

APPENDIX A:

Notation and Mathematical Conventions

In this appendix we collect our notation, and some related mathematical facts and conventions.

A.1 SET NOTATION AND CONVENTIONS

If X is a set and x is an element of X , we write $x \in X$. A set can be specified in the form $X = \{x \mid x \text{ satisfies } P\}$, as the set of all elements satisfying property P . The union of two sets X_1 and X_2 is denoted by $X_1 \cup X_2$, and their intersection by $X_1 \cap X_2$. The empty set is denoted by \emptyset . The symbol \forall means “for all.”

The set of real numbers (also referred to as scalars) is denoted by \mathfrak{R} . The set of extended real numbers is denoted by \mathfrak{R}^* :

$$\mathfrak{R}^* = \mathfrak{R} \cup \{\infty, -\infty\}.$$

We write $-\infty < x < \infty$ for all real numbers x , and $-\infty \leq x \leq \infty$ for all extended real numbers x . We denote by $[a, b]$ the set of (possibly extended) real numbers x satisfying $a \leq x \leq b$. A rounded, instead of square, bracket denotes strict inequality in the definition. Thus $(a, b]$, $[a, b)$, and (a, b) denote the set of all x satisfying $a < x \leq b$, $a \leq x < b$, and $a < x < b$, respectively.

Generally, we adopt standard conventions regarding addition and multiplication in \mathfrak{R}^* , except that we take

$$\infty - \infty = -\infty + \infty = \infty,$$

and we take the product of 0 and ∞ or $-\infty$ to be 0. In this way the sum and product of two extended real numbers is well-defined. Division by 0 or ∞ does not appear in our analysis. In particular, we adopt the following rules in calculations involving ∞ and $-\infty$:

$$\begin{aligned}\alpha + \infty &= \infty + \alpha = \infty, & \forall \alpha \in \mathbb{R}^*, \\ \alpha - \infty &= -\infty + \alpha = -\infty, & \forall \alpha \in [-\infty, \infty), \\ \alpha \cdot \infty &= \infty, \quad \alpha \cdot (-\infty) = \infty, & \forall \alpha \in (0, \infty], \\ \alpha \cdot \infty &= -\infty, \quad \alpha \cdot (-\infty) = -\infty, & \forall \alpha \in [-\infty, 0), \\ 0 \cdot \infty &= \infty \cdot 0 = 0 = 0 \cdot (-\infty) = (-\infty) \cdot 0, & -(-\infty) = \infty.\end{aligned}$$

Under these rules, the following laws of arithmetic are still valid within \mathbb{R}^* :

$$\begin{aligned}\alpha_1 + \alpha_2 &= \alpha_2 + \alpha_1, & (\alpha_1 + \alpha_2) + \alpha_3 &= \alpha_1 + (\alpha_2 + \alpha_3), \\ \alpha_1 \alpha_2 &= \alpha_2 \alpha_1, & (\alpha_1 \alpha_2) \alpha_3 &= \alpha_1 (\alpha_2 \alpha_3).\end{aligned}$$

We also have

$$\alpha(\alpha_1 + \alpha_2) = \alpha\alpha_1 + \alpha\alpha_2$$

if either $\alpha \geq 0$ or else $(\alpha_1 + \alpha_2)$ is not of the form $\infty - \infty$.

Inf and Sup Notation

The *supremum* of a nonempty set $X \subset \mathbb{R}^*$, denoted by $\sup X$, is defined as the smallest $y \in \mathbb{R}^*$ such that $y \geq x$ for all $x \in X$. Similarly, the *infimum* of X , denoted by $\inf X$, is defined as the largest $y \in \mathbb{R}^*$ such that $y \leq x$ for all $x \in X$. For the empty set, we use the convention

$$\sup \emptyset = -\infty, \quad \inf \emptyset = \infty.$$

If $\sup X$ is equal to an $\bar{x} \in \mathbb{R}^*$ that belongs to the set X , we say that \bar{x} is the *maximum point* of X and we write $\bar{x} = \max X$. Similarly, if $\inf X$ is equal to an $\bar{x} \in \mathbb{R}^*$ that belongs to the set X , we say that \bar{x} is the *minimum point* of X and we write $\bar{x} = \min X$. Thus, when we write $\max X$ (or $\min X$) in place of $\sup X$ (or $\inf X$, respectively), we do so just for emphasis: we indicate that it is either evident, or it is known through earlier analysis, or it is about to be shown that the maximum (or minimum, respectively) of the set X is attained at one of its points.

A.2 FUNCTIONS

If f is a function, we use the notation $f : X \mapsto Y$ to indicate the fact that f is defined on a nonempty set X (its *domain*) and takes values in a set Y (its *range*). Thus when using the notation $f : X \mapsto Y$, we implicitly assume that X is nonempty. We will often use the *unit function* $e : X \mapsto \mathbb{R}$, defined by

$$e(x) = 1, \quad \forall x \in X.$$

Given a set X , we denote by $\mathcal{R}(X)$ the set of real-valued functions $J : X \mapsto \mathbb{R}$, and by $\mathcal{E}(X)$ the set of all extended real-valued functions $J : X \mapsto \mathbb{R}^*$. For any collection $\{J_\gamma \mid \gamma \in \Gamma\} \subset \mathcal{E}(X)$, parameterized by the elements of a set Γ , we denote by $\inf_{\gamma \in \Gamma} J_\gamma$ the function taking the value $\inf_{\gamma \in \Gamma} J_\gamma(x)$ at each $x \in X$.

For two functions $J_1, J_2 \in \mathcal{E}(X)$, we use the shorthand notation $J_1 \leq J_2$ to indicate the pointwise inequality

$$J_1(x) \leq J_2(x), \quad \forall x \in X.$$

We use the shorthand notation $\inf_{i \in I} J_i$ to denote the function obtained by pointwise infimum of a collection $\{J_i \mid i \in I\} \subset \mathcal{E}(X)$, i.e.,

$$\left(\inf_{i \in I} J_i \right)(x) = \inf_{i \in I} J_i(x), \quad \forall x \in X.$$

We use similar notation for sup.

Given subsets $S_1, S_2, S_3 \subset \mathcal{E}(X)$ and mappings $T_1 : S_1 \mapsto S_3$ and $T_2 : S_2 \mapsto S_1$, the *composition* of T_1 and T_2 is the mapping $T_1 T_2 : S_2 \mapsto S_3$ defined by

$$(T_1 T_2 J)(x) = (T_1(T_2 J))(x), \quad \forall J \in S_2, x \in X.$$

In particular, given a subset $S \subset \mathcal{E}(X)$ and mappings $T_1 : S \mapsto S$ and $T_2 : S \mapsto S$, the composition of T_1 and T_2 is the mapping $T_1 T_2 : S \mapsto S$ defined by

$$(T_1 T_2 J)(x) = (T_1(T_2 J))(x), \quad \forall J \in S, x \in X.$$

Similarly, given mappings $T_k : S \mapsto S$, $k = 1, \dots, N$, their composition is the mapping $(T_1 \cdots T_N) : S \mapsto S$ defined by

$$(T_1 T_2 \cdots T_N J)(x) = (T_1(T_2(\cdots(T_N J))))(x), \quad \forall J \in S, x \in X.$$

In our notation involving compositions we minimize the use of parentheses, as long as clarity is not compromised. In particular, we write $T_1 T_2 J$ instead of $(T_1 T_2 J)$ or $(T_1 T_2)J$ or $T_1(T_2 J)$, but we write $(T_1 T_2 J)(x)$ to indicate the value of $T_1 T_2 J$ at $x \in X$.

If X and Y are nonempty sets, a mapping $T : S_1 \mapsto S_2$, where $S_1 \subset \mathcal{E}(X)$ and $S_2 \subset E(Y)$, is said to be *monotone* if for all $J, J' \in S_1$,

$$J \leq J' \quad \Rightarrow \quad TJ \leq TJ'.$$

Sequences of Functions

For a sequence of functions $\{J_k\} \subset \mathcal{E}(X)$ that converges pointwise, we denote by $\lim_{k \rightarrow \infty} J_k$ the pointwise limit of $\{J_k\}$. We denote by $\limsup_{k \rightarrow \infty} J_k$ (or $\liminf_{k \rightarrow \infty} J_k$) the pointwise limit superior (or inferior, respectively) of $\{J_k\}$. If $\{J_k\} \subset \mathcal{E}(X)$ converges pointwise to J , we write $J_k \rightarrow J$. Note that we reserve this notation for pointwise convergence. To denote convergence with respect to a norm $\|\cdot\|$, we write $\|J_k - J\| \rightarrow 0$.

A sequence of functions $\{J_k\} \subset \mathcal{E}(X)$ is said to be *monotonically nonincreasing* (or *monotonically nondecreasing*) if $J_{k+1} \leq J_k$ for all k (or $J_{k+1} \geq J_k$ for all k , respectively). Such a sequence always has a (pointwise) limit within $\mathcal{E}(X)$. We write $J_k \downarrow J$ (or $J_k \uparrow J$) to indicate that $\{J_k\}$ is monotonically nonincreasing (or monotonically nonincreasing, respectively) and that its limit is J .

Let $\{J_{mn}\} \subset \mathcal{E}(X)$ be a double indexed sequence, which is monotonically nonincreasing separately for each index in the sense that

$$J_{(m+1)n} \leq J_{mn}, \quad J_{m(n+1)} \leq J_{mn}, \quad \forall m, n = 0, 1, \dots$$

For such sequences, a useful fact is that

$$\lim_{m \rightarrow \infty} \left(\lim_{n \rightarrow \infty} J_{mn} \right) = \lim_{m \rightarrow \infty} J_{mm}.$$

There is a similar fact for monotonically nondecreasing sequences.

Expected Values

Given a random variable w defined over a probability space Ω , the expected value of w is defined by

$$E\{w\} = E\{w^+\} + E\{w^-\},$$

where w^+ and w^- are the positive and negative parts of w ,

$$w^+(\omega) = \max \{0, w(\omega)\}, \quad w^-(\omega) = \min \{0, w(\omega)\}.$$

In this way, taking also into account the rule $\infty - \infty = \infty$, the expected value $E\{w\}$ is well-defined if Ω is finite or countably infinite. In more general cases, $E\{w\}$ is similarly defined by the appropriate form of integration, and more detail will be given at specific points as needed.

APPENDIX B:

Contraction Mappings

B.1 CONTRACTION MAPPING FIXED POINT THEOREMS

The purpose of this appendix is to provide some background on contraction mappings and their properties. Let Y be a real vector space with a norm $\|\cdot\|$, i.e., a real-valued function satisfying for all $y \in Y$, $\|y\| \geq 0$, $\|y\| = 0$ if and only if $y = 0$, and

$$\|ay\| = |a|\|y\|, \quad \forall a \in \mathbb{R}, \quad \|y + z\| \leq \|y\| + \|z\|, \quad \forall y, z \in Y.$$

Let \overline{Y} be a closed subset of Y . A function $F : \overline{Y} \mapsto \overline{Y}$ is said to be a *contraction mapping* if for some $\rho \in (0, 1)$, we have

$$\|Fy - Fz\| \leq \rho\|y - z\|, \quad \forall y, z \in \overline{Y}.$$

The scalar ρ is called the *modulus of contraction* of F .

Example B.1 (Linear Contraction Mappings in \mathbb{R}^n)

Consider the case of a linear mapping $F : \mathbb{R}^n \mapsto \mathbb{R}^n$ of the form

$$Fy = b + Ay,$$

where A is an $n \times n$ matrix and b is a vector in \mathbb{R}^n . Let $\sigma(A)$ denote the *spectral radius* of A (the largest modulus among the moduli of the eigenvalues of A). Then it can be shown that A is a *contraction mapping with respect to some norm if and only if $\sigma(A) < 1$* .

Specifically, given $\epsilon > 0$, there exists a norm $\|\cdot\|_s$ such that

$$\|Ay\|_s \leq (\sigma(A) + \epsilon)\|y\|_s, \quad \forall y \in \mathbb{R}^n. \quad (\text{B.1})$$

Thus, if $\sigma(A) < 1$ we may select $\epsilon > 0$ such that $\rho = \sigma(A) + \epsilon < 1$, and obtain the contraction relation

$$\|Fy - Fz\|_s = \|A(y - z)\|_s \leq \rho\|y - z\|_s, \quad \forall y, z \in \Re^n. \quad (\text{B.2})$$

The norm $\|\cdot\|_s$ can be taken to be a weighted Euclidean norm, i.e., it may have the form $\|y\|_s = \|My\|$, where M is a square invertible matrix, and $\|\cdot\|$ is the standard Euclidean norm, i.e., $\|x\| = \sqrt{x'x}$. †

Conversely, if Eq. (B.2) holds for some norm $\|\cdot\|_s$ and all real vectors y, z , it also holds for all complex vectors y, z with the squared norm $\|c\|_s^2$ of a complex vector c defined as the sum of the squares of the norms of the real and the imaginary components. Thus from Eq. (B.2), by taking $y - z = u$, where u is an eigenvector corresponding to an eigenvalue λ with $|\lambda| = \sigma(A)$, we have $\sigma(A)\|u\|_s = \|Au\|_s \leq \rho\|u\|_s$. Hence $\sigma(A) \leq \rho$, and it follows that if F is a contraction with respect to a given norm, we must have $\sigma(A) < 1$.

A sequence $\{y_k\} \subset Y$ is said to be a *Cauchy sequence* if $\|y_m - y_n\| \rightarrow 0$ as $m, n \rightarrow \infty$, i.e., given any $\epsilon > 0$, there exists N such that $\|y_m - y_n\| \leq \epsilon$ for all $m, n \geq N$. The space Y is said to be *complete* under the norm $\|\cdot\|$ if every Cauchy sequence $\{y_k\} \subset Y$ is convergent, in the sense that for some $y \in Y$, we have $\|y_k - y\| \rightarrow 0$. Note that a Cauchy sequence is always bounded. Also, a Cauchy sequence of real numbers is convergent, implying that the real line is a complete space and so is every real finite-dimensional vector space. On the other hand, an infinite dimensional space may not be complete under some norms, while it may be complete under other norms.

When Y is complete and \overline{Y} is a closed subset of Y , an important property of a contraction mapping $F : \overline{Y} \mapsto \overline{Y}$ is that it has a unique fixed point within \overline{Y} , i.e., the equation

$$y = Fy$$

has a unique solution $y^* \in \overline{Y}$, called the *fixed point of F* . Furthermore, the sequence $\{y_k\}$ generated by the iteration

$$y_{k+1} = Fy_k$$

† We may show Eq. (B.1) by using the Jordan canonical form of A , which is denoted by J . In particular, if P is a nonsingular matrix such that $P^{-1}AP = J$ and D is the diagonal matrix with $1, \delta, \dots, \delta^{n-1}$ along the diagonal, where $\delta > 0$, it is straightforward to verify that $D^{-1}P^{-1}APD = \hat{J}$, where \hat{J} is the matrix that is identical to J except that each nonzero off-diagonal term is replaced by δ . Defining $\hat{P} = PD$, we have $A = \hat{P}\hat{J}\hat{P}^{-1}$. Now if $\|\cdot\|$ is the standard Euclidean norm, we note that for some $\beta > 0$, we have $\|\hat{J}z\| \leq (\sigma(A) + \beta\delta)\|z\|$ for all $z \in \Re^n$ and $\delta \in (0, 1]$. For a given $\delta \in (0, 1]$, consider the weighted Euclidean norm $\|\cdot\|_s$ defined by $\|y\|_s = \|\hat{P}^{-1}y\|$. Then we have for all $y \in \Re^n$,

$$\|Ay\|_s = \|\hat{P}^{-1}Ay\| = \|\hat{P}^{-1}\hat{P}\hat{J}\hat{P}^{-1}y\| = \|\hat{J}\hat{P}^{-1}y\| \leq (\sigma(A) + \beta\delta)\|\hat{P}^{-1}y\|,$$

so that $\|Ay\|_s \leq (\sigma(A) + \beta\delta)\|y\|_s$, for all $y \in \Re^n$. For a given $\epsilon > 0$, we choose $\delta = \epsilon/\beta$, so the preceding relation yields Eq. (B.1).

converges to y^* , starting from an arbitrary initial point y_0 .

Proposition B.1: (Contraction Mapping Fixed-Point Theorem) Let Y be a complete vector space and let \overline{Y} be a closed subset of Y . Then if $F : \overline{Y} \mapsto \overline{Y}$ is a contraction mapping with modulus $\rho \in (0, 1)$, there exists a unique $y^* \in \overline{Y}$ such that

$$y^* = Fy^*.$$

Furthermore, the sequence $\{F^k y\}$ converges to y^* for any $y \in \overline{Y}$, and we have

$$\|F^k y - y^*\| \leq \rho^k \|y - y^*\|, \quad k = 1, 2, \dots$$

Proof: Let $y \in \overline{Y}$ and consider the iteration $y_{k+1} = Fy_k$ starting with $y_0 = y$. By the contraction property of F ,

$$\|y_{k+1} - y_k\| \leq \rho \|y_k - y_{k-1}\|, \quad k = 1, 2, \dots,$$

which implies that

$$\|y_{k+1} - y_k\| \leq \rho^k \|y_1 - y_0\|, \quad k = 1, 2, \dots.$$

It follows that for every $k \geq 0$ and $m \geq 1$, we have

$$\begin{aligned} \|y_{k+m} - y_k\| &\leq \sum_{i=1}^m \|y_{k+i} - y_{k+i-1}\| \\ &\leq \rho^k (1 + \rho + \dots + \rho^{m-1}) \|y_1 - y_0\| \\ &\leq \frac{\rho^k}{1 - \rho} \|y_1 - y_0\|. \end{aligned}$$

Therefore, $\{y_k\}$ is a Cauchy sequence in \overline{Y} and must converge to a limit $y^* \in \overline{Y}$, since Y is complete and \overline{Y} is closed. We have for all $k \geq 1$,

$$\|Fy^* - y^*\| \leq \|Fy^* - y_k\| + \|y_k - y^*\| \leq \rho \|y^* - y_{k-1}\| + \|y_k - y^*\|$$

and since y_k converges to y^* , we obtain $Fy^* = y^*$. Thus, the limit y^* of y_k is a fixed point of F . It is a unique fixed point because if \tilde{y} were another fixed point, we would have

$$\|y^* - \tilde{y}\| = \|Fy^* - F\tilde{y}\| \leq \rho \|y^* - \tilde{y}\|,$$

which implies that $y^* = \tilde{y}$.

To show the convergence rate bound of the last part, note that

$$\|F^k y - y^*\| = \|F^k y - F y^*\| \leq \rho \|F^{k-1} y - y^*\|.$$

Repeating this process for a total of k times, we obtain the desired result.
Q.E.D.

The convergence rate exhibited by $F^k y$ in the preceding proposition is said to be *geometric*, and $F^k y$ is said to converge to its limit y^* *geometrically*. This is in reference to the fact that the error $\|F^k y - y^*\|$ converges to 0 faster than some geometric progression ($\rho^k \|y - y^*\|$ in this case).

In some contexts of interest to us one may encounter mappings that are not contractions, but become contractions when iterated a finite number of times. In this case, one may use a slightly different version of the contraction mapping fixed point theorem, which we now present.

We say that a function $F : Y \mapsto Y$ is an *m-stage contraction mapping* if there exists a positive integer m and some $\rho < 1$ such that

$$\|F^m y - F^m y'\| \leq \rho \|y - y'\|, \quad \forall y, y' \in Y,$$

where F^m denotes the composition of F with itself m times. Thus, F is an *m-stage contraction* if F^m is a contraction. Again, the scalar ρ is called the modulus of contraction. We have the following generalization of Prop. B.1.

Proposition B.2: (m-Stage Contraction Mapping Fixed-Point Theorem) Let Y be a complete vector space and let \overline{Y} be a closed subset of Y . Then if $F : \overline{Y} \mapsto \overline{Y}$ is an *m-stage contraction mapping* with modulus $\rho \in (0, 1)$, there exists a unique $y^* \in \overline{Y}$ such that

$$y^* = Fy^*.$$

Furthermore, $\{F^k y\}$ converges to y^* for any $y \in \overline{Y}$.

Proof: Since F^m maps \overline{Y} into \overline{Y} and is a contraction mapping, by Prop. B.1, it has a unique fixed point in \overline{Y} , denoted y^* . Applying F to both sides of the relation $y^* = F^m y^*$, we see that Fy^* is also a fixed point of F^m , so by the uniqueness of the fixed point, we have $y^* = Fy^*$. Therefore y^* is a fixed point of F . If F had another fixed point, say \tilde{y} , then we would have $\tilde{y} = F^m \tilde{y}$, which by the uniqueness of the fixed point of F^m implies that $\tilde{y} = y^*$. Thus, y^* is the unique fixed point of F .

To show the convergence of $\{F^k y\}$, note that by Prop. B.1, we have for all $y \in \overline{Y}$,

$$\lim_{k \rightarrow \infty} \|F^{mk} y - y^*\| = 0.$$

Using $F^\ell y$ in place of y , we obtain

$$\lim_{k \rightarrow \infty} \|F^{mk+\ell}y - y^*\| = 0, \quad \ell = 0, 1, \dots, m-1,$$

which proves the desired result. **Q.E.D.**

B.2 WEIGHTED SUP-NORM CONTRACTIONS

In this section, we will focus on contraction mappings within a specialized context that is particularly important in DP. Let X be a set (typically the state space in DP), and let $v : X \mapsto \mathbb{R}$ be a positive-valued function,

$$v(x) > 0, \quad \forall x \in X.$$

Let $\mathcal{B}(X)$ denote the set of all functions $J : X \mapsto \mathbb{R}$ such that $J(x)/v(x)$ is bounded as x ranges over X . We define a norm on $\mathcal{B}(X)$, called the *weighted sup-norm*, by

$$\|J\| = \sup_{x \in X} \frac{|J(x)|}{v(x)}. \quad (\text{B.3})$$

It is easily verified that $\|\cdot\|$ thus defined has the required properties for being a norm. Furthermore, $\mathcal{B}(X)$ is complete under this norm. To see this, consider a Cauchy sequence $\{J_k\} \subset \mathcal{B}(X)$, and note that $\|J_m - J_n\| \rightarrow 0$ as $m, n \rightarrow \infty$ implies that for all $x \in X$, $\{J_k(x)\}$ is a Cauchy sequence of real numbers, so it converges to some $J^*(x)$. We will show that $J^* \in \mathcal{B}(X)$ and that $\|J_k - J^*\| \rightarrow 0$. To this end, it will be sufficient to show that given any $\epsilon > 0$, there exists an integer K such that

$$\frac{|J_k(x) - J^*(x)|}{v(x)} \leq \epsilon, \quad \forall x \in X, k \geq K.$$

This will imply that

$$\sup_{x \in X} \frac{|J^*(x)|}{v(x)} \leq \epsilon + \|J_k\|, \quad \forall k \geq K,$$

so that $J^* \in \mathcal{B}(X)$, and will also imply that $\|J_k - J^*\| \leq \epsilon$, so that $\|J_k - J^*\| \rightarrow 0$. Assume the contrary, i.e., that there exists an $\epsilon > 0$ and a subsequence $\{x_{m_1}, x_{m_2}, \dots\} \subset X$ such that $m_i < m_{i+1}$ and

$$\epsilon < \frac{|J_{m_i}(x_{m_i}) - J^*(x_{m_i})|}{v(x_{m_i})}, \quad \forall i \geq 1.$$

The right-hand side above is less or equal to

$$\frac{|J_{m_i}(x_{m_i}) - J_n(x_{m_i})|}{v(x_{m_i})} + \frac{|J_n(x_{m_i}) - J^*(x_{m_i})|}{v(x_{m_i})}, \quad \forall n \geq 1, i \geq 1.$$

The first term in the above sum is less than $\epsilon/2$ for i and n larger than some threshold; fixing i and letting n be sufficiently large, the second term can also be made less than $\epsilon/2$, so the sum is made less than ϵ - a contradiction. In conclusion, the space $\mathcal{B}(X)$ is complete, so the fixed point results of Props. B.1 and B.2 apply.

In our discussions, unless we specify otherwise, we will assume that $\mathcal{B}(X)$ is equipped with the weighted sup-norm above, where the weight function v will be clear from the context. There will be frequent occasions where the norm will be unweighted, i.e., $v(x) \equiv 1$ and $\|J\| = \sup_{x \in X} |J(x)|$, in which case we will explicitly state so.

Finite-Dimensional Cases

Let us now focus on the finite-dimensional case $X = \{1, \dots, n\}$, in which case $\mathcal{R}(X)$ and $\mathcal{B}(X)$ can be identified with \mathbb{R}^n . We first consider a linear mapping (cf. Example B.1). We have the following proposition.

Proposition B.3: Consider a linear mapping $F : \mathbb{R}^n \mapsto \mathbb{R}^n$ of the form

$$Fy = b + Ay,$$

where A is an $n \times n$ matrix with components a_{ij} , and b is a vector in \mathbb{R}^n . Denote by $|A|$ the matrix whose components are the absolute values of the components of A and let $\sigma(A)$ and $\sigma(|A|)$ denote the spectral radii of A and $|A|$, respectively. Then:

- (a) $|A|$ has a real eigenvalue λ , which is equal to its spectral radius, and an associated nonnegative eigenvector.
- (b) F is a contraction with respect to some weighted sup-norm if and only if $\sigma(|A|) < 1$. In particular, any substochastic matrix P ($p_{ij} \geq 0$ for all i, j , and $\sum_{j=1}^n p_{ij} \leq 1$, for all i) is a contraction with respect to some weighted sup-norm if and only if $\sigma(P) < 1$.
- (c) F is a contraction with respect to the weighted sup-norm

$$\|y\| = \max_{i=1, \dots, n} \frac{|y_i|}{v(i)}$$

if and only if

$$\frac{\sum_{j=1}^n |a_{ij}| v(j)}{v(i)} < 1, \quad i = 1, \dots, n.$$

Proof: (a) This is the Perron-Frobenius Theorem; see e.g., [BeT89], Chapter 2, Prop. 6.6.

(b) This follows from the Perron-Frobenius Theorem; see [BeT89], Ch. 2, Cor. 6.2.

(c) This is proved in more general form in the following Prop. B.4. **Q.E.D.**

Consider next a nonlinear mapping $F : \mathbb{R}^n \mapsto \mathbb{R}^n$ that has the property

$$|Fy - Fz| \leq P |y - z|, \quad \forall y, z \in \mathbb{R}^n,$$

for some matrix P with nonnegative components and $\sigma(P) < 1$. Here, we generically denote by $|w|$ the vector whose components are the absolute values of the components of w , and the inequality is componentwise. Then we claim that F is a contraction with respect to some weighted sup-norm. To see this note that by the preceding discussion, P is a contraction with respect to some weighted sup-norm $\|y\| = \max_{i=1,\dots,n} |y_i|/v(i)$, and we have

$$\frac{(|Fy - Fz|)(i)}{v(i)} \leq \frac{(P|y - z|)(i)}{v(i)} \leq \alpha \|y - z\|, \quad \forall i = 1, \dots, n,$$

for some $\alpha \in (0, 1)$, where $(|Fy - Fz|)(i)$ and $(P|y - z|)(i)$ are the i th components of the vectors $|Fy - Fz|$ and $P|y - z|$, respectively. Thus, F is a contraction with respect to $\|\cdot\|$. For additional discussion of linear and nonlinear contraction mapping properties and characterizations such as the one above, see the book [OrR70].

Linear Mappings on Countable Spaces

The case where X is countable (or, as a special case, finite) is frequently encountered in DP. The following proposition provides some useful criteria for verifying the contraction property of mappings that are either linear or are obtained via a parametric minimization of other contraction mappings.

Proposition B.4: Let $X = \{1, 2, \dots\}$.

(a) Let $F : \mathcal{B}(X) \mapsto \mathcal{B}(X)$ be a linear mapping of the form

$$(FJ)(i) = b_i + \sum_{j \in X} a_{ij} J(j), \quad i \in X,$$

where b_i and a_{ij} are some scalars. Then F is a contraction with modulus ρ with respect to the weighted sup-norm (B.3) if and only if

$$\frac{\sum_{j \in X} |a_{ij}| v(j)}{v(i)} \leq \rho, \quad i \in X. \quad (\text{B.4})$$

(b) Let $F : \mathcal{B}(X) \mapsto \mathcal{B}(X)$ be a mapping of the form

$$(FJ)(i) = \inf_{\mu \in M} (F_\mu J)(i), \quad i \in X,$$

where M is parameter set, and for each $\mu \in M$, F_μ is a contraction mapping from $\mathcal{B}(X)$ to $\mathcal{B}(X)$ with modulus ρ . Then F is a contraction mapping with modulus ρ .

Proof: (a) Assume that Eq. (B.4) holds. For any $J, J' \in \mathcal{B}(X)$, we have

$$\begin{aligned} \|FJ - FJ'\| &= \sup_{i \in X} \frac{\left| \sum_{j \in X} a_{ij} (J(j) - J'(j)) \right|}{v(i)} \\ &\leq \sup_{i \in X} \frac{\sum_{j \in X} |a_{ij}| v(j) \left(|J(j) - J'(j)| / v(j) \right)}{v(i)} \\ &\leq \sup_{i \in X} \frac{\sum_{j \in X} |a_{ij}| v(j)}{v(i)} \|J - J'\| \\ &\leq \rho \|J - J'\|, \end{aligned}$$

where the last inequality follows from the hypothesis.

Conversely, arguing by contradiction, let's assume that Eq. (B.4) is violated for some $i \in X$. Define $J(j) = v(j) \operatorname{sgn}(a_{ij})$ and $J'(j) = 0$ for all $j \in X$. Then we have $\|J - J'\| = \|J\| = 1$, and

$$\frac{|(FJ)(i) - (FJ')(i)|}{v(i)} = \frac{\sum_{j \in X} |a_{ij}| v(j)}{v(i)} > \rho = \rho \|J - J'\|,$$

showing that F is not a contraction of modulus ρ .

(b) Since F_μ is a contraction of modulus ρ , we have for any $J, J' \in \mathcal{B}(X)$,

$$\frac{(F_\mu J)(i)}{v(i)} \leq \frac{(F_\mu J')(i)}{v(i)} + \rho \|J - J'\|, \quad i \in X,$$

so by taking the infimum over $\mu \in M$,

$$\frac{(FJ)(i)}{v(i)} \leq \frac{(FJ')(i)}{v(i)} + \rho \|J - J'\|, \quad i \in X.$$

Reversing the roles of J and J' , we obtain

$$\frac{|(FJ)(i) - (FJ')(i)|}{v(i)} \leq \rho \|J - J'\|, \quad i \in X,$$

and by taking the supremum over i , the contraction property of F is proved.
Q.E.D.

The preceding proposition assumes that $FJ \in \mathcal{B}(X)$ for all $J \in \mathcal{B}(X)$. The following proposition provides conditions, particularly relevant to the DP context, which imply this assumption.

Proposition B.5: Let $X = \{1, 2, \dots\}$, let M be a parameter set, and for each $\mu \in M$, let F_μ be a linear mapping of the form

$$(F_\mu J)(i) = b_i(\mu) + \sum_{j \in X} a_{ij}(\mu) J(j), \quad i \in X.$$

- (a) We have $F_\mu J \in \mathcal{B}(X)$ for all $J \in \mathcal{B}(X)$ provided $b(\mu) \in \mathcal{B}(X)$ and $V(\mu) \in \mathcal{B}(X)$, where

$$b(\mu) = \{b_1(\mu), b_2(\mu), \dots\}, \quad V(\mu) = \{V_1(\mu), V_2(\mu), \dots\},$$

with

$$V_i(\mu) = \sum_{j \in X} |a_{ij}(\mu)| v(j), \quad i \in X.$$

- (b) Consider the mapping F

$$(FJ)(i) = \inf_{\mu \in M} (F_\mu J)(i), \quad i \in X.$$

We have $FJ \in \mathcal{B}(X)$ for all $J \in \mathcal{B}(X)$, provided $b \in \mathcal{B}(X)$ and $V \in \mathcal{B}(X)$, where

$$b = \{b_1, b_2, \dots\}, \quad V = \{V_1, V_2, \dots\},$$

with $b_i = \sup_{\mu \in M} b_i(\mu)$ and $V_i = \sup_{\mu \in M} V_i(\mu)$.

Proof: (a) For all $\mu \in M$, $J \in \mathcal{B}(X)$ and $i \in X$, we have

$$(F_\mu J)(i) \leq |b_i(\mu)| + \sum_{j \in X} |a_{ij}(\mu)| |J(j)/v(j)| v(j)$$

$$\begin{aligned} &\leq |b_i(\mu)| + \|J\| \sum_{j \in X} |a_{ij}(\mu)| v(j) \\ &= |b_i(\mu)| + \|J\| V_i(\mu), \end{aligned}$$

and similarly $(F_\mu J)(i) \geq -|b_i(\mu)| - \|J\| V_i(\mu)$. Thus

$$|(F_\mu J)(i)| \leq |b_i(\mu)| + \|J\| V_i(\mu), \quad i \in X.$$

By dividing this inequality with $v(i)$ and by taking the supremum over $i \in X$, we obtain

$$\|F_\mu J\| \leq \|b_\mu\| + \|J\| \|V_\mu\| < \infty.$$

(b) By doing the same as in (a), but after first taking the infimum of $(F_\mu J)(i)$ over μ , we obtain

$$\|FJ\| \leq \|b\| + \|J\| \|V\| < \infty.$$

Q.E.D.

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