

APPENDIX A:

Notation and Mathematical Conventions

In this appendix we collect our notation, and some related mathematical facts and conventions.

A.1 SET NOTATION AND CONVENTIONS

If X is a set and x is an element of X , we write $x \in X$. A set can be specified in the form $X = \{x \mid x \text{ satisfies } P\}$, as the set of all elements satisfying property P . The union of two sets X_1 and X_2 is denoted by $X_1 \cup X_2$, and their intersection by $X_1 \cap X_2$. The empty set is denoted by \emptyset . The symbol \forall means “for all.”

The set of real numbers (also referred to as scalars) is denoted by \mathbb{R} . The set of extended real numbers is denoted by \mathbb{R}^* :

$$\mathbb{R}^* = \mathbb{R} \cup \{\infty, -\infty\}.$$

We write $-\infty < x < \infty$ for all real numbers x , and $-\infty \leq x \leq \infty$ for all extended real numbers x . We denote by $[a, b]$ the set of (possibly extended) real numbers x satisfying $a \leq x \leq b$. A rounded, instead of square, bracket denotes strict inequality in the definition. Thus $(a, b]$, $[a, b)$, and (a, b) denote the set of all x satisfying $a < x \leq b$, $a \leq x < b$, and $a < x < b$, respectively.

Generally, we adopt standard conventions regarding addition and multiplication in \mathbb{R}^* , except that we take

$$\infty - \infty = -\infty + \infty = \infty,$$

and we take the product of 0 and ∞ or $-\infty$ to be 0. In this way the sum and product of two extended real numbers is well-defined. Division by 0 or ∞ does not appear in our analysis. In particular, we adopt the following rules in calculations involving ∞ and $-\infty$:

$$\alpha + \infty = \infty + \alpha = \infty, \quad \forall \alpha \in \mathbb{R}^*,$$

$$\alpha - \infty = -\infty + \alpha = -\infty, \quad \forall \alpha \in [-\infty, \infty),$$

$$\alpha \cdot \infty = \infty, \quad \alpha \cdot (-\infty) = \infty, \quad \forall \alpha \in (0, \infty],$$

$$\alpha \cdot \infty = -\infty, \quad \alpha \cdot (-\infty) = -\infty, \quad \forall \alpha \in [-\infty, 0),$$

$$0 \cdot \infty = \infty \cdot 0 = 0 = 0 \cdot (-\infty) = (-\infty) \cdot 0, \quad -(-\infty) = \infty.$$

Under these rules, the following laws of arithmetic are still valid within \mathbb{R}^* :

$$\alpha_1 + \alpha_2 = \alpha_2 + \alpha_1, \quad (\alpha_1 + \alpha_2) + \alpha_3 = \alpha_1 + (\alpha_2 + \alpha_3),$$

$$\alpha_1 \alpha_2 = \alpha_2 \alpha_1, \quad (\alpha_1 \alpha_2) \alpha_3 = \alpha_1 (\alpha_2 \alpha_3).$$

We also have

$$\alpha(\alpha_1 + \alpha_2) = \alpha\alpha_1 + \alpha\alpha_2$$

if either $\alpha \geq 0$ or else $(\alpha_1 + \alpha_2)$ is not of the form $\infty - \infty$.

Inf and Sup Notation

The *supremum* of a nonempty set $X \subset \mathbb{R}^*$, denoted by $\sup X$, is defined as the smallest $y \in \mathbb{R}^*$ such that $y \geq x$ for all $x \in X$. Similarly, the *infimum* of X , denoted by $\inf X$, is defined as the largest $y \in \mathbb{R}^*$ such that $y \leq x$ for all $x \in X$. For the empty set, we use the convention

$$\sup \emptyset = -\infty, \quad \inf \emptyset = \infty.$$

If $\sup X$ is equal to an $\bar{x} \in \mathbb{R}^*$ that belongs to the set X , we say that \bar{x} is the *maximum point* of X and we write $\bar{x} = \max X$. Similarly, if $\inf X$ is equal to an $\bar{x} \in \mathbb{R}^*$ that belongs to the set X , we say that \bar{x} is the *minimum point* of X and we write $\bar{x} = \min X$. Thus, when we write $\max X$ (or $\min X$) in place of $\sup X$ (or $\inf X$, respectively), we do so just for emphasis: we indicate that it is either evident, or it is known through earlier analysis, or it is about to be shown that the maximum (or minimum, respectively) of the set X is attained at one of its points.

A.2 FUNCTIONS

If f is a function, we use the notation $f : X \mapsto Y$ to indicate the fact that f is defined on a nonempty set X (its *domain*) and takes values in a set Y (its *range*). Thus when using the notation $f : X \mapsto Y$, we implicitly assume that X is nonempty. We will often use the *unit function* $e : X \mapsto \mathbb{R}$, defined by

$$e(x) = 1, \quad \forall x \in X.$$

Given a set X , we denote by $\mathcal{R}(X)$ the set of real-valued functions $J : X \mapsto \mathbb{R}$, and by $\mathcal{E}(X)$ the set of all extended real-valued functions $J : X \mapsto \mathbb{R}^*$. For any collection $\{J_\gamma \mid \gamma \in \Gamma\} \subset \mathcal{E}(X)$, parameterized by the elements of a set Γ , we denote by $\inf_{\gamma \in \Gamma} J_\gamma$ the function taking the value $\inf_{\gamma \in \Gamma} J_\gamma(x)$ at each $x \in X$.

For two functions $J_1, J_2 \in \mathcal{E}(X)$, we use the shorthand notation $J_1 \leq J_2$ to indicate the pointwise inequality

$$J_1(x) \leq J_2(x), \quad \forall x \in X.$$

We use the shorthand notation $\inf_{i \in I} J_i$ to denote the function obtained by pointwise infimum of a collection $\{J_i \mid i \in I\} \subset \mathcal{E}(X)$, i.e.,

$$\left(\inf_{i \in I} J_i \right)(x) = \inf_{i \in I} J_i(x), \quad \forall x \in X.$$

We use similar notation for sup.

Given subsets $S_1, S_2, S_3 \subset \mathcal{E}(X)$ and mappings $T_1 : S_1 \mapsto S_3$ and $T_2 : S_2 \mapsto S_1$, the *composition* of T_1 and T_2 is the mapping $T_1 T_2 : S_2 \mapsto S_3$ defined by

$$(T_1 T_2 J)(x) = (T_1(T_2 J))(x), \quad \forall J \in S_2, x \in X.$$

In particular, given a subset $S \subset \mathcal{E}(X)$ and mappings $T_1 : S \mapsto S$ and $T_2 : S \mapsto S$, the composition of T_1 and T_2 is the mapping $T_1 T_2 : S \mapsto S$ defined by

$$(T_1 T_2 J)(x) = (T_1(T_2 J))(x), \quad \forall J \in S, x \in X.$$

Similarly, given mappings $T_k : S \mapsto S$, $k = 1, \dots, N$, their composition is the mapping $(T_1 \cdots T_N) : S \mapsto S$ defined by

$$(T_1 T_2 \cdots T_N J)(x) = (T_1(T_2(\cdots (T_N J))))(x), \quad \forall J \in S, x \in X.$$

In our notation involving compositions we minimize the use of parentheses, as long as clarity is not compromised. In particular, we write $T_1 T_2 J$ instead of $(T_1 T_2 J)$ or $(T_1 T_2)J$ or $T_1(T_2 J)$, but we write $(T_1 T_2 J)(x)$ to indicate the value of $T_1 T_2 J$ at $x \in X$.

If X and Y are nonempty sets, a mapping $T : S_1 \mapsto S_2$, where $S_1 \subset \mathcal{E}(X)$ and $S_2 \subset \mathcal{E}(Y)$, is said to be *monotone* if for all $J, J' \in S_1$,

$$J \leq J' \quad \Rightarrow \quad T J \leq T J'.$$

Sequences of Functions

For a sequence of functions $\{J_k\} \subset \mathcal{E}(X)$ that converges pointwise, we denote by $\lim_{k \rightarrow \infty} J_k$ the pointwise limit of $\{J_k\}$. We denote by $\limsup_{k \rightarrow \infty} J_k$ (or $\liminf_{k \rightarrow \infty} J_k$) the pointwise limit superior (or inferior, respectively) of $\{J_k\}$. If $\{J_k\} \subset \mathcal{E}(X)$ converges pointwise to J , we write $J_k \rightarrow J$. Note that we reserve this notation for pointwise convergence. To denote convergence with respect to a norm $\|\cdot\|$, we write $\|J_k - J\| \rightarrow 0$.

A sequence of functions $\{J_k\} \subset \mathcal{E}(X)$ is said to be *monotonically nonincreasing* (or *monotonically nondecreasing*) if $J_{k+1} \leq J_k$ for all k (or $J_{k+1} \geq J_k$ for all k , respectively). Such a sequence always has a (pointwise) limit within $\mathcal{E}(X)$. We write $J_k \downarrow J$ (or $J_k \uparrow J$) to indicate that $\{J_k\}$ is monotonically nonincreasing (or monotonically nondecreasing, respectively) and that its limit is J .

Let $\{J_{mn}\} \subset \mathcal{E}(X)$ be a double indexed sequence, which is monotonically nonincreasing separately for each index in the sense that

$$J_{(m+1)n} \leq J_{mn}, \quad J_{m(n+1)} \leq J_{mn}, \quad \forall m, n = 0, 1, \dots$$

For such sequences, a useful fact is that

$$\lim_{m \rightarrow \infty} \left(\lim_{n \rightarrow \infty} J_{mn} \right) = \lim_{m \rightarrow \infty} J_{mm}.$$

There is a similar fact for monotonically nondecreasing sequences.

Expected Values

Given a random variable w defined over a probability space Ω , the expected value of w is defined by

$$E\{w\} = E\{w^+\} + E\{w^-\},$$

where w^+ and w^- are the positive and negative parts of w ,

$$w^+(\omega) = \max\{0, w(\omega)\}, \quad w^-(\omega) = \min\{0, w(\omega)\}.$$

In this way, taking also into account the rule $\infty - \infty = \infty$, the expected value $E\{w\}$ is well-defined if Ω is finite or countably infinite. In more general cases, $E\{w\}$ is similarly defined by the appropriate form of integration, and more detail will be given at specific points as needed.

APPENDIX B:

Contraction Mappings

B.1 CONTRACTION MAPPING FIXED POINT THEOREMS

The purpose of this appendix is to provide some background on contraction mappings and their properties. Let Y be a real vector space with a norm $\|\cdot\|$, i.e., a real-valued function satisfying for all $y \in Y$, $\|y\| \geq 0$, $\|y\| = 0$ if and only if $y = 0$, and

$$\|ay\| = |a|\|y\|, \quad \forall a \in \mathbb{R}, \quad \|y + z\| \leq \|y\| + \|z\|, \quad \forall y, z \in Y.$$

Let \bar{Y} be a closed subset of Y . A function $F : \bar{Y} \mapsto \bar{Y}$ is said to be a *contraction mapping* if for some $\rho \in (0, 1)$, we have

$$\|Fy - Fz\| \leq \rho\|y - z\|, \quad \forall y, z \in \bar{Y}.$$

The scalar ρ is called the *modulus of contraction* of F .

Example B.1 (Linear Contraction Mappings in \mathbb{R}^n)

Consider the case of a linear mapping $F : \mathbb{R}^n \mapsto \mathbb{R}^n$ of the form

$$Fy = b + Ay,$$

where A is an $n \times n$ matrix and b is a vector in \mathbb{R}^n . Let $\sigma(A)$ denote the *spectral radius* of A (the largest modulus among the moduli of the eigenvalues of A). Then it can be shown that A is a contraction mapping with respect to some norm if and only if $\sigma(A) < 1$.

Specifically, given $\epsilon > 0$, there exists a norm $\|\cdot\|_s$ such that

$$\|Ay\|_s \leq (\sigma(A) + \epsilon)\|y\|_s, \quad \forall y \in \mathbb{R}^n. \quad (\text{B.1})$$

Thus, if $\sigma(A) < 1$ we may select $\epsilon > 0$ such that $\rho = \sigma(A) + \epsilon < 1$, and obtain the contraction relation

$$\|Fy - Fz\|_s = \|A(y - z)\|_s \leq \rho\|y - z\|_s, \quad \forall y, z \in \mathbb{R}^n. \quad (\text{B.2})$$

The norm $\|\cdot\|_s$ can be taken to be a weighted Euclidean norm, i.e., it may have the form $\|y\|_s = \|My\|$, where M is a square invertible matrix, and $\|\cdot\|$ is the standard Euclidean norm, i.e., $\|x\| = \sqrt{x'x}$.[†]

Conversely, if Eq. (B.2) holds for some norm $\|\cdot\|_s$ and all real vectors y, z , it also holds for all complex vectors y, z with the squared norm $\|c\|_s^2$ of a complex vector c defined as the sum of the squares of the norms of the real and the imaginary components. Thus from Eq. (B.2), by taking $y - z = u$, where u is an eigenvector corresponding to an eigenvalue λ with $|\lambda| = \sigma(A)$, we have $\sigma(A)\|u\|_s = \|Au\|_s \leq \rho\|u\|_s$. Hence $\sigma(A) \leq \rho$, and it follows that if F is a contraction with respect to a given norm, we must have $\sigma(A) < 1$.

A sequence $\{y_k\} \subset Y$ is said to be a *Cauchy sequence* if $\|y_m - y_n\| \rightarrow 0$ as $m, n \rightarrow \infty$, i.e., given any $\epsilon > 0$, there exists N such that $\|y_m - y_n\| \leq \epsilon$ for all $m, n \geq N$. The space Y is said to be *complete* under the norm $\|\cdot\|$ if every Cauchy sequence $\{y_k\} \subset Y$ is convergent, in the sense that for some $y \in Y$, we have $\|y_k - y\| \rightarrow 0$. Note that a Cauchy sequence is always bounded. Also, a Cauchy sequence of real numbers is convergent, implying that the real line is a complete space and so is every real finite-dimensional vector space. On the other hand, an infinite dimensional space may not be complete under some norms, while it may be complete under other norms.

When Y is complete and \bar{Y} is a closed subset of Y , an important property of a contraction mapping $F : \bar{Y} \mapsto \bar{Y}$ is that it has a unique fixed point within \bar{Y} , i.e., the equation

$$y = Fy$$

has a unique solution $y^* \in \bar{Y}$, called the *fixed point of F* . Furthermore, the sequence $\{y_k\}$ generated by the iteration

$$y_{k+1} = Fy_k$$

[†] We may show Eq. (B.1) by using the Jordan canonical form of A , which is denoted by J . In particular, if P is a nonsingular matrix such that $P^{-1}AP = J$ and D is the diagonal matrix with $1, \delta, \dots, \delta^{n-1}$ along the diagonal, where $\delta > 0$, it is straightforward to verify that $D^{-1}P^{-1}APD = \hat{J}$, where \hat{J} is the matrix that is identical to J except that each nonzero off-diagonal term is replaced by δ . Defining $\hat{P} = PD$, we have $A = \hat{P}\hat{J}\hat{P}^{-1}$. Now if $\|\cdot\|$ is the standard Euclidean norm, we note that for some $\beta > 0$, we have $\|\hat{J}z\| \leq (\sigma(A) + \beta\delta)\|z\|$ for all $z \in \mathbb{R}^n$ and $\delta \in (0, 1]$. For a given $\delta \in (0, 1]$, consider the weighted Euclidean norm $\|\cdot\|_s$ defined by $\|y\|_s = \|\hat{P}^{-1}y\|$. Then we have for all $y \in \mathbb{R}^n$,

$$\|Ay\|_s = \|\hat{P}^{-1}Ay\| = \|\hat{P}^{-1}\hat{P}\hat{J}\hat{P}^{-1}y\| = \|\hat{J}\hat{P}^{-1}y\| \leq (\sigma(A) + \beta\delta)\|\hat{P}^{-1}y\|,$$

so that $\|Ay\|_s \leq (\sigma(A) + \beta\delta)\|y\|_s$, for all $y \in \mathbb{R}^n$. For a given $\epsilon > 0$, we choose $\delta = \epsilon/\beta$, so the preceding relation yields Eq. (B.1).

converges to y^* , starting from an arbitrary initial point y_0 .

Proposition B.1: (Contraction Mapping Fixed-Point Theorem) Let Y be a complete vector space and let \overline{Y} be a closed subset of Y . Then if $F : \overline{Y} \mapsto \overline{Y}$ is a contraction mapping with modulus $\rho \in (0, 1)$, there exists a unique $y^* \in \overline{Y}$ such that

$$y^* = Fy^*.$$

Furthermore, the sequence $\{F^k y\}$ converges to y^* for any $y \in \overline{Y}$, and we have

$$\|F^k y - y^*\| \leq \rho^k \|y - y^*\|, \quad k = 1, 2, \dots$$

Proof: Let $y \in \overline{Y}$ and consider the iteration $y_{k+1} = Fy_k$ starting with $y_0 = y$. By the contraction property of F ,

$$\|y_{k+1} - y_k\| \leq \rho \|y_k - y_{k-1}\|, \quad k = 1, 2, \dots,$$

which implies that

$$\|y_{k+1} - y_k\| \leq \rho^k \|y_1 - y_0\|, \quad k = 1, 2, \dots$$

It follows that for every $k \geq 0$ and $m \geq 1$, we have

$$\begin{aligned} \|y_{k+m} - y_k\| &\leq \sum_{i=1}^m \|y_{k+i} - y_{k+i-1}\| \\ &\leq \rho^k (1 + \rho + \dots + \rho^{m-1}) \|y_1 - y_0\| \\ &\leq \frac{\rho^k}{1 - \rho} \|y_1 - y_0\|. \end{aligned}$$

Therefore, $\{y_k\}$ is a Cauchy sequence in \overline{Y} and must converge to a limit $y^* \in \overline{Y}$, since Y is complete and \overline{Y} is closed. We have for all $k \geq 1$,

$$\|Fy^* - y^*\| \leq \|Fy^* - y_k\| + \|y_k - y^*\| \leq \rho \|y^* - y_{k-1}\| + \|y_k - y^*\|$$

and since y_k converges to y^* , we obtain $Fy^* = y^*$. Thus, the limit y^* of y_k is a fixed point of F . It is a unique fixed point because if \tilde{y} were another fixed point, we would have

$$\|y^* - \tilde{y}\| = \|Fy^* - F\tilde{y}\| \leq \rho \|y^* - \tilde{y}\|,$$

which implies that $y^* = \tilde{y}$.

To show the convergence rate bound of the last part, note that

$$\|F^k y - y^*\| = \|F^k y - F y^*\| \leq \rho \|F^{k-1} y - y^*\|.$$

Repeating this process for a total of k times, we obtain the desired result.
Q.E.D.

The convergence rate exhibited by $F^k y$ in the preceding proposition is said to be *geometric*, and $F^k y$ is said to converge to its limit y^* *geometrically*. This is in reference to the fact that the error $\|F^k y - y^*\|$ converges to 0 faster than some geometric progression ($\rho^k \|y - y^*\|$ in this case).

In some contexts of interest to us one may encounter mappings that are not contractions, but become contractions when iterated a finite number of times. In this case, one may use a slightly different version of the contraction mapping fixed point theorem, which we now present.

We say that a function $F : Y \mapsto Y$ is an *m-stage contraction mapping* if there exists a positive integer m and some $\rho < 1$ such that

$$\|F^m y - F^m y'\| \leq \rho \|y - y'\|, \quad \forall y, y' \in Y,$$

where F^m denotes the composition of F with itself m times. Thus, F is an *m-stage contraction* if F^m is a contraction. Again, the scalar ρ is called the modulus of contraction. We have the following generalization of Prop. B.1.

Proposition B.2: (*m*-Stage Contraction Mapping Fixed-Point Theorem) Let Y be a complete vector space and let \overline{Y} be a closed subset of Y . Then if $F : \overline{Y} \mapsto \overline{Y}$ is an *m-stage contraction mapping* with modulus $\rho \in (0, 1)$, there exists a unique $y^* \in \overline{Y}$ such that

$$y^* = F y^*.$$

Furthermore, $\{F^k y\}$ converges to y^* for any $y \in \overline{Y}$.

Proof: Since F^m maps \overline{Y} into \overline{Y} and is a contraction mapping, by Prop. B.1, it has a unique fixed point in \overline{Y} , denoted y^* . Applying F to both sides of the relation $y^* = F^m y^*$, we see that $F y^*$ is also a fixed point of F^m , so by the uniqueness of the fixed point, we have $y^* = F y^*$. Therefore y^* is a fixed point of F . If F had another fixed point, say \tilde{y} , then we would have $\tilde{y} = F^m \tilde{y}$, which by the uniqueness of the fixed point of F^m implies that $\tilde{y} = y^*$. Thus, y^* is the unique fixed point of F .

To show the convergence of $\{F^k y\}$, note that by Prop. B.1, we have for all $y \in \overline{Y}$,

$$\lim_{k \rightarrow \infty} \|F^m y - y^*\| = 0.$$

Using $F^\ell y$ in place of y , we obtain

$$\lim_{k \rightarrow \infty} \|F^{mk+\ell} y - y^*\| = 0, \quad \ell = 0, 1, \dots, m-1,$$

which proves the desired result. **Q.E.D.**

B.2 WEIGHTED SUP-NORM CONTRACTIONS

In this section, we will focus on contraction mappings within a specialized context that is particularly important in DP. Let X be a set (typically the state space in DP), and let $v : X \mapsto \mathbb{R}$ be a positive-valued function,

$$v(x) > 0, \quad \forall x \in X.$$

Let $\mathcal{B}(X)$ denote the set of all functions $J : X \mapsto \mathbb{R}$ such that $J(x)/v(x)$ is bounded as x ranges over X . We define a norm on $\mathcal{B}(X)$, called the *weighted sup-norm*, by

$$\|J\| = \sup_{x \in X} \frac{|J(x)|}{v(x)}. \quad (\text{B.3})$$

It is easily verified that $\|\cdot\|$ thus defined has the required properties for being a norm. Furthermore, $\mathcal{B}(X)$ is *complete under this norm*. To see this, consider a Cauchy sequence $\{J_k\} \subset \mathcal{B}(X)$, and note that $\|J_m - J_n\| \rightarrow 0$ as $m, n \rightarrow \infty$ implies that for all $x \in X$, $\{J_k(x)\}$ is a Cauchy sequence of real numbers, so it converges to some $J^*(x)$. We will show that $J^* \in \mathcal{B}(X)$ and that $\|J_k - J^*\| \rightarrow 0$. To this end, it will be sufficient to show that given any $\epsilon > 0$, there exists an integer K such that

$$\frac{|J_k(x) - J^*(x)|}{v(x)} \leq \epsilon, \quad \forall x \in X, k \geq K.$$

This will imply that

$$\sup_{x \in X} \frac{|J^*(x)|}{v(x)} \leq \epsilon + \|J_k\|, \quad \forall k \geq K,$$

so that $J^* \in \mathcal{B}(X)$, and will also imply that $\|J_k - J^*\| \leq \epsilon$, so that $\|J_k - J^*\| \rightarrow 0$. Assume the contrary, i.e., that there exists an $\epsilon > 0$ and a subsequence $\{x_{m_1}, x_{m_2}, \dots\} \subset X$ such that $m_i < m_{i+1}$ and

$$\epsilon < \frac{|J_{m_i}(x_{m_i}) - J^*(x_{m_i})|}{v(x_{m_i})}, \quad \forall i \geq 1.$$

The right-hand side above is less or equal to

$$\frac{|J_{m_i}(x_{m_i}) - J_n(x_{m_i})|}{v(x_{m_i})} + \frac{|J_n(x_{m_i}) - J^*(x_{m_i})|}{v(x_{m_i})}, \quad \forall n \geq 1, i \geq 1.$$

The first term in the above sum is less than $\epsilon/2$ for i and n larger than some threshold; fixing i and letting n be sufficiently large, the second term can also be made less than $\epsilon/2$, so the sum is made less than ϵ - a contradiction. In conclusion, the space $\mathcal{B}(X)$ is complete, so the fixed point results of Props. B.1 and B.2 apply.

In our discussions, unless we specify otherwise, we will assume that $\mathcal{B}(X)$ is equipped with the weighted sup-norm above, where the weight function v will be clear from the context. There will be frequent occasions where the norm will be unweighted, i.e., $v(x) \equiv 1$ and $\|J\| = \sup_{x \in X} |J(x)|$, in which case we will explicitly state so.

Finite-Dimensional Cases

Let us now focus on the finite-dimensional case $X = \{1, \dots, n\}$, in which case $\mathcal{R}(X)$ and $\mathcal{B}(X)$ can be identified with \mathbb{R}^n . We first consider a linear mapping (cf. Example B.1). We have the following proposition.

Proposition B.3: Consider a linear mapping $F : \mathbb{R}^n \mapsto \mathbb{R}^n$ of the form

$$Fy = b + Ay,$$

where A is an $n \times n$ matrix with components a_{ij} , and b is a vector in \mathbb{R}^n . Denote by $|A|$ the matrix whose components are the absolute values of the components of A and let $\sigma(A)$ and $\sigma(|A|)$ denote the spectral radii of A and $|A|$, respectively. Then:

- (a) $|A|$ has a real eigenvalue λ , which is equal to its spectral radius, and an associated nonnegative eigenvector.
- (b) F is a contraction with respect to some weighted sup-norm if and only if $\sigma(|A|) < 1$. In particular, any substochastic matrix P ($p_{ij} \geq 0$ for all i, j , and $\sum_{j=1}^n p_{ij} \leq 1$, for all i) is a contraction with respect to some weighted sup-norm if and only if $\sigma(P) < 1$.
- (c) F is a contraction with respect to the weighted sup-norm

$$\|y\| = \max_{i=1, \dots, n} \frac{|y_i|}{v(i)}$$

if and only if

$$\frac{\sum_{j=1}^n |a_{ij}| v(j)}{v(i)} < 1, \quad i = 1, \dots, n.$$

Proof: (a) This is the Perron-Frobenius Theorem; see e.g., [BeT89], Chapter 2, Prop. 6.6.

(b) This follows from the Perron-Frobenius Theorem; see [BeT89], Ch. 2, Cor. 6.2.

(c) This is proved in more general form in the following Prop. B.4. **Q.E.D.**

Consider next a nonlinear mapping $F : \mathfrak{R}^n \mapsto \mathfrak{R}^n$ that has the property

$$|Fy - Fz| \leq P |y - z|, \quad \forall y, z \in \mathfrak{R}^n,$$

for some matrix P with nonnegative components and $\sigma(P) < 1$. Here, we generically denote by $|w|$ the vector whose components are the absolute values of the components of w , and the inequality is componentwise. Then we claim that F is a contraction with respect to some weighted sup-norm. To see this note that by the preceding discussion, P is a contraction with respect to some weighted sup-norm $\|y\| = \max_{i=1,\dots,n} |y_i|/v(i)$, and we have

$$\frac{(|Fy - Fz|)(i)}{v(i)} \leq \frac{(P |y - z|)(i)}{v(i)} \leq \alpha \|y - z\|, \quad \forall i = 1, \dots, n,$$

for some $\alpha \in (0, 1)$, where $(|Fy - Fz|)(i)$ and $(P |y - z|)(i)$ are the i th components of the vectors $|Fy - Fz|$ and $P |y - z|$, respectively. Thus, F is a contraction with respect to $\|\cdot\|$. For additional discussion of linear and nonlinear contraction mapping properties and characterizations such as the one above, see the book [OrR70].

Linear Mappings on Countable Spaces

The case where X is countable (or, as a special case, finite) is frequently encountered in DP. The following proposition provides some useful criteria for verifying the contraction property of mappings that are either linear or are obtained via a parametric minimization of other contraction mappings.

Proposition B.4: Let $X = \{1, 2, \dots\}$.

(a) Let $F : \mathcal{B}(X) \mapsto \mathcal{B}(X)$ be a linear mapping of the form

$$(FJ)(i) = b_i + \sum_{j \in X} a_{ij} J(j), \quad i \in X,$$

where b_i and a_{ij} are some scalars. Then F is a contraction with modulus ρ with respect to the weighted sup-norm (B.3) if and only if

$$\frac{\sum_{j \in X} |a_{ij}| v(j)}{v(i)} \leq \rho, \quad i \in X. \quad (\text{B.4})$$

(b) Let $F : \mathcal{B}(X) \mapsto \mathcal{B}(X)$ be a mapping of the form

$$(FJ)(i) = \inf_{\mu \in M} (F_\mu J)(i), \quad i \in X,$$

where M is parameter set, and for each $\mu \in M$, F_μ is a contraction mapping from $\mathcal{B}(X)$ to $\mathcal{B}(X)$ with modulus ρ . Then F is a contraction mapping with modulus ρ .

Proof: (a) Assume that Eq. (B.4) holds. For any $J, J' \in \mathcal{B}(X)$, we have

$$\begin{aligned} \|FJ - FJ'\| &= \sup_{i \in X} \frac{\left| \sum_{j \in X} a_{ij} (J(j) - J'(j)) \right|}{v(i)} \\ &\leq \sup_{i \in X} \frac{\sum_{j \in X} |a_{ij}| v(j) \left(|J(j) - J'(j)| / v(j) \right)}{v(i)} \\ &\leq \sup_{i \in X} \frac{\sum_{j \in X} |a_{ij}| v(j)}{v(i)} \|J - J'\| \\ &\leq \rho \|J - J'\|, \end{aligned}$$

where the last inequality follows from the hypothesis.

Conversely, arguing by contradiction, let's assume that Eq. (B.4) is violated for some $i \in X$. Define $J(j) = v(j) \operatorname{sgn}(a_{ij})$ and $J'(j) = 0$ for all $j \in X$. Then we have $\|J - J'\| = \|J\| = 1$, and

$$\frac{|(FJ)(i) - (FJ')(i)|}{v(i)} = \frac{\sum_{j \in X} |a_{ij}| v(j)}{v(i)} > \rho = \rho \|J - J'\|,$$

showing that F is not a contraction of modulus ρ .

(b) Since F_μ is a contraction of modulus ρ , we have for any $J, J' \in \mathcal{B}(X)$,

$$\frac{(F_\mu J)(i)}{v(i)} \leq \frac{(F_\mu J')(i)}{v(i)} + \rho \|J - J'\|, \quad i \in X,$$

so by taking the infimum over $\mu \in M$,

$$\frac{(FJ)(i)}{v(i)} \leq \frac{(FJ')(i)}{v(i)} + \rho \|J - J'\|, \quad i \in X.$$

Reversing the roles of J and J' , we obtain

$$\frac{|(FJ)(i) - (FJ')(i)|}{v(i)} \leq \rho \|J - J'\|, \quad i \in X,$$

and by taking the supremum over i , the contraction property of F is proved.

Q.E.D.

The preceding proposition assumes that $FJ \in \mathcal{B}(X)$ for all $J \in \mathcal{B}(X)$. The following proposition provides conditions, particularly relevant to the DP context, which imply this assumption.

Proposition B.5: Let $X = \{1, 2, \dots\}$, let M be a parameter set, and for each $\mu \in M$, let F_μ be a linear mapping of the form

$$(F_\mu J)(i) = b_i(\mu) + \sum_{j \in X} a_{ij}(\mu) J(j), \quad i \in X.$$

- (a) We have $F_\mu J \in \mathcal{B}(X)$ for all $J \in \mathcal{B}(X)$ provided $b(\mu) \in \mathcal{B}(X)$ and $V(\mu) \in \mathcal{B}(X)$, where

$$b(\mu) = \{b_1(\mu), b_2(\mu), \dots\}, \quad V(\mu) = \{V_1(\mu), V_2(\mu), \dots\},$$

with

$$V_i(\mu) = \sum_{j \in X} |a_{ij}(\mu)| v(j), \quad i \in X.$$

- (b) Consider the mapping F

$$(FJ)(i) = \inf_{\mu \in M} (F_\mu J)(i), \quad i \in X.$$

We have $FJ \in \mathcal{B}(X)$ for all $J \in \mathcal{B}(X)$, provided $b \in \mathcal{B}(X)$ and $V \in \mathcal{B}(X)$, where

$$b = \{b_1, b_2, \dots\}, \quad V = \{V_1, V_2, \dots\},$$

with $b_i = \sup_{\mu \in M} b_i(\mu)$ and $V_i = \sup_{\mu \in M} V_i(\mu)$.

Proof: (a) For all $\mu \in M$, $J \in \mathcal{B}(X)$ and $i \in X$, we have

$$(F_\mu J)(i) \leq |b_i(\mu)| + \sum_{j \in X} |a_{ij}(\mu)| |J(j)/v(j)| v(j)$$

$$\begin{aligned}
&\leq |b_i(\mu)| + \|J\| \sum_{j \in X} |a_{ij}(\mu)| v(j) \\
&= |b_i(\mu)| + \|J\| V_i(\mu),
\end{aligned}$$

and similarly $(F_\mu J)(i) \geq -|b_i(\mu)| - \|J\| V_i(\mu)$. Thus

$$|(F_\mu J)(i)| \leq |b_i(\mu)| + \|J\| V_i(\mu), \quad i \in X.$$

By dividing this inequality with $v(i)$ and by taking the supremum over $i \in X$, we obtain

$$\|F_\mu J\| \leq \|b_\mu\| + \|J\| \|V_\mu\| < \infty.$$

(b) By doing the same as in (a), but after first taking the infimum of $(F_\mu J)(i)$ over μ , we obtain

$$\|FJ\| \leq \|b\| + \|J\| \|V\| < \infty.$$

Q.E.D.

References

- [ABB02] Abounadi, J., Bertsekas, B. P., and Borkar, V. S., 2002. “Stochastic Approximation for Non-Expansive Maps: Q-Learning Algorithms,” *SIAM J. on Control and Opt.*, Vol. 41, pp. 1-22.
- [AnM79] Anderson, B. D. O., and Moore, J. B., 1979. *Optimal Filtering*, Prentice Hall, Englewood Cliffs, N. J.
- [BBB08] Basu, A., Bhattacharyya, and Borkar, V., 2008. “A Learning Algorithm for Risk-Sensitive Cost,” *Math. of OR*, Vol. 33, pp. 880-898.
- [BBD10] Busoniu, L., Babuska, R., De Schutter, B., and Ernst, D., 2010. *Reinforcement Learning and Dynamic Programming Using Function Approximators*, CRC Press, N. Y.
- [Bau78] Baudet, G. M., 1978. “Asynchronous Iterative Methods for Multiprocessors,” *Journal of the ACM*, Vol. 25, pp. 226-244.
- [BeI96] Bertsekas, D. P., and Ioffe, S., 1996. “Temporal Differences-Based Policy Iteration and Applications in Neuro-Dynamic Programming,” *Lab. for Info. and Decision Systems Report LIDS-P-2349*, MIT.
- [BeS78] Bertsekas, D. P., and Shreve, S. E., 1978. *Stochastic Optimal Control: The Discrete Time Case*, Academic Press, N. Y.; may be downloaded from <http://web.mit.edu/dimitrib/www/home.html>
- [BeT89] Bertsekas, D. P., and Tsitsiklis, J. N., 1989. *Parallel and Distributed Computation: Numerical Methods*, Prentice-Hall, Engl. Cliffs, N. J.; may be downloaded from <http://web.mit.edu/dimitrib/www/home.html>
- [BeT91] Bertsekas, D. P., and Tsitsiklis, J. N., 1991. “An Analysis of Stochastic Shortest Path Problems,” *Math. of OR*, Vol. 16, pp. 580-595.
- [BeT96] Bertsekas, D. P., and Tsitsiklis, J. N., 1996. *Neuro-Dynamic Programming*, Athena Scientific, Belmont, MA.
- [BeT08] Bertsekas, D. P., and Tsitsiklis, J. N., 2008. *Introduction to Probability*, 2nd Ed., Athena Scientific, Belmont, MA.
- [BeY07] Bertsekas, D. P., and Yu, H., 2007. “Solution of Large Systems of Equations Using Approximate Dynamic Programming Methods,” *Lab. for Info. and Decision Systems Report LIDS-P-2754*, MIT.
- [BeY09] Bertsekas, D. P., and Yu, H., 2009. “Projected Equation Methods for Approximate Solution of Large Linear Systems,” *J. of Computational and Applied Mathematics*, Vol. 227, pp. 27-50.
- [BeY10] Bertsekas, D. P., and Yu, H., 2010. “Asynchronous Distributed Policy Iteration in Dynamic Programming,” *Proc. of Allerton Conf. on Communication, Control and Computing*, Allerton Park, Ill, pp. 1368-1374.

- [BeY12] Bertsekas, D. P., and Yu, H., 2012. “Q-Learning and Enhanced Policy Iteration in Discounted Dynamic Programming,” *Math. of OR*, Vol. 37, pp. 66-94.
- [BeY16] Bertsekas, D. P., and Yu, H., 2016. “Stochastic Shortest Path Problems Under Weak Conditions,” *Lab. for Information and Decision Systems Report LIDS-2909*, January 2016.
- [Ber71] Bertsekas, D. P., 1971. “Control of Uncertain Systems With a Set-Membership Description of the Uncertainty,” Ph.D. Dissertation, Massachusetts Institute of Technology, Cambridge, MA (available from the author’s website).
- [Ber72] Bertsekas, D. P., 1972. “Infinite Time Reachability of State Space Regions by Using Feedback Control,” *IEEE Trans. Aut. Control*, Vol. AC-17, pp. 604-613.
- [Ber75] Bertsekas, D. P., 1975. “Monotone Mappings in Dynamic Programming,” 1975 IEEE Conference on Decision and Control, pp. 20-25.
- [Ber77] Bertsekas, D. P., 1977. “Monotone Mappings with Application in Dynamic Programming,” *SIAM J. on Control and Opt.*, Vol. 15, pp. 438-464.
- [Ber82] Bertsekas, D. P., 1982. “Distributed Dynamic Programming,” *IEEE Trans. Aut. Control*, Vol. AC-27, pp. 610-616.
- [Ber83] Bertsekas, D. P., 1983. “Asynchronous Distributed Computation of Fixed Points,” *Math. Programming*, Vol. 27, pp. 107-120.
- [Ber87] Bertsekas, D. P., 1987. *Dynamic Programming: Deterministic and Stochastic Models*, Prentice-Hall, Englewood Cliffs, N. J.
- [Ber96] Bertsekas, D. P., 1996. Lecture at NSF Workshop on Reinforcement Learning, Hilltop House, Harper’s Ferry, N. Y.
- [Ber98] Bertsekas, D. P., 1998. *Network Optimization: Continuous and Discrete Models*, Athena Scientific, Belmont, MA.
- [Ber09] Bertsekas, D. P., 2009. *Convex Optimization Theory*, Athena Scientific, Belmont, MA.
- [Ber10] Bertsekas, D. P., 2010. “Williams-Baird Counterexample for Q-Factor Asynchronous Policy Iteration,” http://web.mit.edu/dimitrib/www/Williams-Baird_Counterexample.pdf
- [Ber11a] Bertsekas, D. P., 2011. “Temporal Difference Methods for General Projected Equations,” *IEEE Trans. Aut. Control*, Vol. 56, pp. 2128-2139.
- [Ber11b] Bertsekas, D. P., 2011. “ λ -Policy Iteration: A Review and a New Implementation,” *Lab. for Info. and Decision Systems Report LIDS-P-2874*, MIT; appears in *Reinforcement Learning and Approximate Dynamic Programming for Feedback Control*, by F. Lewis and D. Liu (eds.), IEEE Press, 2012.
- [Ber11c] Bertsekas, D. P., 2011. “Approximate Policy Iteration: A Survey and Some New Methods,” *J. of Control Theory and Applications*, Vol. 9, pp. 310-335; a somewhat expanded version appears as *Lab. for Info. and Decision Systems Report LIDS-2833*, MIT, 2011.
- [Ber12a] Bertsekas, D. P., 2012. *Dynamic Programming and Optimal Control*, Vol. II, 4th Edition: *Approximate Dynamic Programming*, Athena Scientific, Belmont, MA.
- [Ber12b] Bertsekas, D. P., 2012. “Weighted Sup-Norm Contractions in Dynamic Programming: A Review and Some New Applications,” *Lab. for Info. and Decision Systems Report LIDS-P-2884*, MIT.

- [Ber14] Bertsekas, D. P., 2014. “Robust Shortest Path Planning and Semicontractive Dynamic Programming,” Lab. for Information and Decision Systems Report LIDS-P-2915, MIT, Feb. 2014 (revised Jan. 2015 and June 2016); arXiv preprint arXiv:1608.01670; to appear in Naval Research Logistics.
- [Ber15] Bertsekas, D. P., 2015. “Regular Policies in Abstract Dynamic Programming,” Lab. for Information and Decision Systems Report LIDS-P-3173, MIT, May 2015; arXiv preprint arXiv:1609.03115; to appear in SIAM J. on Control and Optimization.
- [Ber16a] Bertsekas, D. P., 2016. “Affine Monotonic and Risk-Sensitive Models in Dynamic Programming,” Lab. for Information and Decision Systems Report LIDS-3204, MIT, June 2016; arXiv preprint arXiv:1608.01393.
- [Ber16b] Bertsekas, D. P., 2016. “Proximal Algorithms and Temporal Differences for Large Linear Systems: Extrapolation, Approximation, and Simulation,” Report LIDS-P-3205, MIT, Oct. 2016; arXiv preprint arXiv:1610.1610.05427.
- [Ber16c] Bertsekas, D. P., 2016. Nonlinear Programming, 3rd Edition, Athena Scientific, Belmont, MA.
- [Ber17a] Bertsekas, D. P., 2017. Dynamic Programming and Optimal Control, Vol. I, 4th Edition, Athena Scientific, Belmont, MA.
- [Ber17b] Bertsekas, D. P., 2017. “Value and Policy Iteration in Deterministic Optimal Control and Adaptive Dynamic Programming,” IEEE Transactions on Neural Networks and Learning Systems, Vol. 28, pp. 500-509.
- [Ber17c] Bertsekas, D. P., 2017. “Stable Optimal Control and Semicontractive Dynamic Programming,” Report LIDS-P-3506, MIT, May 2017; to appear in SIAM J. on Control and Optimization.
- [Ber17d] Bertsekas, D. P., 2017. “Proper Policies in Infinite-State Stochastic Shortest Path Problems”, Report LIDS-P-3507, MIT, May 2017; to appear in IEEE Transactions on Aut. Control.
- [Ber17e] Bertsekas, D. P., 2017. “Proximal Algorithms and Temporal Differences for Solving Fixed Point Problems,” to appear in Computational Optimization and Applications J.
- [Bla65] Blackwell, D., 1965. “Positive Dynamic Programming,” Proc. Fifth Berkeley Symposium Math. Statistics and Probability, pp. 415-418.
- [BoM99] Borkar, V. S., Meyn, S. P., 1999. “Risk Sensitive Optimal Control: Existence and Synthesis for Models with Unbounded Cost,” SIAM J. Control and Opt., Vol. 27, pp. 192-209.
- [BoM00] Borkar, V. S., Meyn, S. P., 1990. “The O.D.E. Method for Convergence of Stochastic Approximation and Reinforcement Learning,” SIAM J. Control and Opt., Vol. 38, pp. 447-469.
- [BoM02] Borkar, V. S., Meyn, S. P., 2002. “Risk-Sensitive Optimal Control for Markov Decision Processes with Monotone Cost,” Math. of OR, Vol. 27, pp. 192-209.
- [Bor98] Borkar, V. S., 1998. “Asynchronous Stochastic Approximation,” SIAM J. Control Opt., Vol. 36, pp. 840-851.
- [Bor08] Borkar, V. S., 2008. Stochastic Approximation: A Dynamical Systems Viewpoint, Cambridge Univ. Press, N. Y.

- [CFH07] Chang, H. S., Fu, M. C., Hu, J., Marcus, S. I., 2007. *Simulation-Based Algorithms for Markov Decision Processes*, Springer, N. Y.
- [CaM88] Carraway, R. L., and Morin, T. L., 1988. "Theory and Applications of Generalized Dynamic Programming: An Overview," *Computers and Mathematics with Applications*, Vol. 16, pp. 779-788.
- [CaR13] Canbolat, P. G., and Rothblum, U. G., 2013. "(Approximate) Iterated Successive Approximations Algorithm for Sequential Decision Processes," *Annals of Operations Research*, Vol. 208, pp. 309-320.
- [Cao07] Cao, X. R., 2007. *Stochastic Learning and Optimization: A Sensitivity-Based Approach*, Springer, N. Y.
- [ChM69] Chazan D., and Miranker, W., 1969. "Chaotic Relaxation," *Linear Algebra and Applications*, Vol. 2, pp. 199-222.
- [ChS87] Chung, K.-J., and Sobel, M. J., 1987. "Discounted MDPs: Distribution Functions and Exponential Utility Maximization," *SIAM J. Control and Opt.*, Vol. 25, pp. 49-62.
- [CoM99] Coraluppi, S. P., and Marcus, S. I., 1999. "Risk-Sensitive and Minimax Control of Discrete-Time, Finite-State Markov Decision Processes," *Automatica*, Vol. 35, pp. 301-309.
- [DFV00] de Farias, D. P., and Van Roy, B., 2000. "On the Existence of Fixed Points for Approximate Value Iteration and Temporal-Difference Learning," *J. of Optimization Theory and Applications*, Vol. 105, pp. 589-608.
- [DeM67] Denardo, E. V., and Mitten, L. G., 1967. "Elements of Sequential Decision Processes," *J. Indust. Engrg.*, Vol. 18, pp. 106-112.
- [DeR79] Denardo, E. V., and Rothblum, U. G., 1979. "Optimal Stopping, Exponential Utility, and Linear Programming," *Math. Programming*, Vol. 16, pp. 228-244.
- [Den67] Denardo, E. V., 1967. "Contraction Mappings in the Theory Underlying Dynamic Programming," *SIAM Review*, Vol. 9, pp. 165-177.
- [Der70] Derman, C., 1970. *Finite State Markovian Decision Processes*, Academic Press, N. Y.
- [DuS65] Dubins, L., and Savage, L. M., 1965. *How to Gamble If You Must*, McGraw-Hill, N. Y.
- [FeM97] Fernandez-Gaucherand, E., and Marcus, S. I., 1997. "Risk-Sensitive Optimal Control of Hidden Markov Models: Structural Results," *IEEE Trans. Aut. Control*, Vol. AC-42, pp. 1418-1422.
- [Fei02] Feinberg, E. A., 2002. "Total Reward Criteria," in E. A. Feinberg and A. Schwartz, (Eds.), *Handbook of Markov Decision Processes*, Springer, N. Y.
- [FiV96] Filar, J., and Vrieze, K., 1996. *Competitive Markov Decision Processes*, Springer, N. Y.
- [FIM95] Fleming, W. H., and McEneaney, W. M., 1995. "Risk-Sensitive Control on an Infinite Time Horizon," *SIAM J. Control and Opt.*, Vol. 33, pp. 1881-1915.
- [Gos03] Gosavi, A., 2003. *Simulation-Based Optimization: Parametric Optimization Techniques and Reinforcement Learning*, Springer, N. Y.
- [GuS17] Guillot, M., and Stauffer, G., 2017. "The Stochastic Shortest Path Problem: A Polyhedral Combinatorics Perspective," *Univ. of Grenoble Report*.

- [HCP99] Hernandez-Lerma, O., Carrasco, O., and Perez-Hernandez. 1999. "Markov Control Processes with the Expected Total Cost Criterion: Optimality, Stability, and Transient Models," *Acta Appl. Math.*, Vol. 59, pp. 229-269.
- [Hay08] Haykin, S., 2008. *Neural Networks and Learning Machines*, (3rd Edition), Prentice-Hall, Englewood-Cliffs, N. J.
- [HeL99] Hernandez-Lerma, O., and Lasserre, J. B., 1999. *Further Topics on Discrete-Time Markov Control Processes*, Springer, N. Y.
- [HeM96] Hernandez-Hernandez, D., and Marcus, S. I., 1996. "Risk Sensitive Control of Markov Processes in Countable State Space," *Systems and Control Letters*, Vol. 29, pp. 147-155.
- [HiW05] Hinderer, K., and Waldmann, K.-H., 2005. "Algorithms for Countable State Markov Decision Models with an Absorbing Set," *SIAM J. of Control and Opt.*, Vol. 43, pp. 2109-2131.
- [HoM72] Howard, R. S., and Matheson, J. E., 1972. "Risk-Sensitive Markov Decision Processes," *Management Science*, Vol. 8, pp. 356-369.
- [JBE94] James, M. R., Baras, J. S., Elliott, R. J., 1994. "Risk-Sensitive Control and Dynamic Games for Partially Observed Discrete-Time Nonlinear Systems," *IEEE Trans. Aut. Control*, Vol. AC-39, pp. 780-792.
- [JaC06] James, H. W., and Collins, E. J., 2006. "An Analysis of Transient Markov Decision Processes," *J. Appl. Prob.*, Vol. 43, pp. 603-621.
- [Jac73] Jacobson, D. H., 1973. "Optimal Stochastic Linear Systems with Exponential Performance Criteria and their Relation to Deterministic Differential Games," *IEEE Transactions on Automatic Control*, Vol. AC-18, pp. 124-131.
- [Kal60] Kalman, R. E., 1960. "Contributions to the Theory of Optimal Control," *Bol. Soc. Mat. Mexicana*, Vol. 5, pp. 102-119.
- [Kuc72] Kucera, V., 1972. "The Discrete Riccati Equation of Optimal Control," *Kybernetika*, Vol. 8, pp. 430-447.
- [Kuc73] Kucera, V., 1973. "A Review of the Matrix Riccati Equation," *Kybernetika*, Vol. 9, pp. 42-61.
- [LaR95] Lancaster, P., and Rodman, L., 1995. *Algebraic Riccati Equations*, Clarendon Press, Oxford, UK.
- [Mey07] Meyn, S., 2007. *Control Techniques for Complex Networks*, Cambridge Univ. Press, N. Y.
- [Mit64] Mitten, L. G., 1964. "Composition Principles for Synthesis of Optimal Multistage Processes," *Operations Research*, Vol. 12, pp. 610-619.
- [Mit74] Mitten, L. G., 1964. "Preference Order Dynamic Programming," *Management Science*, Vol. 21, pp. 43 - 46.
- [Mor82] Morin, T. L., 1982. "Monotonicity and the Principle of Optimality," *J. of Math. Analysis and Applications*, Vol. 88, pp. 665-674.
- [OrR70] Ortega, J. M., and Rheinboldt, W. C., 1970. *Iterative Solution of Nonlinear Equations in Several Variables*, Academic Press, N. Y.
- [PaB99] Patek, S. D., and Bertsekas, D. P., 1999. "Stochastic Shortest Path Games," *SIAM J. on Control and Opt.*, Vol. 36, pp. 804-824.
- [Pal67] Pallu de la Barriere, R., 1967. *Optimal Control Theory*, Saunders, Phila; republished by Dover, N. Y., 1980.

- [Pat01] Patek, S. D., 2001. "On Terminating Markov Decision Processes with a Risk Averse Objective Function," *Automatica*, Vol. 37, pp. 1379-1386.
- [Pat07] Patek, S. D., 2007. "Partially Observed Stochastic Shortest Path Problems with Approximate Solution by Neuro-Dynamic Programming," *IEEE Trans. on Systems, Man, and Cybernetics Part A*, Vol. 37, pp. 710-720.
- [Pli78] Pliska, S. R., 1978. "On the Transient Case for Markov Decision Chains with General State Spaces," in *Dynamic Programming and its Applications*, by M. L. Puterman (ed.), Academic Press, N. Y.
- [Pow07] Powell, W. B., 2007. *Approximate Dynamic Programming: Solving the Curses of Dimensionality*, J. Wiley and Sons, Hoboken, N. J; 2nd ed., 2011.
- [Put94] Puterman, M. L., 1994. *Markovian Decision Problems*, J. Wiley, N. Y.
- [Rei16] Reissig, G., 2016. "Approximate Value Iteration for a Class of Deterministic Optimal Control Problems with Infinite State and Input Alphabets," *Proc. 2016 IEEE Conf. on Decision and Control*, pp. 1063-1068.
- [Roc70] Rockafellar, R. T., 1970. *Convex Analysis*, Princeton Univ. Press, Princeton, N. J.
- [Ros67] Rosenfeld, J., 1967. "A Case Study on Programming for Parallel Processors," Research Report RC-1864, IBM Res. Center, Yorktown Heights, N. Y.
- [Rot79] Rothblum, U. G., 1979. "Iterated Successive Approximation for Sequential Decision Processes," in *Stochastic Control and Optimization*, by J. W. B. van Overhagen and H. C. Tijms (eds), Vrije University, Amsterdam.
- [Rot84] Rothblum, U. G., 1984. "Multiplicative Markov Decision Chains," *Math. of OR*, Vol. 9, pp. 6-24.
- [ScL12] Scherrer, B., and Lesner, B., 2012. "On the Use of Non-Stationary Policies for Stationary Infinite-Horizon Markov Decision Processes," *NIPS 2012 - Neural Information Processing Systems*, South Lake Tahoe, Ne.
- [Sch75] Schal, M., 1975. "Conditions for Optimality in Dynamic Programming and for the Limit of n -Stage Optimal Policies to be Optimal," *Z. Wahrscheinlichkeitstheorie und Verw. Gebiete*, Vol. 32, pp. 179-196.
- [Sch11] Scherrer, B., 2011. "Performance Bounds for Lambda Policy Iteration and Application to the Game of Tetris," Report RR-6348, INRIA, France; *J. of Machine Learning Research*, Vol. 14, 2013, pp. 1181-1227.
- [Sch12] Scherrer, B., 2012. "On the Use of Non-Stationary Policies for Infinite-Horizon Discounted Markov Decision Processes," *INRIA Lorraine Report*, France.
- [Sha53] Shapley, L. S., 1953. "Stochastic Games," *Proc. Nat. Acad. Sci. U.S.A.*, Vol. 39.
- [Sob75] Sobel, M. J., 1975. "Ordinal Dynamic Programming," *Management Science*, Vol. 21, pp. 967-975.
- [Str66] Strauch, R., 1966. "Negative Dynamic Programming," *Ann. Math. Statist.*, Vol. 37, pp. 871-890.
- [SuB98] Sutton, R. S., and Barto, A. G., 1998. *Reinforcement Learning*, MIT Press, Cambridge, MA.
- [Sze98a] Szepesvari, C., 1998. *Static and Dynamic Aspects of Optimal Sequential Decision Making*, Ph.D. Thesis, Bolyai Institute of Mathematics, Hungary.

- [Sze98b] Szepesvari, C., 1998. "Non-Markovian Policies in Sequential Decision Problems," *Acta Cybernetica*, Vol. 13, pp. 305-318.
- [Sze10] Szepesvari, C., 2010. *Algorithms for Reinforcement Learning*, Morgan and Claypool Publishers, San Francisco, CA.
- [TBA86] Tsitsiklis, J. N., Bertsekas, D. P., and Athans, M., 1986. "Distributed Asynchronous Deterministic and Stochastic Gradient Optimization Algorithms," *IEEE Trans. Aut. Control*, Vol. AC-31, pp. 803-812.
- [ThS10a] Thiery, C., and Scherrer, B., 2010. "Least-Squares λ -Policy Iteration: Bias-Variance Trade-off in Control Problems," in *ICML'10: Proc. of the 27th Annual International Conf. on Machine Learning*.
- [ThS10b] Thiery, C., and Scherrer, B., 2010. "Performance Bound for Approximate Optimistic Policy Iteration," Technical Report, INRIA, France.
- [Tsi94] Tsitsiklis, J. N., 1994. "Asynchronous Stochastic Approximation and Q-Learning," *Machine Learning*, Vol. 16, pp. 185-202.
- [VVL13] Vrabie, V., Vamvoudakis, K. G., and Lewis, F. L., 2013. *Optimal Adaptive Control and Differential Games by Reinforcement Learning Principles*, The Institution of Engineering and Technology, London.
- [VeP87] Verdu, S., and Poor, H. V., 1987. "Abstract Dynamic Programming Models under Commutativity Conditions," *SIAM J. on Control and Opt.*, Vol. 25, pp. 990-1006.
- [Wat89] Watkins, C. J. C. H., *Learning from Delayed Rewards*, Ph.D. Thesis, Cambridge Univ., England.
- [Whi80] Whittle, P., 1980. "Stability and Characterization Conditions in Negative Programming," *Journal of Applied Probability*, Vol. 17, pp. 635-645.
- [Whi82] Whittle, P., 1982. *Optimization Over Time*, Wiley, N. Y., Vol. 1, 1982, Vol. 2, 1983.
- [Whi90] Whittle, P., 1990. *Risk-Sensitive Optimal Control*, Wiley, Chichester.
- [WiB93] Williams, R. J., and Baird, L. C., 1993. "Analysis of Some Incremental Variants of Policy Iteration: First Steps Toward Understanding Actor-Critic Learning Systems," Report NU-CCS-93-11, College of Computer Science, Northeastern University, Boston, MA.
- [Wil71] Willems, J., 1971. "Least Squares Stationary Optimal Control and the Algebraic Riccati Equation," *IEEE Trans. on Automatic Control*, Vol. 16, pp. 621-634.
- [YuB10] Yu, H., and Bertsekas, D. P., 2010. "Error Bounds for Approximations from Projected Linear Equations," *Math. of OR*, Vol. 35, pp. 306-329.
- [YuB12] Yu, H., and Bertsekas, D. P., 2012. "Weighted Bellman Equations and their Applications in Dynamic Programming," Lab. for Info. and Decision Systems Report LIDS-P-2876, MIT.
- [YuB13a] Yu, H., and Bertsekas, D. P., 2013. "Q-Learning and Policy Iteration Algorithms for Stochastic Shortest Path Problems," *Annals of Operations Research*, Vol. 208, pp. 95-132.
- [YuB13b] Yu, H., and Bertsekas, D. P., 2013. "On Boundedness of Q-Learning Iterates for Stochastic Shortest Path Problems," *Math. of OR*, Vol. 38, pp. 209-227.

- [YuB15] Yu, H., and Bertsekas, D. P., 2015. "A Mixed Value and Policy Iteration Method for Stochastic Control with Universally Measurable Policies," *Math. of OR*, Vol. 40, pp. 926-968.
- [Yu11] Yu, H., 2011. "Stochastic Shortest Path Games and Q-Learning," Lab. for Info. and Decision Systems Report LIDS-P-2875, MIT.
- [Yu15] Yu, H., 2015. "On Convergence of Value Iteration for a Class of Total Cost Markov Decision Processes," *SIAM J. on Control and Optimization*, Vol. 53, pp. 1982-2016.
- [Zac64] Zachrisson, L. E., 1964. "Markov Games," in *Advances in Game Theory*, by M. Dresher, L. S. Shapley, and A. W. Tucker, (eds.), Princeton Univ. Press, Princeton, N. J., pp. 211-253.

INDEX

A

Abstraction, 31
Affine monotonic model, 19, 170, 198, 204, 213
Aggregation, 20
Aggregation, distributed, 23
Aggregation, multistep, 28
Aggregation equation, 26
Aggregation probability, 21
Approximate DP, 25, 31
Approximation models, 24, 49
Asynchronous algorithms, 23, 31, 77, 98, 195, 205, 243
Asynchronous convergence theorem, 80, 100
Asynchronous policy iteration, 23, 84, 89, 92, 94, 95, 98, 195, 205
Asynchronous value iteration, 77, 98, 195, 205, 243

B

Bellman's equation, 6, 40, 107, 136, 221, 223, 230, 234, 277, 280, 297, 299, 302
Blackmailer's dilemma, 115
Box condition, 81

C

Cauchy sequence, 326
Complete space, 326
Composition of mappings, 323
Continuous-state optimal control, 191, 205, 210, 211, 257, 260, 266, 291, 307, 315
Contraction assumption, 8, 41
Contraction mapping, 8, 33, 319, 325, 329
Contraction mapping fixed-point theorem, 41, 325, 327, 328, 330-333
Contractive models, 29, 39
Controllability, 118, 204, 272, 281, 307
Convergent models, 262, 307
Cost function, 40, 127

D

Disaggregation probability, 20
Discounted MDP, 12, 260, 316
Distributed aggregation, 23, 24
Distributed computation, 23, 31

E

ϵ -optimal policy, 43, 218, 222, 225, 228, 239, 263, 274
Error amplification, 55
Error bounds, 45, 47, 50, 54, 59, 62, 71
Euclidean norm, 326
Exponential cost model, 171, 173, 176, 205

F

Finite-horizon problems, 219
First passage problem, 16
Fixed point, 325

G

Games, dynamic, 13, 95
Gauss-Seidel method, 78, 98
Geometric convergence rate, 328

H

Hard aggregation, 21

I

Imperfect state information, 206
Improper policy, 16, 112, 113, 164, 181
Interpolated mappings, 319
Interpolation, 95, 203

J, K

L

λ -aggregation, 27
 λ -policy iteration, 27, 63, 76, 97, 146, 245, 305
LSPE(λ), 27
LSTD(λ), 27
Least squares approximation, 55
Limited lookahead policy, 47

Linear contraction mappings, 325, 331
 Linear-quadratic problems, 118, 189, 282, 307, 311

M

MDP, 10, 12,
 Markovian decision problem, *see* MDP
 Mathematical programming, 103, 148, 309
 Minimax problems, 15, 95, 179, 198
 Modulus of contraction, 325
 Monotone mapping, 323
 Monotone decreasing model, 226, 304
 Monotone fixed point iterations, 316, 317
 Monotone increasing model, 225, 255, 304
 Monotonicity assumption, 7, 40, 126
 Multiplicative model, 18, 171
 Multistage lookahead, 49
 Multistep aggregation, 28
 Multistep mapping, 27, 33, 34, 36
 Multistep methods, 27, 33, 34

N

N -stage optimal policy, 218
 Negative cost DP model, 32, 226, 304
 Neural networks, 25
 Neuro-dynamic programming, 25
 Noncontractive model, 30, 217
 Nonmonotonic-contractive model, 74, 101
 Nonstationary policy, 40, 44

O

ODE approach, 98
 Oblique projection, 28
 Observability, 118, 204, 281, 307
 Optimality conditions, 42, 131, 150, 166, 168, 176, 187, 194, 220, 236, 256, 277, 280, 297

P

p - ϵ -optimality, 274
 p -stable policy, 270
 Parallel computation, 78
 Partially asynchronous algorithms, 80
 Periodic policies, 50, 96, 99

Perturbations, 155, 169, 190, 212, 213, 270, 293, 313
 Policy, 5, 40
 Policy, contractive, 174
 Policy evaluation, 56, 63, 64, 84
 Policy improvement, 56, 84, 137
 Policy iteration, 9, 56, 84, 89, 136, 191, 242, 246, 247, 282, 285
 Policy iteration, approximate, 59, 103
 Policy iteration, asynchronous, 84, 90, 94, 99, 196, 205
 Policy iteration, constrained, 23
 Policy iteration, convergence, 56
 Policy iteration, modified, 96
 Policy iteration, optimistic, 63, 65, 70, 85, 89, 94, 95, 96, 144, 244, 290
 Policy iteration, perturbations, 158, 169, 204, 287
 Policy, noncontractive, 174
 Policy, terminating, 181, 193, 272
 Positive cost DP model, 32, 226, 304
 Projected Bellman equation, 25
 Projected equation, 25
 Proper policy, 16, 111, 113, 164, 181, 293, 307, 316
 Proximal algorithm, 26, 245, 248
 Proximal mapping, 27, 34, 245, 248

Q

Q-factor, 89
 Q-learning, 98

R

Reachability, 310
 Reduced space implementation, 93
 Regular, *see* S -regular
 Reinforcement learning, 25
 Risk-sensitive model, 18
 Robust SSP, 179, 205
 Rollout, 25

S

SSP problems, 15, 113, 162, 203, 204, 247, 291, 307
 S -irregular policy, 106, 128, 149, 155
 S -regular collection, 249
 S -regular policy, 106, 128
 Search problems, 155
 Self-learning, 25

Semi-Markov problem, 13
 Seminorm projection, 28
 Semicontinuity conditions, 165
 Semicontractive model, 29, 106, 125, 203
 Shortest path problem, 15, 17, 111, 161, 267, 291, 312
 Simulation, 28, 31, 78, 98
 Spectral radius, 325
 Stable policies, 119, 261, 266, 270, 273, 282, 307
 Stationary policy, 40, 44
 Stochastic shortest path problems, see SSP problems
 Stopping problems, 90, 94, 283
 Strong PI property, 140
 Strong SSP conditions, 165
 Synchronous convergence condition, 81

T

TD(λ), 27
 Temporal differences, 26, 27, 245
 Terminating policy, 193, 210, 211, 269
 Totally asynchronous algorithms, 81
 Transient programming problem, 16

U

Uniform fixed point, 89
 Uniformly N -stage optimal policy, 219
 Uniformly proper policy, 301, 307, 316
 Unit function, 323

V

Value iteration, 9, 52, 53, 77, 98, 134, 166, 168, 176, 178, 187, 191, 194, 195, 205, 240, 243, 252, 255, 258, 261, 266, 277, 279, 280, 297, 302, 304, 317, 318
 Value iteration, asynchronous, 77, 98, 195, 205, 243

W

Weak PI property, 138
 Weak SSP conditions, 167
 Weighted Bellman equation, 38
 Weighted Euclidean norm, 25, 326
 Weighted multistep mapping, 38
 Weighted sup norm, 41, 329
 Weighted sup-norm contraction, 90, 96, 329
 Well-behaved region, 131, 251

X, Y

Z

Zero-sum games, 13, 95

