, j$

7+

$dp[i][j]$

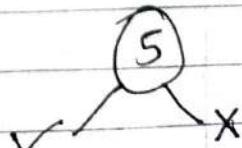
RC

$$wt = 13 + 5$$

$w = 7$

$$val = 14 + 5$$

$\max(\text{val}[n-1] + \text{dp}[i-1][j - wt[i-1]],$
 $\text{dp}[i-1][j])$



Date _____
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<p>Recursive</p> <pre> if (wt[n-1] <= w) return max(val[n-1] + Knapsack(wt, val, w-wt[n-1], n), knapsack(wt, val, w, n-1)) else if (wt[n-1] > w) return knapsack(wt, val, w, n-1) </pre>	<p>Top down</p> <pre> if (wt[n-1] <= w) t[n][w] = max(val[n-1] + t[w-wt[n-1]][n-1], t[n-1][w]) else t[n][w] = t[n-1][w] </pre>
---	---

<p>pseudo code</p> <pre> Top - down approach int t[n+1][w+1]; for (int i=1; i < n+1; i++) for (int j=1; j < w+1; j++) if (wt[i-1] <= j) t[i][j] = max(val[i-1] + t[i-1][j-wt[i-1]], t[i-1][j]) else t[i][j] = t[i-1][j] return t[n][w]; </pre>	<p style="text-align: center;">0 1 2 3</p> <p style="text-align: right;">$w = 7$</p> <p style="text-align: center;">wt = 1 3 4 5 val = 1 4 5 7 n = 4</p> <p style="text-align: right;">$w = 0 1 2 3 4 5 6 7$</p> <p style="text-align: right;">$0 1 2 3 4 5 6 7$</p>
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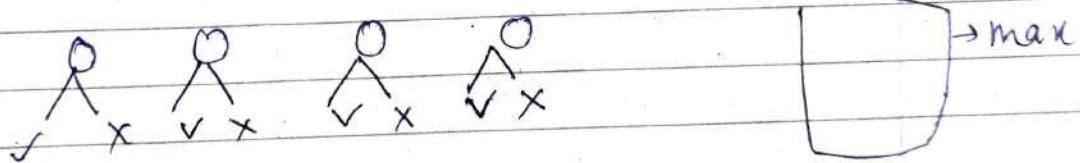


Identification of Knapsack problem

- 1) Subset sum problem
- 2) Equal sum partition
- 3) Count of subset sum.
- 4) Minimum subset sum diff
- 5) Target sum
- 6) No of subset with a given diff.

Ip: Item array : | | | |

w: Capacity



1. Subset Sum problem

arr [] : 2 3 7 8 10

sum =

- a) Problem statement
- b) similarity with knapsack
- c) Code variation

D) Problem statement : find if there is a subset present in an array with given sum.

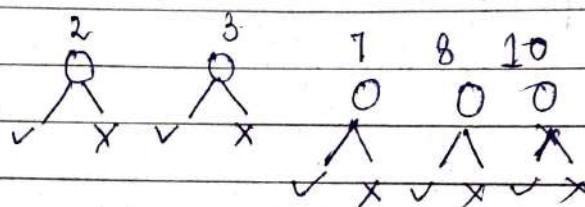
arr []: 2 3 7 8 10

Sum : 11

2) Similarity

item array [] : → 2 3 7 8 10

weight : capacity → 11



3) Code variation

$t[n+1][w+1]$

$t[5+1][11+1]$

sum

$t[6][12]$

Initialization:

	0	1	2	3	4	5	6	7	8	9	10	11
0	T	F	F	F	F	F	F	F	F	F	F	F
1	T											
2	T											
3	T											
4	T											
5	T											

\rightarrow True/false

arr []:
sum : 0
arr []: 1
sum = 0

arr []: 2
sum = 0

arr []:
sum = 2



when array is empty sum can't be anything

$arr[]$: no elements

sum: 1 → not possible

$t[n+1][sum+1]$

Initialisation: $\text{for } (\text{int } i = 1 \dots n+1)$
 $\quad \text{for } (\text{int } j = 1 \dots m+1)$

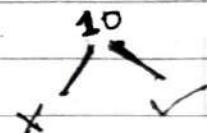
$\text{if } (i == 0)$

$t[i][j] = \text{false}$.

2378!0

$\text{if } (j == 0)$

$t[i][j] = \text{True}$



knapSack $\rightarrow arr$

$\text{if } (wt[i-1] \leq j)$

$t[i][j] = \max(\text{val}[i-1] + t[i-1])$

↓
there is
no max
in true
or false.

else

$t[i][j] = t[i-1][j]$

$t[i][j][0]$

$= t[i-1][j][0]$
 $= t[i-2][j][0]$
 $= t[i-3][j][0]$

$= t[i][j][0]$

Subset sum

$\text{if } (arr[i-1] \leq j)$

$t[i][j] = t[i-1][j-wt[i-1]]$

||

$t[i-1][j]$

else

$t[i][j] = t[i-1][j]$

return $t[n][sum]$;

2, 3, 8
 $\{3, 8\} \rightarrow \text{true } \checkmark$
 $\{3\} \rightarrow \text{false } \times$

2. Equal Sum Partition Problem

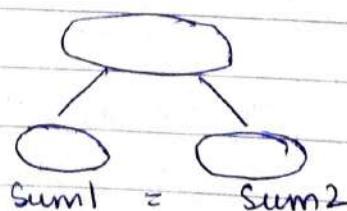
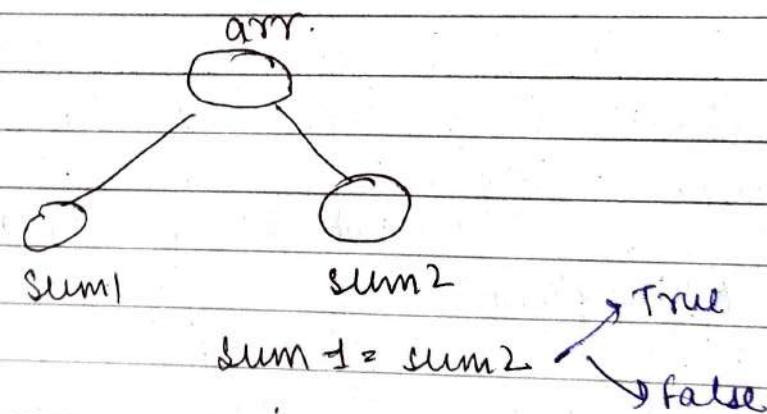
- 1) Problem statement
- 2) Subset sum similarity
- 3) Odd/Even significance
- 4) Code variation

1) Problem statement

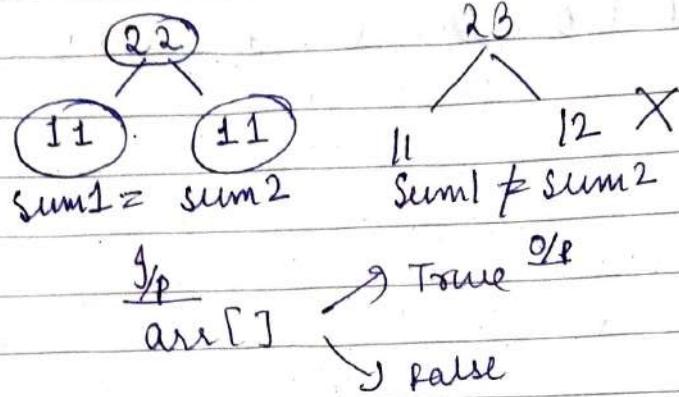
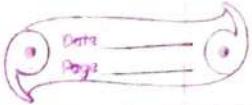
$$\text{arr}[] = \{1, 5, 11, 5\}$$

$\therefore \text{O/P} : \text{T/F}$

Is it possible to divide the array such that both the subset gives an equal sum

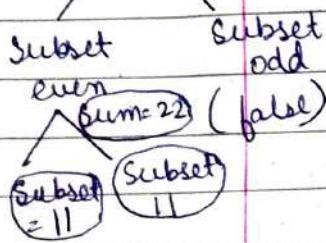


$\text{sum1} = \text{sum2} \rightarrow$ It is equal when the no is even. Sum of all the array elements should be even to change it into equal parts.



`for (int i = 0; i < size; i++) {`

$\sum = \sum + arr[i];$



if ($\sum \% 2 \neq 0$) (sum is odd)
return false

else if ($\sum \% 2 == 0$)

*→ we need to find one subset with sum 11 the
next subset would automatically be 11.

return subsetsum(arr, sum/2);

3. Code:

ip → arr[], n.

int sum = 0;

for (int i = 0; i < n; i++)
 $\sum += arr[i];$

if ($\sum \% 2 \neq 0$)

return false

else

return subsetsum(arr, sum/2);

3. Count of Subsets sum with a given sum.

S/p:

$arr[] = 2 \ 3 \ 5 \ 6 \ 8 \ 10$

sum = 10.

Flow

- 1) Problem Statement
- 2) Similarity to subset sum
- 3) Code variation
- 4) Return Type.

✓ ↘

Initialisation Code

1) Problem Statement

$arr[] = 2 \ 3 \ 5 \ 6 \ 8 \ 10$

Sum : 10

O/p = 3.

$\{2, 8\} \rightarrow$ Yes/True (in subset sum)

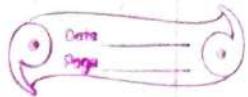
$$\left. \begin{array}{l} \{2, 8\} = 10 \\ \{5, 2, 3\} = 10 \\ \{10\} = 10 \end{array} \right\} \text{count} = 3$$

2) Similarity

2, 8

Yes No

return count



3. Code variation

Subset sum

Count
int

due to ↙ bool
T/F

False → 0 (no of subset
0)

True → null
subset (no of subset
1)

0	0	0	0	0	0	-
1						
1						

if ($\text{arr}[i-1] \leq j$) +
 $\text{dp}[i][j] = \text{dp}[i-1][j] + \text{dp}[i-1][j-\text{arr}[i-1]]$

else

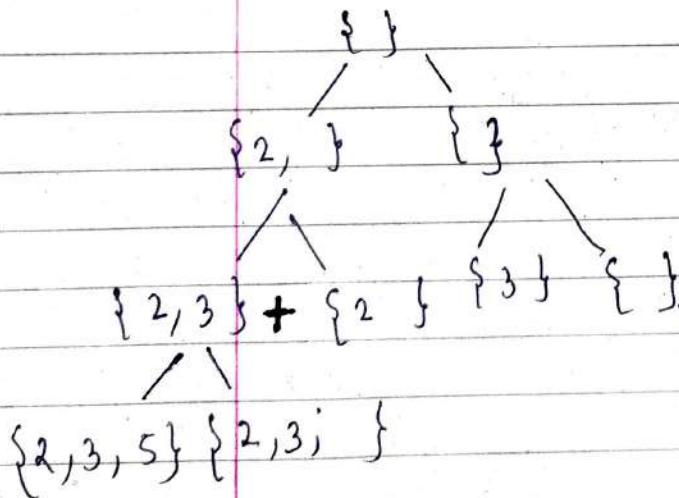
$\text{dp}[i][j] = \text{dp}[i-1][j]$

we will add all the subset
so arr would be changed
to + - True & false case
we can use ~~arr~~ arr.

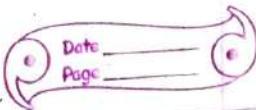
if ($\text{arr}[i-1] \leq j$)
 $\text{dp}[i][j] = \text{dp}[i-1][j] + \text{dp}[i-1][j-\text{arr}[i-1]]$

else

$\text{dp}[i][j] = \text{dp}[i-1][j]$



if ($\text{arr}[i-1] \leq j$)
 $\text{dp}[i][j] = \text{dp}[i-1][j - \text{arr}[i-1]] + \text{dp}[i-1][j]$
 else
 $\text{dp}[i][j] = \text{dp}[i-1][j]$



$\text{arr}[] = [3 5 6 8 10]$

Sum = 10

O/p = 3

1, 3, 2, 5 5

O/p $\rightarrow \{3, 2\}$
 $\{5\}$

$\text{arr}[0] \leq 2^1$

~~10000~~

$$\text{dp}[N+1][\text{sum}+1] = \text{dp}[6+1][10+1]$$

$$= \text{dp}[7][11]$$

$i \neq j \leq 2^{20}$

Sum \rightarrow

$+1 \quad \{ \quad \} \quad -1$
 $+3 \quad \{ \quad \} \quad \{ \quad \}$
 $\{1, 3\} \quad \{1\} \quad \{3\} \quad \{ \}$

	0	1	2	3	4	5	6	7	8	9	10
$i = 1$	0	1	0	0	0	0	0	0	0	0	0
$j = 2$	1	1	0	1							
$\text{arr}[0] \leq 2^2$	2	1									
$i = 2$	3	1									
$\text{dp}[0][0]$	4	1									
$\text{dp}[0][2]$	5	1									
	6	1									

Minimum $\sum_{i=1}^n |S_i - S_j|$

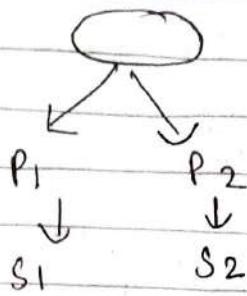
- 1) Problem statement
- 2) Similarity
- 3) Solve using its previous concept

+

- 4) Problem statement

$\text{arr}[]: [16 | 11 | 5]$

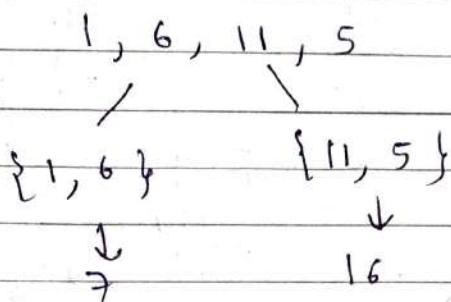
O/P: 1



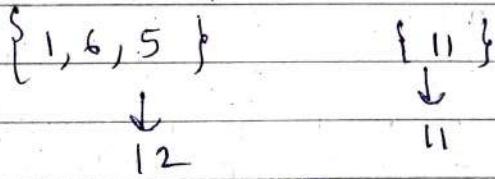
Equal sum: $S_1 - S_2 = 0$

Num^m subset: $S_1 - S_2 = \min$

$\text{abs}(S_1 - S_2) = \min$ (should be min)



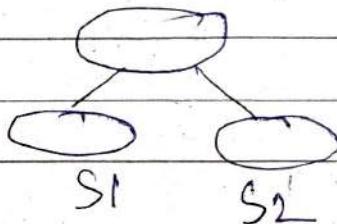
$16 - 7 = 9 \rightarrow$ minimize of
find could
it be done
in a better
way.



+ $12 - 11 = 1 \rightarrow$ Can't be minimized further
→ O/P

2) Similarity

It's similar to equal sum partition.



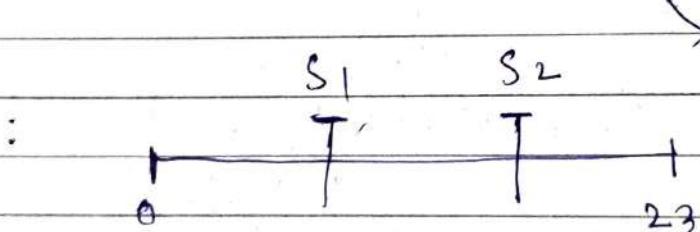
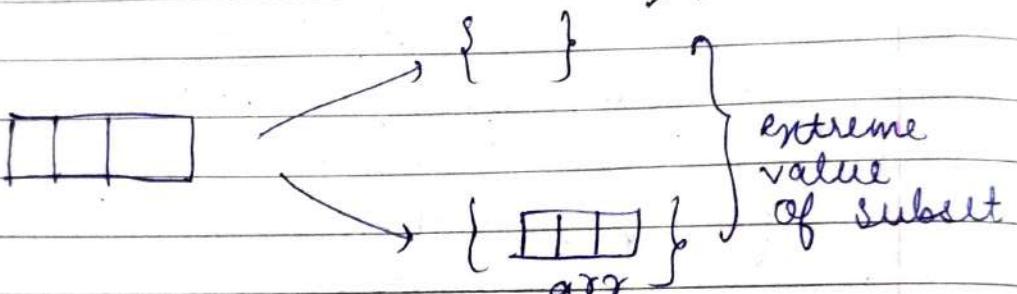
We need to find S_1 & S_2 .

$\text{arr}[] [1 | 6 | 11 | 5 | \text{?}]$

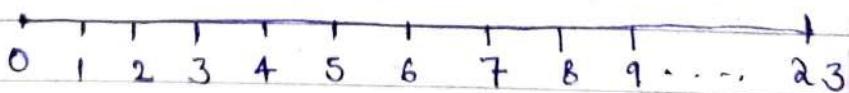


We can find the range of S_1 & S_2 .

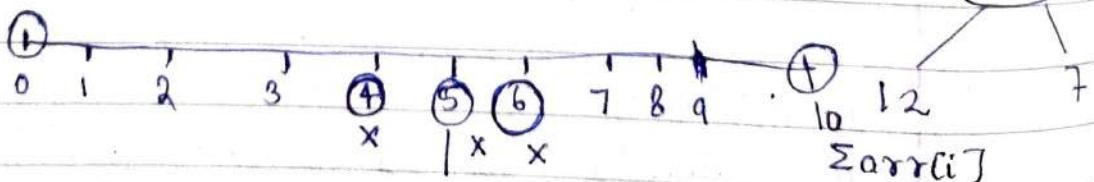
$$S_1 = 0 \quad (0+0+0+0)$$



$\text{arr}[] : \{1, 2, 7\}$



Sum line

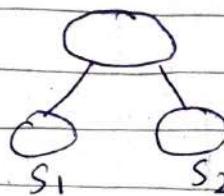
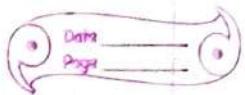


S_1 / S_2

can't be

5

$$S_1 / S_2 = \{0, 1, 2, 3, 7, 8, 9\}$$

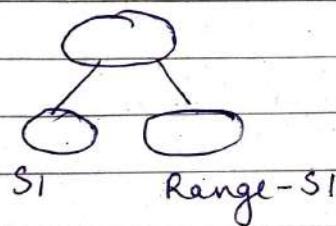


$$S_1 + S_2 = \sum arr[i]$$

$$S_1 / S_2 = \{ 0, 1, 2, 3, | 7, 8, 9, 10 \}$$

S_1
 (S_1) Range - S_1
 (S_2)

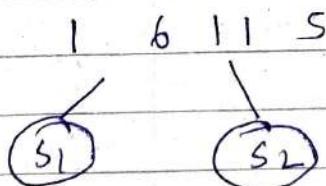
$$S_1 - S_2 = \text{minimize}$$



$$\text{abs}(.S_2 - S_1) \text{ or } \text{abs}(S_1 - S_2)$$

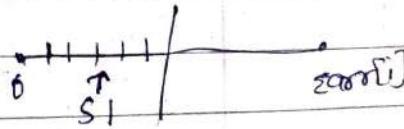
$$\begin{aligned} \text{abs}(S_2 - S_1) &= \text{minimize} \\ \text{abs}(\text{Range} - S_1 - S_1) &= \text{minimize. } \text{abs}(\text{Range} - 2S_1) \end{aligned} \quad \xrightarrow{\text{minimize}}$$

Sum up

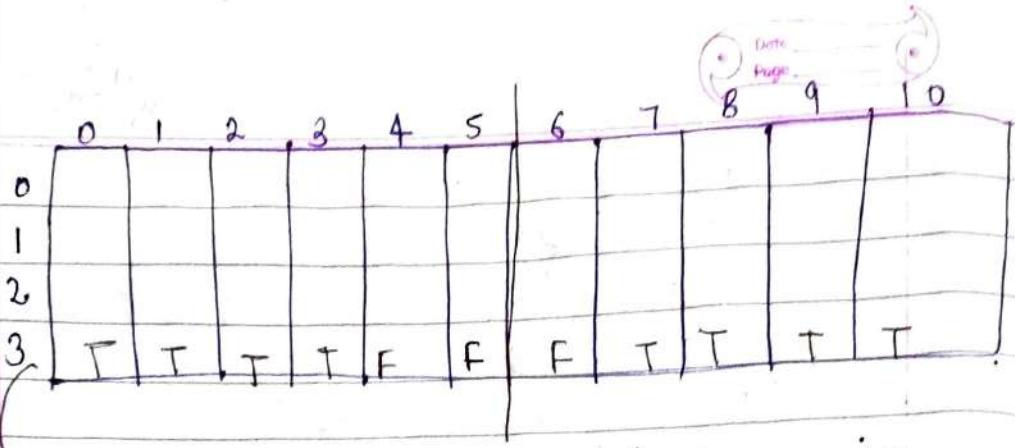


$$\left. \begin{array}{l} S_1 - S_2 \\ S_2 - S_1 \end{array} \right\} \text{minimize}$$

$$\begin{aligned} \downarrow \\ ((\text{Range} - S_1) - S_1) \\ (\text{Range} - 2S_1) \end{aligned}$$



Range/2



→ we will push the last row in vector till half way.

0	1	2	3
---	---	---	---

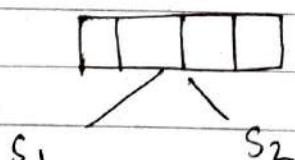
```
int min = INT_MAX;
```

```
for (int i = 0; i < v.size(); i++) {
```

```
    mn = min (mn, Range - 2v[i])
```

```
return mn;
```

Concept

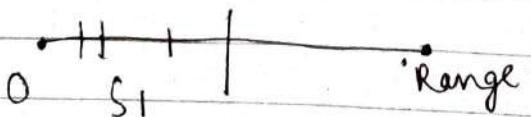


$$(S_2 - S_1) = \min$$

↓
smaller

$$(S_1) \quad (\text{Range} - S_1)$$

$$\text{Range} - 2S_1 \rightarrow \min^m$$



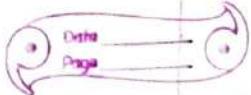
Code

```
subsetSum( int arr[], int range ) {
```

```
{
```

```
last row filled in vec till
```

0	1	0	0		halfway
---	---	---	---	--	---------



```
int mindiff (int arr[], int n) {
```

```
    int sum = 0;
```

```
    for (int i = 0; i < n; i++) {
```

```
        sum += arr[i];
```

```
}
```

```
    bool dp[n+1][sum+1];
```

```
    for (int i = 0; i < n+1; i++) {
```

```
        for (int j = 0; j < sum+1; j++) {
```

```
            if (i == 0) dp[i][j] = false;
```

~~```
 if (j == 0) dp[i][j] = true;
```~~

```
}
```

```
 }
```

```
 for (int i = 1; i < n+1; i++) {
```

```
 for (int j = 1; j < sum+1; j++) {
```

```
 if (arr[i-1] <= j)
```

```
 dp[i][j] = dp[i-1][j - arr[i-1]] || dp[i-1][j];
```

```
 else
```

```
 dp[i][j] = dp[i-1][j];
```

```
}
```

```
 }
```

```
 int diff = INT_MAX;
```

```
 for (int j = sum/2; j >= 0; j--) {
```

```
 if (dp[n][j] == true) {
```

```
 diff = min(sum - 2*j, diff)
```

~~```
        }
```~~

```
}
```

```
return diff;
```

```
}
```

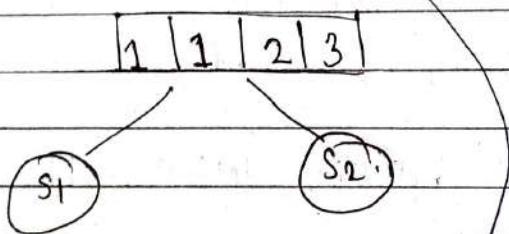
Count the number of subset with
a given diff

- 1) Problem statement
- 2) will try to reduce the actual statement
- 3) solve it using already solved problem.

D) Problem statement

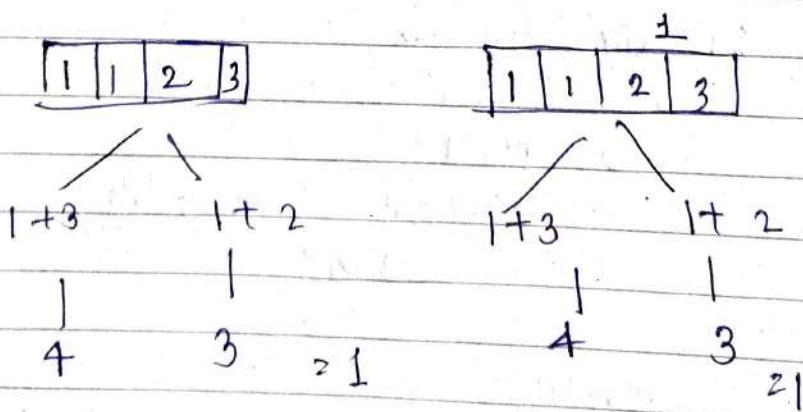
arr[] : [1 | 1 | 2 | 3]

Diff : 1

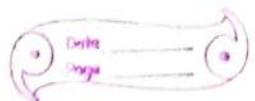


$$S_1 - S_2 = \text{diff}$$

return the count of
subset with diff"



Final result:
 $4 - 3 = 1$



| | | | |
|---|---|---|---|
| 1 | 1 | 2 | 3 |
|---|---|---|---|

s_1 s_2

(diff) $\rightarrow s = \text{sum of arr}$

$$\text{sum}(s_1) - \text{sum}(s_2) = \text{diff}$$

$$\text{sum}(s_1) + \text{sum}(s_2) = s.$$

$$2s_1 = \text{diff} + \text{sum(arr)}$$

$$s_1 = \frac{\text{diff} + \text{sum(arr)}}{2}$$

$$= \frac{1+7}{2} = \frac{8}{2} = 4$$

$$s_1 = 4$$

$$s_2 = s_1 - \text{diff}$$

$$= 4 - 1$$

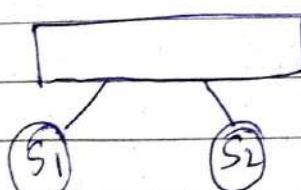
$$\text{Count} = ?$$

$$= 3.$$

② Find count when sum of $s_1 = 4$

no of count of \rightarrow count of
Subset with subset
given diff sum.

Sum up.



$$2s_1 = \text{diff} + \text{sum(arr)}$$

$$s_1 = \frac{\text{diff} + \text{sum(arr)}}{2}$$

$$s_1 - s_2 = \text{diff}$$

$$s_1 + s_2 = \text{sum(arr)}$$

$\text{int sum} = \frac{\text{diff} + \text{sum}(arr)}{2}$

$\text{return countofsubsetsum}(arr, \text{sum});$

~~int countofsubsetwithdiff (int arr[], int n) {
int sum = 0; for (int i=0; i<n; i++) sum += arr[i];
int diff = sum - arr[0];
for (int i=1; i<n; i++)
arr[i] = arr[i] - arr[0];
for (int i=1; i<n; i++)
arr[i] = arr[i] / 2;~~

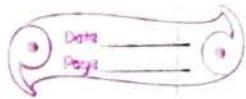
$\text{int arrsum} = 0$
 $\text{for (int } i=0; i < n; i++) \{$
 $\quad \text{arrsum} += \text{arr}[i];$
 $\}$

$\text{int sum} = 0;$

$\text{Sum} = \frac{\text{diff} + \text{arrsum}}{2}.$

$\text{return countofsubsetsum}(arr, \text{sum}, n);$
 $\}$

~~int countofsubsetsum (int arr[], int sum, int n) {
int dp[n+1][sum+1];
for (int i=0; i<n+1; i++)
for (int j=0; j<sum+1; j++) {
if (i == 0) dp[i][j] = 0;
if (j == 0) dp[i][j] = 1;~~



```

for (int i = 0; i < n+1; i++) {
    for (int j = 1; j < sum+1; j++) {
        if (arr[i-1] <= j) {
            dp[i][j] = dp[i-1][j - arr[i-1]] + dp[i-1][j];
        } else {
            dp[i][j] = dp[i-1][j];
        }
    }
    return dp[n][sum];
}

```

Target Sum.

arr :

| | | | |
|---|---|---|---|
| 1 | 1 | 2 | 3 |
|---|---|---|---|

 ↓
 sum : 1 +/ - +c : [0, 0, 0, 0, 0, 0, 1]
 target = 1

Op : 3

Target value = 2

$$+1 -1 -2 +3 = 1$$

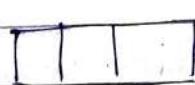
$$-1 +1 -2 +3 = 1 \quad \{0, 2\}$$

$$+1 +1 +2 -3 = 1$$

(+0, 2) (-0, 2)

arr [] → [+ - + +]

{0, 0, 2} count = 2



{+0, +0, 2} {+0, -0, 2} {-0, 0, 2}

S1 S2

$$S1 - S2 = \text{diff}$$

{-0, 0, 2}

= 4 (count)

So answer is increased
by a power of 2
no of zeros
no of zeros

$$\begin{array}{ccccccc}
 & + & - & - & + \\
 1 & | & 1 & 2 & 3 \\
 & / & \searrow & & \\
 +1 & +3 & & -1 & -2
 \end{array}$$

$$(1+3) - (1+2)$$

$S_1 - S_2 \rightarrow$ count of subset with given diff.

Count of Subsets with Difference K

```
int findTargetSumWays(vector<int>& nums, int s) {
    int cnt = 0, sum = 0;
```

```
int n = nums.size();
```

```
for (int i = 0; i < nums.size(); i++) {
```

```
    sum = sum + nums[i];
```

```
    if (nums[i] == 0)
```

```
        cnt = cnt + 1;
```

Cnt the
no. of zeroes
in subset

```
S = abs(s);
```

```
if (S > sum) || (S + sum) % 2 != 0)
```

```
return 0;
```

```
int s = (S + sum) / 2;
```

```
int dp[n+1][s+1];
```

```
for (int i = 0; i < n+1; i++)
```

```
    for (int j = 0; j < s+1; j++)
```

```
        if (i == 0) dp[i][j] = 0;
```

```
        if (j == 0) dp[i][j] = 1;
```

```
for (int i=1; i<n+1; i++) {
```

```
    for (int j=1; j<s+1; j++) {
```

if (num[i-1] == 0)

$dp[i][j] = dp[i-1][j];$

else if (num[i-1] > j)

$dp[i][j] = dp[i-1][j];$

else

$dp[i][j] = dp[i-1][j - num[i-1]] +$
 $dp[i-1][j]$

.

).

return pow(2, n) * dp[n][s];

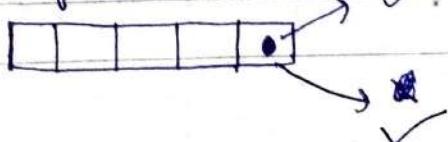
}.

13. Unbounded Knapsack :

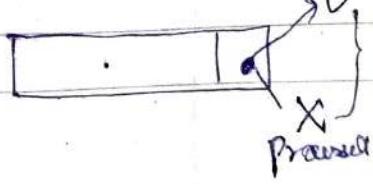
Related Problems

- 1) Road cutting
 - 2) Coin change I (Max no of ways)
 - 3) Coin change II (Min no of ways)
 - 4) Max m: Ribbon cut
- (Variations
of
unbounded
knapsack)

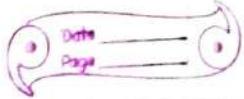
Unbounded
(multiple occurrence
of same item)



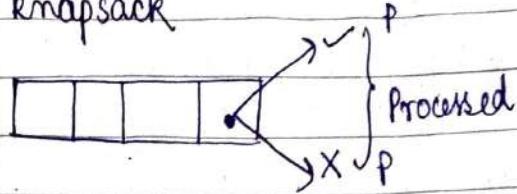
Knapsack
(only one occurrence)



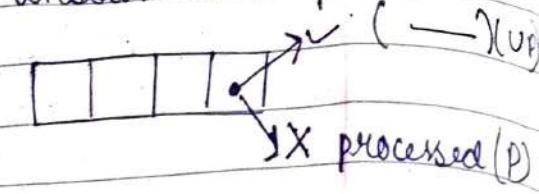
Present



knapsack



Unbounded knapsack



Multiple occurrences

✓ (can visit many times)

Comparison betw.

Knapsack

| | 0 | 1 | 2 | 3 | 4 |
|---|---|---|---|---|---|
| 0 | 0 | 0 | 0 | 0 | 0 |
| 1 | 0 | | | | |
| 2 | 0 | | | | |

Unbounded

| | 0 | 1 | 2 | 3 | 4 |
|---|---|---|---|---|---|
| 0 | 0 | 0 | 0 | 0 | 0 |
| 1 | 0 | | | | |
| 2 | 0 | | | | |

$t[n+1][w+1]$

if ($wt[i-1] \leq j$) {

$$t[i][j] = \max(t[i-1][j], t[i-1][j-wt[i-1]] + t[i-1][j]);$$

else {

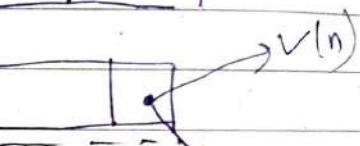
$$t[i][j] = t[i-1][j];$$

if ($wt[i-1] \leq j$) {

$$t[i][j] = \max(t[i-1][j], t[i-1][j-wt[i-1]] + t[i-1][j]);$$

else {

$$t[i][j] = t[i-1][j];$$



14. Road Cutting Problem

length [] : 1 2 3 4 5 6 7 8

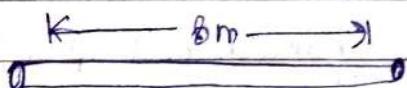
peice [] :

| | | | | | | | |
|---|---|---|---|----|----|----|----|
| 1 | 5 | 8 | 9 | 10 | 17 | 17 | 20 |
|---|---|---|---|----|----|----|----|

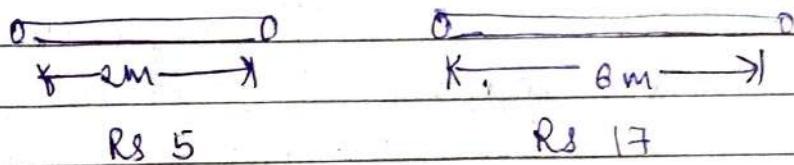
N : 8

N : B

1) Problem statement



length of rod is given. we need to cut it in such a way the profit is max^m.



$$\begin{aligned}\text{Total} &= 5 + 17 \\ &= \text{Rs } 22\end{aligned}$$

Knap Sack

wt
val
W

length = 1 to N if length array is not given
 Peice =

| | | | |
|--|--|--|--|
| | | | |
|--|--|--|--|

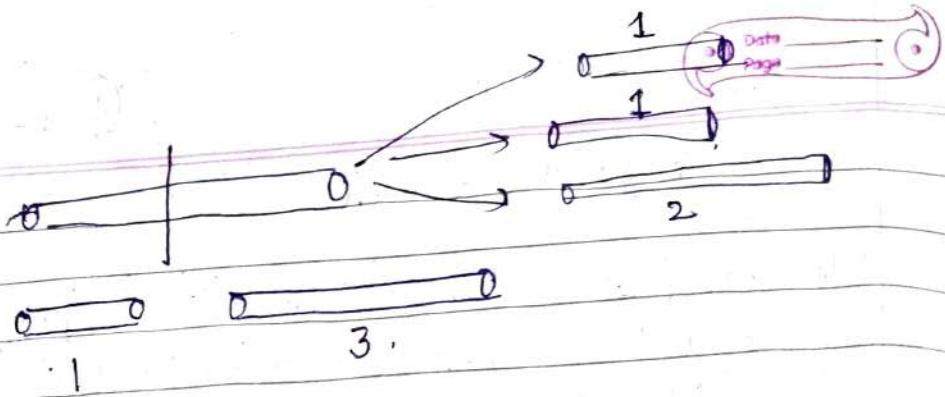
 make it & fill it will
 N = 8 1 to N.

knapsack
problem

vals]
W []
W

Unbounded knapsack

price []
length []
N



Code variation

$t[N+1][N+1]$

if ($\text{length}[i-1] \leq j$)

$dp[i][j] = \max(\text{price}[i-1] + dp[i][j - \text{length}[i-1]],$
 $dp[i-1][j])$

else

$dp[i][j] = dp[i-1][j];$

Sometimes our size changes therefore we need to find the size of array.

$dp[\text{size}+1][N+1].$

| | | length → | | | | | |
|-----------|---|----------|---|---|---|---|---|
| | | 0 | 0 | 0 | 0 | 0 | 0 |
| ↓
size | 0 | | | | | | |
| | 0 | | | | | | |

Coinchange Problem

Max^m no of
ways

Min^m no
of coins

Max^m no of ways.

Problem

statement

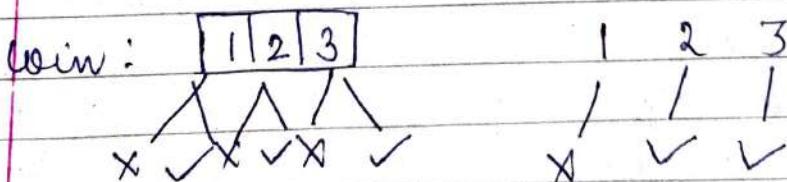
coin [] : [1 | 2 | 3] → (unlimited)
sum: 5.

Use the coin array unlimited time & get the sum 5. Now find the max^m no of ways in which we can get 5.

| | OP |
|--------|------------------------|
| 5 ways | { 2 + 3 5. |
| | } 1 + 2 + 2 5. |
| | 1 + 1 + 3 5. |
| | 1 + 1 + 1 + 1 + 1 5. |
| | 1 + 1 + 1 + 2 5 |

Q Why knapsack

→ Every coin has a choice to include or not



wt [] -

item

val [] x (when one array is given)

Matching

$w[] \rightarrow \text{coin}[]$
 $w \rightarrow \text{sum}$

| | | |
|---|---|---|
| 1 | 2 | 3 |
|---|---|---|

sum5 : $\textcircled{1} + \textcircled{1} + 3$

one item used
many times

If taking one item many times does the sum allows then it is unbounded knapsack.

Subset sum

| | | | |
|--------|-------------|---|---|
| 1 | 2 | 3 | 5 |
| Sum: 8 | → T
↓ F. | | |

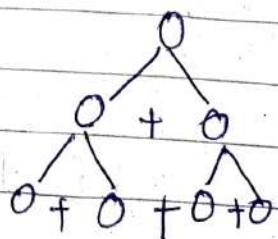
Subset sum

```

if curr[i-1] <= j {
    t[i][j] = t[i-1][j] || t[i-1][j - arr[i-1]]
} else {
    t[i][j] = t[i-1][j];   for count
}                                         we remove
                                         || by +.
                                         ↓
                                         for count
                                         we remove
                                         || by +.

```

count / no of ways



We need to find maxim. no of ways.

Matching \rightarrow Knapsack $\xrightarrow{0-1}$
 $\xrightarrow{\text{Unbounded}}$

wt \rightarrow coin

W \rightarrow sum.

if ($\text{coin}[i-1] \leq j$)

$$t[i][j] = t[i-1][j] + t[i][j - \text{coin}[i-1]]$$

else

$$t[i][j] = t[i-1][j]$$

| | 0 | 1 | 2 | 3 | ... | sum \rightarrow |
|-------------------|---|---|---|---|-----|-------------------|
| 0 | 1 | 0 | 0 | 0 | 0 | 0 |
| 1 | 1 | | | | | |
| 2 | 1 | | | | | |
| size \downarrow | 3 | 1 | | | | |

Coin change-II (Min^m no of coins)

$$\text{coin}[] = [1 \ 2 \ 3]$$

$$\text{sum} = 5$$

$$\begin{aligned}
 2+3 &\rightarrow 5 & \rightarrow 2 \text{ coins} \\
 1+2+2 &\rightarrow 5 & \rightarrow 3 \text{ coins} \\
 1+1+1+2 &\rightarrow 5 & \rightarrow 4 \text{ coins} \\
 1+1+3 &\rightarrow 5 & \rightarrow 3 \text{ coins} \\
 1+1+1+1+1 &\rightarrow 5 & \rightarrow 5 \text{ coins}
 \end{aligned}
 \left. \begin{array}{l} \text{find min}^m \\ \text{no. of coins} \\ \text{to make} \\ \text{sum as 5.} \end{array} \right\}$$

$\text{coin}[] = [1 \ 2 \ 3]$

$$\text{sum} = 5$$

$$\text{O/p : } 2 \quad (2+3=5)$$

Initialisation: $t[n+1][w+1]$

\downarrow
 $t[n+1][\text{sum}+1]$

$\text{wt} \rightarrow \text{coin}[j] = n$

$\text{val} \rightarrow x$

$w \rightarrow \text{sum}$

$$n^2 3 \\ \text{sum} = 5$$

\downarrow
size
(n)

$$\text{sum} = 3 \\ \text{coin}[] = [1 \ 2 \ 3]$$

Initialisation



+ Twist

$\text{coin}[]: \text{empty}$

$\text{sum} : 1$

$\rightarrow \text{INT_MAX}$ (∞ coins)

$\text{coin}[]: \text{empty}$

$\text{sum} : 0$

$\rightarrow \text{INT_MAX - 1}$



$\text{coin}[]: 1$

$\text{sum} : 0$

$\} 0 \text{ coins}$

$\text{coin}[]: 1 \ 2$

$\text{sum} : 0$

$\} 0 \text{ coins}$

for guess

$$\text{sum} = 5$$

$$\text{arr} = [3, 5, 2]$$

when

$$\text{arr}[3] = 3$$

$$\text{sum} = 4$$

Domino
Principle

$$\frac{3}{3} = 1$$

$$\frac{4}{3} = \text{INT_MAX}$$

| | 0 | 1 | 2 | 3 | 4 | 5 |
|---|---|---|---|---|---|---|
| 0 | | | | | | |
| 1 | | | | | | |
| 2 | 0 | | | | | |
| 3 | | | | | | |
| 4 | | | | | | |

$$\frac{4}{3} = \text{INT_MAX}$$

$$\frac{j}{\text{arr}[0]} = \frac{3}{3} = 1$$

For second row

```
for(int i=1; j < sum+1; i++) {  
    if (j % arr[0] == 0)  
        t[i][j] = j / arr[0]
```

else

$$t[i][j] = \text{INT_MAX}-1$$

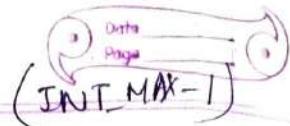
$$\frac{j}{\text{arr}[0]} = \frac{3}{3} = 1$$

$$\frac{4}{3} = \text{INT_MAX}$$

size will always
remain 1 so
we will
use arr[0] as
it is the second
row

Code variation

1st row \rightarrow INT-MAX 1st col \rightarrow INT-MAX

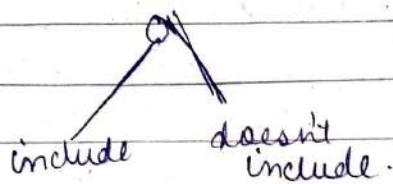


if (~~coln~~ $i-1 \leq j$)

$$t[i][j] = \min(t[i-1][j], t[i][j - \text{coins}[i-1]])$$

else

$$t[i][j] = t[i-1][j]$$



```
int minCoins( int coins[], int M, int V ) {
```

 m = coins.size() v = sum.

```
    int dp[M+1][v+1];
```

```
    for (int i=0; i < M+1; i++) {
```

```
        for (int j=0; j < v+1; j++) {
```

 if (j == 0) dp[i][j] = 0;

 if (i == 0) dp[i][j] = INT_MAX-1;

```
        }
```

```
    for (int i=1, j=1; j < v+1; j++) {
```

 if (j * coins[0] == 0) dp[i][j] = j / coins[0];

 else

 dp[i][j] = INT_MAX-1;

```
    for (int i=2; i < M+1; i++) {
```

```
        for (int j=1; j < v+1; j++) {
```

 if (coins[i-1] <= j)

 dp[i][j] = min(dp[i-1][j], 1 + dp[i-1][j - coins[i-1]]);

 else dp[i][j] = dp[i-1][j];

 if (dp[M][V] == INT_MAX-1) return -1;

 return dp[M][V];

```
}
```



Largest Common Subsequence.

- 1) Longest common substring
- 2) Print LCS
- 3) Shortest common supersequence
- 4) Print SCS
- 5) Min^m no of insertion and deletion $a \rightarrow b$
- 6) Largest repeating subsequence
- 7) length of largest subsequence of a which is a substring is b.
- 8) Subsequence pattern matching
- 9) Count how many times a appear subsequence in b
- 10) largest Palindromic subsequence
- 11) largest palindromic substring
- 12) Count of palindromic substring
- 13) minimum no of deletion in a string to make it a palindrome
- 14) Minimum no of insertion in a string to make it a palindrome

Largest Common Subsequence (Recursive)

Problem Statement

X : @⑥ c② g④ h
Y : @③ b④ e⑤ d⑥ f⑦ h⑧ u.

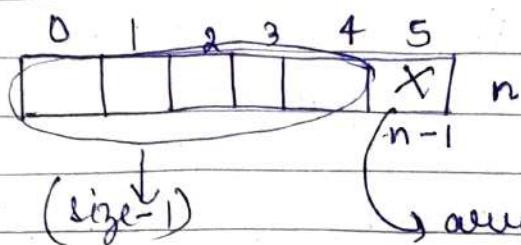
String X = abcdg

String Y = abedfhu

→ abdh
→ 4 (length of string)

diffn. → longest common subsequence - abdh.
contr. → longest common substring - ab

Recursive approach: Base condⁿ + Choice diagram + i/p small
 \downarrow
 (x, y)



fun(x, y)
 \downarrow
 $\text{fun}(x, y)$
 \downarrow
 x smaller

Base condⁿ: Think of the smallest valid input.

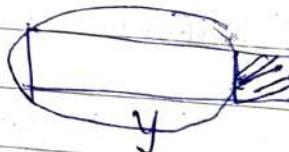
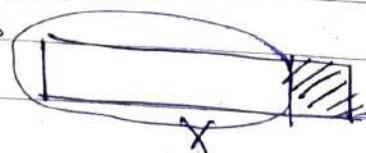
$x: \square \rightarrow n$
 $y: \square \rightarrow m$

$n=0 \quad m=0 \quad LCS=0$
 (empty string)

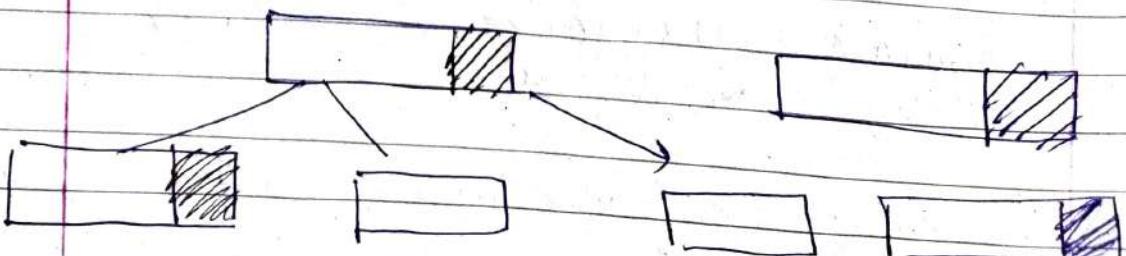
if ($n = 0$ || $m = 0$)
 return 0

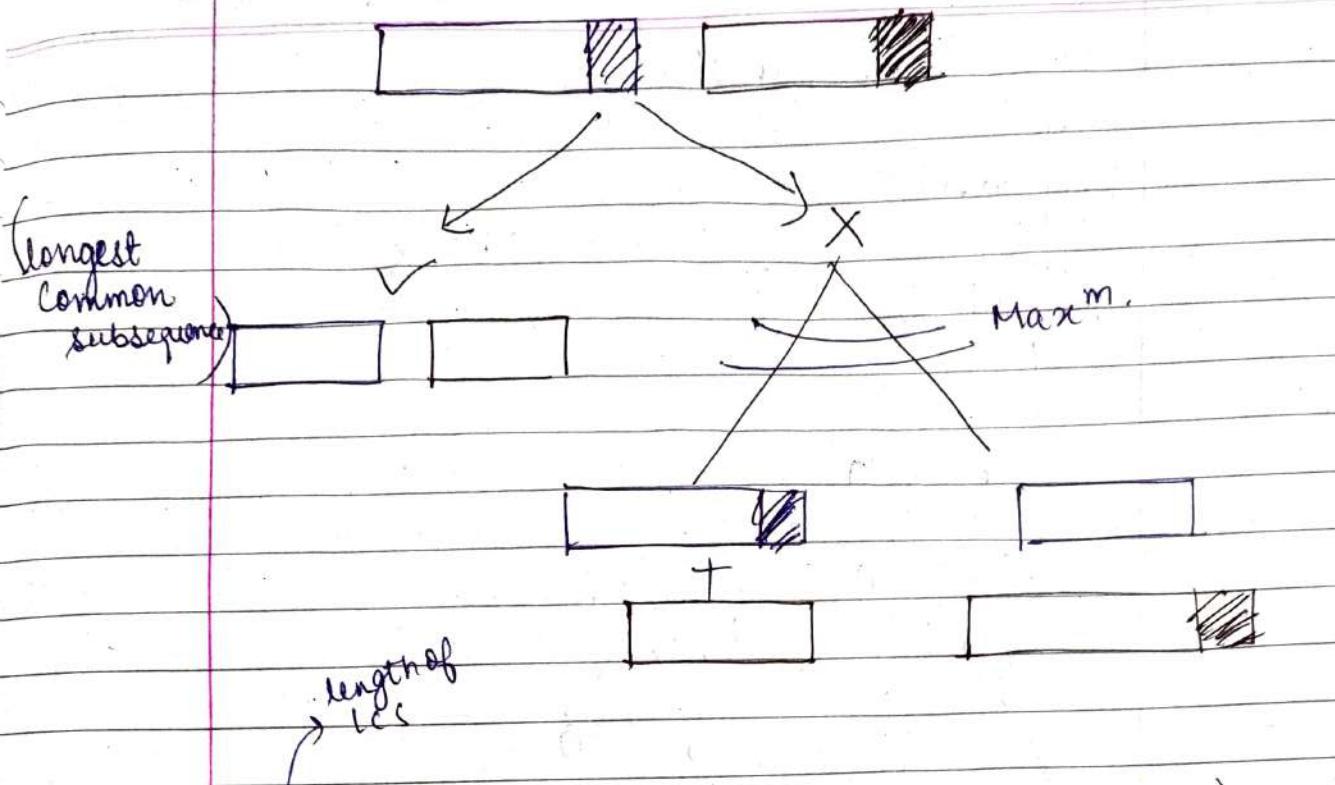
Choice diagram:
 $x: abc\text{dgh}$
 $y: abedf\text{hee}$

① when last char matches



② when last char don't match



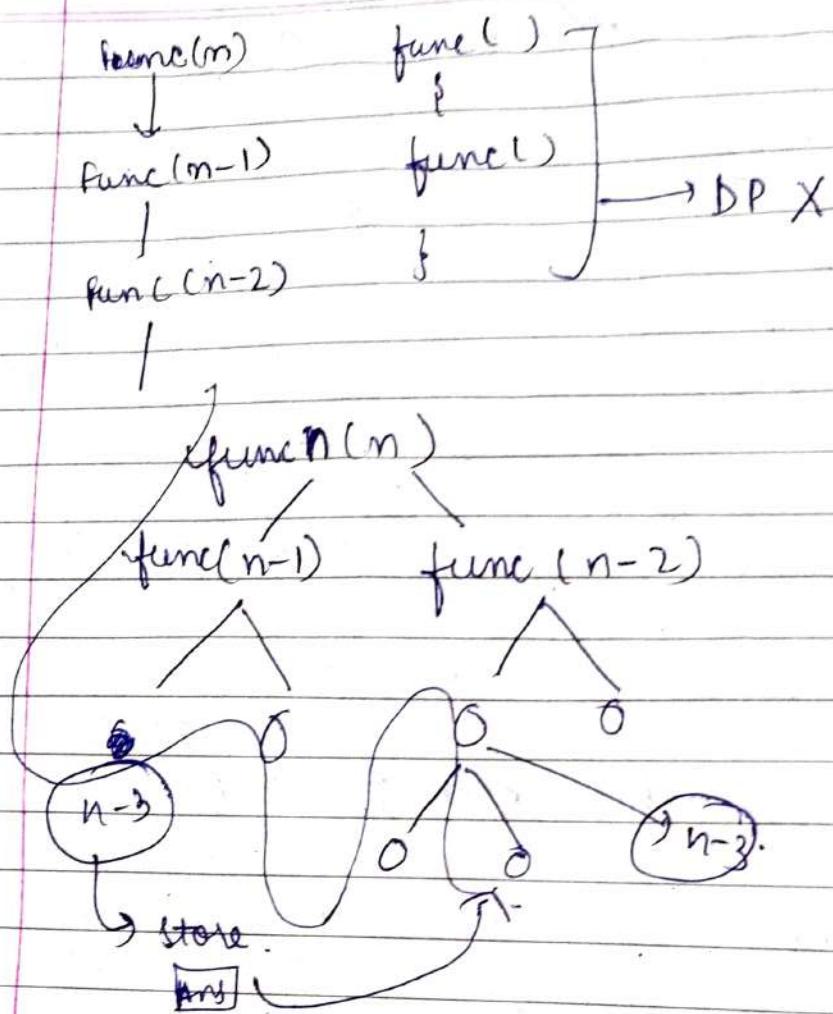


```

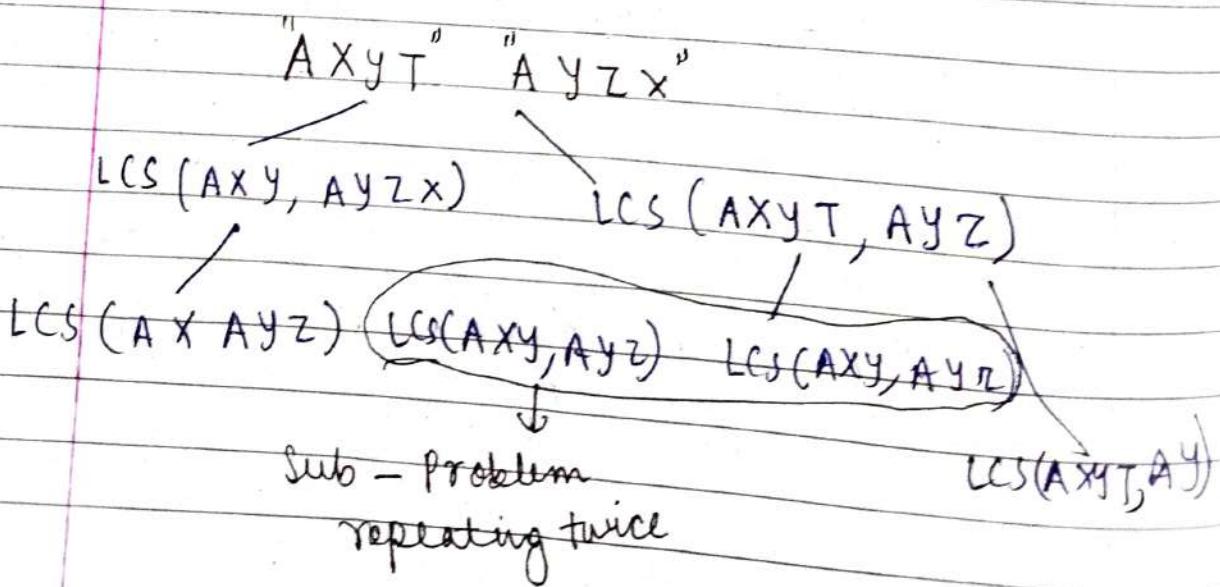
int lcs (string X, string Y, int n, int m) {
    if (n == 0 || m == 0) return 0;
    if (X[n-1] == Y[m-1])
        return 1 + lcs(X, Y, n-1, m-1);
    else
        return max (lcs(X, Y, n, m-1),
                    lcs(X, Y, n-1, m));
}

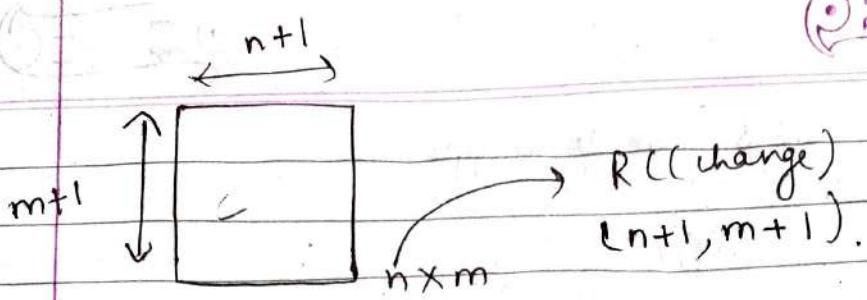
```

LCS memoization (bottom up approach)



R.C + table
 $(\boxed{\quad \quad \quad})$





int t[100][100];

~~int lcs()~~

int main()

memset(t, -1, sizeof(t));

lcs()

}

| | | | |
|----|----|----|----|
| 1 | -1 | -1 | -1 |
| -1 | 1 | -1 | -1 |
| -1 | -1 | 1 | -1 |

int lcs(string x, string y, int m, int n)

if (n == 0 || m == 0)

return 0;

if (t[m][n] != -1)

return t[m][n];

if (x[m-1] == y[n-1])

return t[m][n] = 1 + lcs(x, y, m-1, n-1);

else

return t[m][n] = max(lcs(x, y, m, n-1),

lcs(x, y, m-1, n));

:

exponential $\rightarrow O(n^2)$

LCS Top-down Approach

$x : @ \underline{a} \underline{b} \underline{c} d a f$
 $y : @ c b \underline{c} f$

O/P \rightarrow 4 (abcf)

| | | length Y | | | | | |
|------------|---|----------|---|---|---|---|---|
| | | 0 | 1 | 2 | 3 | 4 | 5 |
| length X ↓ | 0 | 0 | 0 | 0 | 0 | 0 | 0 |
| | 1 | 0 | | | | | |
| | 2 | 0 | | | | | |
| | 3 | 0 | | | | | |

Base cond'n \rightarrow initialization
 (R.S) (Top down)

int LCS(string X, string Y, int n, int m) {
~~if (n <= 0 || m <= 0)~~

```
for (int i = 0; i < m + 1; i++)
    for (int j = 0; j < n + 1; j++)
        if (i == 0 || j == 0)
            dp[i][j] = 0
```

```
for (int i = 1; i < m + 1; i++)
    for (int j = 1; j < n + 1; j++)
        if (X[i - 1] == Y[j - 1])
            dp[i][j] = 1 + dp[i - 1][j - 1]
        else
            dp[i][j] = max(dp[i - 1][j], dp[i][j - 1])
```

dp[i][j] = max(dp[i - 1][j], dp[i][j - 1])

between $dp[m][n]$;

Largest common Substring

I/p a : $\rightarrow \text{a b c d e}$
 b : $\rightarrow \text{a b f c e}$

O/p $\rightarrow 2 (ab) \quad ab, c, e$

longest = ab

a b c d e

a b f c e

② o. x. x 1

m length = 0

| | | | | | | |
|---|---|---|---|---|--|--------------------|
| | 0 | 0 | 0 | 0 | | dis continuity = 0 |
| n | 0 | | | | | |
| | 0 | | | | | |
| | 0 | | | | | |

if ($a[i-1] == b[j-1]$)

$dp[i][j] = dp[i-1][j-1] + 1$

else

$dp[i][j] = 0$

// for max^m val.

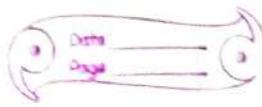
int maxi = 0;

for (i = 0 \rightarrow n+1)

 for (j = 0 \rightarrow m+1)

 max1 = max(maxi, dp[i][j])

return maxi;



Point LCS b/w 2 string

a : @ c b @ f
 b : @ b c d a f

O/P \rightarrow abc f

| | | | | | | |
|---|---|---|---|---|---|---|
| | a | b | c | d | a | f |
| a | | | | | | |
| b | | | | | | |
| c | | | | | | |
| f | | | | | | |

| | Ø | a | b | c | d | a | f |
|---|---|---|---|---|---|---|---|
| Ø | 0 | 0 | 0 | 0 | 0 | 0 | b |
| a | 0 | 1 | 1 | 1 | 1 | 1 | 1 |
| c | 0 | 1 | 1 | 2 | 2 | 2 | 2 |
| b | 0 | 1 | 2 | 2 | 2 | 2 | 2 |
| d | 0 | 1 | 2 | 3 | 3 | 3 | 3 |
| f | 0 | 1 | 2 | 3 | 3 | 3 | 4 |

Now, we need to print LCS.

"fcba"
 abc f ← reverse

if equal $i, j \rightarrow (i--, j--)$

if not equal $i, j \rightarrow \max((i-1, j), (i, j-1))$

when equal

$(i--, j--)$

4

when not equal

$(\max \text{ between } a \& b)$

a

5

4

add in string

nothing to be added

int $i = m, j = n;$

String $s = " "$;

while ($i > 0 \&& j > 0$)

{ a b
if ($t[i-1] == t[j-1]$)

{

$s.push_back(t[i-1])$

$i--;$

$j--;$

}

else {

if ($t[i][j-1] > t[i-1][j]$)

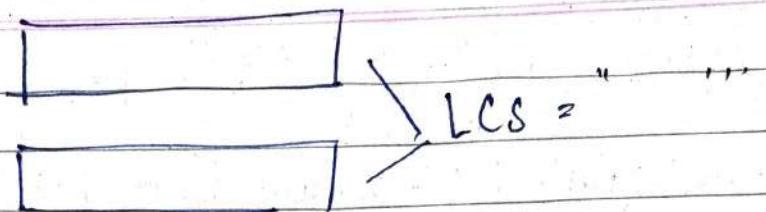
$j--;$

else

$i--;$

}

reverse ($s.begin(), s.end()$) ;



if $a[i-1] == b[j-1]$
 S.push_back($a[i-1]$)
 $i--$
 $j--$

else

None in the direction of maximum
 reverse (s.begin(), s.end())

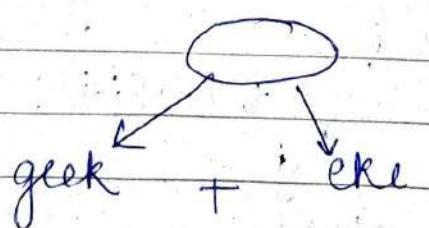
24. Shortest Common Supersequence

a: "geek."

b: "eke"

1) Problem statement :

→ 2 strings are given. We need to merge the string such that we get both the subsequence.
 Mix 2 sequence to make it a supersequence.



Op

(geek|eke), (geek|e)

find shortest

Subsequence \rightarrow Sequence of events

$s_1 \quad s_2 \quad s_3$

Order should be maintained
but don't need to be
continuous always

a : A G G T A B

b : G X T X A Y B

Super sequence : A G I G I T G I X A B T X A Y B

A G I G I X T X A Y B \rightarrow shortest supersequence

\swarrow \searrow
(A G I G I T A B) (G I X T X A Y B)

Give the length in O/P.

The one letter which is common write once

A G G I T A B
G I X T X A Y B
 \downarrow

G I T A B is common in both
i.e LCS.

Now thinking the brute force approach we
can merge both the string & then subtract
G I T A B.

a: AGIGTAB

b: GXTXAYB

Worst case: $a+b = AGIGTAB GXTXAYB$

length of SS

$a+b$

AGIGTABGXTXAYB - GTAB



AGIGXTXAYB

Shortest length = $(m+n) - LCS$

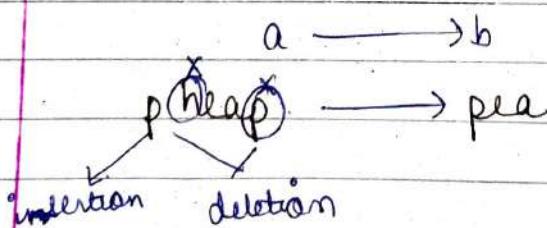
25. Minimum no of insertion & deletion to convert a string a to string b.

a: heap

$\%P = 1$ (Insertion)

b: pea

$\%P = 2$ (deletion)

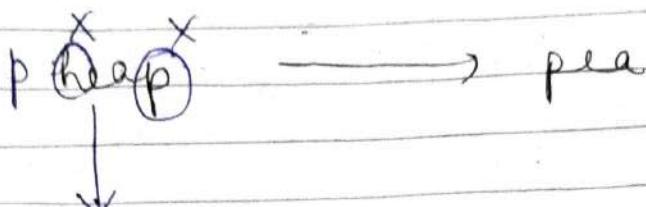


When to use LCS.

| | I/P | Q | O/P | | | |
|------------|-----|----------------|-----|--|--|--|
| LCS | a: | LCS | int | | | |
| | b: | | | | | |
| Given Ques | a: | insertion/ | int | | | |
| | b: | deletion | | | | |

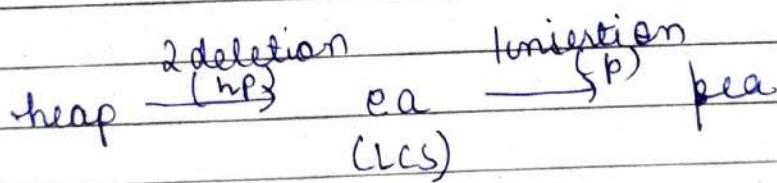
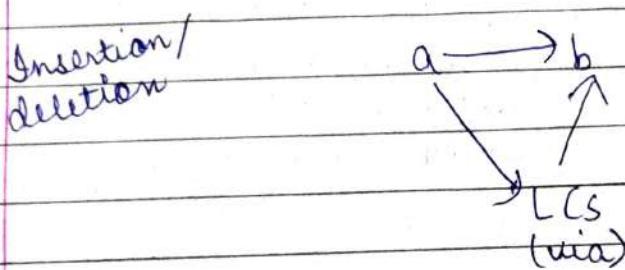
} → Pattern matching algorithm

Convert heap to pea.



"ea" is the LCS of both string

heap → pea



$$\text{No of deletion} = a \text{ length} - \text{LCS}$$

$$\text{No of insertion} = b \text{ length} - \text{LCS}$$

26. Longest Palindromic Subsequence

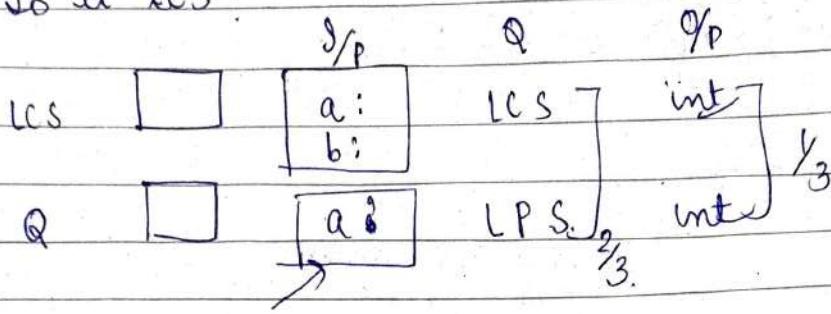
Problem Statement : S : agbcba

$$O/P = 5$$

$S \leftarrow \text{longest}$ [abcba]
 [bcb]
 [b] ✓

O/p \rightarrow 5 (abcbba)

2) Is it LCS?



LCS?

LCS S/P a, b
 LPS S/P a b = func(a)

(hidden
string/redundant)

"agbcba"
 ↓

LPS



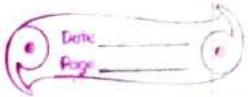
a \rightarrow agbcba

b \rightarrow reverse(a) abc bga.

↓
 LCS

@bCba@ @
 @g bCba@ > abcba

LPS(a) \equiv LCS (a, reverse(a))
 LPS(agbcba) \rightarrow return(agbcba,



28. Minimum no. of deletion in a string to make it palindrome

a : "agbcba"

$$g_p = 1$$

S : agbcba

| | | | | | |
|---|---|---|---|---|---|
| a | g | b | c | b | a |
|---|---|---|---|---|---|

↓ No of deletions (Minimize)

→ New string
(palindrome)

~~agbcba~~
bcb (palindrome) ageba (notpalindrome)
3

agbcba
 bcb (3) (LPS) C (5) (LPS) abcba (1) (LPS)
 Min^m.

\uparrow length of LPS \swarrow no of deletions \downarrow

LPS

\downarrow min^m no of deletions.

\rightarrow Palindromic subsequence.

length of LPS α $\frac{1}{\text{no of deletions}}$

(longest palindromic subsequence)

agbcba

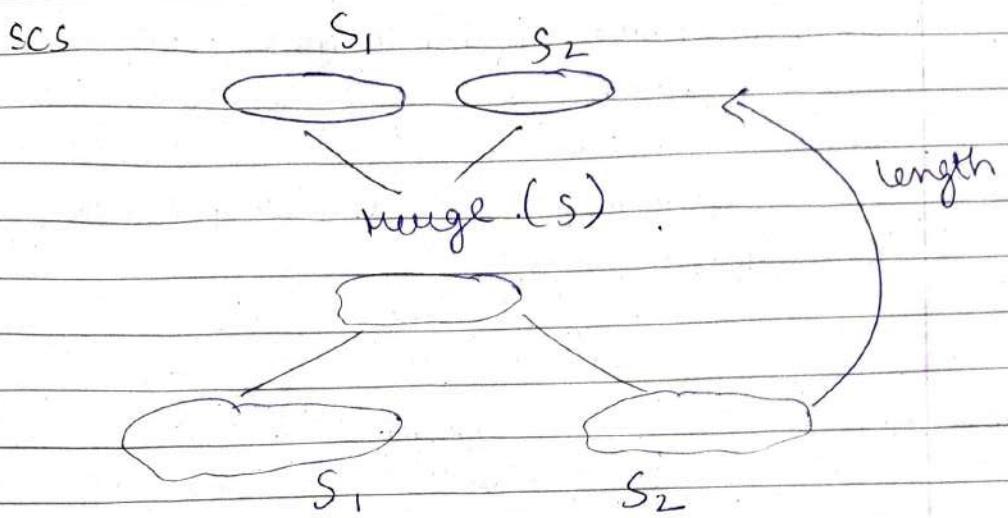
\downarrow
LCS(s, reverse(s))

\downarrow
abcbab (LPS)

$$\text{Min no of deletion} = S - \frac{\text{LPS length}}{\text{(length)}}$$

29. Print shortest common subsequence

i/p a: acbcbf
b: abcdaf
o/p: acbcdaf



worst case $a c b c f$ $a b c d a f$

$a c b c d a f \rightarrow$ print the whole string

$m+n$
 $a c b c f a b c d a f$

now

$m+n - LCS$

LCS $\xrightarrow{\text{similar}}$ SCS

Print LCS Print SCS

| ϕ | a | b | c | d | a | t | \rightarrow LCS |
|--------|---|---|---|---|---|---|-------------------|
| a | 0 | 1 | 1 | 1 | 1 | 1 | |
| c | 0 | 1 | 1 | 2 | 2 | 2 | |
| b | 0 | 1 | 2 | 2 | 2 | 2 | |
| c | 0 | 1 | 2 | 3 | 3 | 3 | |
| f | 0 | 1 | 2 | 3 | 3 | 3 | 4 |

LCS → common ni hai to move kar jao
 common hai to print karo.

SCS → common hai to ek baar print kro
 else print karo.

LCS.

```
String s = " ";
while (i > 0 && j > 0) {
    if (a[i-1] == b[j-1]) {
        s.push_back(a[i])
        i--;
        j--;
    }
    else {
        if (t[i][j-1] > t[i-1][j])
            j--;
        else if (t[i-1][j] > t[i][j-1])
            i--;
    }
}
```

SCS

```
String s = " ";
while (i > 0 && j > 0) {
    if (a[i-1] == b[j-1]) {
        s.push_back(a[i])
        i--;
        j--;
    }
    else {
        if (t[i][j-1] > t[i-1][j])
            s.push_back(b[j-1]);
        else if (t[i-1][j] > t[i][j-1])
            s.push_back(a[i-1]);
    }
}
while (i > 0)
    s.push_back(a[i-1]);
i--;
while (j > 0)
    s.push_back(b[j-1]);
j--;
```

- i) Now in SCS code we include the letter on moving at $i-1$ see $j-1$.
- ii) Now in $(i > 0 \&\& j > 0)$ we need to change because in case of LCS we don't need to stop at the topmost but in SCS we need to.

In case
of LCS

| | ϕ | a | b | f |
|--------|--------|---|---|---|
| ϕ | 0 | 0 | 0 | 0 |
| a | 0 | | | |
| c | 0 | 0 | | |
| b | 0 | | | |
| d | 0 | | | |

| | ϕ | a | b | f |
|--------|--------|---|---|---|
| ϕ | 0 | 0 | 0 | 0 |
| a | 0 | | | |
| f | 0 | | | |
| b | 0 | | | |
| d | 0 | | | |

Some
case
stop
here.

"ac",] lcs ("")

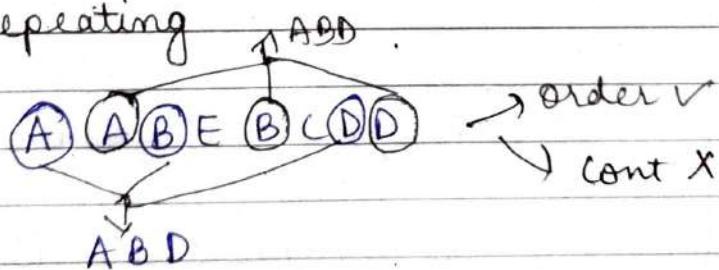
"ac",] scs (ab)

longest repeating subsequence

Str = "A A B E B C D D"

%p = 3

Problem Statement - find a subsequence which is repeating



ABD (2x)] longest \rightarrow "ABD"
AB (2x) %p = 3

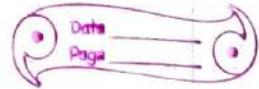
S = A A B E B C D D

E → 3
C → 5

A F 2 3 4 5 6 6
A K B E B C D D \rightarrow LCS = A A B (E) B C D D.

A A B B D D
once ↙ → ans

E & C are occurring once.
letter at the same index don't take



$E \rightarrow 3 \xrightarrow{i} j \quad i = j X$
 $C \rightarrow 5$

$A \rightarrow 0, 1$
 $A \rightarrow 0, 1$ $\xrightarrow{\text{diff index}}$
 $B \rightarrow 2, 4$
 $B \rightarrow 2, 4$

$\xrightarrow{\text{same index}}$

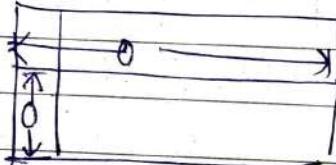
Sum up:

So don't take A of same index take from other index.

$a = s$

$s \rightarrow \text{empty set} \quad \text{LCS } (i = j)$
 $b = s$ \times

where



We can't take same letter for both the string during
(LCS i.e. $i = j$)

if $(a[i-1] = b[j-1] \text{ & } i = j)$

$$+ [i][j] = t[i-1][j-1] + 1$$

else

$$+ [i][j] = \max (+ [i][j-1], + [i-1][j])$$

31. Sequence Pattern Matching

$a = "Axy"$
 $b = "ADXCPY"$

O/P \rightarrow T/F

Problem statement: Is string "A" a subsequence of "B"

a : AXy

b : ADXCPY

b : A D X C P Y

a : AXy

O/P → True.

LCS : AXy

| S/P | | | |
|-----|--------|----------|-----|
| LCS | a
b | Sequence | int |
| Q | a
b | Sequence | T/F |

S/P

b : ADXCPY

b : A D X C P Y

a : AXy

a : AXyZ

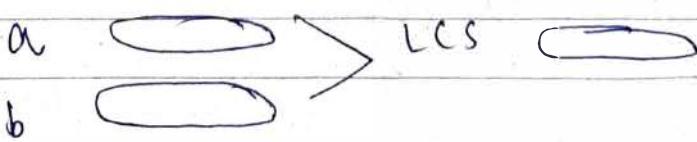
LCS : AXy (LCS == a)

LCS : A X y

(LCS == a)

if (LSS == A) return true
else
return false

Can it be done using length ?



a: $A \times Y$ (3)

b: $A \times D \times C \times P \times Z$ (6)

$LCS = 0$ to $\min(m, n)$

$LCS = 2$

if ($LCS == a.length()$)

return true

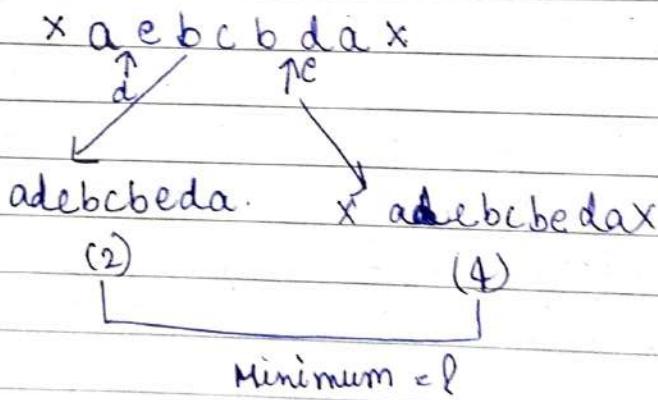
else

return false.

32. Min^m no of insertions in a string to make it palindrome

S/p: S: "aebcbda"

Q/p $\rightarrow 2$



Min^m no of deletion to make a string palindrome.

a e b e b d a

a e a
(4)

a b c b a
(2)

$s.length() \rightarrow LPS.$ (deletion)

$S : ^a e b c b d a'$

Palindromic
string

$\boxed{a} \ e \ \boxed{b \ c \ b} \ d \ \boxed{a}$

$e \times \left\{ \begin{array}{l} \text{delete} \\ \text{d} \times \left\{ \begin{array}{l} e \& d \end{array} \right. \end{array} \right. \right.$

$a \ e \checkmark \quad \left\{ \begin{array}{l} \text{insert e \& d} \\ \text{d} \checkmark \quad \left\{ \begin{array}{l} \text{make their} \\ \text{pair} \end{array} \right. \end{array} \right. \right.$

no of insertions = no of deletions
||

$s.length - LPS.$

Deletion \rightarrow single (Pair) $\rightarrow X$
Insertion \rightarrow Single (Pair) $\rightarrow \checkmark$

33. MCM (Matrix chain Multiplication)

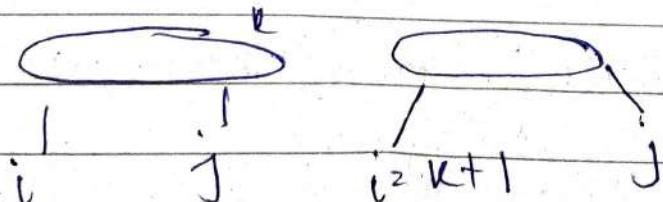
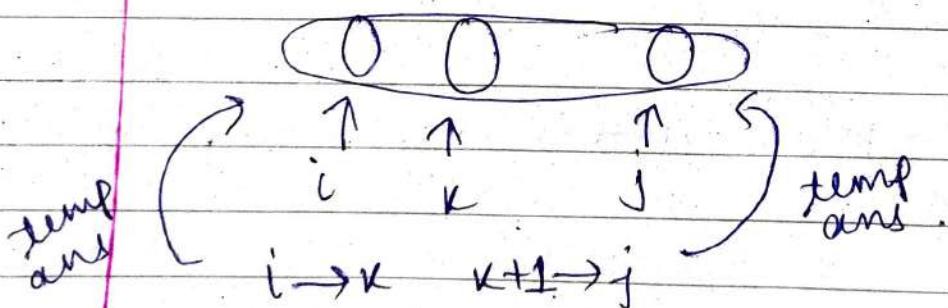
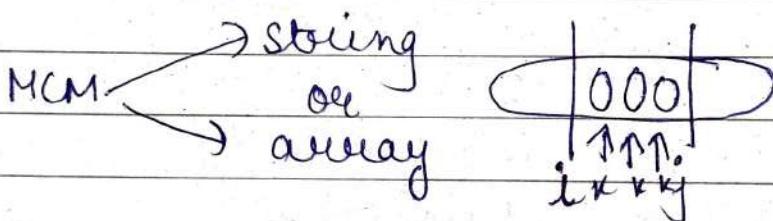
- 1) MCM
- 2) Painting MCM
- 3) Evaluate Exp. to True / Boolean Parenthesization
- 4) Min / Max value of an Expr.
- 5) Palindrome partitioning
- 6) Scramble string
- 7) Egg Dropping problem

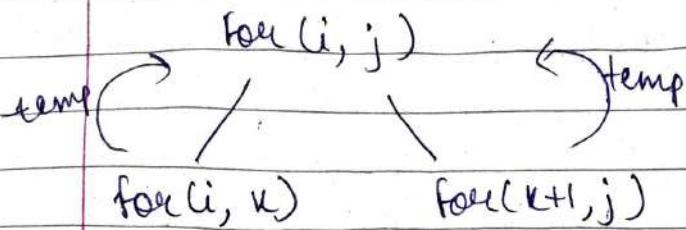
MCM format

- 1) MCM
- 2) Printing MCM
- 3) Evaluate Expr to True/ Boolean parenthesization
- 4) Min / Max value of an Expr
- 5) Palindrome partitioning
- 6) Scramble string
- 7) Egg dropping problem

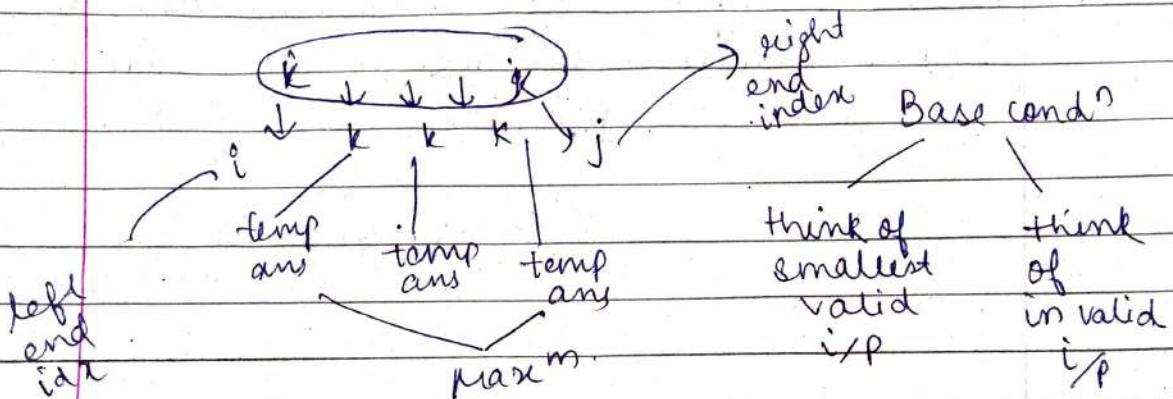
Identify \rightarrow MCM \rightarrow basic format

Identification + Format





$\text{ans} \leftarrow f(\text{temp ans})$



int solve (int arr[], int i, int j)

{

if ($i > j$) \rightarrow this may be diff accn'-le
return 0;

for (int k = i ; k < j ; k++)

// Calc. temp ans

temp ans = solve(arr, i, k) +
solve(arr, k+1, j);

ans = func(temp ans)

MCM (Recursive)

Problem

Statement : $\text{arr}[] = \{40, 20, 30, 10, 30\}$

$\overset{m \times n}{A_1}, \quad \overset{m \times m}{A_2}, \quad \overset{n \times n}{A_3}, \quad A_4$

Some matrix are given like A_1, A_2, A_3, A_4 we have to multiply the matrix to reduce the cost (no of multiplications).

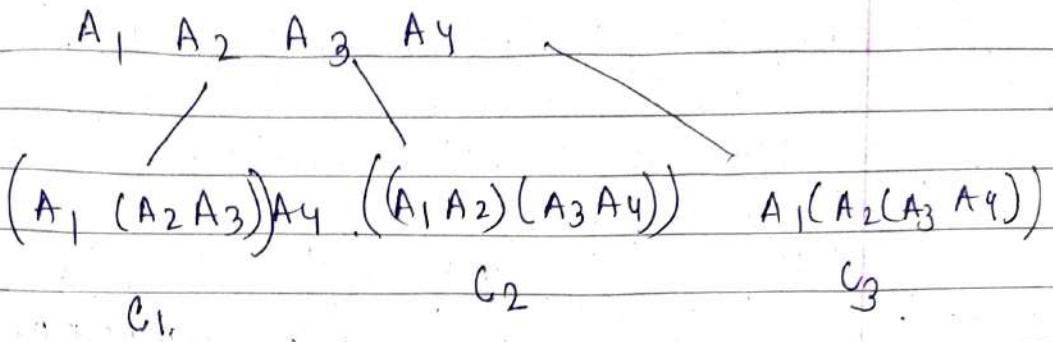
$b = c$ (then only we can multiply)

$$\begin{array}{c}
 [] \\
 \downarrow \\
 a \times b
 \end{array}
 \quad
 \begin{array}{c}
 [] \\
 \downarrow \\
 c
 \end{array}
 \quad
 \begin{array}{c}
 [] \\
 \downarrow \\
 d \times e
 \end{array}$$

$$\begin{array}{ccc}
 2 \times 3 & & 3 \times 6 \\
 \downarrow & \downarrow & \\
 2 & 3 & 6
 \end{array}
 = 36 \text{ cost (no of multiplications)}$$

A_1, A_2, A_3, A_4

dimension [] = {x, y, z, w}



A → 10 * 30

B → 30 * 5

C → 5 * 60

$$(A \setminus B)C = 10 * 30 * 30 * 5$$

~~10 * 500~~ ~~500~~ 0. ~~10 * 500~~ ~~500~~ 0.

$$(A B) C = \left(10 * \overbrace{30 * 30}^{} * 5 \right) S * 60.$$

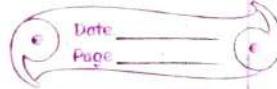
$$A_1 A_2 A_3 \quad A_4 = 10 * 5 * 30 * 5 * 60$$

↓ do parenthesis
in such a way
cost is the least

$A_1 (A_2 A_3) A_4$ C1
 $(A_1 A_2) (A_3 A_4)$ C2
 $A_1 (A_2 (A_3 A_4))$ C3

min cost
 $30 * 5 \ 5 * 60$
 $B \ C \rightarrow 30 * 5 * 60$
 $A \ B \ C \rightarrow 10 * 30 \ 30 * 60$
 $10 * 30 * 60$
 Cost

Bracket lagane par aleg aleg cost aa rahi
hai.



$$arr : \{ 40, 20, 30, 10, 30 \}$$

arr size n

n-1 matrix x

$$A_1 \rightarrow 40 * 20$$

$$A_2 \rightarrow 20 * 30$$

$$A_3 \rightarrow 30 * 10$$

$$A_4 \rightarrow 10 * 30$$

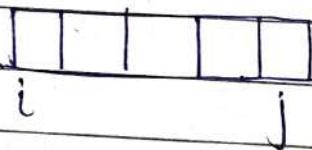
$$A_1 * A_2 * A_3 * A_4$$

Cost (min^m)

$$A_i \rightarrow arr[i-1] * arr[i] \quad (\text{no of multiplication should be less})$$

$$A[i] = arr[0] + arr[1]$$

Next step → Format



$$A_1 \ A_2 \ A_3 \ A_4$$

↓

$$(A_1) \ (A_2 \ A_3 \ A_4)$$

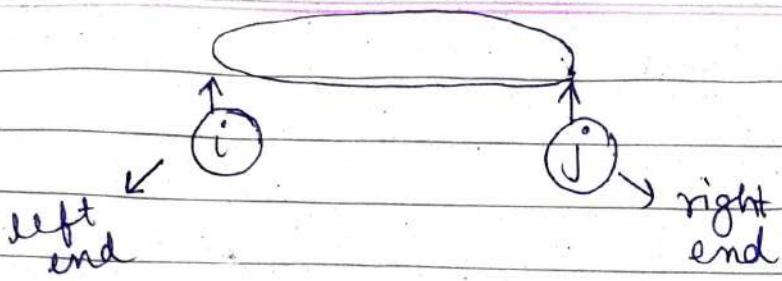
↑

↓

Min cost + min cost temp ans

$$(A_1 \ A_2 \ A_3) (A_4)$$

put brackets on different places to get the answer
slide k at different positions.



| | | | | |
|----|----|----|----|----|
| 40 | 20 | 30 | 10 | 30 |
|----|----|----|----|----|

i i j

$$A_i \rightarrow arr[i-1] * arr[i]$$

when $i=0$ $A_i \rightarrow arr[-1] * arr[0]$, X

when $i=1$ $A_i \rightarrow arr[0] * arr[1]$, ✓

$$A_j \rightarrow arr[j-1] * arr[j]$$

$$arr[3] * arr[4]$$

$i > j \rightarrow \text{size} = 0$

$i = j \rightarrow \text{size} = 1$

$$\begin{matrix} i = 1 \\ j = n-1 \end{matrix} \quad \} \text{ return cost}$$

$\frac{(i-1)}{n}$
middle

int solve (int arr[], int i, int j)

{

if ($i \geq j$)

return 0;

int mn = INT_MAX;

for (int k = i ; k <= j-1 ; k++) {

int temp ans = solve (arr, i, k) +
solve (arr, k+1, j)

+ $arr[i-1] * arr[k] * arr[j]$

if (temp ans)

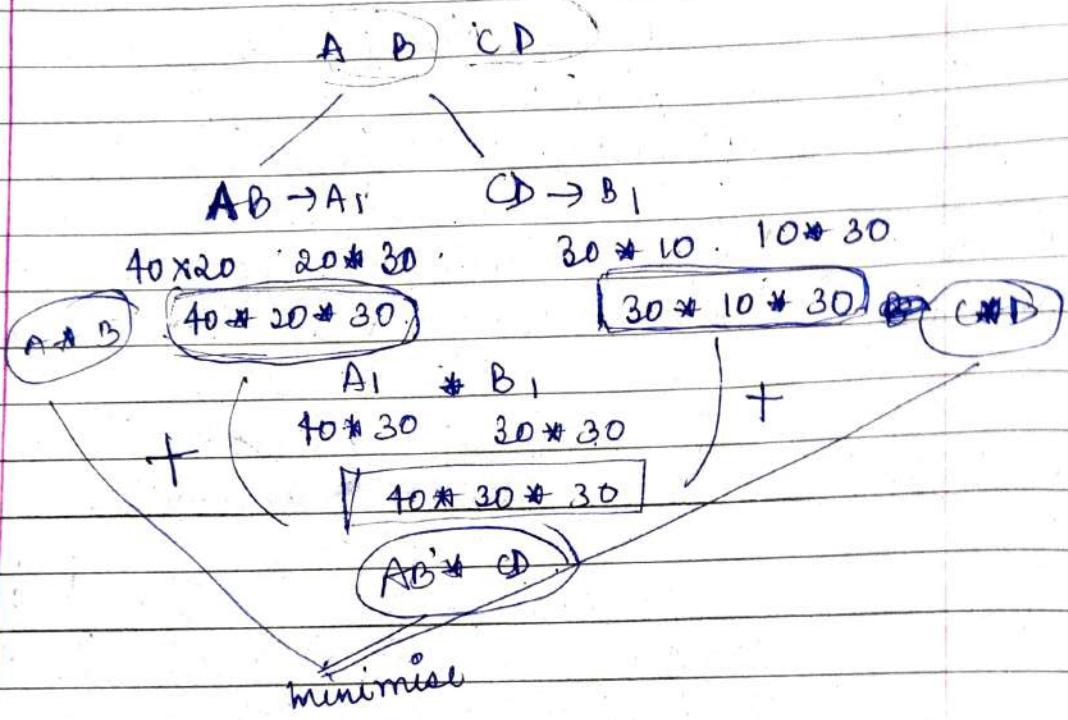
{

mn = temp ans;

}

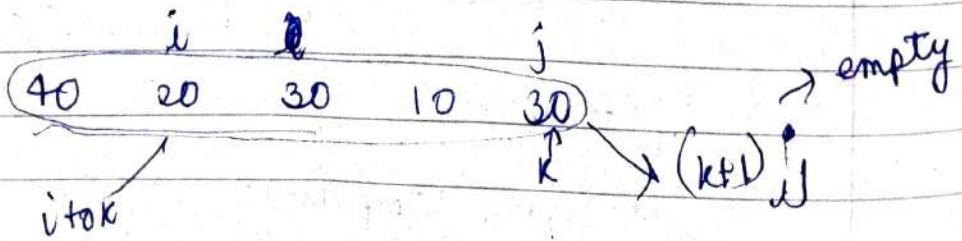
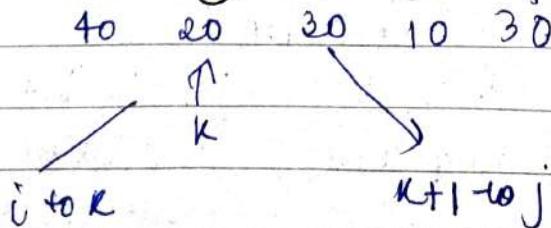
between mn

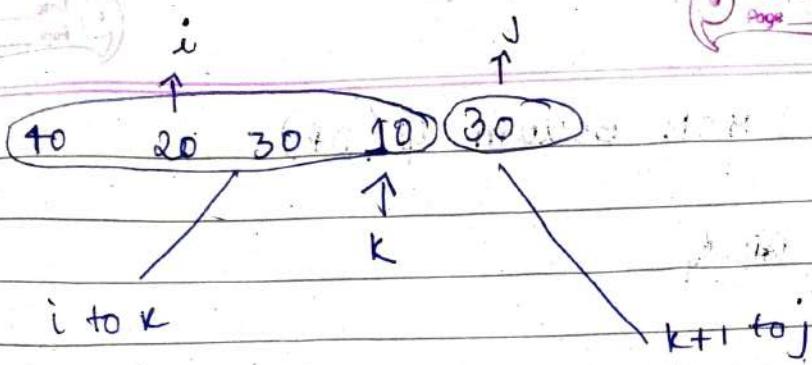
f



Steps for NCM

- 1) find i & j value
- 2) find eight base condn.
- 3) Move K \rightarrow i to j. (find K-loop scheme)
- 4) calculate cost for temp ans.



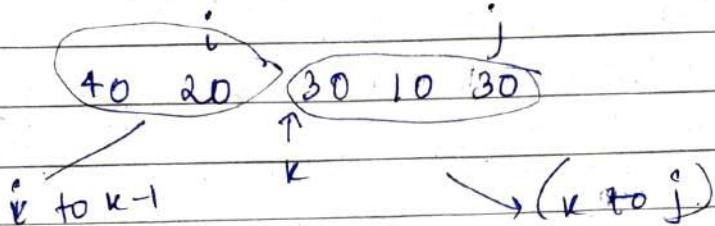


$$40 * 20$$

$$20 * 30$$

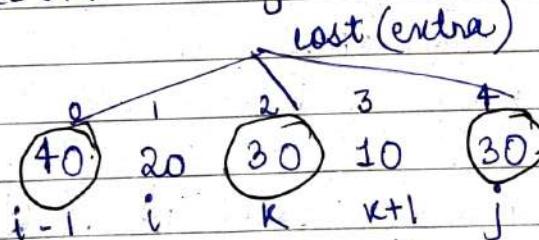
$$30 * 10$$

$$k = i \quad k = j - 1 \quad (i \rightarrow k \quad k+1 \rightarrow j)$$



2 schemes.

$$\begin{array}{ll} k = i \rightarrow k = j - 1 & i \rightarrow k \quad k + 1 \rightarrow j \\ k = i + 1 \rightarrow k = j & i \rightarrow k - 1 \quad k \rightarrow j \end{array}$$



for ($i \rightarrow k$)

$$40 * 20 * 20 * 30$$

solve($i \rightarrow k$)

$$40 * 20 * 30$$

for ($k+1 \rightarrow j$)

$$30 * 10 \quad 10 * 30$$

$$30 * 10 * 30$$

solve ($k+1 \rightarrow j$)

$$40 * 20 * 30 * 30$$

arr[i-1]

$\leftarrow 40 * 30 * 30 \rightarrow arr[j]$

$\rightarrow arr[k]$

dimension arr: []

MCM Bottom Up (DP)

~~Top Down~~

(dimension) arr[] : []

int dp[100][100]

memset (dp, -1, sizeof(dp))

int solve (int arr[], int i, int j)

if (i >= j)
~~dp[i][j] = 0~~
 return 0;

if (dp[i][j] == -1)
 return dp[i][j];

int mn = INT_MAX;
 for (int k = i ; k < j ; k++) {

int temp_ans = solve (arr, i, k) +
 solve (arr, k+1, j) +
 arr[i-1] * arr[k] * arr[j];

if (temp_ans < mn)
 mn = temp_ans;
 }

dp[i][j] = mn;
 return mn;



Palindrome Partitioning

```
int dp[100][100];
```

class Solution {

public:

```
int solve(int arr[], int i, int j) {
```

```
if (i >= j)
```

```
return 0;
```

```
if (dp[i][j] == -1)
```

```
return dp[i][j];
```

```
int mn = INT_MAX;
```

```
for (int k = i; k <= j - 1; k++) {
```

```
int temp = solve(arr, i, k) + solve(arr, k + 1, j)
```

```
+ arr[i - 1] * arr[k] * arr[j]
```

```
mn = min (temp, mn);
```

```
}
```

```
dp[i][j] = mn;
```

```
return dp[i][j];
```

```
}
```

```
int matrixmult (int N, int arr[]) {
```

```
memset (dp, -1, sizeof (dp));
```

~~int ans =~~ solve (arr, 1, N - 1);

```
return ans;
```

```
}
```

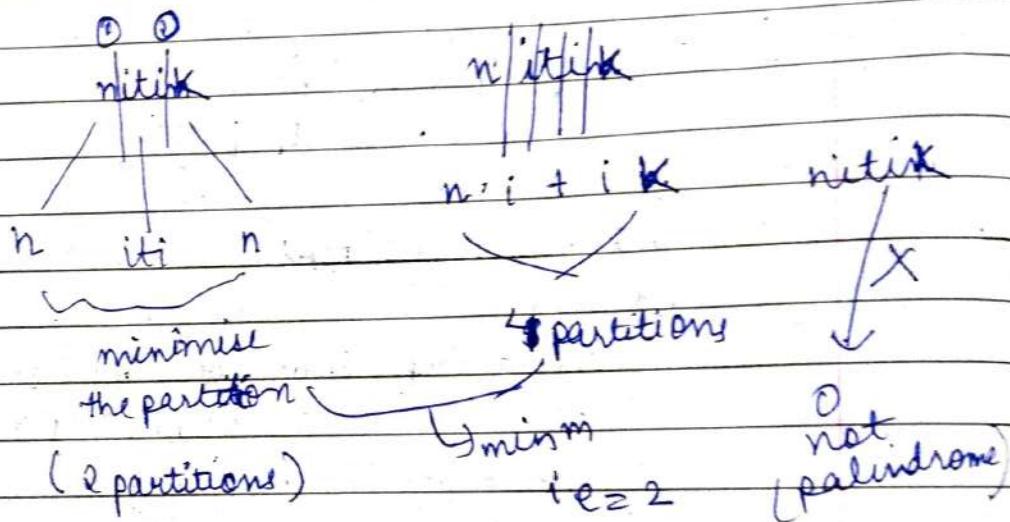
```
,
```

Palindrome partitioning

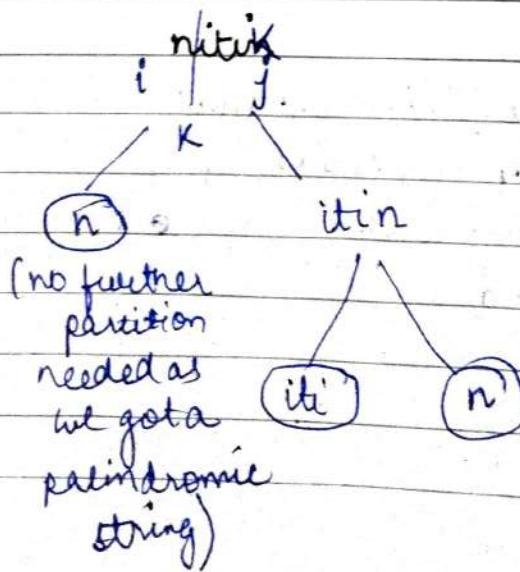
D) Problem statement

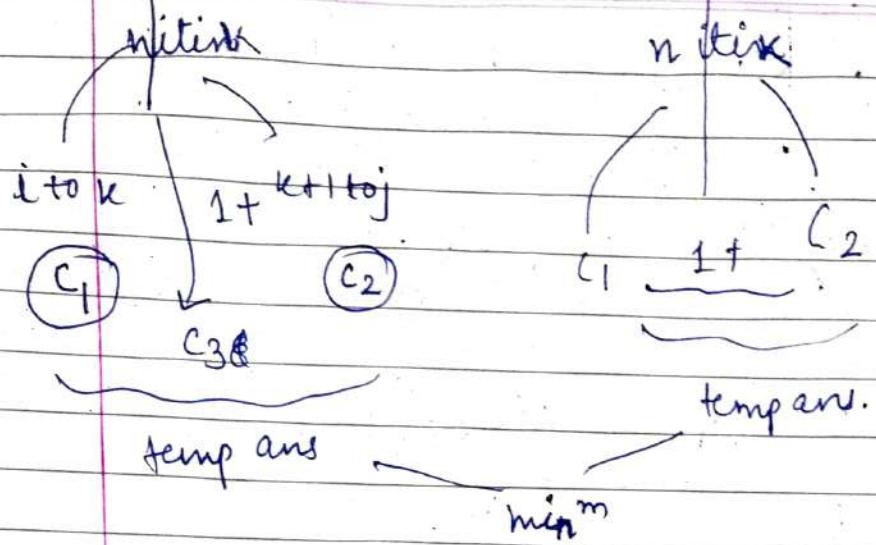
S: nitik → it is ~~not~~ a whole
 is a palindrome
 Q: 2

divide string in such a way all the strings are
 palindrome



worst case partition: ~~2~~ n-1





Format

1. Find i & j
2. Find B.C.
3. Find K loop (scheme)

4. $\text{temp ans} \rightarrow \min^m$

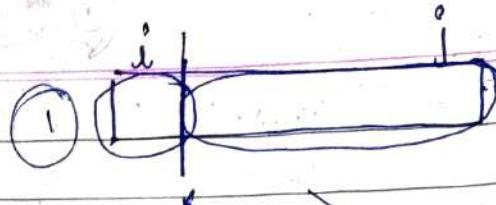
Recursive code

B.C $i = j$ $\text{size} = 1$ $\left\{ \begin{array}{l} \text{if size is not given or 0 or equal} \\ \text{no of partition would be} \\ 0 \text{ as string is empty} \end{array} \right.$

$i > j$ $\text{size} = 0$ $\left\{ \begin{array}{l} \text{if size is not given or 0 or equal} \\ \text{no of partition would be} \\ 0 \text{ as string is empty} \end{array} \right.$

$\text{is palindrome}(s, i, j)$ $\left\{ \begin{array}{l} \text{& if palindrome partition} \\ \text{should be 0.} \end{array} \right.$

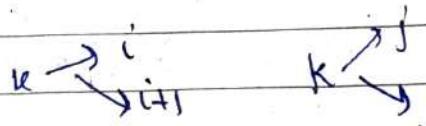
loop



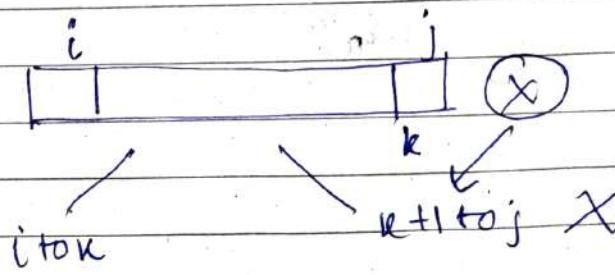
$k+1 \rightarrow j$

$$k=i \quad k=j-1$$

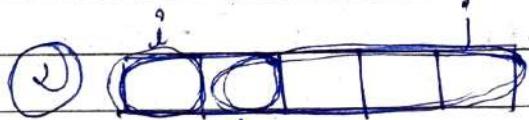
$i \rightarrow k \quad k+1 \rightarrow j$



$i \rightarrow k \quad k \rightarrow j$



$k+1 \rightarrow j \quad X$



$k(i+1)$

~~$i \rightarrow k-1$~~

$k \rightarrow j$

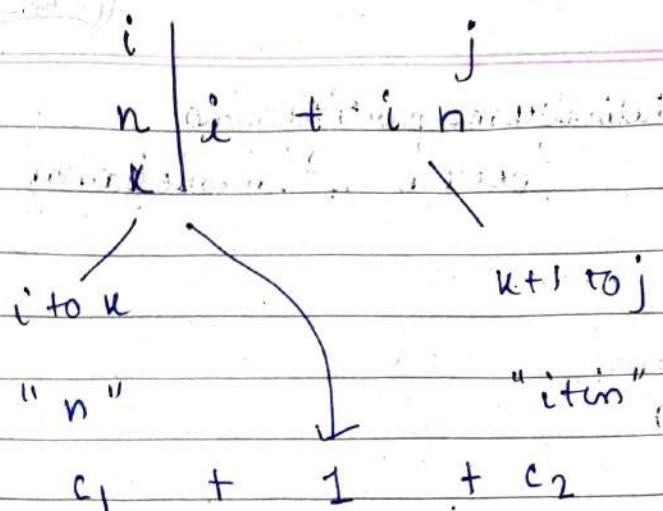
$$k = i+1$$

$$k = j$$

- 1) $k=i \quad k=j-1 \quad i \rightarrow k \quad k+1 \rightarrow j$
- 2) $k=i+1 \quad k=j \quad i \rightarrow k-1 \quad k \rightarrow j$

for (int $k=i$; $k=j-1$; $k++$)
{

 int temp = solve(s, i, k) + solve(s, k+1, j)
 + 1



} ans = min (ans, temp)

} return ans;

Final code

```
int solve (string s, int i, int j) {
    if (i >= j)
        return 0;
```

```
    if (isPalindrome (s, i, j) == True)
        return 0;
```

```
    int mn = INT-MAX;
```

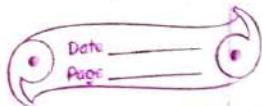
```
    for (int k = i ; k <= j-1 ; k++) {
```

```
        int temp = 1 + solve (s, i, k) + solve (s, k+1, j);
```

```
        mn = min (mn, temp);
```

```
}  
return mn;
```

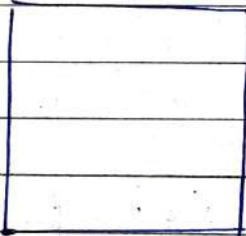
memset(\rightarrow only allow 2 values)



Palindrome partitioning (bottom up) (memoization)

i/p = nitin nitik
o/p = 0 2.

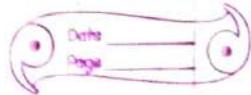
- 1) Initialise the matrix with -1.
- 2) Check if the value is $\begin{cases} -1 & \text{value isn't evaluated} \\ \neq -1 & \text{value is evaluated so we stored} \end{cases}$
Return the value.



i & j
changes.

\checkmark
variables
which
changes
basically

```
int dp [100][100];  
memset (dp, -1, sizeof(dp));  
int solve ( string s, int i, int j ) {  
    if ( i >= j )  
        return 0;  
  
    if ( is palindrome(s, i, j) )  
        return 0;
```



```
if (dp[i][j] != -1)
    return dp[i][j];
int mn = INT_MAX;
for (int k = i ; k <= j-1 ; k++) {
    int temp = solve(s, i, k) + solve(s, k+1, j) + 1
    mn = min (mn, temp);
}
dp[i][j] = mn;
return mn;
}
```

```
int palindromePart ( int string s) {
    int n = s.length() - 1;
    int ans = solve(s, 0, n);
    return ans;
}
```

```
bool ispalindrome (string s, int i, int j) {
```

```
if (i == j)
    return true;
```

```
if (i > j)
    return true;
```

```
while (i < j)
    if (s[i] != s[j])
        return false
    else
```

```
    i++;
    j--;
```

GEEKSFORGEeks = ✓

Interviewbit = ✗

Further Optimization

Why the above code is not most optimized?

Since we are calling

`int temp = 1 + solve(s, i, k) + solve(s, k+1, j)`

^{RC}
(left) , ^{RL}
(right)

There is a possibility, might be one of the RC is solved or called.

The code will remain same but the diff* is

`int temp = 1 + solve(s, i, k) + solve(s, k+1, j)`

`if (+[i][k]) = -1)`

`left = +[i][k].`

`else`

`left = solve(s, i, k),`

`+[i][k] = left`

`if (+[k+1][j]) = -1)`

`right = + [k+1][j]`



else

```
    right = solve(s, k+1, j)
    t[k+1:j] = right
```

```
int temp = 1 + left + right;
```

evaluate expression to true
Boolean parenthesization

String : True F and T.
Op : & ^

problem : A string is given : string might have some statement characters like T → true & → false



How to put bracket such that the expression evaluates to true.

Find the no of ways in which when bracket is put it evaluates to True.

Ques: "T | F & T ∨ F"

no of ways

$$\begin{array}{ccccc}
 (()) & () & \xrightarrow{\quad} & T & \\
 () & () & \xrightarrow{\quad} & F & \\
 () & () & \xrightarrow{\quad} & T & \\
 () & () & \xrightarrow{\quad} & T &
 \end{array}
 \text{3 ways.}$$

In MCM we put brackets for min^m cost &
in this also we do the same.

(T | F & T ∨ F)

$\xrightarrow{k-1} \xrightarrow{k} \xrightarrow{k+1}$

We need to break
bracket on
operator

Exprⁿ: operator Exprⁿ:

4 steps:

- 1) find i & j
- 2) find base condⁿ
- 3) Find k loop
- 4) temp ans & funcⁿ

↓
Main ans.

$i \quad j$
 $T \mid F \& T \wedge F$

1) $i = 0$
 $j = \text{st.length}() - 1$

2) BC $i \quad j$
 $(T \text{ or } F \text{ and } T) \text{xor}(F)$
 $i \rightarrow k-1 \quad k \quad k+1 \rightarrow j$

(left) Exprⁿ¹ XOR Exprⁿ² (right)
 $i \rightarrow k-1 \quad k \quad k+1 \rightarrow j$

$$T \wedge T = \text{False}$$

$$F \wedge F = \text{False}$$

$$T \wedge F = \text{True}$$

$$F \wedge T = \text{True}$$

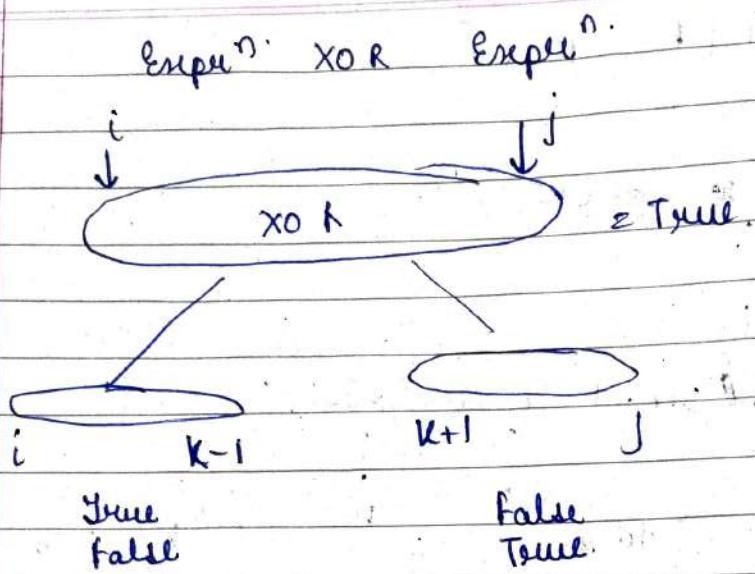
$$\text{no of true ways} = T \text{if } \text{left} + F \text{if } \text{right}$$

$T \wedge F = T$
 $F \wedge T = T$

$\text{no of ways} \rightarrow 2$

$\text{Expr}^n \quad \text{XOR} \quad \text{Expr}^n$

* $\leftarrow \text{no of ways}$
 false



int solve (string s, int i, int j, boolean ^T_{or}_{isFalse})

Base condⁿ

$i \quad j$
T or F and T XOR F
T IF & T AF

boolean isTrue = F

if ($i > j$) return false.

if ($i == j$)

if ($isTrue == \text{True}$)

return $s[i] == 'T'$

else

return $s[i] == 'F'$

loop

$$\begin{array}{c}
 i \quad j \\
 T \mid F \quad \& \quad T \wedge F \\
 \uparrow \quad \uparrow \quad \uparrow \\
 k = i+1 \quad k = k+2 \quad k = j-1
 \end{array}$$

for (int $k = i+1$; $k \leq j-1$; $k = k+2$)

 int ans = 0;

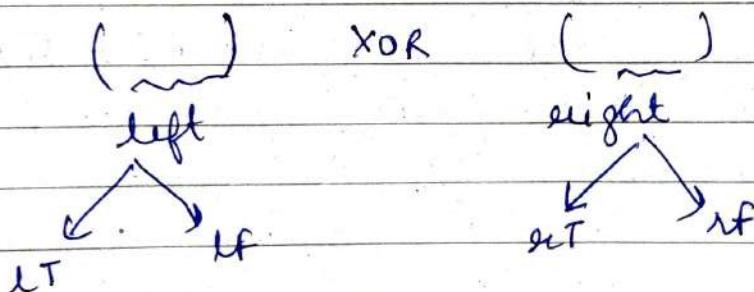
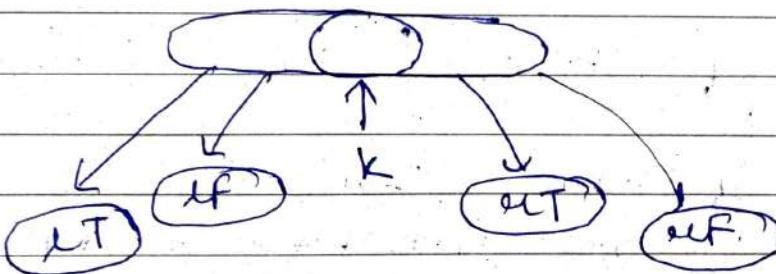
 int LT = (S, i, k-1, T)

 int LF = (S, i, k-1, F)

 int aLT = (S, k+1, j, T)

 int aLF = (S, k+1, j, F)

} temp ans



$$ans = LT * aLF + LF * aLT$$

if ($S[k] == '8'$)

{

if ($iTrue == \text{True}$)

ans = ans + LT * aT

else {

ans = ans + LF * aT + LT * aF + LF * aF;

}

else if ($S[k] == '1'$) {

if ($iTrue == \text{True}$)

ans = ans + LT * aT + LT * aF +
LF * aLT..

else :

ans = ans + LF * aF.

}

else if ($S[k] == '1'$)

{

if ($iTrue == \text{True}$)

ans = ans + LF * aT + LT * aF

else

ans = ans + LT * aT + LF * aF

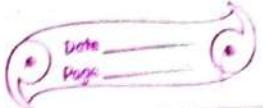
}

return ans;

}

whole code

```
int solve (string s , int i , int j , bool isTrue ) {  
    if (i > j) {  
        return false ; }  
    if (i == j) {  
        if (isTrue == true)  
            return s[i] == 'T' ;  
        else  
            return s[i] == 'F' ;  
    }  
    for ( int k = i+1 ; k <= j-1 ; k+=2) {  
        int ans = 0 ;  
        int LT = solve (s , i , k-1 , T ) ;  
        int LF = solve (s , i , k-1 , F ) ;  
        int RT = solve (s , k+1 , j , T ) ;  
        int RF = solve (s , k+1 , j , F ) ;  
  
        if (s[k] == '&') {  
            if (isTrue == true)  
                ans = ans + LT * RT ;  
            else  
                ans = ans + LF * RT + LT * RF + LF * RF ;  
        }  
        else if ( s[k] == '|') {  
            if (isTrue == true)  
                ans += LT * RT + LT * RF + LF * RT ;  
            else  
                ans += LF * RF ;  
        }  
    }
```

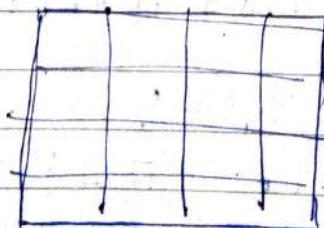


```
else if ( S[k] == 'N' ) {  
    if ( isTrue == True )  
        ans += LF * RT + LT * RF ;  
    else  
        ans += LT * RT + LF * RF ;  
}  
}  
return ans ;  
}
```

Evaluate Expression to True Boolean
Parenthesization - (memoization)
(Bottom Up - DP)
BD dp

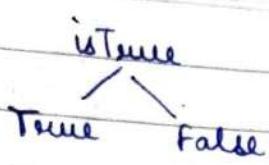
Recursive \rightarrow Recursive call (R.C.)
Top down \rightarrow Table
Bottom up \rightarrow R.C + table

P.S. \rightarrow put brackets in such a way that the expression evaluates to true.

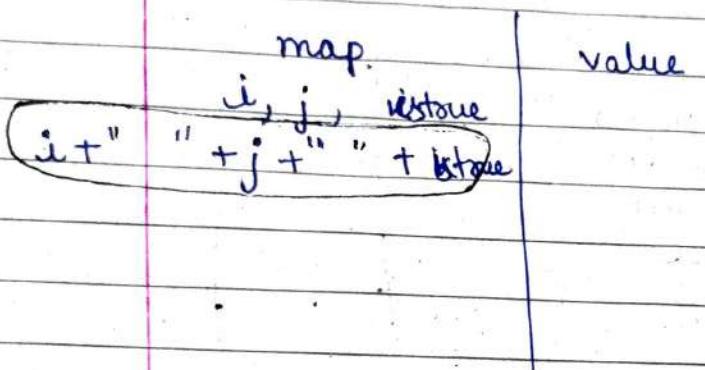


$i * j * \text{isTrue}$
(3d) k

int t [100][100][2]



More better way: we can use map



unordered_map<string, int> mp;

```

int main() {
    mp.clear()
    solve()
}

```

After Base cond:-

```

string temp = to_string(i);
temp.push_back(' ');
temp.append(to_string(j));
temp.push_back(' ');
temp.append(to_string(isTrue));

```

```

if(mp.find(temp) != mp.end()){
    return mp[temp];
}

```

return mp[temp] = ans;

b.

Egg Dropping Problem

I/p : e = 3

f = 5

O/p : 3 → minimize no
of attempts
in worst
case.

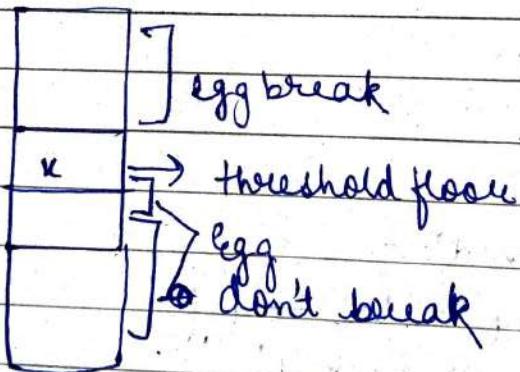
Problem

statement :

| |
|---|
| 5 |
| 4 |
| 3 |
| 2 |
| 1 |

0 0 0

~~~~~  
Eggs.



We are in a building, we need to find the  
minm no of attempts to find the critical  
floor.



|   |
|---|
| 5 |
| 4 |
| 3 |
| 2 |
| 1 |

0 0 0

3 eggs

safe strategy  $\rightarrow$  worst case (1 egg) so drop from the last it won't break until threshold.

|   |
|---|
| 5 |
| 4 |
| 3 |
| 2 |
| 1 |

→ eggbreak (threshold floor)

|   |
|---|
| 7 |
| 6 |
| 5 |
| 4 |
| 3 |
| 2 |
| 1 |

worst-case  $\rightarrow$  minimize no of attempts to find threshold floor.  
 e = 3.

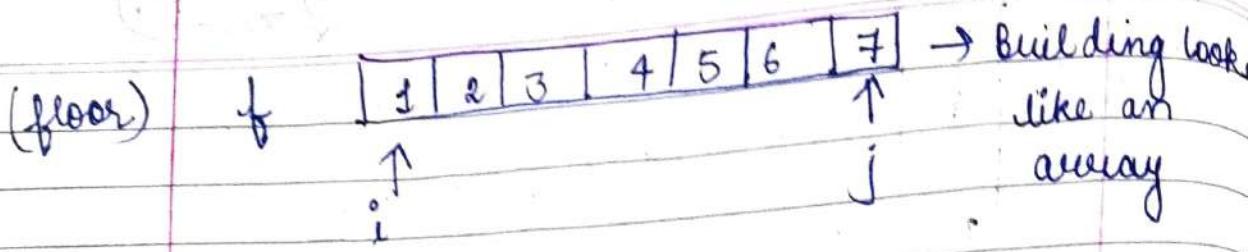
9/10 e = 3

f = 5

0/1 = 3 attempts

|   |
|---|
| 5 |
| 4 |
| 3 |
| 2 |
| 1 |

egg break.



from where to take  $k$ ?

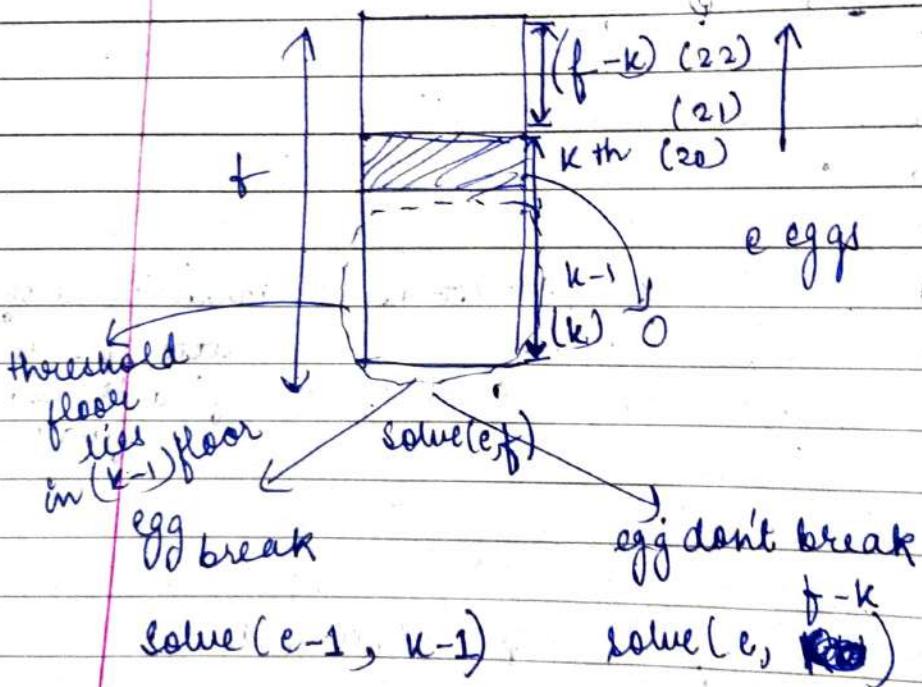
→ We will check for all  $k$  values.

for  $k = 1$        $k = f$        $k++$

base cond<sup>n</sup>. (smallest valid i/p)

$e = 1$  return f

$f = 0/1$  return f



int solve (int e, int f)

8

```
if (f == 0 || f == 1)  
    return f;
```

if ( $e = 1$ )  
return f;

int mn = INT\_MAX;

```
for (int k=1; k<=f; k++)
```

۸

`int temp = 1 + max( solve(e-1, n-1),  
 solve(e, f-n) );`

$$mn = \min(mn, \text{temp});$$

1

return mn;

۲

## Egg Dropping Memoization

$$| \angle = T \angle = 30$$

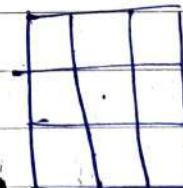
$$1 \leq e \leq 1000$$

$$1 \leq f \leq \$100$$

1) globally  
defined

```
int dp[100][100];
```

```
memset(dp, -1, sizeof(dp));
```



ext

```
int solve ( int e , int f ) {
```

if ( $e == 1$ )

return f;

if ( $f = 0$  ||  $f = -1$ )

return;

```

if ( $t[e][f] \neq -1$ )
    return  $t[e][f]$ ;
int mn = INT_MAX;
for (int k = i; k <= f; k++) {
    int temp =  $\max(\text{solve}(e-1, k-1),$ 
                $\text{solve}(e, f-k)) + 1$ ;
    mn = min(mn, temp);
}
return mn;  $t[e][f] = mn$ ;

```

### Further Optimization

Inside the loop

```

if ( $t[e-1][k-1] \neq -1$ )
    int low =  $t[e-1][k-1]$ 
else
    low = solve(e-1, k-1)
     $t[e-1][k-1] = low$ ;

if ( $t[e][f-k] \neq -1$ )
    int high =  $t[e][f-k]$ ;
else
    high = solve(e, f-k);
     $t[e][f-k] = high$ ;

int temp =  $1 + (\text{low}, \text{high})$ ;

```