

# **DATA STRUCTURES USING “C”**

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**LECTURE NOTES**

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# SYLLABUS

## BE 2106 DATA STRUCTURE (3-0-0)

### **Module – I**

Introduction to data structures: storage structure for arrays, sparse matrices, Stacks and Queues: representation and application. Linked lists: Single linked lists, linked list representation of stacks and Queues. Operations on polynomials, Double linked list, circular list.

### **Module – II**

Dynamic storage management-garbage collection and compaction, infix to post fix conversion, postfix expression evaluation. Trees: Tree terminology, Binary tree, Binary search tree, General tree, B+ tree, AVL Tree, Complete Binary Tree representation, Tree traversals, operation on Binary tree-expression Manipulation.

### **Module –III**

Graphs: Graph terminology, Representation of graphs, path matrix, BFS (breadth first search), DFS (depth first search), topological sorting, Warshall's algorithm (shortest path algorithm.) Sorting and Searching techniques – Bubble sort, selection sort, Insertion sort, Quick sort, merge sort, Heap sort, Radix sort. Linear and binary search methods, Hashing techniques and hash functions.

### **Text Books:**

1. Gilberg and Forouzan: "Data Structure- A Pseudo code approach with C" by Thomson publication
2. "Data structure in C" by Tanenbaum, PHI publication / Pearson publication.
3. Pai: "Data Structures & Algorithms; Concepts, Techniques & Algorithms "Tata McGraw Hill.

### **Reference Books:**

1. "Fundamentals of data structure in C" Horowitz, Sahani & Freed, Computer Science Press.
2. "Fundamental of Data Structure" ( Schaums Series) Tata-McGraw-Hill.

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# **Module-1**

## **Lecture-01**

### **Introduction to Data structures**

In computer terms, a data structure is a Specific way to store and organize data in a computer's memory so that these data can be used efficiently later. Data may be arranged in many different ways such as the logical or mathematical model for a particular organization of data is termed as a data structure. The variety of a particular data model depends on the two factors -

- Firstly, it must be loaded enough in structure to reflect the actual relationships of the data with the real world object.
- Secondly, the formation should be simple enough so that anyone can efficiently process the data each time it is necessary.

### **Categories of Data Structure:**

The data structure can be sub divided into major types:

- Linear Data Structure
- Non-linear Data Structure

### **Linear Data Structure:**

A data structure is said to be linear if its elements combine to form any specific order. There are basically two techniques of representing such linear structure within memory.

- First way is to provide the linear relationships among all the elements represented by means of linear memory location. These linear structures are termed as arrays.
- The second technique is to provide the linear relationship among all the elements represented by using the concept of pointers or links. These linear structures are termed as linked lists.

The common examples of linear data structure are:

- Arrays
- Queues
- Stacks
- Linked lists

### **Non linear Data Structure:**

This structure is mostly used for representing data that contains a hierarchical relationship among various elements.

Examples of Non Linear Data Structures are listed below:

- Graphs
- family of trees and
- table of contents

**Tree:** In this case, data often contain a hierarchical relationship among various elements. The data structure that reflects this relationship is termed as rooted tree graph or a tree.

**Graph:** In this case, data sometimes hold a relationship between the pairs of elements which is not necessarily following the hierarchical structure. Such data structure is termed as a Graph.

**Array** is a container which can hold a fix number of items and these items should be of the same type. Most of the data structures make use of arrays to implement their algorithms. Following are the important terms to understand the concept of Array.

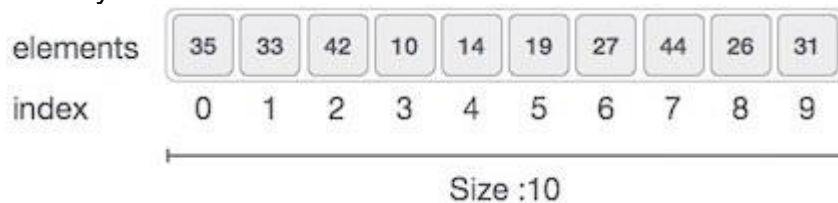
- **Element** – Each item stored in an array is called an element.
- **Index** – Each location of an element in an array has a numerical index, which is used to identify the element.

#### **Array Representation:(Storage structure)**

Arrays can be declared in various ways in different languages. For illustration, let's take C array declaration.



Arrays can be declared in various ways in different languages. For illustration, let's take C array declaration.



As per the above illustration, following are the important points to be considered.

- Index starts with 0.
- Array length is 10 which means it can store 10 elements.
- Each element can be accessed via its index. For example, we can fetch an element at index 6 as 9.

#### **Basic Operations**

Following are the basic operations supported by an array.

- **Traverse** – print all the array elements one by one.
- **Insertion** – Adds an element at the given index.
- **Deletion** – Deletes an element at the given index.
- **Search** – Searches an element using the given index or by the value.
- **Update** – Updates an element at the given index.

In C, when an array is initialized with size, then it assigns defaults values to its elements in following order.

Data Type	Default Value
bool	false

char	0
int	0
float	0.0
double	0.0f
void	
wchar_t	0

### Insertion Operation

Insert operation is to insert one or more data elements into an array. Based on the requirement, a new element can be added at the beginning, end, or any given index of array.

Here, we see a practical implementation of insertion operation, where we add data at the end of the array –

### Algorithm

Let **LA** be a Linear Array (unordered) with **N** elements and **K** is a positive integer such that **K ≤ N**. Following is the algorithm where ITEM is inserted into the **K<sup>th</sup>** position of LA –

1. Start
2. Set J = N
3. Set N = N+1
4. Repeat steps 5 and 6 while J >= K
5. Set LA[J+1] = LA[J]
6. Set J = J-1
7. Set LA[K] = ITEM
8. Stop

### Example

Following is the implementation of the above algorithm –

```
#include <stdio.h>

main() {
    int LA[] = {1,3,5,7,8};
    int item = 10, k = 3, n = 5;
    int i = 0, j = n;
    printf("The original array elements are :\n");
    for(i = 0; i<n; i++) {
        printf("LA[%d] = %d \n", i, LA[i]);
    }
}
```

```

n = n + 1;
while( j >= k) {
    LA[j+1] = LA[j];
    j = j - 1;
}
LA[k] = item;
printf("The array elements after insertion :\n");
for(i = 0; i<n; i++) {
    printf("LA[%d] = %d \n", i, LA[i]);
}
}

```

When we compile and execute the above program, it produces the following result –  
**Output**

```

The original array elements are :
LA[0] = 1
LA[1] = 3
LA[2] = 5
LA[3] = 7
LA[4] = 8
The array elements after insertion :
LA[0] = 1
LA[1] = 3
LA[2] = 5
LA[3] = 10
LA[4] = 7
LA[5] = 8

```

### Deletion Operation

Deletion refers to removing an existing element from the array and re-organizing all elements of an array.

### Algorithm

Consider **LA** is a linear array with **N** elements and **K** is a positive integer such that **K ≤ N**. Following is the algorithm to delete an element available at the **K<sup>th</sup>** position of LA.

1. Start
2. Set J = K
3. Repeat steps 4 and 5 while J < N
4. Set LA[J] = LA[J + 1]
5. Set J = J+1
6. Set N = N-1
7. Stop

### Example



Following is the implementation of the above algorithm –

```
#include <stdio.h>

void main() {
    int LA[] = {1,3,5,7,8};
    int k = 3, n = 5;
    int i, j;
    printf("The original array elements are :\n");
    for(i = 0; i<n; i++) {
        printf("LA[%d] = %d \n", i, LA[i]);
    }

    j = k;
    while( j < n) {
        LA[j-1] = LA[j];
        j = j + 1;
    }
    n = n -1;
    printf("The array elements after deletion :\n");
    for(i = 0; i<n; i++) {
        printf("LA[%d] = %d \n", i, LA[i]);
    }
}
```

When we compile and execute the above program, it produces the following result –

### **Output**

```
The original array elements are :
LA[0] = 1
LA[1] = 3
LA[2] = 5
LA[3] = 7
LA[4] = 8
The array elements after deletion :
LA[0] = 1
LA[1] = 3
LA[2] = 7
LA[3] = 8
```

## Lecture-02

### Search Operation

You can perform a search for an array element based on its value or its index.

### Algorithm

Consider **LA** is a linear array with **N** elements and **K** is a positive integer such that **K ≤ N**. Following is the algorithm to find an element with a value of **ITEM** using sequential search.

1. Start
2. Set J = 0
3. Repeat steps 4 and 5 while J < N
4. IF LA[J] is equal ITEM THEN GOTO STEP 6
5. Set J = J + 1
6. PRINT J, ITEM
7. Stop

### Example

Following is the implementation of the above algorithm –

```
#include <stdio.h>

void main() {
    int LA[] = {1,3,5,7,8};
    int item = 5, n = 5;
    int i = 0, j = 0;
    printf("The original array elements are :\n");
    for(i = 0; i < n; i++) {
        printf("LA[%d] = %d \n", i, LA[i]);
    }
    while( j < n){
        if( LA[j] == item ) {
            break;
        }
        j = j + 1;
    }
    printf("Found element %d at position %d\n", item, j+1);
}
```

When we compile and execute the above program, it produces the following result –

### Output

```
The original array elements are :
LA[0] = 1
LA[1] = 3
LA[2] = 5
```

LA[3] = 7  
LA[4] = 8  
Found element 5 at position 3

### Update Operation

Update operation refers to updating an existing element from the array at a given index.

### Algorithm

Consider **LA** is a linear array with **N** elements and **K** is a positive integer such that **K ≤ N**. Following is the algorithm to update an element available at the **K<sup>th</sup>** position of LA.

1. Start
2. Set LA[K-1] = ITEM
3. Stop

### Example

Following is the implementation of the above algorithm –

```
#include <stdio.h>

void main() {
    int LA[] = {1,3,5,7,8};
    int k = 3, n = 5, item = 10;
    int i, j;
    printf("The original array elements are :\n");
    for(i = 0; i < n; i++) {
        printf("LA[%d] = %d \n", i, LA[i]);
    }

    LA[k-1] = item;
    printf("The array elements after updation :\n");
    for(i = 0; i < n; i++) {
        printf("LA[%d] = %d \n", i, LA[i]);
    }
}
```

When we compile and execute the above program, it produces the following result –

### Output

The original array elements are :  
LA[0] = 1  
LA[1] = 3  
LA[2] = 5  
LA[3] = 7  
LA[4] = 8

The array elements after updation :

LA[0] = 1

LA[1] = 3

LA[2] = 10

LA[3] = 7

LA[4] = 8

## Lecture-03

### **Sparse Matrix and its representations**

A [matrix](#) is a two-dimensional data object made of m rows and n columns, therefore having total m x n values. If most of the elements of the matrix have **0 value**, then it is called a sparse matrix.

#### **Why to use Sparse Matrix instead of simple matrix ?**

- **Storage:** There are lesser non-zero elements than zeros and thus lesser memory can be used to store only those elements.
- **Computing time:** Computing time can be saved by logically designing a data structure traversing only non-zero elements..

Example:

0 0 3 0 4

0 0 5 7 0

0 0 0 0 0

0 2 6 0 0

Representing a sparse matrix by a 2D array leads to wastage of lots of memory as zeroes in the matrix are of no use in most of the cases. So, instead of storing zeroes with non-zero elements, we only store non-zero elements. This means storing non-zero elements with **triples- (Row, Column, value)**.

Sparse Matrix Representations can be done in many ways following are two common representations:

1. Array representation
2. Linked list representation

#### **Method 1: Using Arrays**

```
#include<stdio.h>
int main()
{
    // Assume 4x5 sparse matrix
    int sparseMatrix[4][5] =
    {
        {0 , 0 , 3 , 0 , 4 },
        {0 , 0 , 5 , 7 , 0 },
        {0 , 0 , 0 , 0 , 0 },
        {0 , 2 , 6 , 0 , 0 }
    };

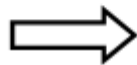
    int size = 0;
    for (int i = 0; i < 4; i++)
        for (int j = 0; j < 5; j++)
            if (sparseMatrix[i][j] != 0)
                size++;
    int compactMatrix[3][size];
    // Making of new matrix
```

```

int k = 0;
for (int i = 0; i < 4; i++)
    for (int j = 0; j < 5; j++)
        if (sparseMatrix[i][j] != 0)
        {
            compactMatrix[0][k] = i;
            compactMatrix[1][k] = j;
            compactMatrix[2][k] = sparseMatrix[i][j];
            k++;
        }
for (int i=0; i<3; i++)
{
    for (int j=0; j<size; j++)
        printf("%d ", compactMatrix[i][j]);
    printf("\n");
}
return 0;
}

```

0	0	3	0	4
0	0	5	7	0
0	0	0	0	0
0	2	6	0	0



Row	0	0	1	1	3	3
Column	2	4	2	3	1	2
Value	3	4	5	7	2	6

## Lecture-04

### STACK

A stack is an Abstract Data Type (ADT), commonly used in most programming languages. It is named stack as it behaves like a real-world stack, for example – a deck of cards or a pile of plates, etc.

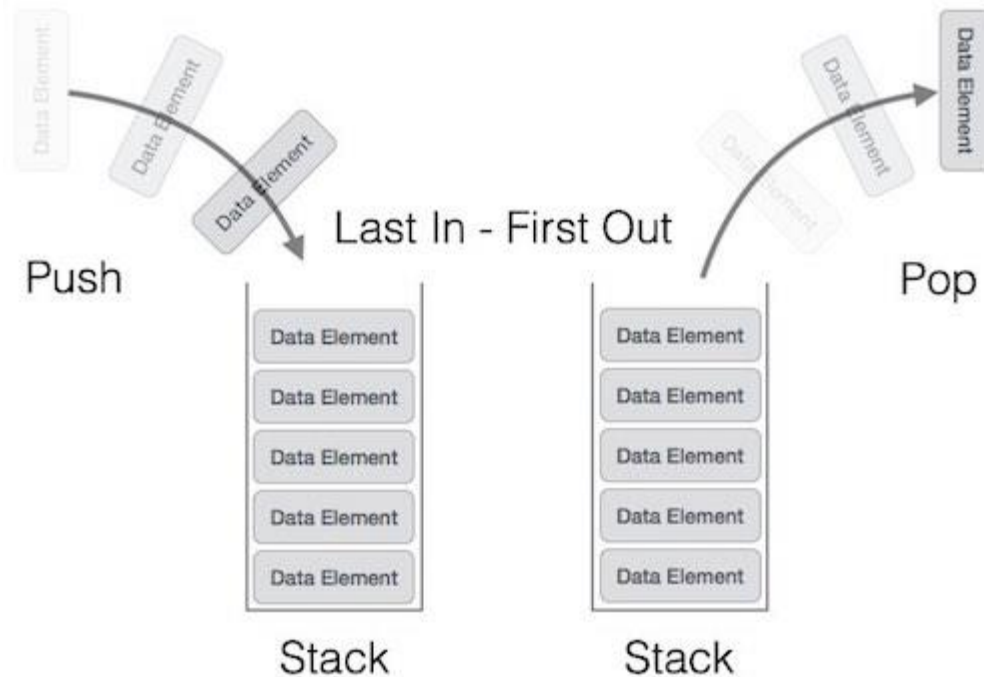


A real-world stack allows operations at one end only. For example, we can place or remove a card or plate from the top of the stack only. Likewise, Stack ADT allows all data operations at one end only. At any given time, we can only access the top element of a stack.

This feature makes it LIFO data structure. LIFO stands for Last-in-first-out. Here, the element which is placed (inserted or added) last, is accessed first. In stack terminology, insertion operation is called **PUSH** operation and removal operation is called **POP** operation.

### Stack Representation

The following diagram depicts a stack and its operations –



A stack can be implemented by means of Array, Structure, Pointer, and Linked List. Stack can either be a fixed size one or it may have a sense of dynamic resizing. Here, we are going to implement stack using arrays, which makes it a fixed size stack implementation.

### Basic Operations

Stack operations may involve initializing the stack, using it and then de-initializing it. Apart from these basic stuffs, a stack is used for the following two primary operations –

- **push()** – Pushing (storing) an element on the stack.

- **pop()** – Removing (accessing) an element from the stack.

When data is PUSHed onto stack.

To use a stack efficiently, we need to check the status of stack as well. For the same purpose, the following functionality is added to stacks –

- **peek()** – get the top data element of the stack, without removing it.
- **isFull()** – check if stack is full.
- **isEmpty()** – check if stack is empty.

At all times, we maintain a pointer to the last PUSHed data on the stack. As this pointer always represents the top of the stack, hence named **top**. The **top** pointer provides top value of the stack without actually removing it.

First we should learn about procedures to support stack functions –

### **peek()**

Algorithm of peek() function –

```
begin procedure peek
    return stack[top]
end procedure
```

Implementation of peek() function in C programming language –

### **Example**

```
int peek() {
    return stack[top];
}
```

### **isfull()**

Algorithm of isfull() function –

```
begin procedure isfull

    if top equals to MAXSIZE
        return true
    else
        return false
    endif

end procedure
```

Implementation of isfull() function in C programming language –

### **Example**

```
bool isfull() {
    if(top == MAXSIZE)
        return true;
    else
        return false;
}
```



## isempty()

Algorithm of isempty() function –

```
begin procedure isempty
```

```
    if top less than 1
        return true
    else
        return false
    endif
```

```
end procedure
```

Implementation of isempty() function in C programming language is slightly different. We initialize top at -1, as the index in array starts from 0. So we check if the top is below zero or -1 to determine if the stack is empty. Here's the code –

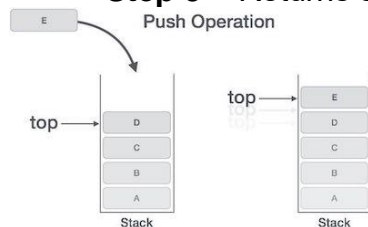
### Example

```
bool isempty() {
    if(top == -1)
        return true;
    else
        return false;
}
```

## Push Operation

The process of putting a new data element onto stack is known as a Push Operation. Push operation involves a series of steps –

- **Step 1** – Checks if the stack is full.
- **Step 2** – If the stack is full, produces an error and exit.
- **Step 3** – If the stack is not full, increments **top** to point next empty space.
- **Step 4** – Adds data element to the stack location, where top is pointing.
- **Step 5** – Returns success.



If the linked list is used to implement the stack, then in step 3, we need to allocate space dynamically.

Algorithm for PUSH Operation

A simple algorithm for Push operation can be derived as follows –

```
begin procedure push: stack, data
```

```
    if stack is full
```

```

    return null
endif

top ← top + 1
stack[top] ← data

```

end procedure

Implementation of this algorithm in C, is very easy. See the following code –

### Example

```

void push(int data) {
    if(!isFull()) {
        top = top + 1;
        stack[top] = data;
    } else {
        printf("Could not insert data, Stack is full.\n");
    }
}

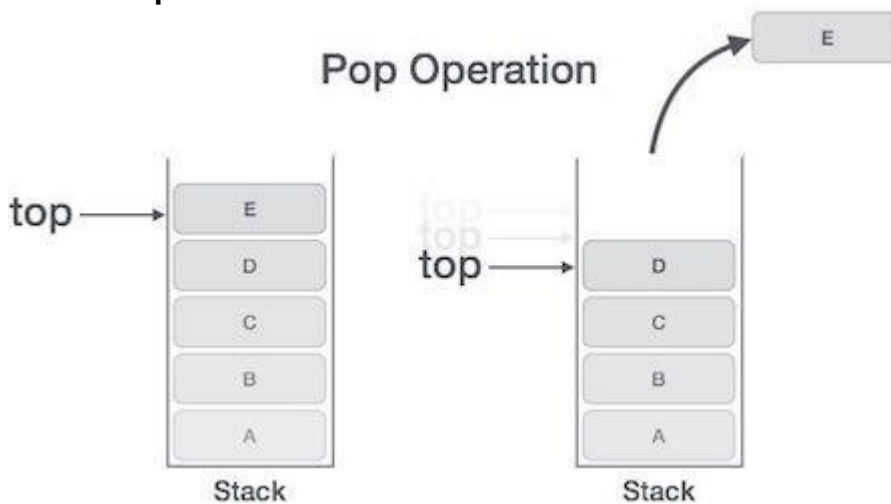
```

### Pop Operation

Accessing the content while removing it from the stack, is known as a Pop Operation. In an array implementation of pop() operation, the data element is not actually removed, instead **top** is decremented to a lower position in the stack to point to the next value. But in linked-list implementation, pop() actually removes data element and deallocates memory space.

A Pop operation may involve the following steps –

- **Step 1** – Checks if the stack is empty.
- **Step 2** – If the stack is empty, produces an error and exit.
- **Step 3** – If the stack is not empty, accesses the data element at which **top** is pointing.
- **Step 4** – Decreases the value of top by 1.
- **Step 5** – Returns success.



Algorithm for Pop Operation

A simple algorithm for Pop operation can be derived as follows –

```
begin procedure pop: stack
```

```
    if stack is empty  
        return null  
    endif
```

```
    data ← stack[top]  
    top ← top - 1  
    return data
```

```
end procedure
```

Implementation of this algorithm in C, is as follows –

### **Example**

```
int pop(int data) {  
  
    if(!isempty()) {  
        data = stack[top];  
        top = top - 1;  
        return data;  
    } else {  
        printf("Could not retrieve data, Stack is empty.\n");  
    }  
}
```

## Lecture-05

### Stack Applications

Three applications of stacks are presented here. These examples are central to many activities that a computer must do and deserve time spent with them.

1. Expression evaluation
2. Backtracking (game playing, finding paths, exhaustive searching)
3. Memory management, run-time environment for nested language features.

### Expression evaluation

In particular we will consider arithmetic expressions. Understand that there are boolean and logical expressions that can be evaluated in the same way. Control structures can also be treated similarly in a compiler.

This study of arithmetic expression evaluation is an example of problem solving where you solve a simpler problem and then transform the actual problem to the simpler one.

Aside: *The NP-Complete problem*. There are a set of apparently intractable problems: finding the shortest route in a graph (Traveling Salesman Problem), bin packing, linear programming, etc. that are similar enough that if a polynomial solution is ever found (exponential solutions abound) for one of these problems, then the solution can be applied to all problems.

### Infix, Prefix and Postfix Notation

We are accustomed to write arithmetic expressions with the operation between the two operands:  $a+b$  or  $c/d$ . If we write  $a+b*c$ , however, we have to apply precedence rules to avoid the ambiguous evaluation (add first or multiply first?).

There's no real reason to put the operation between the variables or values. They can just as well precede or follow the operands. You should note the advantage of prefix and postfix: the need for precedence rules and parentheses are eliminated.

Infix	Prefix	Postfix
$a + b$	$+ a b$	$a b +$
$a + b * c$	$+ a * b c$	$a b c * +$
$(a + b) * (c - d)$	$* + a b - c d$	$a b + c d - *$
$b * b - 4 * a * c$		
$40 - 3 * 5 + 1$		

Postfix expressions are easily evaluated with the aid of a stack.

---

### Infix, Prefix and Postfix Notation KEY

Infix	Prefix	Postfix
$a + b$	$+ a b$	$a b +$

$a + b * c$	$+ a * b c$	$a b c * +$
$(a + b) * (c - d)$	$* + a b - c d$	$a b + c d - *$
$b * b - 4 * a * c$	$- * b b * * 4 a c$	$b b * 4 a * c * -$
$40 - 3 * 5 + 1 = 26$	$+ - 40 * 3 5 1$	$40 3 5 * - 1 +$

### Postfix Evaluation Algorithm

Assume we have a string of operands and operators, an informal, by hand process is

1. Scan the expression left to right
2. Skip values or variables (operands)
3. When an operator is found, apply the operation to the preceding two operands
4. Replace the two operands and operator with the calculated value (three symbols are replaced with one operand)
5. Continue scanning until only a value remains--the result of the expression

The time complexity is  $O(n)$  because each operand is scanned once, and each operation is performed once.

A more formal algorithm:

create a new stack

```
while(input stream is not empty){
    token = getNextToken();
    if(token instanceof operand){
        push(token);
    } else if (token instanceof operator)
        op2 = pop();
        op1 = pop();
        result = calc(token, op1, op2);
        push(result);
    }
}
```

return pop();

Demonstration with  $2\ 3\ 4\ +\ *\ 5\ -$

### Infix transformation to Postfix

This process uses a stack as well. We have to hold information that's expressed inside parentheses while scanning to find the closing ')'. We also have to hold information on operations that are of lower precedence on the stack. The algorithm is:

1. Create an empty stack and an empty postfix output string/stream
2. Scan the infix input string/stream left to right
3. If the current input token is an operand, simply append it to the output string (note the examples above that the operands remain in the same order)
4. If the current input token is an operator, pop off all operators that have equal or higher

precedence and append them to the output string; push the operator onto the stack. The order of popping is the order in the output.

5. If the current input token is '(', push it onto the stack
6. If the current input token is ')', pop off all operators and append them to the output string until a '(' is popped; discard the '('.
7. If the end of the input string is found, pop all operators and append them to the output string.

This algorithm doesn't handle errors in the input, although careful analysis of parenthesis or lack of parenthesis could point to such error determination.

Apply the algorithm to the above expressions.

### **Backtracking**

Backtracking is used in algorithms in which there are steps along some path (state) from some starting point to some goal.

- Find your way through a maze.
- Find a path from one point in a graph (roadmap) to another point.
- Play a game in which there are moves to be made (checkers, chess).

In all of these cases, there are choices to be made among a number of options. We need some way to remember these decision points in case we want/need to come back and try the alternative

Consider the maze. At a point where a choice is made, we may discover that the choice leads to a dead-end. We want to retrace back to that decision point and then try the other (next) alternative.

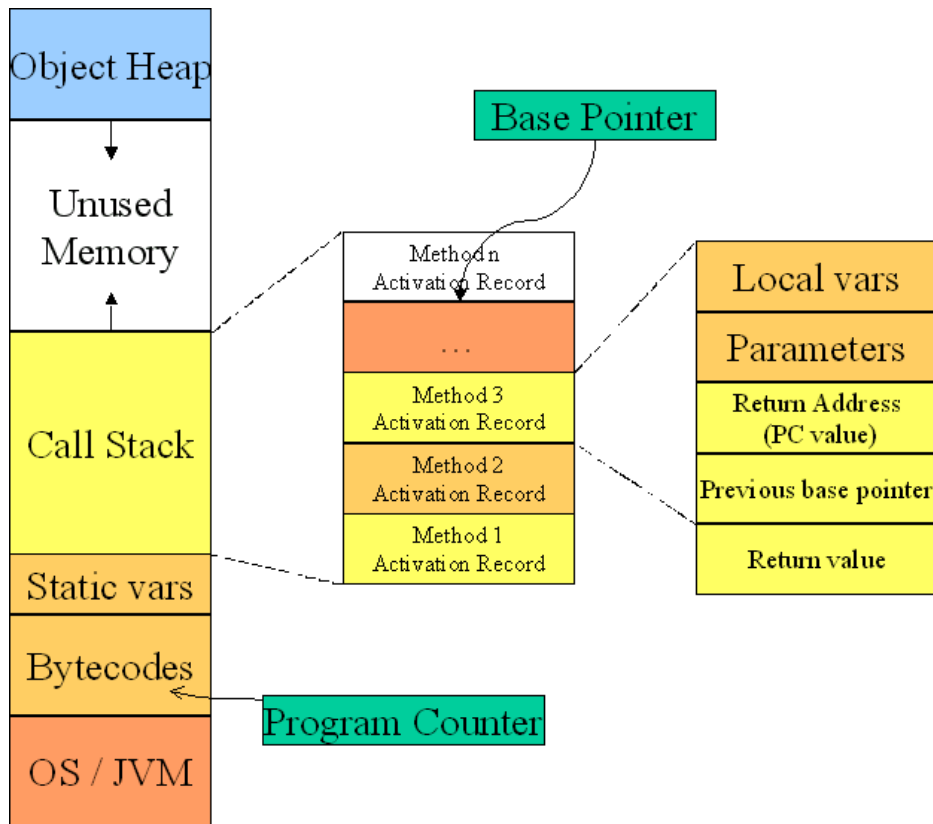
Again, stacks can be used as part of the solution. Recursion is another, typically more favored, solution, which is actually implemented by a stack.

### **Memory Management**

Any modern computer environment uses a stack as the primary memory management model for a running program. Whether it's native code (x86, Sun, VAX) or JVM, a stack is at the center of the run-time environment for Java, C++, Ada, FORTRAN, etc.

The discussion of JVM in the text is consistent with NT, Solaris, VMS, Unix runtime environments.

Each program that is running in a computer system has its own memory allocation containing the typical layout as shown below.



### Call and return process

When a method/function is called

1. An activation record is created; its size depends on the number and size of the local variables and parameters.
2. The Base Pointer value is saved in the special location reserved for it
3. The Program Counter value is saved in the Return Address location
4. The Base Pointer is now reset to the new base (top of the call stack prior to the creation of the AR)
5. The Program Counter is set to the location of the first bytecode of the method being called
6. Copies the calling parameters into the Parameter region
7. Initializes local variables in the local variable region

While the method executes, the local variables and parameters are simply found by adding a constant associated with each variable/parameter to the Base Pointer.

When a method returns

1. Get the program counter from the activation record and replace what's in the PC
2. Get the base pointer value from the AR and replace what's in the BP
3. Pop the AR entirely from the stack.

## Lecture-06

### QUEUE

Queue is an abstract data structure, somewhat similar to Stacks. Unlike stacks, a queue is open at both its ends. One end is always used to insert data (enqueue) and the other is used to remove data (dequeue). Queue follows First-In-First-Out methodology, i.e., the data item stored first will be accessed first.



A real-world example of queue can be a single-lane one-way road, where the vehicle enters first, exits first. More real-world examples can be seen as queues at the ticket windows and bus-stops.

### Queue Representation

As we now understand that in queue, we access both ends for different reasons. The following diagram given below tries to explain queue representation as data structure –



As in stacks, a queue can also be implemented using Arrays, Linked-lists, Pointers and Structures. For the sake of simplicity, we shall implement queues using one-dimensional array.

### Basic Operations

Queue operations may involve initializing or defining the queue, utilizing it, and then completely erasing it from the memory. Here we shall try to understand the basic operations associated with queues –

- **enqueue()** – add (store) an item to the queue.
- **dequeue()** – remove (access) an item from the queue.

Few more functions are required to make the above-mentioned queue operation efficient. These are –

- **peek()** – Gets the element at the front of the queue without removing it.
- **isfull()** – Checks if the queue is full.
- **isempty()** – Checks if the queue is empty.

In queue, we always dequeue (or access) data, pointed by **front** pointer and while enqueueing (or storing) data in the queue we take help of **rear** pointer.

Let's first learn about supportive functions of a queue –

### **peek()**



This function helps to see the data at the **front** of the queue. The algorithm of peek() function is as follows –

#### Algorithm

```
begin procedure peek
    return queue[front]
end procedure
```

Implementation of peek() function in C programming language –

#### Example

```
int peek() {
    return queue[front];
}
```

#### isfull()

As we are using single dimension array to implement queue, we just check for the rear pointer to reach at MAXSIZE to determine that the queue is full. In case we maintain the queue in a circular linked-list, the algorithm will differ. Algorithm of isfull() function –

#### Algorithm

```
begin procedure isfull

    if rear equals to MAXSIZE
        return true
    else
        return false
    endif

end procedure
```

Implementation of isfull() function in C programming language –

#### Example

```
bool isfull() {
    if(rear == MAXSIZE - 1)
        return true;
    else
        return false;
}
```

#### isempty()

Algorithm of isempty() function –

#### Algorithm

```
begin procedure isempty

    if front is less than MIN OR front is greater than rear
        return true
    endif

end procedure
```

```

else
    return false
endif

```

end procedure

If the value of **front** is less than MIN or 0, it tells that the queue is not yet initialized, hence empty.

Here's the C programming code –

### Example

```

bool isempty() {
    if(front < 0 || front > rear)
        return true;
    else
        return false;
}

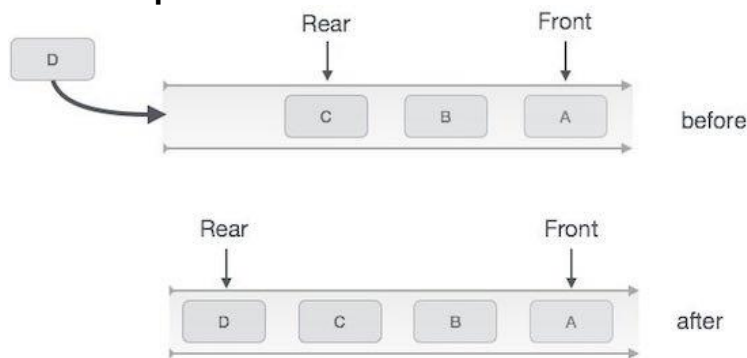
```

### Enqueue Operation

Queues maintain two data pointers, **front** and **rear**. Therefore, its operations are comparatively difficult to implement than that of stacks.

The following steps should be taken to enqueue (insert) data into a queue –

- **Step 1** – Check if the queue is full.
- **Step 2** – If the queue is full, produce overflow error and exit.
- **Step 3** – If the queue is not full, increment **rear** pointer to point the next empty space.
- **Step 4** – Add data element to the queue location, where the rear is pointing.
- **Step 5** – return success.



### Queue Enqueue

Sometimes, we also check to see if a queue is initialized or not, to handle any unforeseen situations.

Algorithm for enqueue operation

```

procedure enqueue(data)

```

```

    if queue is full
        return overflow

```

```
endif
```

```
rear ← rear + 1
```

```
queue[rear] ← data
```

```
return true
```

```
end procedure
```

Implementation of enqueue() in C programming language –

### Example

```
int enqueue(int data)
```

```
if(isfull())
```

```
return 0;
```

```
rear = rear + 1;
```

```
queue[rear] = data;
```

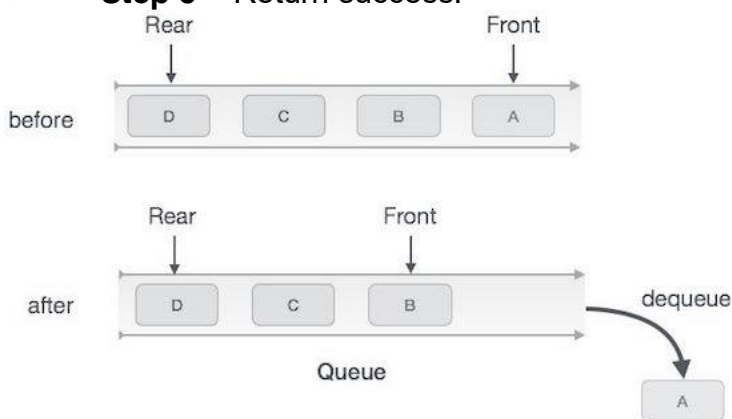
```
return 1;
```

```
end procedure
```

### Dequeue Operation

Accessing data from the queue is a process of two tasks – access the data where **front** is pointing and remove the data after access. The following steps are taken to perform **dequeue** operation –

- **Step 1** – Check if the queue is empty.
- **Step 2** – If the queue is empty, produce underflow error and exit.
- **Step 3** – If the queue is not empty, access the data where **front** is pointing.
- **Step 4** – Increment **front** pointer to point to the next available data element.
- **Step 5** – Return success.



### Queue Dequeue

Algorithm for dequeue operation

```
procedure dequeue
```

```
if queue is empty
    return underflow
end if
```

```
data = queue[front]
front ← front + 1
return true
```

```
end procedure
```

Implementation of dequeue() in C programming language –

### **Example**

```
int dequeue() {
    if(isempty())
        return 0;

    int data = queue[front];
    front = front + 1;

    return data;
}
```

## Lecture-07

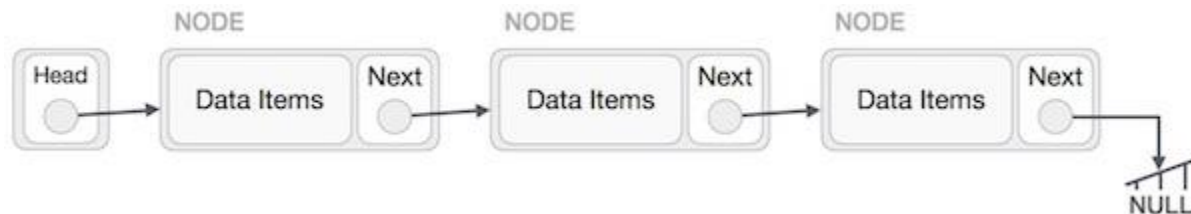
### LINKED LIST

A linked list is a sequence of data structures, which are connected together via links. Linked List is a sequence of links which contains items. Each link contains a connection to another link. Linked list is the second most-used data structure after array. Following are the important terms to understand the concept of Linked List.

- **Link** – Each link of a linked list can store a data called an element.
- **Next** – Each link of a linked list contains a link to the next link called Next.
- **LinkedList** – A Linked List contains the connection link to the first link called First.

#### Linked List Representation

Linked list can be visualized as a chain of nodes, where every node points to the next node.



As per the above illustration, following are the important points to be considered.

- Linked List contains a link element called first.
- Each link carries a data field(s) and a link field called next.
- Each link is linked with its next link using its next link.
- Last link carries a link as null to mark the end of the list.

#### Types of Linked List

Following are the various types of linked list.

- **Simple Linked List** – Item navigation is forward only.
- **Doubly Linked List** – Items can be navigated forward and backward.
- **Circular Linked List** – Last item contains link of the first element as next and the first element has a link to the last element as previous.

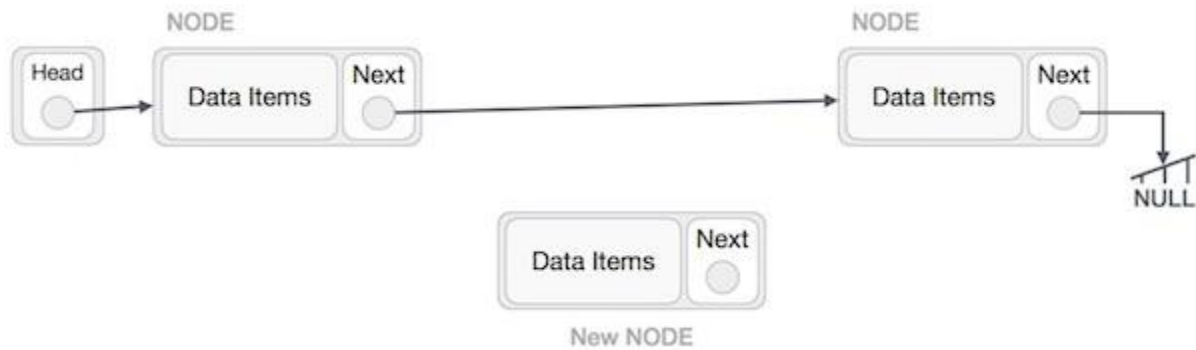
#### Basic Operations

Following are the basic operations supported by a list.

- **Insertion** – Adds an element at the beginning of the list.
- **Deletion** – Deletes an element at the beginning of the list.
- **Display** – Displays the complete list.
- **Search** – Searches an element using the given key.
- **Delete** – Deletes an element using the given key.

#### Insertion Operation

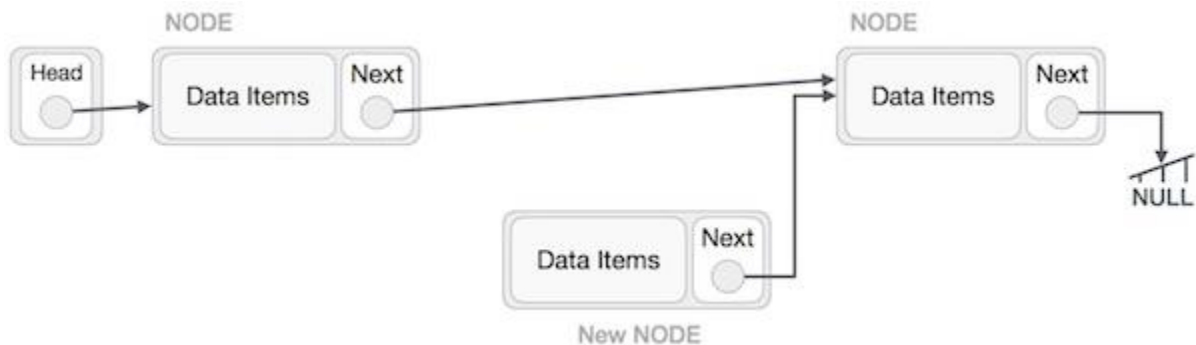
Adding a new node in linked list is a more than one step activity. We shall learn this with diagrams here. First, create a node using the same structure and find the location where it has to be inserted.



Imagine that we are inserting a node **B** (NewNode), between **A** (LeftNode) and **C** (RightNode). Then point B.next to C –

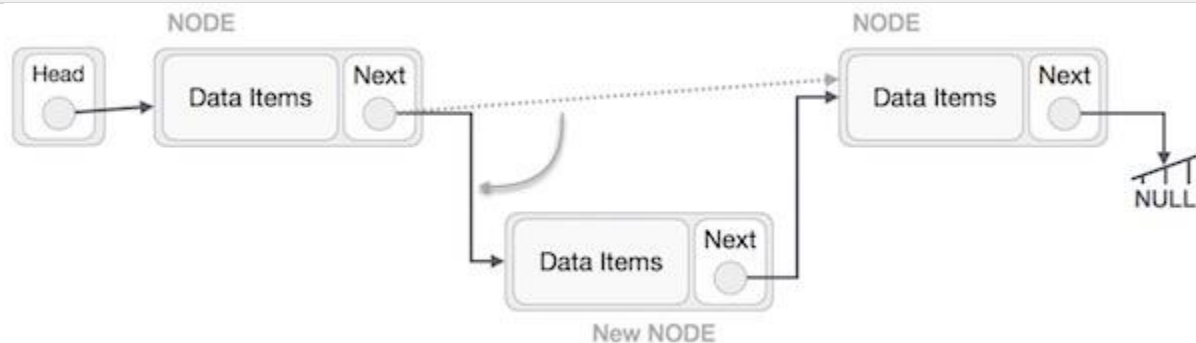
NewNode.next -> RightNode;

It should look like this –



Now, the next node at the left should point to the new node.

LeftNode.next -> NewNode;



This will put the new node in the middle of the two. The new list should look like this –



Similar steps should be taken if the node is being inserted at the beginning of the list. While inserting it at the end, the second last node of the list should point to the new node and the new node will point to NULL.