Inter-process communication (IPC)

Kartik Gopalan
Chapter 2 of Tanenbaum's book
Chapter 4 and 5 of OSTEP book

In this lecture, we cover

- Pipes
- · Signals
- · Shared Memory

Other forms of IPC

·Parent-child

- · Command-line arguments,
- wait(...), waitpid(...)
- exit(...)
- ·Reading/modifying common files
 - · Servers commonly use 'pid' file to determine other active servers.
- ·Semaphores (separate topic)
 - · Locking and event signaling mechanism between processes

·Sockets

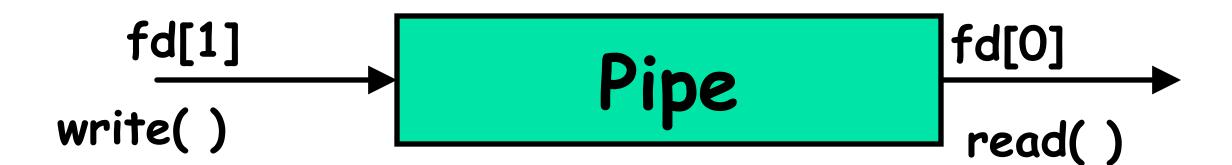
- Bi-directional
- Not just across the network, but also between processes.

<u>Pipes</u>

Pipe Abstraction

· Write to one end, read from another

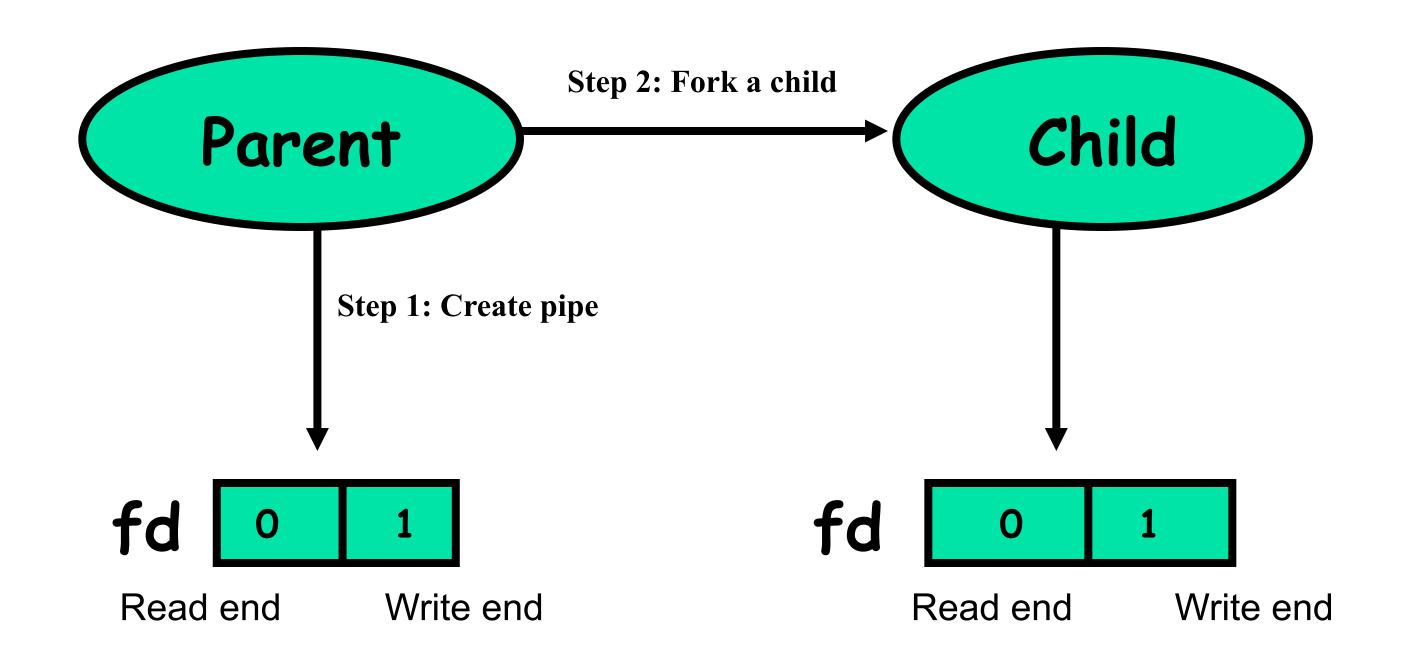
• pipe()



Pipe provides a byte-stream abstraction

- You can read and write at arbitrary byte boundaries.
 - •E.g. Byte lengths sequence written
 - •10, 10, 10, 10
 - •byte lengths sequence read
 - •5, 15, 15, 5
- As opposed to message abstraction, which provides explicit message boundaries.
 - •E.g. network packets

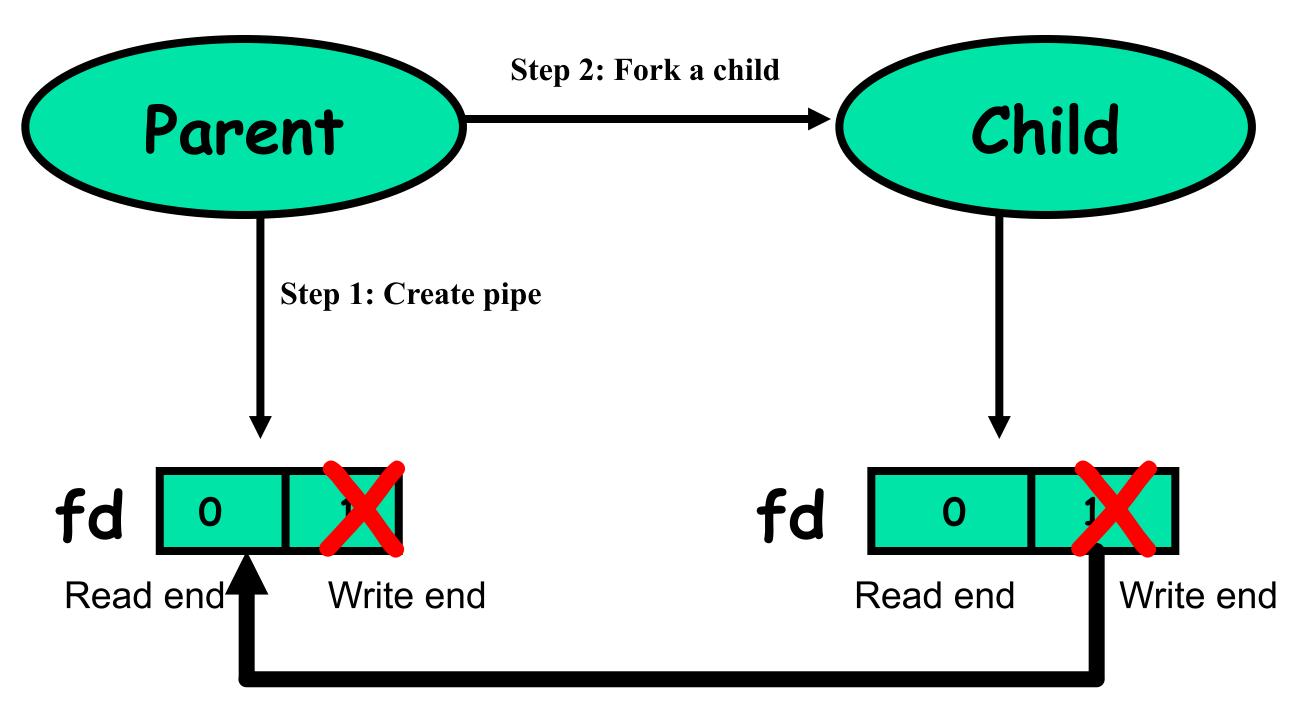
Parent-child communication using pipe



Here's an example.

https://oscourse.github.io/examples/pipe1.c

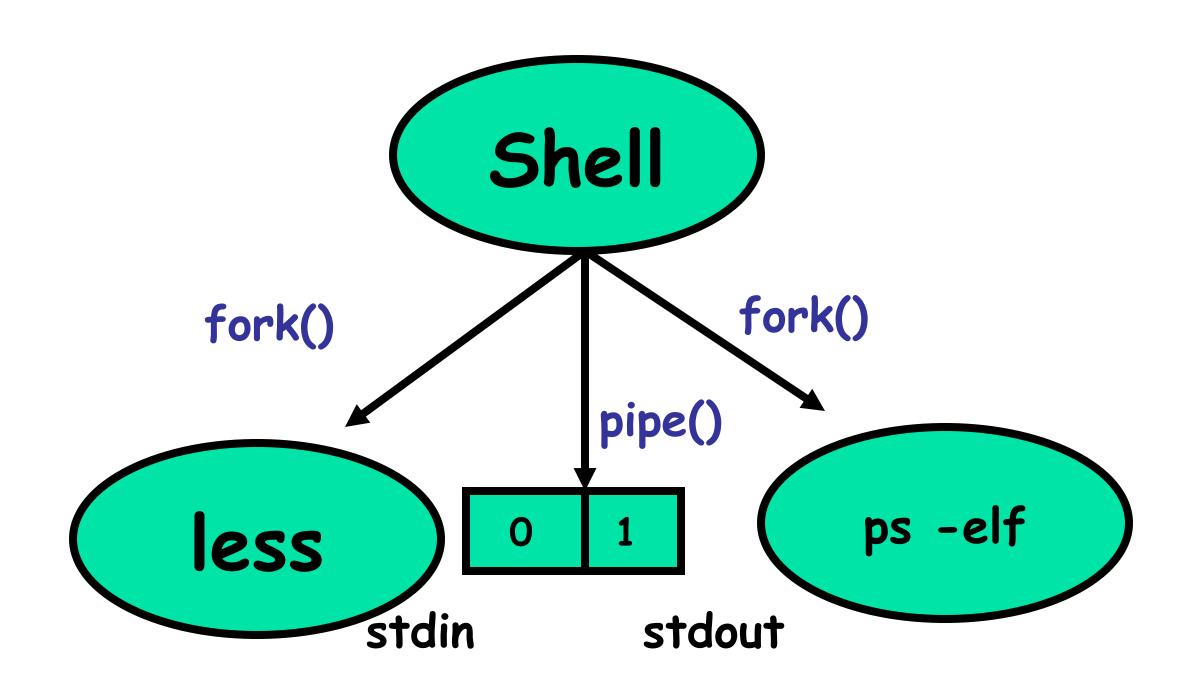
Parent-child communication using pipe



Here's an example.

https://oscourse.github.io/examples/pipe1.c

Filters in shell command-line ps -elf | less

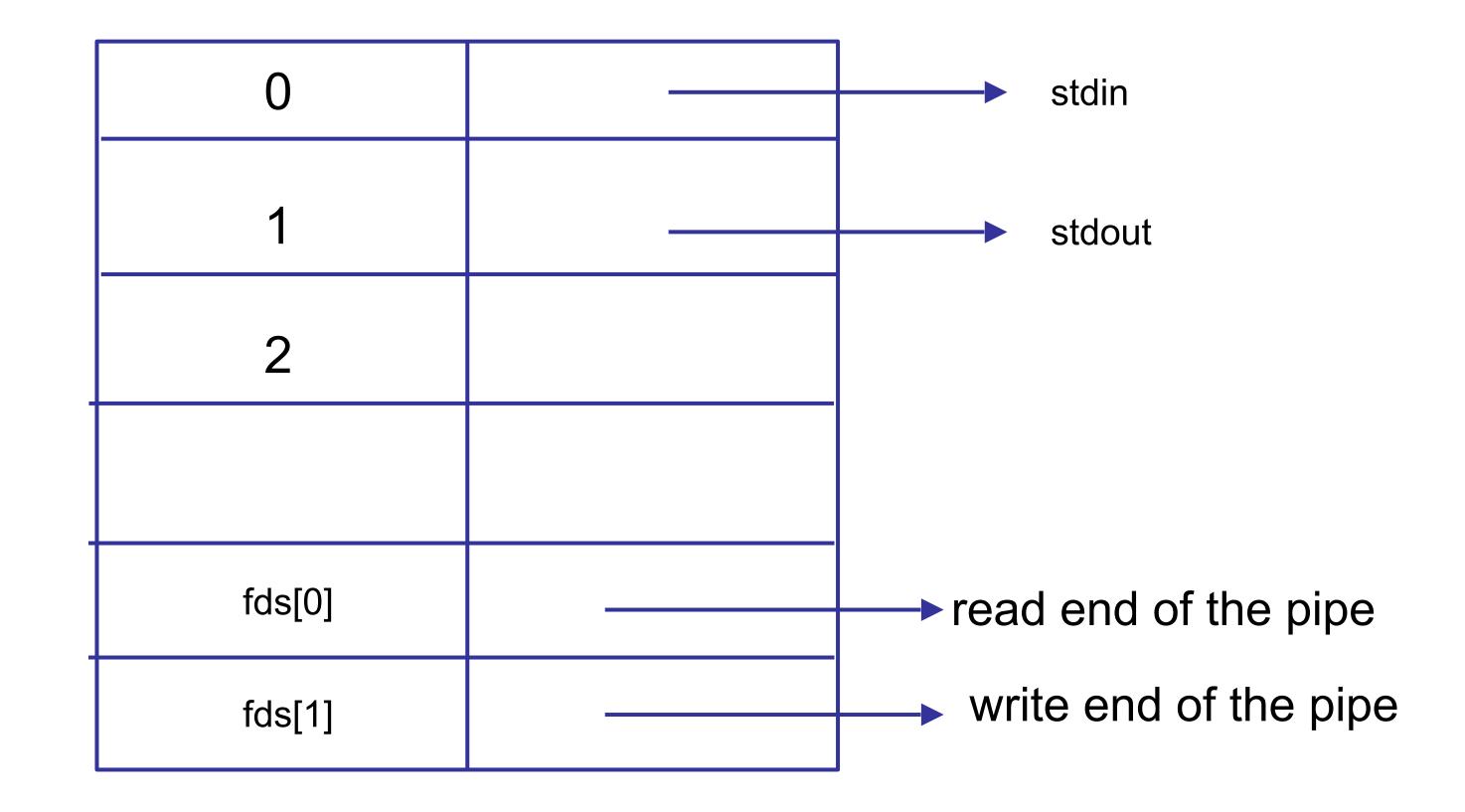


Here's an example.

https://oscourse.github.io/examples/pipe2.c

File-Descriptor Table

- Each process has a file-descriptor table
- One entry for each open file
- "File" = regular files, stdin, stdout, pipes, I/O devices etc.



Being careful with read()/write()

•read(fds[0], buf, 6);

- Doesn't mean read will return with 6 bytes of data!
- · It could be less. Why?

·Some reasons

- read() could reach end of input stream (EOF).
- · Other endpoint may abruptly close the connection
- · read() could return on a signal.
- ·So you must handle errors with every system call

Error handling...

·You must

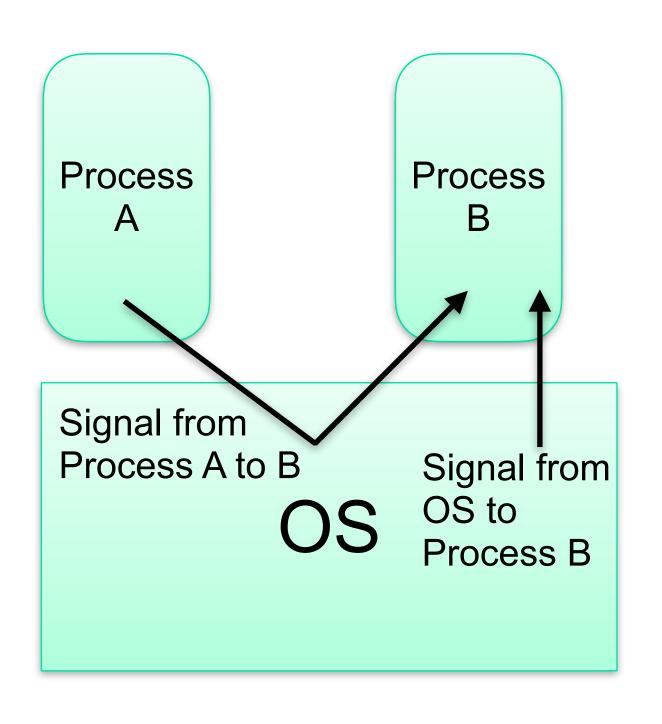
- · First check the return value of every read(...)/write(...) system call.
- · Then either...
- Wait to read/write more data
 OR
- · Handle any error conditions

```
More convinient to write a wrapper function
/* Write "n" bytes to a descriptor. */
ssize_t writen(int fd, const void *vptr, size_t n)
                 nleft;
     size_t
               nwritten;
     size_t
     const char
                   *ptr;
     ptr = vptr;
     nleft = n;
     while (nleft > 0) {
         if ((nwritten = write(fd, ptr, nleft))<=0){
             if (errno == EINTR)
              nwritten = 0; /* call write() again*/
             else return(-1); /* error */
          nleft -= nwritten;
          ptr += nwritten;
     return(n);
```

Signals

Signals Overview

- Signal is a notification to a process that an event has occurred.
 - Could come from another process or from the OS
- Type of event determined by the type of signal
- Try listing all signal types using \$\ \text{kill} \ -\ \lambda\$
- Some interesting signals
 - SIGCHLD, SIGKILL, SIGSTOP



Handling Signals

- Signals can be <u>caught</u> i.e. an action (or handler) can be associated with them
 - SIGKILL and SIGSTOP cannot be caught.
- Actions can be customized using
 - sigaction(...)
 - which associates a <u>signal handler</u> with the signal.
- <u>Default</u> action for most signals is to terminate the process
 - Except SIGCHLD and SIGURG are ignored by default.
- Unwanted signals can be <u>ignored</u>
 - Except SIGKILL or SIGSTOP
- Here's an example.
 - https://oscourse.github.io/examples/signals ex.c

Process in Execution

Instruction 1 Instruction 2

. . . .

Instruction K

Signal received
Save execution state
Jump to signal handler
Execute handler code

time

.

return from handler

(restore saved state)
Instruction K+1
Instruction K+2

.

More on SIGCHLD

- Sent to parent when a child process terminates or stops.
- If act.sa_handler is SIG_IGN
 - SIGCHLD will be ignored (default behavior)
- If act.sa_flags is SA_NOCLDSTOP
 - SIGCHLD won't be generated when children stop
- act.sa flags is SA NOCLDWAIT
 - children of the calling process will not be transformed into zombies when they terminate.
- These need to be set in sigaction() before parent calls fork()

Handling child's exit without blocking on wait()

- · Parent could install a signal handler for SIGCHLD
- · Call wait (...) /waitpid (...) inside the signal handler

```
void handle_sigchld(int signo) {
    pid_t pid;
    int stat;

pid = wait(&stat); //returns without blocking
    printf("child %d terminated\n", pid);
}
```

- · Here's an example.
 - •https://oscourse.github.io/examples/sigchld.c

More information...

• Check 'man sigaction(...)'

- Understand what happens when signal is delivered in the middle of a system call?
 - Different OSes have different behavior.

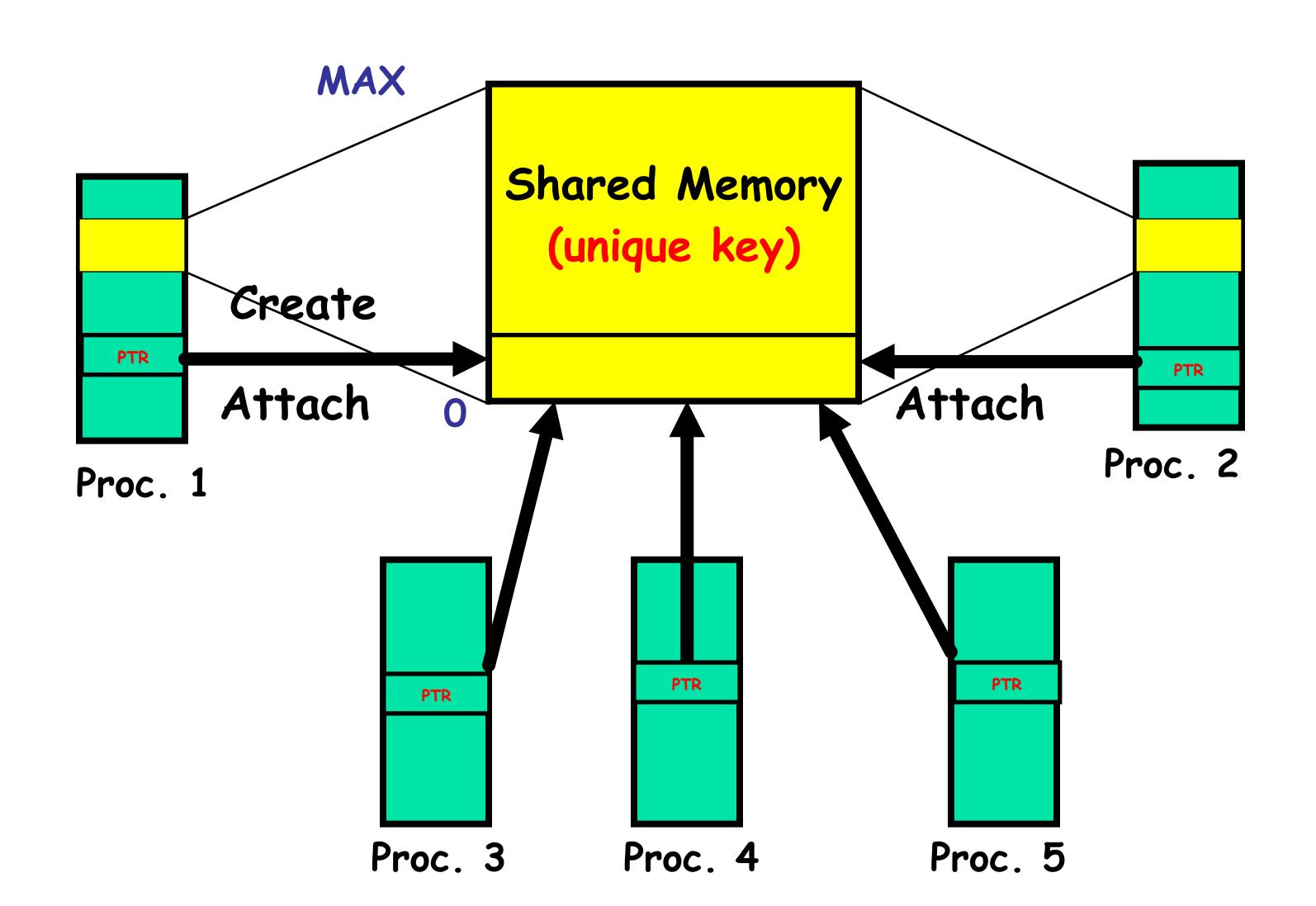
- Google for keywords "Unix Signals"
 - Tons of useful links

Shared Memory

• Man pages: shmget, shmat, shmdt, shmctl, semget, semop, semctl

Shared Memory

Common chunk of read/write memory among processes



Creating Shared Memory

```
int shmget(key t key, size t size, int shmflg);
Example:
key t key;
int shmid;
key = ftok("<somefile>", 'A');
shmid = shmget(key, 1024, 0644 | IPC CREAT);
Here's an example.
https://oscourse.github.io/examples/shm_create.c
```

Attach and Detach Shared Memory

```
void *shmat(int shmid, void *shmaddr, int shmflg);
int shmdt(void *shmaddr);
 Example:
   key t key;
   int shmid;
   char *data;
   key = ftok("<somefile>", 'A');
   shmid = shmget(key, 1024, 0644);
   data = shmat(shmid, (void *)0, 0);
   // read or write something to data here.
   shmdt (data);
 Here's an example.
 https://oscourse.github.io/examples/shm_attach.c
```

Deleting Shared Memory

```
int shmctl(int shmid, int cmd, struct shmid_ds *buf);
shmctl(shmid, IPC_RMID, NULL);
```

Example:

https://oscourse.github.io/examples/shm_delete.c

Command-line IPC control

·ipcs

· Lists all IPC objects owned by the user

•ipcrm

· Removes specific IPC object

References

·Unix man pages

- "Advanced Programming in Unix Environment" by Richard Stevens
 - http://www.kohala.com/start/apue.html