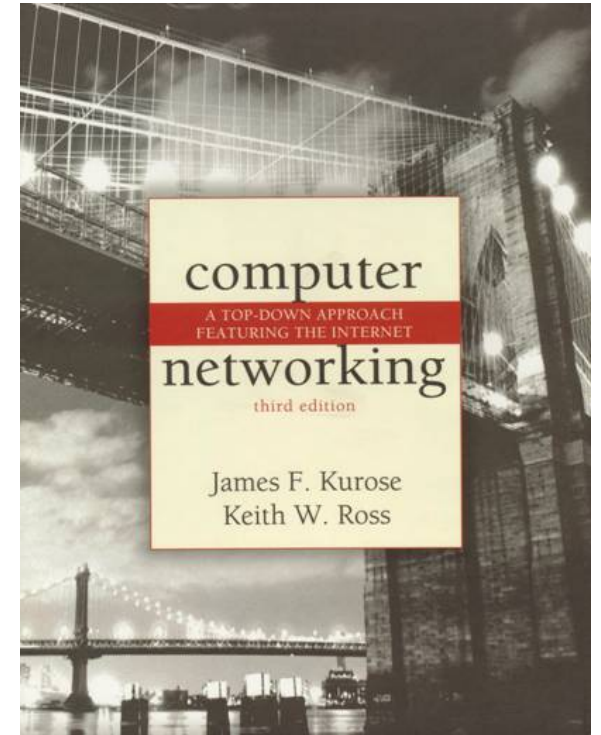


Switching and Routing



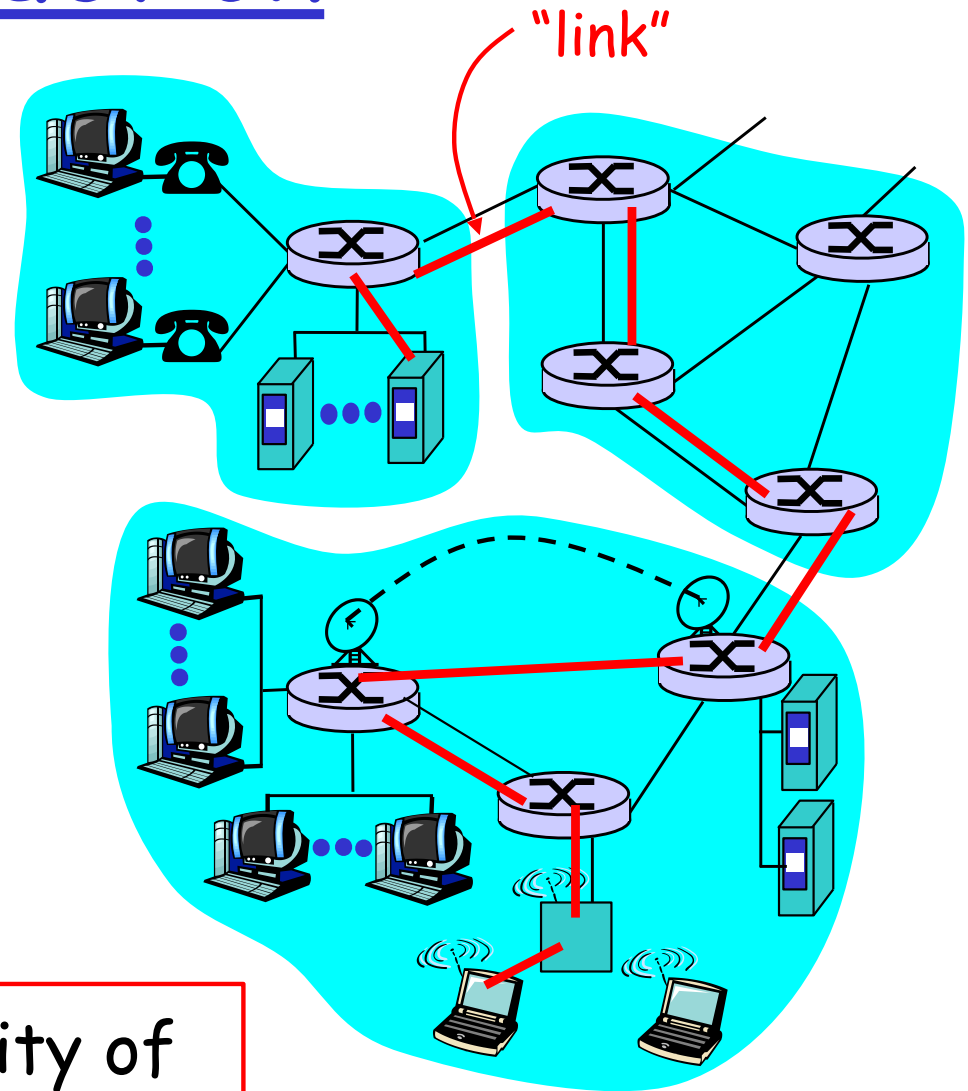
Computer Networking:
A Top Down Approach
Featuring the Internet,
Jim Kurose, Keith Ross
Addison-Wesley

Link Layer: Introduction

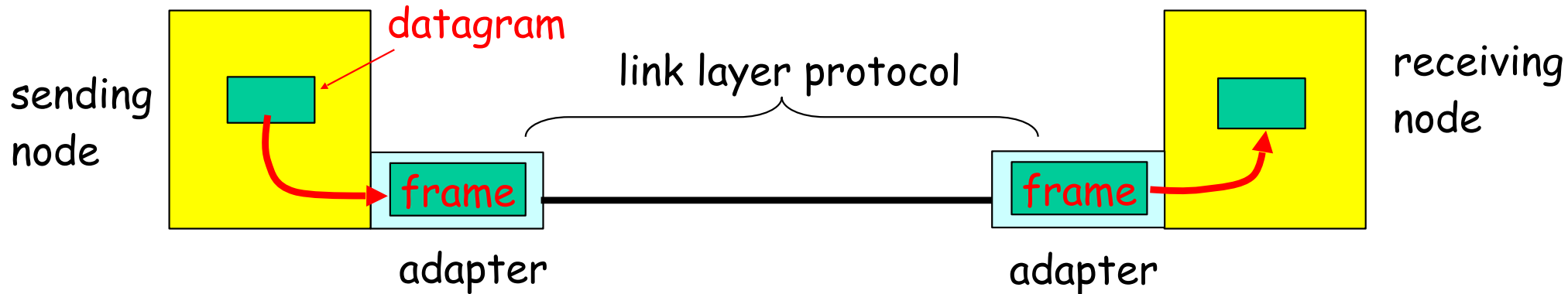
Some terminology:

- ❑ hosts and routers are **nodes**
- ❑ communication channels that connect adjacent nodes along communication path are **links**
 - wired links
 - wireless links
 - LANs
- ❑ layer-2 packet is a **frame**, encapsulates datagram

data-link layer has responsibility of transferring datagram from one node to adjacent node over a link



Adaptors Communicating



- ❑ link layer implemented in “adaptor” (aka NIC)
 - Ethernet card, PCMCIA card, 802.11 card
- ❑ sending side:
 - encapsulates datagram in a frame
 - adds error checking bits, rdt, flow control, etc.

- ❑ receiving side
 - looks for errors, rdt, flow control, etc
 - extracts datagram, passes to rcving node
- ❑ adapter is semi-autonomous
- ❑ link & physical layers

MAC Addresses

❑ IP address:

- network-layer address
- used to get datagram to destination IP subnet

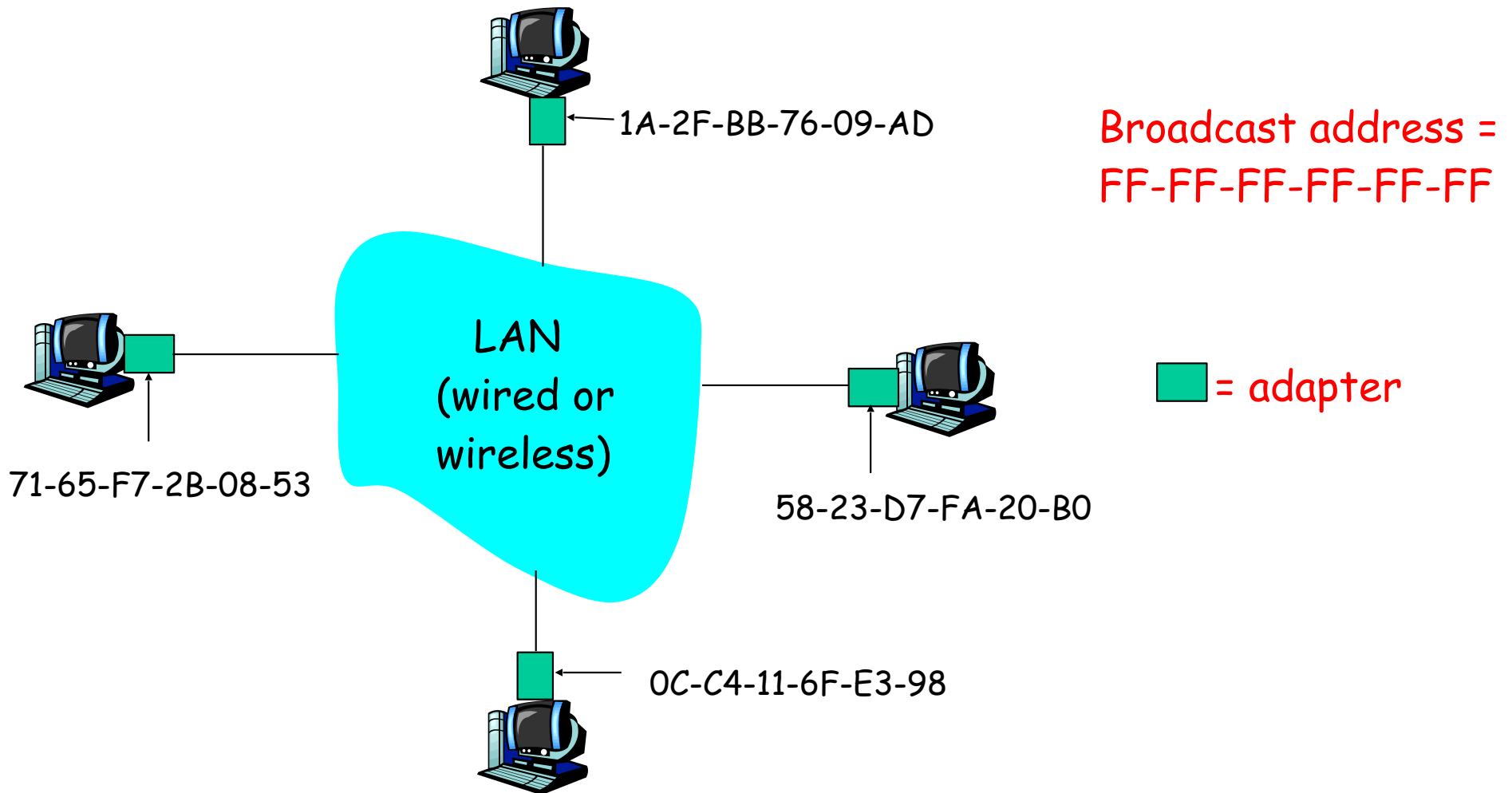
❑ MAC (or LAN or physical or Ethernet) address:

❑ MAC = Media Access Control

- used to get datagram from one interface to another physically-connected interface (same network)
- 48 bit MAC address (for most LANs)
burned in the adapter ROM

MAC Addresses(more)

Each adapter on LAN has unique 6-byte MAC address



MAC Address (more)

- ❑ MAC address allocation administered by IEEE
- ❑ manufacturer buys portion of MAC address space (to assure uniqueness)
- ❑ Analogy:
 - (a) MAC address: like Social Security Number
 - (b) IP address: like postal address
- ❑ MAC flat address → portability
 - can move LAN card from one LAN to another
- ❑ IP hierarchical address NOT portable
 - depends on IP subnet to which node is attached

ARP: Address Resolution Protocol

Question: how to determine MAC address of a node knowing its IP address?

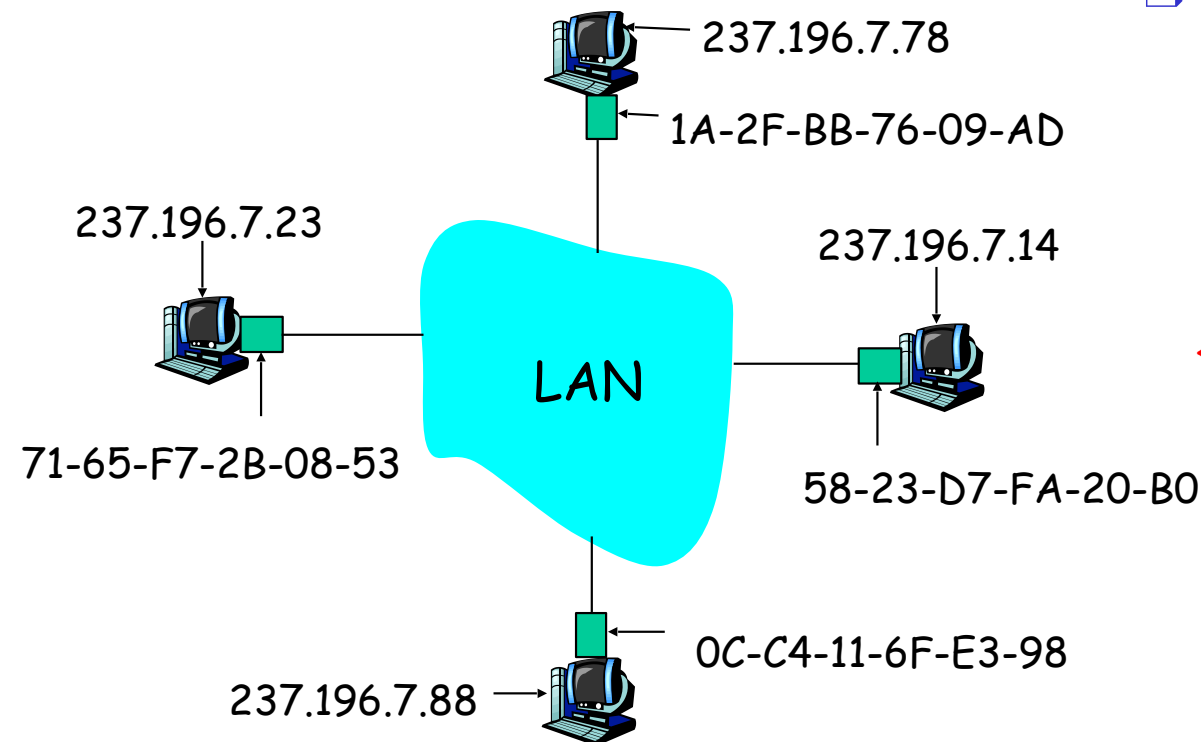
- Each IP node (Host, Router) on LAN has **ARP** table

- ARP Table

- IP-MAC address mappings for some LAN nodes

< IP address; MAC address; TTL >

- TTL (Time To Live): time after which address mapping will be forgotten (typically 20 min)



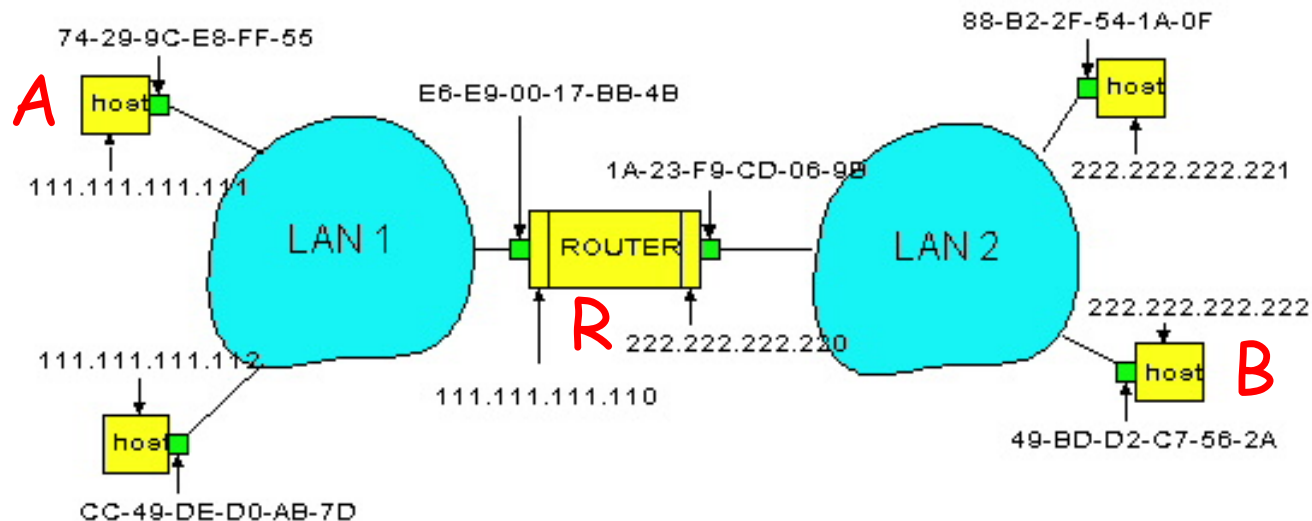
ARP protocol in the Same LAN

- ❑ A wants to send datagram to B, and B's MAC address not in A's ARP table.
- ❑ A **broadcasts** ARP query packet, containing B's IP address
 - "Who has B's MAC address?"
 - Dest MAC address = FF-FF-FF-FF-FF-FF
 - all machines on LAN receive ARP query
- ❑ B receives ARP packet, replies to A with its (B's) MAC address
 - frame sent to A's MAC address (unicast)
- ❑ A caches (saves) IP-to-MAC address pair in its ARP table until information becomes old (times out)
 - soft state: information that times out (goes away) unless refreshed
- ❑ ARP is "plug-and-play":
 - nodes create their ARP tables without intervention from net administrator
- ARP Spoofing: Malicious nodes can lie about another node's MAC address to receive their traffic.

Routing to another LAN

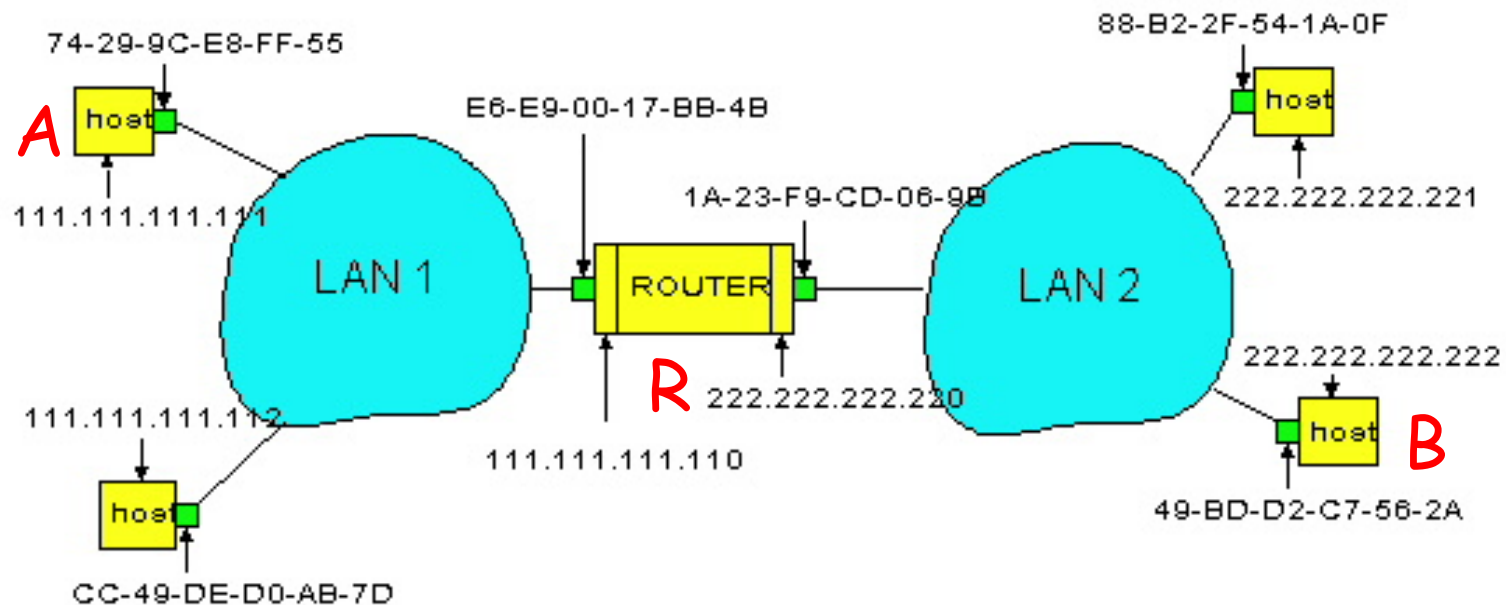
walkthrough: **send datagram from A to B via R**

assume A know's B IP address



- ❑ Two ARP tables in router R, one for each IP network (LAN)
- ❑ In routing table at source Host, find router 111.111.111.110
- ❑ In ARP table at source, find MAC address E6-E9-00-17-BB-4B, etc

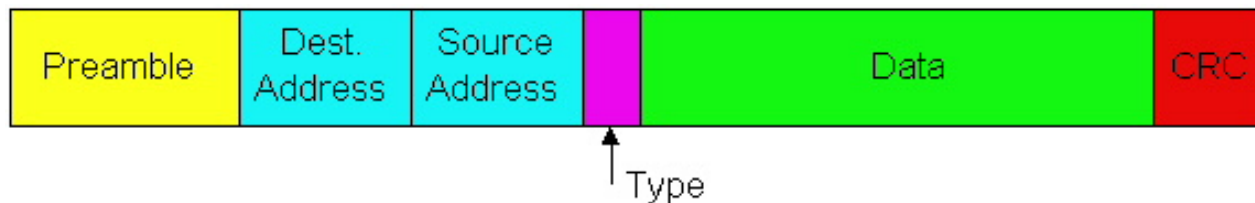
- ❑ A creates datagram with source A, destination B
- ❑ A uses ARP to get R's MAC address for 111.111.111.110
- ❑ A creates link-layer frame with R's MAC address as dest, frame contains A-to-B IP datagram
- ❑ A's adapter sends frame
- ❑ R's adapter receives frame
- ❑ R removes IP datagram from Ethernet frame, sees its destined to B
- ❑ R uses ARP to get B's MAC address
- ❑ R creates frame containing A-to-B IP datagram sends to B



Ethernet

Ethernet Frame Structure

Sending adapter encapsulates IP datagram (or other network layer protocol packet) in **Ethernet frame**



- ❑ **Preamble:** 8 bytes used for synchronizing sender and receiver.
- ❑ **Addresses:** 6 bytes
 - if adapter receives frame with matching destination address, or with broadcast address (eg ARP packet), it passes data in frame to net-layer protocol
 - otherwise, adapter discards frame
- ❑ **Type:** indicates the higher layer protocol (mostly IP but others may be supported such as Novell IPX and AppleTalk)
- ❑ **CRC: (cyclic redundancy checksum)**
 - ❑ checked at receiver, if error is detected, the frame is simply dropped

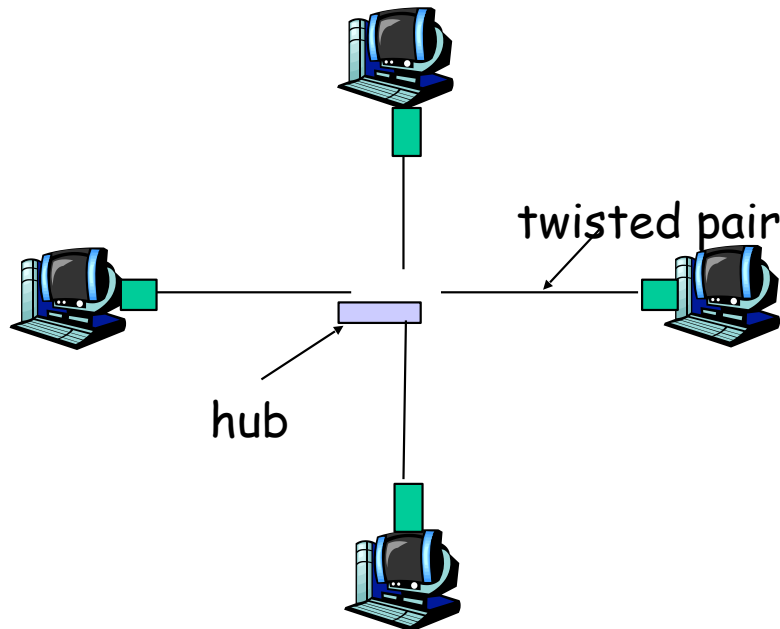
Ethernet uses CSMA/CD

- ❑ Adapter transmits whenever data is ready, but listens first.
- ❑ Adapter doesn't transmit if it senses that some other adapter is transmitting, that is, **carrier sense**
- ❑ Transmitting adapter aborts when it senses that another adapter is transmitting, that is, **collision detection**
- ❑ Before attempting a retransmission, adapter waits a random time, that is, **random access**

Hubs

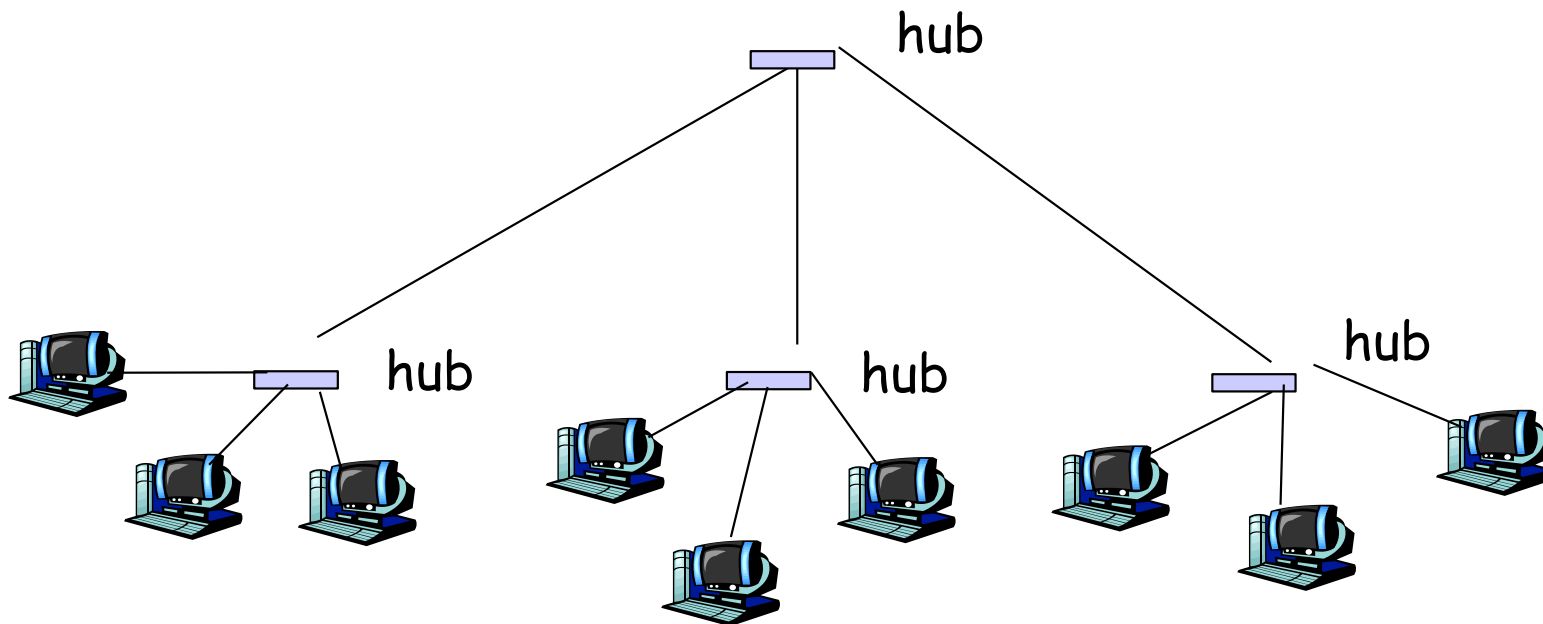
Hubs are essentially physical-layer repeaters:

- bits coming from one link go out all other links
- at the same rate
- no frame buffering
- no CSMA/CD at hub: adapters detect collisions
- provides net management functionality



Interconnecting with hubs

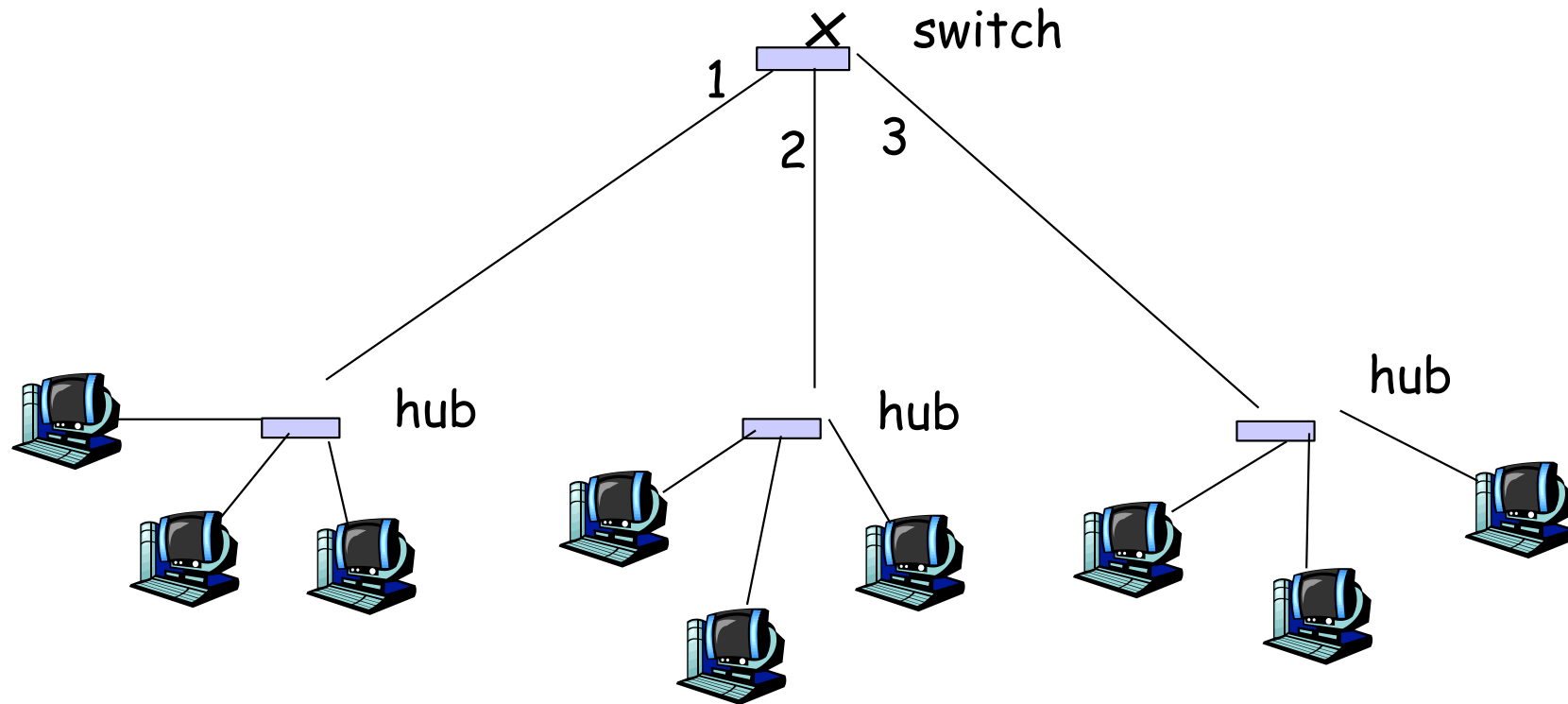
- ❑ Backbone hub interconnects LAN segments
- ❑ Extends max distance between nodes
- ❑ But individual segment collision domains become one large collision domain



Switch

- ❑ Link layer device
 - stores and forwards Ethernet frames
 - examines frame header and **selectively** forwards frame based on MAC dest address
 - when frame is to be forwarded on segment, uses CSMA/CD to access segment
- ❑ transparent
 - hosts are unaware of presence of switches
- ❑ plug-and-play, self-learning
 - switches do not need to be configured

Forwarding



- How do determine onto which LAN segment to forward frame?

Self learning

- ❑ Also called “backward learning” or “transparent bridging”
- ❑ A switch has a **switch table**
- ❑ entry in switch table:
 - [MAC Address, Interface, Time Stamp]
 - stale entries in table dropped (TTL can be 60 min)
- ❑ switch **learns** which hosts can be reached through which interfaces
 - when frame received, switch “learns” location of sender: incoming LAN segment
 - records sender/location pair in switch table

Filtering/Forwarding

When switch receives a frame:

index switch table using MAC dest address

if entry found for destination

then{

if dest on segment from which frame arrived

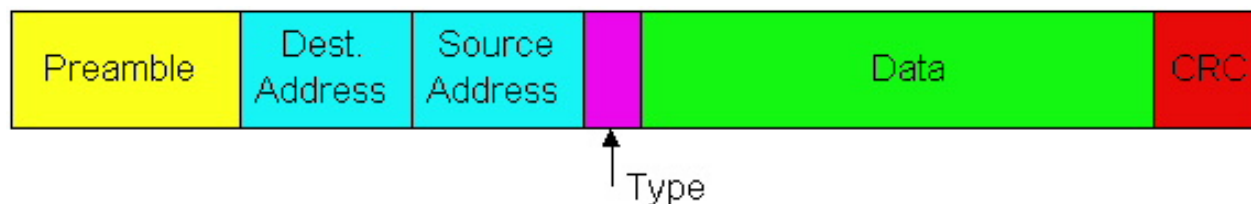
then drop the frame

else forward the frame on interface indicated

}

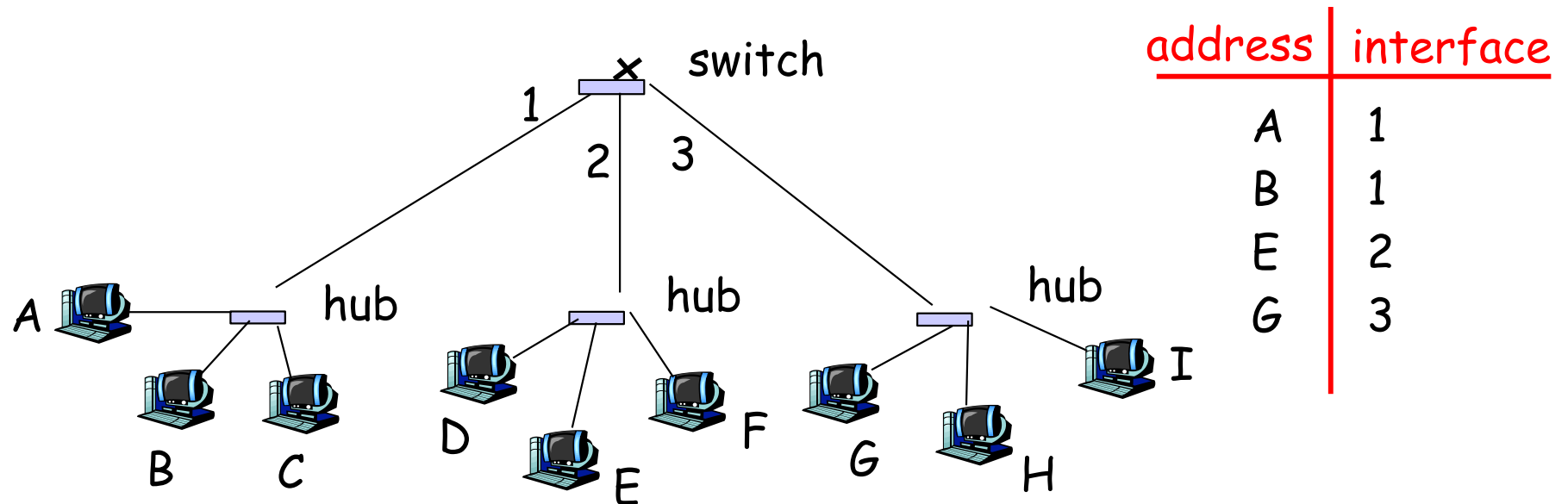
else flood

forward on all but the interface
on which the frame arrived



Switch example

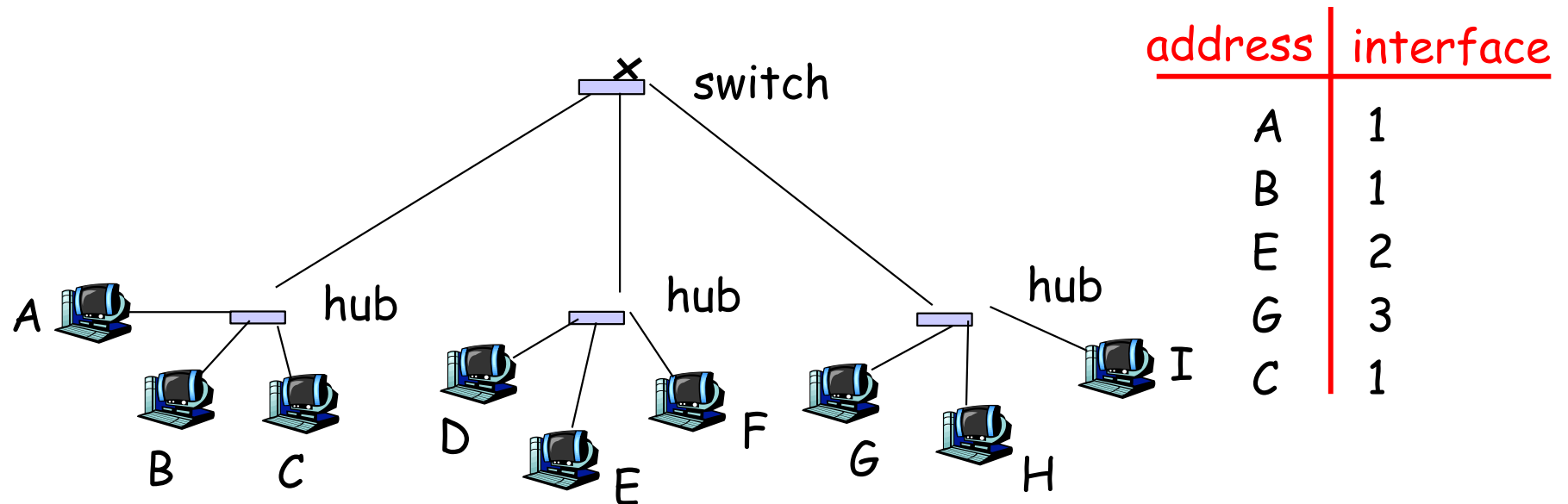
Suppose C sends frame to D



- ❑ Switch receives frame from from C
 - notes in bridge table that C is on interface 1
 - because D is not in table, switch forwards frame into interfaces 2 and 3
- ❑ frame received by D

Switch example

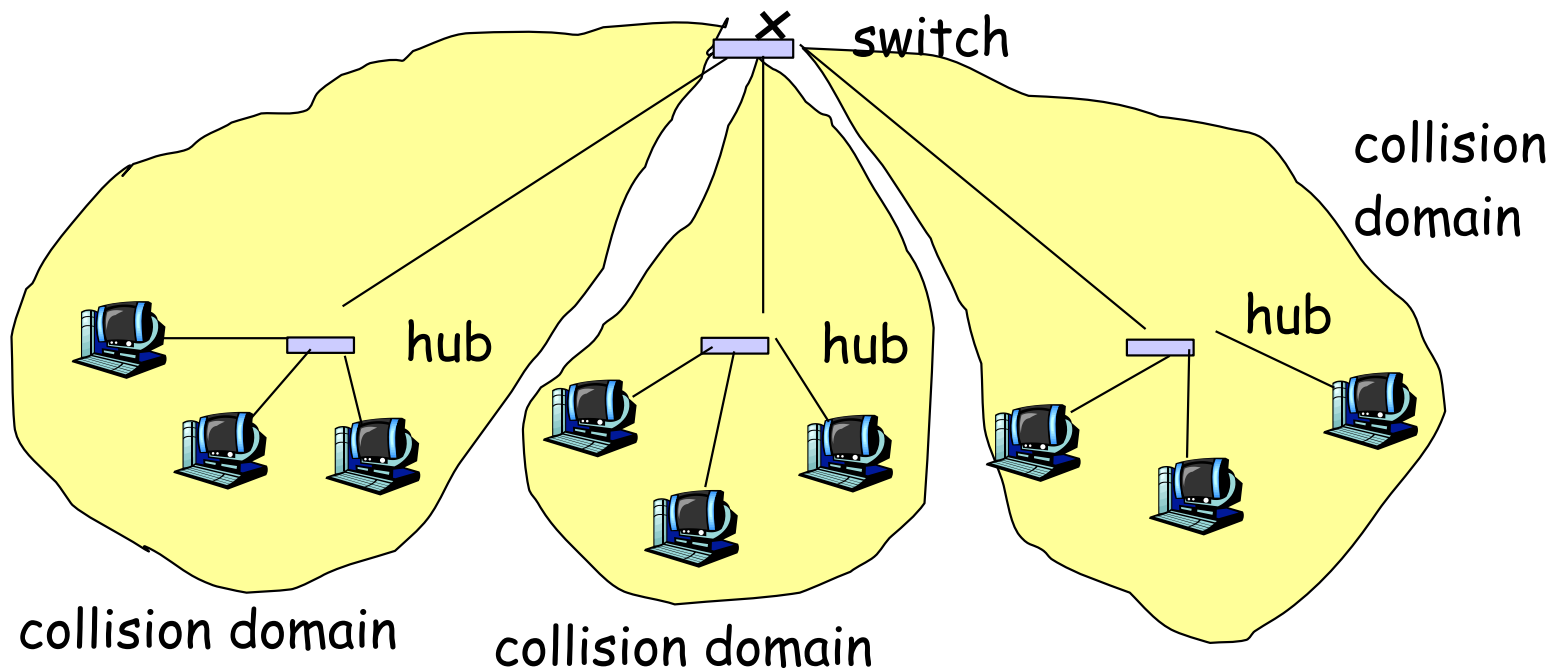
Suppose D replies back with frame to C.



- ❑ Switch receives frame from from D
 - notes in bridge table that D is on interface 2
 - because C is in table, switch forwards frame only to interface 1
- ❑ frame received by C

Switch: traffic isolation

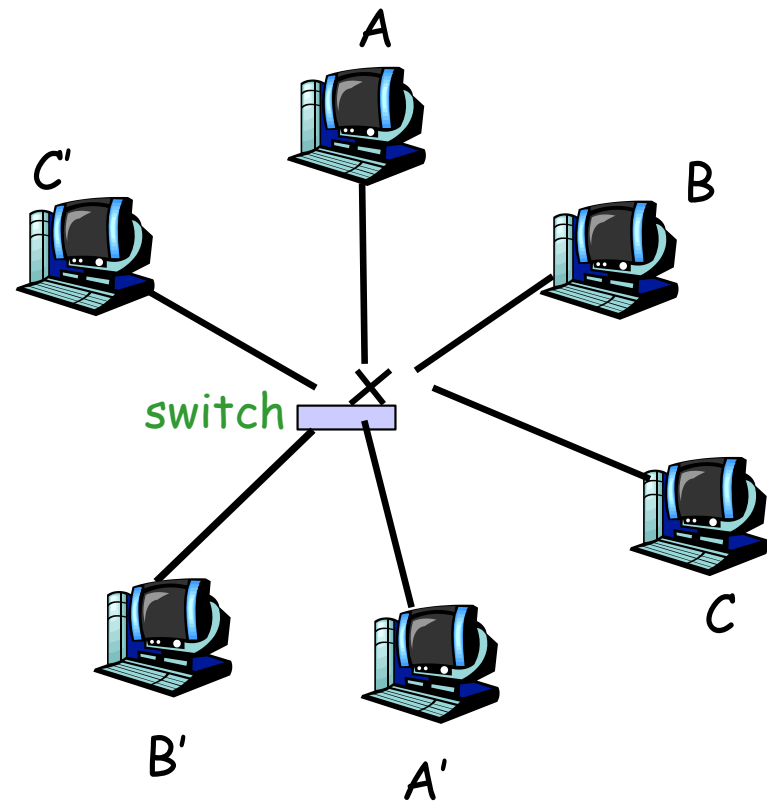
- switch installation breaks subnet into LAN segments
- switch **filters** packets:
 - same-LAN-segment frames not usually forwarded onto other LAN segments
 - segments become separate **collision domains**



Switches: dedicated access

- ❑ Switch with many interfaces
- ❑ Hosts have direct connection to switch
- ❑ No collisions; full duplex

Switching: A-to-A' and B-to-B' simultaneously, no collisions

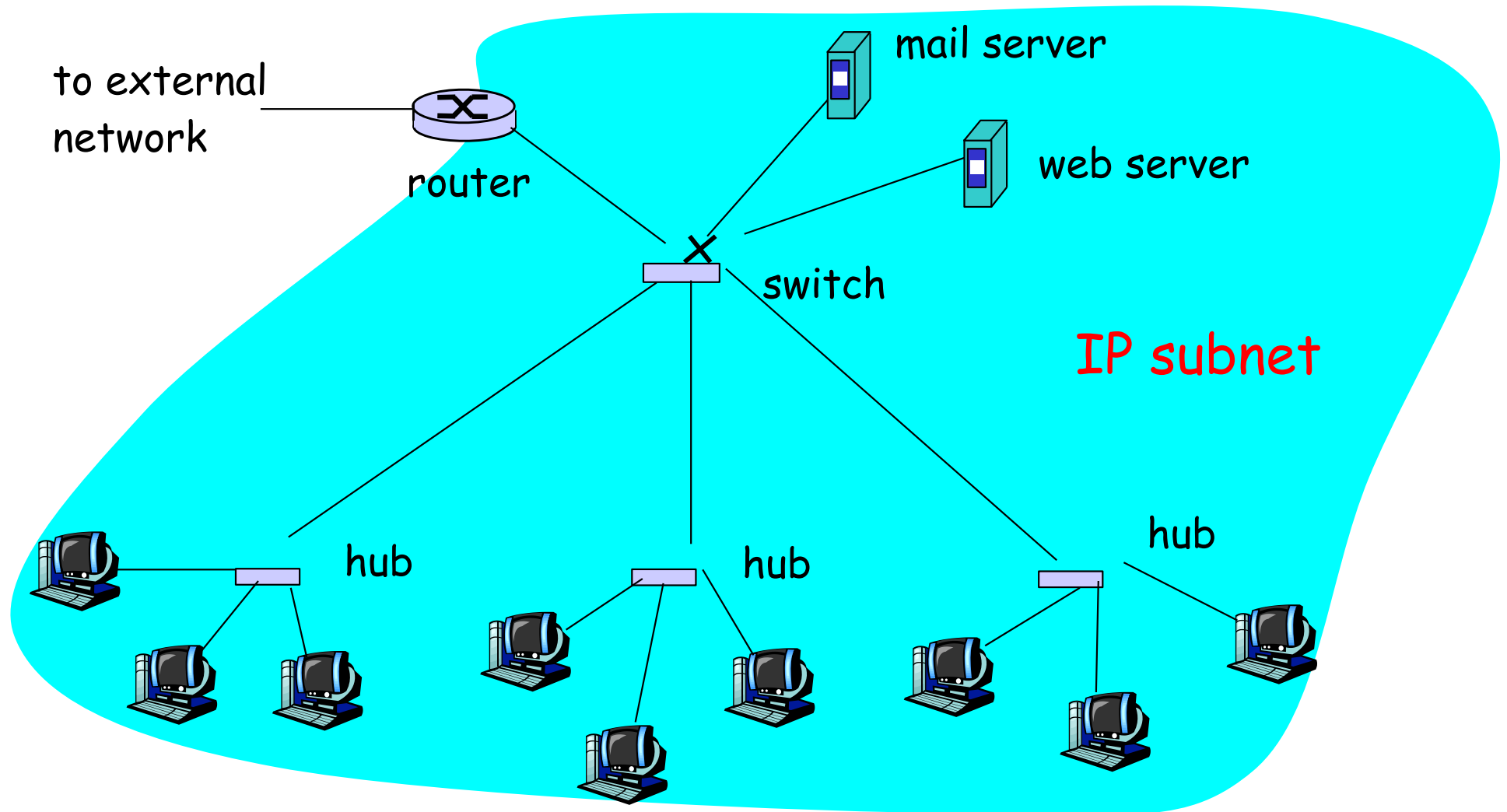


More on Switches

- ❑ **cut-through switching:** frame forwarded from input to output port without first collecting entire frame
 - slight reduction in latency

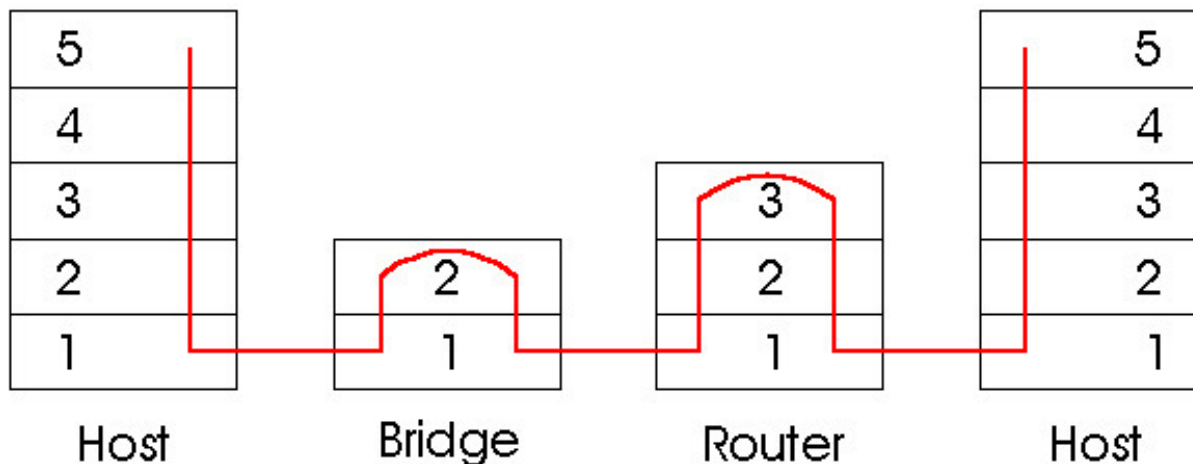
- ❑ combinations of shared/dedicated, 10/100/1000 Mbps interfaces

Institutional network



Switches vs. Routers

- ❑ both store-and-forward devices
 - routers: network layer devices (examine network layer headers)
 - switches are link layer devices
- ❑ routers maintain routing tables, implement routing algorithms
- ❑ switches maintain switch tables, implement filtering, learning algorithms



Summary comparison

	<u>hubs</u>	<u>routers</u>	<u>switches</u>
traffic isolation	no	yes	yes
plug & play	yes	no	yes
optimal routing	no	yes	no
cut through	yes	no	yes