MA5233 Computational Mathematics

Lecture 10: Runge-Kutta Methods

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Introduction

The aim of this lecture is to develop numerical methods for solving problems of the following form.

Def: Ordinary differential equation (ODE)

Given $f: \mathbb{R}^n \to \mathbb{R}^n$, $y_0 \in \mathbb{R}^n$ and T > 0, find $y: [0, T) \to \mathbb{R}^n$ such that

$$y(0) = y_0$$
 and $\dot{y}(t) = f(y(t))$ for all $t \in [0, T)$.

 $\dot{y}=y'=rac{dy}{dt}$ is a shorthand notation for taking derivatives of a function $y:\mathbb{R} o\mathbb{R}^n$.

Outlook

The following slides illustrate the above definition by discussing three example ODEs. These ODEs will reappear later in this lecture, so pay attention even if you are already familiar with ODEs.

Example 1

Consider the problem of finding $y:[0,\infty)\to\mathbb{R}$ such that

$$y(0) = y_0$$
 and $\dot{y}(t) = \lambda y(t)$

for some given $y_0, \lambda \in \mathbb{R}$.

The solution to this problem is given by

$$y(t) = y_0 \exp(\lambda t)$$

because this function satisfies

$$y(0) = y_0 \exp(\lambda 0) = y_0$$
 and $\dot{y}(t) = y_0 \exp(\lambda t) \lambda = \lambda y(t)$.

Example 2

Consider the problem of finding $y:[0,\frac{1}{v_0})\to\mathbb{R}$ such that

$$y(0) = y_0$$
 and $\dot{y}(t) = y(t)^2$

for some given $y_0 \in \mathbb{R}$.

The solution to this problem is given by

$$y(t) = \frac{y_0}{1 - y_0 t}$$

because this function satisfies

$$y(0) = \frac{y_0}{1 - y_0 \cdot 0} = y_0$$
 and $\dot{y}(t) = \frac{y_0^2}{(1 - y_0 t)^2} = y(t)^2$.

Note that this y(t) is undefined for $t = \frac{1}{y_0}$. This shows that unlike the solution of Example 1, the solution to this ODE has a finite domain of definition $[0, \frac{1}{y_0})$.

The formula for y(t) is well defined for $t > \frac{1}{y_0}$, but this part of the solution has no mathematical meaning since it is disconnected from the starting point t = 0.

Example 3

Consider the problem of finding $x:[0,\infty)\to\mathbb{R}$ such that

$$x(0) = 1,$$
 $\dot{x}(0) = 0$ and $\ddot{x} = -x.$

The solution to this equation is given by

$$x(t) = \cos(t)$$

since this function satisfies $x(0) = \cos(0) = 1$, $\dot{x}(0) = -\sin(0) = 0$ and $\ddot{x}(t) = -\cos(t) = -x(t)$.

 $\ddot{x} = -x$ is not of the form $\dot{y} = f(y)$, but it can easily be reduced to this form: if we set

$$y = \begin{pmatrix} x \\ \dot{x} \end{pmatrix}$$
 and $f(y) = \begin{pmatrix} y_2 \\ -y_1 \end{pmatrix}$,

then this y(t) and f(y) indeed satisfy

$$\dot{y} = \begin{pmatrix} \dot{x} \\ \ddot{x} \end{pmatrix} = \begin{pmatrix} y_2 \\ -y_1 \end{pmatrix} = f(y).$$

This shows that considering ODEs involving only first-order derivatives is enough to cover ODEs involving arbitrarily high derivatives.

Real-world example: Newton's law of motion

Consider a point particle of mass m > 0 subject to a force $F(x) \in \mathbb{R}^3$ which depends on the location $x \in \mathbb{R}^3$ of the particle.

Newton's law of motion then states that the trajectory $x(t) \in \mathbb{R}^3$ of this particle satisfies

$$m\ddot{x}(t) = F(x(t)).$$

This equation is again not of the form $\dot{y} = f(y)$ but can reduced to this form by setting

$$y = \begin{pmatrix} x \\ \dot{x} \end{pmatrix}, \qquad f(y) = \begin{pmatrix} y[2] \\ \frac{1}{m} F(y[1]) \end{pmatrix}$$

such that

$$\dot{y} = \begin{pmatrix} \dot{x} \\ \ddot{x} \end{pmatrix} = \begin{pmatrix} \dot{x} \\ \frac{1}{m}F(x) \end{pmatrix} = \begin{pmatrix} y[2] \\ \frac{1}{m}F(y[1]) \end{pmatrix} = f(y).$$

Time-dependent ODEs

Many textbooks allow the ODE-defining function f(y,t) to also depend on the time variable t. I omit this generality here because it is rarely needed in applications and it significantly complicates the analysis.

ODEs vs PDEs

ODEs are similar to PDEs in that the problem is to find a function y(t) given an equation in terms of y(t) and its derivatives which is to hold at every point in the domain.

Historically / formally, the defining property of an ODE is that the unknown y(t) depends on only a single variable and hence the derivatives in the differential equations are *ordinary* derivatives.

By contrast, the unknown $u(x_1, ..., x_d)$ in a PDE depends on several variables and hence the derivatives are *partial* derivatives.

The terms "ODE" and "PDE" are hardly ever used in this way anymore, however.

In modern terminology, the defining property of an ODE is that fixed values for y(t) and its derivatives are specified at a single point t_0 . For this reason, ODEs are also called *initial value problems*.

The defining property of a PDE is that fixed values of u(x) and its derivatives are specified at two or more points $x \in \partial \Omega$. For this reason, PDEs are also called *boundary value problems*.

Example: ODEs vs PDEs

The one-dimensional Poisson equation -u''(x) = f(x) is an ODE in the formal sense because there is only a single independent variable x. However, this equation is usually called a PDE because it is almost always paired with boundary conditions rather than initial conditions and hence it is much more similar to e.g. the higher-dimensional Poisson equation $-\Delta u = f$ than to Newton's law of motion $m\ddot{x} = F(x)$.

Discussion

As mentioned earlier, our goal in this lecture is to develop numerical algorithms which approximate the map

$$(f(y), y_0, T) \mapsto y(t)$$
 such that $y(0) = y_0, \dot{y}(t) = f(y(t)).$

However, it is advisable to first establish that this map is well defined and well conditioned, since otherwise we might end up constructing (and trusting!) numerical solutions $\tilde{y}(t)$ which have no meaning.

It turns out that the key condition guaranteeing the existence of solutions y(t) is Lipschitz continuity as defined on the next slide.

Def: (Global) Lipschitz continuity

A function $f: D \to \mathbb{R}^n$ with $D \subset \mathbb{R}^n$ is called *(globally) Lipschitz* continuous with Lipschitz constant L > 0 if for all $y_1, y_2 \in D$ we have

$$||f(y_1)-f(y_2)|| \leq L ||y_1-y_2||.$$

Note that whether a function f(y) is Lipschitz continuous does not depend on the choice of norm due to norm equivalence in finite dimensions. However, the Lipschitz constant L may depend on the choice of norm.

Def: Local Lipschitz continuity

A function $f: \mathbb{R}^n \to \mathbb{R}^n$ is called *locally Lipschitz continuous* if for every $y_0 \in \mathbb{R}^n$ there exists a pair $\delta, L > 0$ such that for all $y_1, y_2 \in \mathbb{R}^n$ we have

$$||y_k - y_0|| \le \delta$$
 \Longrightarrow $||f(y_1) - f(y_2)|| \le L ||y_1 - y_2||.$

Discussion

Lipschitz continuity of a function f(y) is usually most conveniently verified using the following result.

Corollary

Assume $D \subset \mathbb{R}^n$ is convex and $f: D \to \mathbb{R}^n$ is differentiable everywhere in D. Then, f(y) is locally Lipschitz continuous, and it is globally Lipschitz continuous if $\|\nabla f\|$ is bounded.

Proof. Immediate corollary of the result on the next slide.

Lemma: Lipschitz constants and derivatives

Assume $D \subset \mathbb{R}^n$ is convex and $f: D \to \mathbb{R}^n$ has a bounded derivative. Then,

$$||f(y_1) - f(y_2)|| \le L ||y_1 - y_2||$$
 where $L = \sup_{y \in D} ||\nabla f(y)||$

Proof. According to the chain rule, we have that

$$\frac{d}{dt} \Big(f \big(y_1 + t \, (y_2 - y_1) \big) \Big) = \nabla f \big(y_1 + t \, (y_2 - y_1) \big) \, (y_2 - y_1)$$

and hence we conclude using the fundamental theorem of calculus that

$$||f(y_1) - f(y_2)|| = \left\| \int_0^1 \nabla f(y_1 + t(y_2 - y_1)) (y_2 - y_1) dt \right\|$$

$$\leq \int_0^1 \left\| \nabla f(y_1 + t(y_2 - y_1)) \right\| \left\| y_2 - y_1 \right\| dt$$

$$\leq \left(\sup_{y \in D} \|\nabla f(y)\| \right) \left\| y_2 - y_1 \right\|.$$

Discussion

Armed with the definition of Lipschitz continuity, we can now establish that the ODE function $(f(y), y_0, T) \mapsto y(t)$ is indeed well defined if f(y) is Lipschitz continuous.

Picard-Lindelöf theorem

Assume $f: \mathbb{R}^n \to \mathbb{R}^n$ is Lipschitz continuous. Then, there exists a unique function $y: [0, \infty) \to \mathbb{R}^n$ such that

$$y(0) = y_0$$
 and $\dot{y}(t) = f(y(t))$ for all $t \in [0, \infty)$.

Proof. Beyond the scope of this module.

Example

Consider the function $f(y) = \lambda y$ from the example on slide 3. Since $f'(y) = \lambda$ is bounded for all y, this function is globally Lipschitz continuous and hence the solution $y(t) = y_0 \exp(\lambda t)$ exists for all $t \ge 0$.

Discussion

The above result assumes global Lipschitz continuity of f(y) but in returns guarantees that the solution is well defined for all times $t \geq 0$. If f(y) is only local Lipschitz, then the following result establishes that the solution y(t) is well defined at least on some interval $[0,T] \subset [0,\infty)$.

Picard-Lindelöf theorem. local version

Assume $f: \mathbb{R}^n \to \mathbb{R}^n$ is locally Lipschitz continuous. Then, there exists a T>0 and a unique function $y:[0,T)\to \mathbb{R}^n$ such that

$$y(0) = y_0$$
 and $\dot{y}(t) = f(y(t))$ for all $t \in [0, T)$.

Proof. Beyond the scope of this module.

Example

Consider the function $f(y)=y^2$ from the example on slide 4. Since f'(y)=2y exists but is unbounded \mathbb{R} , this function is locally Lipschitz but not globally Lipschitz. Correspondingly, the solution $y(t)=\frac{y_0}{1-y_0t}$ is defined only on the finite interval $[0,\frac{1}{y_0})$ but not on all of $[0,\infty)$.

Discussion

In order to approximate the ODE map $(f(y), y_0, T) \mapsto y(t)$ numerically, we need this map to be not only well defined but also continuous with respect to y_0 ; if this is not the case, then any small perturbation in y_0 (e.g. floating-point rounding) may lead to arbitrarily large errors in the solution y(t).

We will see on the next slide that continuity is a simple consequence of the following auxiliary result.

Gronwall's inequality

$$\dot{y}(t) \le \lambda y(t) \implies y(t) \le \exp(\lambda t) y(0).$$

Proof. Consider $z(t) = \exp(-\lambda t) y(t)$. Then, z(0) = y(0) and

$$\dot{z}(t) = -\lambda \exp(-\lambda t) y(t) + \exp(-\lambda t) \dot{y}(t)$$

$$\leq -\lambda \exp(-\lambda t) y(t) + \exp(-\lambda t) \lambda y(t) = 0;$$

hence $z(t) \le y(0)$ and thus $y(t) \le y(0) \exp(\lambda t)$.

Thm: Lipschitz continuity of the ODE map

Assume $f: \mathbb{R}^n \to \mathbb{R}^n$ is Lipschitz continuous with Lipschitz constant L, and assume $y_1, y_2: [0, T) \to \mathbb{R}^n$ satisfy $\dot{y}_k = f(y_k)$. Then,

$$\|y_1(t) - y_2(t)\| \le \exp(Lt) \|y_1(0) - y_2(0)\|$$
 for any $t < T$.

Proof. We have for any $y: \mathbb{R} \to \mathbb{R}^n$ that

$$\frac{d}{dt}\|y(t)\| = \lim_{\tilde{t}\to t} \frac{\|y(\tilde{t})\| - \|y(t)\|}{\tilde{t}-t} \le \lim_{\tilde{t}\to t} \frac{\|y(\tilde{t})-y(t)\|}{\tilde{t}-t} = \|\dot{y}(t)\|.$$

Combining the above with the Lipschitz continuity of f(y), we obtain

$$\frac{d}{dt} \|y_2(t) - y_1(t)\| \le \|\dot{y}_2(t) - \dot{y}_1(t)\|
= \|f(y_2(t)) - f(y_1(t))\|
\le L \|y_2(t) - y_1(t)\|$$

from which the claim follows by Gronwall's inequality.

Discussion

The bound on the previous slide is a two-sided coin:

- ▶ Good: error at time *t* is proportional to error at time 0.
- ▶ Bad: constant of proportionality is exp(Lt).

The exponential function has a noticeable two-stage character:

- ▶ For $t \leq \frac{1}{L}$, exp(Lt) is about 1.
- ► For $t \gtrsim \frac{1}{L}$, exp(Lt) grows very quickly.

In practice, this means that there is often a characteristic time-scale $\frac{1}{L}$ beyond which numerical simulations become unreliable.

For example, we can predict the weather fairly accurately for the next 2-3 days, but predictions beyond one week are virtually impossible.

Solving ODEs using quadrature

We have now established that if f(y) is Lipschitz continuous, then $\dot{y} = f(y)$ has a unique solution and this solution is a Lipschitz continuous function of the initial conditions y_0 .

Let us now move on to discuss numerical methods for approximating y(t).

According to the fundamental theorem of calculus, we have

$$\begin{array}{c} y(0) = y_0 \\ \dot{y} = f(y) \end{array} \right\} \qquad \Longleftrightarrow \qquad y(t) = y_0 + \int_0^t f(y(\tau)) d\tau.$$

Given a quadrature rule $(\theta_k, w_k)_{k=1}^s$ for [0, 1], we can hence compute a numerical approximation to y(t) using the formula

$$y(t) \approx y_0 + \sum_{k=1}^{s} f(y(\theta_k t)) w_k t.$$

The quadrature points θ_k and weights w_k are multiplied by t because we need to scale the quadrature rule from [0,1] to [0,t].

However, we need to be careful because we do not know y(t) for t > 0 and hence we cannot evaluate $f(y(\theta_k t))$ for quadrature points $\theta_k > 0$.

Solving ODEs using quadrature (continued)

The above suggests that we solve ODEs using the 1-point quadrature rule

$$heta_1=0, \qquad w_1=1 \qquad \longleftrightarrow$$

This yields the following ODE solver.

Def: Euler step

Approximate the solution to $\dot{y} = f(y)$ using

$$y(t) \approx y(0) + f(y(0)) t.$$

Discussion

The error of this solver is easily estimated using Taylor's theorem.

Thm: Error estimate for Euler step

Assume f(y) is differentiable, y(t) satisfies $\dot{y}=f(y)$, and t>0. Then, there exists $\xi\in(0,t)$ such that

$$y(t) \ = \ \underbrace{y(0) + f(y(0))\,t}_{\text{Euler step}} \ + \ \underbrace{\frac{1}{2}\,\nabla f(\xi)\,\dot{y}(\xi)\,t^2}_{\text{error}}.$$

Proof. According to Taylor's theorem, there exists $\xi \in (0, t)$ such that

$$y(t) = y(0) + \dot{y}(0) t + \frac{1}{2} \ddot{y}(\xi) t^{2}.$$

The claim follows by noting that since $\dot{y}(t) = f(y(t))$, we have

$$\dot{y}(0) = f(y(0))$$
 and $\ddot{y}(\xi) = \nabla f(y(\xi)) \dot{y}(\xi)$.

Discussion

The error estimate

$$y(t) = \underbrace{y(0) + f(y(0)) t}_{\text{Euler step}} + \underbrace{\frac{1}{2} \nabla f(\xi) \dot{y}(\xi) t^2}_{\text{error}}$$

gives us two important pieces of information.

The error is smaller the closer f(y) and y(t) are to being constant (i.e. the smaller their derivative).

This is not surprising since the quadrature rule $\theta_1 = 0$, $w_1 = 1$ is exact for constant integrands.

ightharpoonup The error is smaller for smaller times t.

This is again not surprising because for smaller t, the integral $\int_0^t f(y(\tau)) d\tau$ is both smaller and the integrand $f(y(\tau))$ is closer to being constant over the interval [0, t].

Both the ODE function f(y) and the time interval [0, T] are imposed on us by the application; hence the Euler step presented above gives us no means to improve the accuracy by investing more compute time.

This circumstance can be remedied using the composition idea presented on the next slide.

Def: Euler's method

Approximate the solution to $\dot{y} = f(y)$ using

$$\tilde{y}(0) = y(0), \qquad \tilde{y}(t_k) = \tilde{y}(t_{k-1}) + f(\tilde{y}(t_{k-1}))(t_k - t_{k-1})$$

where $(t_k)_k$ denotes a given temporal mesh.

Def: Temporal mesh

An ordered sequence of points $0 = t_0 < t_1 < \dots$

Numerical demonstration

See euler_step(), integrate() and example().

Discussion

We intuitively expect that composing Euler steps into Euler's method yields a numerical approximation $\tilde{y}(t)$ which converges to the exact solution y(t) as the mesh width $\max_k |t_k - t_{k-1}|$ goes to 0.

I will next show that this is indeed the case, and I will do so using the notions of exact and Euler time propagators introduced on the next slide.

Def: (Exact) time propagator

The *(exact) time propagator* associated with $f: \mathbb{R}^n \to \mathbb{R}^n$ is the function

$$\Phi: \mathbb{R}^n \times \mathbb{R} \to \mathbb{R}^n, \qquad \Phi(y_0, t) = y(t)$$

where y(t) is the solution to $\dot{y} = f(y)$, $y(0) = y_0$.

Def: Euler time propagator

The *Euler time propagator* associated with $f: \mathbb{R}^n \to \mathbb{R}^n$ is the function

$$\tilde{\Phi}: \mathbb{R}^n \times \mathbb{R} \to \mathbb{R}^n, \qquad \tilde{\Phi}(y_0, t) = y_0 + f(y_0) t.$$

Discussion

It will be important in the following that the exact time propagator $\Phi(y_0, t)$ is a Lipschitz continuous function of y_0 . We have already seen on slide 17 that this is indeed the case, so all we have to do is to translate that result into the new notation.

Thm: Lipschitz continuity of the exact time propagator

If f(y) is Lipschitz continuous with Lipschitz constant L, then $\Phi(y_0, t)$ is Lipschitz continuous in y_0 with Lipschitz constant $\exp(Lt)$, i.e.

$$\|\Phi(y_1,t)-\Phi(y_2,t)\| \le \exp(Lt)\|y_1-y_2\|.$$

Proof. Immediate consequence of the result on slide 17.

Shorthand notations

In order to keep the statement and proof of the error estimate for Euler's method manageable, I will use the shorthand notations

$$y_k = y(t_k), \qquad \tilde{y}_k = \tilde{y}(t_k),$$

and

$$\Phi_k(y_0) = \Phi(y_0, t_k - t_{k-1}), \qquad \tilde{\Phi}_k(y_0) = \tilde{\Phi}(y_0, t_k - t_{k-1}).$$

As usual, $(t_k)_k$ denotes some temporal mesh specified by the context. With this, we now have all the pieces in place show that Euler's method indeed converges.

Thm: Error estimate for Euler's method

Denote by y(t) the solution to $\dot{y}=f(y)$ and by $\tilde{y}(t)$ the numerical approximation obtained by Euler's method with mesh $(t_k)_{k=0}^n$. Then,

$$\|\tilde{y}_n - y_n\| \le \sum_{k=1}^n \exp(L(t_n - t_k)) \|\tilde{\Phi}_k(\tilde{y}_{k-1}) - \Phi_k(\tilde{y}_{k-1})\|,$$

assuming f(y) is Lipschitz continuous with Lipschitz constant L.

Proof.

$$\begin{split} \|\tilde{y}_{n} - y_{n}\| &= \|\tilde{\Phi}_{n}(\tilde{y}_{n-1}) - \Phi_{n}(y_{n-1})\| & \text{ (definition of time propagators)} \\ \text{ (triangle ineq.)} &\leq \|\tilde{\Phi}_{n}(\tilde{y}_{n-1}) - \Phi_{n}(\tilde{y}_{n-1})\| + \|\Phi_{n}(\tilde{y}_{n-1}) - \Phi_{n}(y_{n-1})\| \\ \text{ (Φ Lipschitz)} &\leq \|\tilde{\Phi}_{n}(\tilde{y}_{n-1}) - \Phi_{n}(\tilde{y}_{n-1})\| + \exp(L\left(t_{n} - t_{n-1}\right))\|\tilde{y}_{n-1} - y_{n-1}\| \\ \text{ (recursion)} &\leq \sum_{k=1}^{n} \exp(L\left(t_{n} - t_{k}\right))\|\tilde{\Phi}_{k}(\tilde{y}_{k-1}) - \Phi_{k}(\tilde{y}_{k-1})\|. \end{split}$$

Terminology

 $\|\tilde{\Phi}_k(\tilde{y}_{k-1}) - \Phi_k(\tilde{y}_{k-1})\|$ is called the *local*, *consistency* or *truncation* error of the numerical time propagator $\tilde{\Phi}$.

Remark

The bound

$$\|\tilde{y}_{n} - y_{n}\| \le \sum_{k=1}^{n} \exp(L(t_{n} - t_{k})) \|\tilde{\Phi}_{k}(\tilde{y}_{k-1}) - \Phi_{k}(\tilde{y}_{k-1})\|$$

can be put into words as follows.

The error at the final time t_n is upper-bounded by the sum of the errors introduced in the time steps $(t_{k-1} \to t_k)_{k=1}^n$ multiplied by the Lipschitz constant of $\Phi(y,t)$ derived on slide 17.

Discussion

The above bound shows that if

$$\|\tilde{\Phi}_k(\tilde{y}_{k-1}) - \Phi_k(\tilde{y}_{k-1})\| \to 0$$
 in the limit $|t_k - t_{k-1}| \to 0$, (1)

then

$$\max_{k} |t_k - t_{k-1}| \to 0$$
 implies that $\|\tilde{y}_n - y_n\| \to 0$.

We have already seen on slide 21 that (1) is indeed the case, so once again all we have to is to translate the old result into the new notation.

Lemma: Consistency of Euler time propagator

The Euler rule time propagator $\tilde{\Phi}(y_0,t)$ satisfies

$$\|\tilde{\Phi}(y_0,t)-\Phi(y_0,t)\|=O(t^2)$$
 for $t\to 0$.

Proof. Immediate consequence of the result on slide 21.

Thm: Convergence of Euler's method

Denote by y(t) the solution to $\dot{y}=f(y)$ and by $\tilde{y}(t)$ the numerical approximation obtained using Euler's method with the equispaced temporal mesh $(t_k=\frac{k}{n}\,T)_{k=0}^n$.

Then,

$$\left\| \widetilde{y}(T) - y(T) \right\| = O\left(\exp(LT) \frac{T^2}{n} \right)$$
 for both $n, T \to \infty$,

assuming f(y) is Lipschitz continuous with Lipschitz constant L.

Proof.
$$\|\tilde{y}_n - y_n\| \le \sum_{k=1}^n \exp(L(t_n - t_k)) \|\tilde{\Phi}_k(\tilde{y}_{k-1}) - \Phi_k(\tilde{y}_{k-1})\|$$
$$\le n \exp(L(t_n - t_0)) O\left(\left(\frac{T}{n}\right)^2\right)$$
$$= \exp(L(t_n - t_0)) O\left(\frac{T^2}{n}\right)$$

Numerical demonstration

See convergence().

Remark 1

Note that Euler's method is second-order consistent but only first-order convergent. The heuristic reason for this is that Euler's method makes n errors of magnitude $O(n^{-2})$; hence the total error is $n O(n^{-2}) = O(n^{-1})$.

Remark 2

The above convergence estimate indicates that the number of steps required to reach a fixed error tolerance

$$\|\tilde{y}(T) - y(T)\| = O\left(\exp(LT)\frac{T^2}{n}\right) \le \tau$$

is
$$n = O(\tau^{-1} T^2 \exp(LT))$$
.

This once again indicates that it is difficult to solve ODEs on time-scales larger than one over the Lipschitz constant.

See nsteps() for numerical demonstration.

Discussion

The theorem on slide 30 establishes that Euler's method

$$\tilde{y}(0) = y(0), \qquad \tilde{y}(t_k) = \tilde{y}(t_{k-1}) + f(\tilde{y}(t_{k-1}))(t_k - t_{k-1})$$

converges to the exact solution y(t), but it also shows that the rate of convergence is only

$$\|\tilde{y}(T) - y(T)\| = O(n^{-1}).$$

This slow convergence is ultimately due to the fact that Euler's method is based on the very poor quadrature approximation

$$\int_0^t f(y(\tau)) d\tau \approx f(y(0)) t;$$

hence the question arises whether we can improve the speed of convergence by choosing a more accurate quadrature rule.

Discussion (continued)

The reason why we settled for the "left-point rule"

$$heta_1=0, \qquad w_1=1 \qquad \longleftrightarrow$$

on slide 20 was that we assumed that we do not know y(t) at times t > 0 and hence we cannot have quadrature points $\theta_k > 0$.

However, this constraint no longer applies now that we know how to (approximately) propagate y(0) into the future using Euler's step.

For example, we can use Euler's step to obtain an estimate

$$\tilde{y}_2(t) = y(0) + f(y(0)) t,$$

and then we can use a trapezoidal rule approximation

$$\tilde{y}(t) = y(0) + \left(f(y(0)) + f(\tilde{y}_2(t))\right)\frac{t}{2},$$

to improve this estimate. This idea is known as the trapezoidal rule.

Def: Trapezoidal rule

Approximate the solution to $\dot{y} = f(y)$ using $\tilde{y}(0) = y(0)$ and

$$\begin{split} \tilde{y}_{2}(t_{k}) &= \tilde{y}(t_{k-1}) + f(\tilde{y}(t_{k-1}))(t_{k} - t_{k-1}), \\ \tilde{y}(t_{k}) &= \tilde{y}(t_{k-1}) + \left(f(\tilde{y}(t_{k-1})) + f(\tilde{y}_{2}(t_{k}))\right) \frac{t_{k} - t_{k-1}}{2}, \end{split}$$

where $(t_k)_k$ denotes a given temporal mesh.

Terminology: Quadrature rules vs. ODE solvers

It is common to refer to both the above ODE solver as well as the quadrature rule from which it is derived as "trapezoidal rule". This is not a problem since it is always clear from context whether the ODE solver or the quadrature rule is meant.

Notation: Distinguishing approximations to the same value

Note that the trapezoidal rule computes two different numerical approximation $\tilde{y}_2(t_k)$ and $\tilde{y}(t_k)$ to the same exact value $y(t_k)$. Here and throughout this lecture, I will distinguish these approximations by writing $\tilde{y}_i(t_k)$, $i=1,2,3,\ldots$, for the auxiliary approximations, and $\tilde{y}(t_k)$ without a subscript for the final approximation.

Implementation of the trapezoidal rule

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See trapezoidal_step() and integrate().
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Note that f1 = t * f(y0) is used both in

$$f2 = t * f(y0 + f1)$$
 and $y0 + (f1 + f2)/2$.

Introducing f1 hence avoids having to evaluate t * f(y0) twice.

Terminology: Bootstrapping

Recall that we derived the trapezoidal rule ODE solver by combining the trapezoidal quadrature rule

$$\int_0^t f(y(\tau)) d\tau \approx \left(f(y(0)) + f(y(t)) \right) \frac{t}{2}$$

with an Euler step to produce an estimate for y(t).

This idea of using a simple procedure to fill in some blanks in a more complicated procedure appears repeatedly in applied mathematics, statistics and computer science and is known as *bootstrapping*.

This terminology is derived from the idiom "to pull oneself over the fence by one's bootstrap" which is used to mock claims of having achieved obviously impossible feats.

Example: "I completed the Computational Mathematics assignment in just one hour" - "Yeah right, and I pulled myself over the fence by my bootstraps."

Picture of a bootstrap: https://i.stack.imgur.com/qjyj6.png

Convergence of the trapezoidal rule

We have derived on slide 27 the estimate

$$\|\tilde{y}_n - y_n\| \le \sum_{k=1}^n \exp(L(t_n - t_k)) \|\tilde{\Phi}_k(\tilde{y}_{k-1}) - \Phi_k(\tilde{y}_{k-1})\|$$

which bounds the total error in terms of the local errors. Looking back at the proof, we realise that the same bounds also holds for the trapezoidal rule if we just replace the Euler time propagator

$$\tilde{\Phi}(y_0, t) = y_0 + f(y_0) t$$

with the following.

Def: Trapezoidal rule time propagator

The trapezoidal rule time propagator associated with $f: \mathbb{R}^n \to \mathbb{R}^n$ is

$$\tilde{\Phi}(y_0, t) = y_0 + (f(y_0) + f(y_2)) \frac{t}{2}$$
 where $y_2 = y_0 + f(y_0) t$,

or equivalently,

$$\tilde{\Phi}(y_0, t) = y_0 + f(y_0) \frac{t}{2} + f(y_0 + f(y_0) t) \frac{t}{2}.$$

Convergence of the trapezoidal rule (continued)

All we have to do to obtain a convergence estimate for the trapezoidal rule method is hence to estimate the order of convergence of the trapezoidal rule consistency error.

Lemma: Consistency of trapezoidal rule

The trapezoidal rule time propagator $\tilde{\Phi}(y_0, t)$ satisfies

$$\|\tilde{\Phi}(y_0,t) - \Phi_k(y_0,t)\| = O(t^3)$$
 for $t \to 0$.

Proof. The proof amounts to showing that the Taylor series of $\tilde{y}(t) = \tilde{\Phi}(y_0, t)$ agrees with the Taylor series of $y(t) = \Phi(y_0, t)$ up to (and including) the second-order term.

I therefore compute

$$\tilde{y}(t) = y_0 + f(y_0) \frac{t}{2} + f(y_0 + f(y_0) t) \frac{t}{2}
\dot{\tilde{y}}(t) = f(y_0) \frac{1}{2} + f'(y_0 + f(y_0) t) f(y_0) \frac{t}{2} + f(y_0 + f(y_0) t) \frac{1}{2}
\ddot{\tilde{y}}(t) = f''(y_0 + f(y_0) t) f(y_0)^2 \frac{t}{4} + 2 f'(y_0 + f(y_0) t) f(y_0) \frac{1}{2}$$

and conclude that we indeed have

$$\ddot{y}(0) = y_0 = y(0), \quad \dot{\ddot{y}}(0) = f(y_0) = \dot{y}(0), \quad \ddot{\ddot{y}}(0) = f'(y_0) f(y_0) = \ddot{y}(0).$$

Remark

The above proof is analogous to the proof on slide 21 which showed that Euler's step is second-order consistent.

Thm: Convergence of the trapezoidal rule

Denote by y(t) the solution to $\dot{y}=f(y)$ and by $\tilde{y}(t)$ the numerical approximation obtained using the trapezoidal rule with the equispaced temporal mesh $(t_k=\frac{k}{n}\,T)_{k=0}^n$.

Then,

$$\left\| \widetilde{y}(T) - y(T)
ight\| = O \left(\exp(LT) \, rac{T^3}{n^2}
ight) \qquad ext{for both } n, \, T o \infty,$$

assuming f(y) is Lipschitz continuous with Lipschitz constant L.

Proof. Analogous to the one on slide 30.

Numerical demonstration

See convergence().

Euler's method vs. the trapezoidal rule

The above result shows that the trapezoidal rule indeed achieves our goal of improving on the $O(n^{-1})$ convergence of Euler's method.

However, comparing the time propagation equations

Euler:
$$\tilde{y}(t_k) = \tilde{y}(t_{k-1}) + f(\tilde{y}(t_{k-1}))(t_k - t_{k-1}),$$

Trapezoidal: $\tilde{y}(t_k) = \tilde{y}(t_{k-1}) + \left(f(\tilde{y}(t_{k-1})) + f(\tilde{y}_2(t_k))\right) \frac{t_k - t_{k-1}}{2}$

where $\tilde{y}_2(t_k) = \tilde{y}(t_{k-1}) + f(\tilde{y}(t_{k-1}))(t_k - t_{k-1}),$

we also observe that the improved speed of convergence of the trapezoidal rule comes at the price of more computations per time step. It is therefore not clear whether the trapezoidal rule is indeed faster than Euler's method.

Euler's method vs. the trapezoidal rule (continued)

We can answer this question by observing that if we have a convergence estimate

$$\|\tilde{y}(T) - y(T)\| = O(n^{-p})$$

and need to achieve an error tolerance

$$\|\tilde{y}(T) - y(T)\| \leq \tau,$$

then we need to take $O(\tau^{-1/p})$ time steps and hence we incur an $O(w \tau^{-1/p})$ runtime where w denotes the runtime per time step.

This shows that increasing the order of convergence p at the expense of a larger runtime per time step w pays off whenever the error tolerance τ is small enough.

Discussion

Given the above, one may wonder whether we can invest even more work per time step to achieve an even better order of convergence and hence a smaller runtime in the limit $\tau \to 0$.

The answer is yes, and in fact we can do so simply by continuing the bootstrapping procedure which lead us to the trapezoidal rule.

Algorithms derived from this idea are known as Runge-Kutta methods.

Def: Runge-Kutta method

Approximate the solution to $\dot{y} = f(y)$ using

1:
$$\tilde{y}(0) = y(0)$$
,

2: **for** $k = 1, 2, \dots$ **do** (time stepping loop)

3: **for**
$$i = 1, ..., s$$
 do (stage loop)

4:
$$\tilde{y}_i = \tilde{y}(t_{k-1}) + \sum_{j=1}^s f(\tilde{y}_j) W_{ij}(t_k - t_{k-1}) \approx y(t_{k-1} + \theta_i(t_k - t_{k-1}))$$

5: end for

6:
$$\tilde{y}(t_k) = \tilde{y}(t_{k-1}) + \sum_{i=1}^{s} f(\tilde{y}_i) w_j (t_k - t_{k-1})$$

7: end for

where

- \triangleright $(t_k)_k$ denotes a given temporal mesh,
- $(\theta_j, W_{ij})_{j=1}^s$ with $i \in \{1, \dots, s\}$ denotes a sequence of quadrature rules for the intervals $[0, \theta_i]$, and
- \blacktriangleright $(\theta_i, w_i)_{i=1}^s$ is a quadrature rule for [0, 1],
- ▶ $s \in \mathbb{N}$ is called the *number of stages* of the Runge-Kutta scheme.

Discussion

The parameters $\theta \in \mathbb{R}^s$, $W \in \mathbb{R}^{s \times s}$ and $w \in \mathbb{R}^s$ introduced in the above definition are most conveniently represented in a tabular form known as *Butcher tableau*.

Def: Butcher tableau

$$\begin{array}{c|cccc} \theta_1 & W_{11} & \dots & W_{1s} \\ \vdots & \vdots & \ddots & \vdots \\ \theta_s & W_{s1} & \dots & W_{ss} \\ \hline & w_1 & \dots & w_s \end{array}$$

The examples on the next slide might help you to better understand the definition of Runge-Kutta schemes and Butcher tableaus.

Example 1: Euler's method

Tableau Interpretation
$$\begin{array}{c|c}
0 & & \tilde{y}_1(t_{k-1}) = \tilde{y}(t_{k-1}) \\
\hline
1 & & \tilde{y}(t_k) = \tilde{y}(t_{k-1}) + f(\tilde{y}_1)(t_k - t_{k-1})
\end{array}$$

Example 2: Trapezoidal rule

Tableau Interpretation

0 |
$$\tilde{y}_1(t_{k-1}) = \tilde{y}(t_{k-1})$$

1 | 1 | $\tilde{y}_2(t_k) = \tilde{y}(t_{k-1}) + f(\tilde{y}_1)(t_k - t_{k-1})$
 $\tilde{y}_2(t_k) = \tilde{y}(t_{k-1}) + f(\tilde{y}_1)\frac{t_k - t_{k-1}}{2} + f(\tilde{y}_2)\frac{t_k - t_{k-1}}{2}$

Remark

Note that the quadrature weight matrix $W \in \mathbb{R}^{s \times s}$ is strictly lower triangular for both the Euler method and trapezoidal rule.

This is for a good reason, for if this was not the case, then we would not be able to compute the auxiliary variables using simple forward substitution.

Example

Consider the Runge-Kutta scheme

$$\frac{1}{1} \frac{1}{1} \qquad \frac{\tilde{y}_1(t_k) = \tilde{y}(t_{k-1}) + f(\tilde{y}_1(t_k))(t_k - t_{k-1})}{\tilde{y}(t_k) = \tilde{y}(t_{k-1}) + f(\tilde{y}_1(t_k))(t_k - t_{k-1})}$$

Computing $\tilde{v}_1(t_k)$ would require solving

$$\tilde{y}_1(t_k) = \tilde{y}(t_{k-1}) + f(\tilde{y}_1(t_k))(t_k - t_{k-1}).$$

Depending on the ODE, this could be a high-dimensional nonlinear system of equations!

Discussion

The above definition allows us to turn any set of parameters θ , W and w into a Runge-Kutta method, but of course not all methods constructed in this way are equally interesting.

Specifically, there is no point considering a particular Runge-Kutta method if there is another method with a comparable runtime per time step but better accuracy.

Combining this insight with the observation that the runtime per time step is mainly determined by the number of stages s, we conclude that a particular Runge-Kutta method is only interesting if its order of convergence is as large as possible for the given s.

Convergence theory of Runge-Kutta methods

Determining Runge-Kutta methods with maximal order of convergence is straightforward.

- Like Euler's method and the trapezoidal rule, Runge-Kutta methods construct an approximate solution $\tilde{y}(t)$ by repeatedly applying a particular numerical time propagator $\tilde{\Phi}(y_0, t)$;
- lacktriangle Any Runge-Kutta solution $ilde{y}(t_n)$ hence satisfies the bound

$$\|\tilde{y}(t_n) - y(t_n)\| \le \sum_{k=1}^n \exp(L(t_n - t_k)) \|\tilde{\Phi}_k(\tilde{y}_{k-1}) - \Phi_k(\tilde{y}_{k-1})\|$$

derived on slide 27.

The order of convergence is thus maximised if we choose θ , W and w such that the Taylor series of the Runge-Kutta time propagator $\tilde{\Phi}(y_0,t)$ agrees with the one of the exact time propagator $\Phi(y_0,t)$ as far as possible.

Convergence theory of Runge-Kutta methods (continued)

The only obstacle in the above procedure is that computing derivatives of $\Phi(y_0, t)$ and $\tilde{\Phi}(y_0, t)$ becomes more and more tedious as the order of the derivative and the number of stages of the Runge-Kutta method increase.

Fortunately, you will hardly ever have to do these computations yourself, because other people have already done them and they have catalogued their findings for us. See

https://en.wikipedia.org/wiki/List_of_Runge-Kutta_methods.

In practice, determining a suitable Runge-Kutta method is hence as simple as skimming the appropriate Wikipedia list.

Example

Runge-Kutta ODE solvers are named after two German mathematicians who proposed the following scheme and showed that it is a fourth-order method.

I demonstrate in rk4_step() how to implement this scheme and in convergence() that it is indeed a fourth-order method.

Number of stages vs. order of convergence

We have now seen three particular Runge-Kutta methods.

Method	# stages	Convergence
Euler	1	$O(n^{-1})$
Trapezoidal	2	$O(n^{-2})$
RK4	4	$O(n^{-4})$

Based on this table, you might expect but that optimal s-stage Runge-Kutta schemes achieve $O(n^{-s})$ convergence, but this is not true; it has been shown that the optimal order of convergence p as a function of the number of stages s is given by

For this reason, Runge-Kutta schemes with s > 4 stages are rarely used.

Adaptive Runge-Kutta methods

So far, I have used equispaced temporal meshes $(t_k = \frac{k}{n} T)_{k=0}^n$ in my numerical experiments and when estimating the rate of convergence of Runge-Kutta methods.

Equispaced meshes are convenient due to their simplicity, but they do not deliver optimal performance: we have seen that the global error is upper-bounded by

$$\|\tilde{y}(t_n) - y(t_n)\| \le \sum_{k=1}^n \exp(L(t_n - t_k)) \|\tilde{\Phi}_k(\tilde{y}_{k-1}) - \Phi_k(\tilde{y}_{k-1})\|$$

and that the local error satisfies

$$\|\tilde{\Phi}_k(\tilde{y}_{k-1}) - \Phi_k(\tilde{y}_{k-1})\| \le \frac{1}{p!} \|\tilde{y}^{(p)}(t) - y^{(p)}(t)\|_{[t_{k-1},t_k]} (t_k - t_{k-1})^p,$$

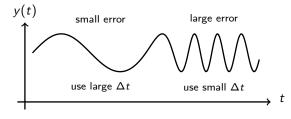
where p denotes the order of consistency (e.g. p=2 for Euler's method and p=3 for the trapezoidal rule).

Adaptive Runge-Kutta methods (continued)

If we choose the same time step t_k-t_{k-1} for all k, then the local error will therefore be small in areas where $\|\tilde{y}^{(p)}(t)-y^{(p)}(t)\|$ is small and large in areas where $\|\tilde{y}^{(p)}(t)-y^{(p)}(t)\|$ is large.

Since the global error is just a weighted sum of the local errors, we therefore conclude that we could have increased the step size $t_k - t_{k-1}$ in areas where $\|\tilde{y}^{(p)}(t) - y^{(p)}(t)\|$ is small without significantly increasing the overall error since the overall error is dominated by the errors in the areas where $\|\tilde{y}^{(p)}(t) - y^{(p)}(t)\|$ is large.

I assume here that $t_n \lesssim \frac{1}{L}$ such that all exponential prefactors are O(1). This is justified since we have already seen that there is little hope to compute $\tilde{y}(t_n)$ accurately if $t_n \gg \frac{1}{L}$.



Adaptive Runge-Kutta methods (continued)

The above indicates that we would ideally choose the temporal mesh $(t_k)_{k=0}^n$ such that $(t_k - t_{k-1})^p$ is roughly inversely proportional to the local curvature $\|\tilde{y}^{(p)} - y^{(p)}\|_{[t_{k-1},t_k]}$.

Calling $y^{(\rho)}(t)$ the *curvature* of y(t) is not standard terminology. I use "curvature" in an informal sense here to indicate that $y^{(\rho)}(t)$ measures some sense of "bendedness" of y(t).

See adaptive_example_1() for an illustration of this idea.

Determining such a mesh requires two procedures.

- A way to estimate the local curvature $\|\tilde{y}^{(p)}(t) y^{(p)}(t)\|$, or equivalently the local error $\|\tilde{\Phi}_k(\tilde{y}_{k-1}) \Phi_k(\tilde{y}_{k-1})\|$.
- ▶ A rule for turning this information into step sizes $t_k t_{k-1}$.

Of course, both of these procedures should be computationally feasible, i.e. they should assume that we can only evaluate f(y) at a finite number of points but have no other information about f(y) (or y(t) for t>0).

Unfortunately, it is not possible to provide a rigorous bound on the local error under these assumptions, see next slide.

"Thm:" Bounding the local error is impossible

No algorithm can compute a bound on $\|\tilde{\Phi}(y_0, t) - \Phi(y_0, t)\|$ using only a finite number of evaluations of f(y).

"Proof." Consider an arbitrary algorithm which bounds $\|\tilde{\Phi}(y_0,t)-\Phi(y_0,t)\|$ by sampling f(y) at a collection of points y_k . If we pass this algorithm a function f(y) such that $f(y_k)=0$ for all k but f(y)>0 otherwise, then the algorithm must assume that f(y)=0 for all y and hence $\tilde{\Phi}(y_0,t)=\Phi(y_0,t)=y_0$ for all t, but actually we have $\tilde{\Phi}(y_0,t),\Phi(y_0,t)>y_0$ and in general $\tilde{\Phi}(y_0,t)\neq\Phi(y_0,t)$. Hence this algorithm would predict that the local error is zero when actually it is positive.

I have put "Thm" and "Proof" in quotation marks because neither the statement nor the proof are fully rigorous. They convey the correct message, however, and that is what matters here.

Estimating the local error

Since computing a rigorous bound on the local error is impossible, we must instead rely on heuristics.

The most commonly employed heuristic is to use two Runge-Kutta methods of different orders of convergence and treat the difference between their solutions as an approximation of the error in the less accurate solution. This heuristic is justified by the following result.

Thm: Estimating the local error by comparing methods

Let $\tilde{\Phi}_1(y_0,t)$, $\tilde{\Phi}_2(y_0,t)$ be two Runge-Kutta time propagators such that the local errors satisfy

$$\left\|\tilde{\Phi}_{\ell}(y_0,t) - \Phi(y_0,t)\right\| = O(t^{p_{\ell}})$$

for two different exponents $p_1 < p_2$.

We then have

$$\left\|\tilde{\Phi}_{1}(y_{0},t)-\tilde{\Phi}_{2}(y_{0},t)\right\|=\left\|\tilde{\Phi}_{1}(y_{0},t)-\Phi(y_{0},t)\right\|+O(t^{p_{1}+1}).$$

Proof. See next slide.

Proof. According to Taylor's theorem, we have

$$\begin{split} &\tilde{\Phi}_{1}(y_{0},t)=y(0)+\dot{y}(0)\,t+\ldots+\frac{y^{(\rho_{1})}(0)}{(\rho_{1}-1)!}\,t^{\rho_{1}-1}+C\,t^{\rho_{1}}+O\big(t^{\rho_{1}+1}\big)\\ &\tilde{\Phi}_{2}(y_{0},t)=y(0)+\dot{y}(0)\,t+\ldots+\frac{y^{(\rho_{1})}(0)}{(\rho_{1}-1)!}\,t^{\rho_{1}-1}+\frac{y^{(\rho_{1})}(0)}{\rho_{1}!}\,t^{\rho_{1}}+O\big(t^{\rho_{1}+1}\big)\\ &\tilde{\Phi}(y_{0},t)=y(0)+\dot{y}(0)\,t+\ldots+\frac{y^{(\rho_{1})}(0)}{(\rho_{1}-1)!}\,t^{\rho_{1}-1}+\frac{y^{(\rho_{1})}(0)}{\rho_{1}!}\,t^{\rho_{1}}+O\big(t^{\rho_{1}+1}\big) \end{split}$$

for some constant $C \neq \frac{y^{(p_1)}(0)}{p_1!}$. We therefore conclude that

$$\tilde{\Phi}_{1}(y_{0}, t) - \tilde{\Phi}_{2}(y_{0}, t) = \left(C - \frac{y^{(p_{1})}(0)}{p_{1}!}\right) t^{p_{1}} + O(t^{p_{1}+1})
= \tilde{\Phi}_{1}(y_{0}, t) - \tilde{\Phi}(y_{0}, t) + O(t^{p_{1}+1}).$$

Embedded Runge-Kutta methods

In addition to being exact to leading order, estimating local errors by comparing Runge-Kutta methods has the advantage that it often incurs little or no additional cost.

For example, the trapezoidal rule step

$$\tilde{y}_2(t) = y(0) + f(y(0)) t,$$

$$\tilde{y}(t) = y(0) + \left(f(y(0)) + f(\tilde{y}_2(t)) \right) \frac{t}{2}$$

already computes the Euler solution $\tilde{y}_2(t)$ and thus we can compute

$$\|\tilde{y}_2(t) - \tilde{y}(t)\| \approx \|\tilde{y}_2(t) - y(t)\|$$

without any additional cost compared to just evaluating the trapezoidal solution $\tilde{y}(t)$, see embedded_ET_step().

Pairs of Runge-Kutta methods where the lower-order method uses only intermediate quantities of the higher-order method are called *embedded*.

Embedded Runge-Kutta methods of various orders are known, see

Choosing the step size

Now that we know how to estimate the local error, we can address the problem of choosing the step sizes $t_k - t_{k-1}$ such that the local errors are approximately the same everywhere, i.e. such that

$$\|\tilde{\Phi}_{1,k}(\tilde{y}_{k-1}) - \Phi_k(\tilde{y}_{k-1})\| \approx \tau \text{ independent of } k.$$

This problem can be solved in many ways, but the most efficient is to note that

$$\|\tilde{\Phi}_{1,k}(\tilde{y}_{k-1}) - \Phi_k(\tilde{y}_{k-1})\| = O((t_k - t_{k-1})^p)$$

implies that

$$\|\tilde{\Phi}_{1,k}(\tilde{y}_{k-1}) - \Phi(\tilde{y}_{k-1})\| \approx C (t_k - t_{k-1})^p$$

for some C > 0 as long as $t_k - t_{k-1} > 0$ is small enough, and thus

$$\|\tilde{\Phi}_{1,k}(\tilde{y}_{k-1}) - \Phi_k(\tilde{y}_{k-1})\| \approx \tau \qquad \iff \qquad t_k - t_{k-1} = \left(\frac{\tau}{C}\right)^{1/p}.$$

This formula for $t_k - t_{k-1}$ can be evaluated as soon as we have an estimate for C, and such an estimate can be obtained by solving for C the error estimate for the previous time step, which is

$$\|\tilde{\Phi}_{1,k-1}(\tilde{y}_{k-2}) - \Phi_{2,k-1}(\tilde{y}_{k-2})\| \approx C(t_{k-1} - t_{k-2})^{p_1}.$$

Choosing the step size (continue)

Putting all of the above together, we obtain the formula

$$t_{k}-t_{k-1}=\left(t_{k-1}-t_{k-2}\right)\left(\frac{\tau}{\left\|\tilde{\Phi}_{1,k-1}(\tilde{y}_{k-2})-\Phi_{2,k-1}(\tilde{y}_{k-2})\right\|}\right)^{1/p_{1}}.$$
 (2)

This formula provides a first practical rule for determining a temporal mesh $(t_k)_{k=0}^n$ such that all local errors are approximately equal to τ , but it is usually not quite accurate enough on its own.

Most adaptive ODE solvers therefore check whether the above choice for the step width t_k-t_{k-1} indeed yields

$$\|\tilde{\Phi}_{1,k}(\tilde{y}_{k-1}) - \tilde{\Phi}_{2,k}(\tilde{y}_{k-1})\| \leq \tau,$$

and if not then (2) is iterated until we find a step which such that the local error tolerance is satisfied. Furthermore, since rejecting and recomputing steps is expensive, we usually multiply (2) by a "safety factor" (e.g. 0.9) to increase our chances of accepting the proposed step size $t_k - t_{k-1}$.

Code describes this better than words, so I strongly recommend that you have a look at propagate_adaptively().

Adaptive step selection applied to $\dot{y} = -y$

adaptive_example_1() demonstrates that the heuristics encoded in propagate_adaptively() work fairly well for this particular example.

Encouraged by this finding, I apply the same code to the simple ODE $\dot{y}=-y$, y(0)=1 in adaptive_rk_example_2().

Since the exact solution $y(t)=\exp(-t)$ becomes increasingly flat for $t\to\infty$, we would expect that our adaptive step size selection routine would choose step sizes t_k-t_{k-1} which grow beyond bounds for $t\to\infty$. However, we empirically observe that the step size grows only until $t\approx 15$ and remains approximately constant afterwards, see stepsize().

Choosing finite step sizes even though the solution y(t) is not noticeably changing indicates that our algorithm probably does more work than what would be necessary.

Let us therefore investigate why our step selection routine refuses to increase the step size beyond some upper bound $\Delta t_{\rm max}$.

Runge-Kutta methods applied to $\dot{y} = -y$

A partial explanation for why the time steps $t_k - t_{k-1}$ eventually stop increasing is that the numerical solution $\tilde{y}(t)$ in general satisfies only

$$\lim_{t \to \infty} \widetilde{y}(t) pprox 0$$
 but not necessarily $\lim_{t \to \infty} \widetilde{y}(t) = 0$.

The numerical solution $\tilde{y}(t)$ thus gets "stuck" at a point $\tilde{y}(\infty) \neq 0$ where the curvature $|f'(\tilde{y}(\infty))|$ is small but nonzero, and correspondingly also the step size $t_k - t_{k-1}$ gets "stuck" at a large but finite value.

We can see this by running adaptive_rk_example_2() again but setting logy = true.

Runge-Kutta methods applied to $\dot{y} = -y$ (continued)

An even better understanding for why the step size remains bounded can be obtained by computing explicit formulae for the Euler and trapezoidal rule propagators.

► Euler:
$$\tilde{y}(t) = y_0 + f(y_0) t = y_0 - y_0 t = (1 - t) y_0.$$

► Trapezoidal: $\tilde{y}(t) = y_0 + \left(f(y_0) + f(y_0 + f(y_0) t)\right) \frac{t}{2}$

$$= y_0 + \left(-y_0 - \left(y_0 - y_0 t\right)\right) \frac{t}{2}$$

$$= \left(1 - t + \frac{t^2}{2}\right) y_0.$$

This shows that after k steps with constant step size Δt , the Euler and trapezoidal rule solutions are given by

$$\tilde{y}(k\,\Delta t)=R(-\Delta t)^k\,y(0),$$

where

$$R(z) = \begin{cases} 1+z & \text{for Euler's method, and} \\ 1+z+\frac{z^2}{2} & \text{for the trapezoidal rule.} \end{cases}$$

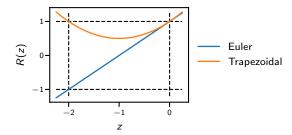
Runge-Kutta methods applied to $\dot{y} = -y$ (continued)

We conclude from the above that the Runge-Kutta solutions $\tilde{y}(t)$

- grow exponentially if $|R(-\Delta t)| > 1$,
- lacktriangle are constant or oscillate if $|R(-\Delta t)|=1$, and
- decay exponentially if $|R(-\Delta t)| < 1$.

In the case $\dot{y}=-y$, we know that the exact solution $y(t)=\exp(-t)$ decays exponentially and hence we want our numerical solutions to do the same.

Combining the above with the plots of R(z) shown below, we conclude that this is the case if and only if $0 < \Delta t < 2$.



We may summarise our findings from the last two slides as follows.

Thm: Step size constraint for $\dot{y} = -y$

Euler's method and the trapezoidal rule applied to $\dot{y}=-y$ with an equispaced temporal mesh $(t_k=k\,\Delta t)_{k=1}^\infty$ produce numerical solutions $\tilde{y}(t)$ such that $\lim_{t\to\infty}\tilde{y}(t)=0$ if and only if $\Delta t<2$.

Terminology: Step size constraint and largest admissible step size Bounds of the form $\Delta t < \Delta t_{\rm max}$ are called *step size constraints*, and the upper bound $\Delta t_{\rm max}$ is called the *largest admissible step size*.

Runge-Kutta methods applied to $\dot{y} = -y$ (continued)

Running stepsize() with show_ref = true, we observe that our adaptive step size routine was actually able to detect the step size constraint and choose t_k-t_{k-1} correspondingly.

Furthermore, stability_example() shows that $\tilde{y}(t)$ indeed converges to 0 only if $\Delta t < 2$.

Discussion

The last four slides discussed a particular issue arising when Runge-Kutta methods are applied to the ODE $\dot{y}=-y$.

However, we already know that

$$\dot{y} = -y \implies y(t) = y(0) \exp(-t),$$

so you may wonder why I spent so much time discussing how to numerically solve an ODE whose exact solution is already known.

The reason for this is the result presented on the next slide.

Def: Fixed point / steady state

A point $y_F \in \mathbb{R}^n$ is said to be a *fixed point* or a *steady state* of the ODE $\dot{y} = f(y)$ if $f(y_F) = 0$.

Thm: Linearisation of ODEs

Assume y(t) solves the ODE $\dot{y} = f(y)$ which has a fixed point $y_F \in \mathbb{R}^n$, and assume the Jacobian $\nabla f(y_F)$ has the eigenvalue decomposition

$$\nabla f(y_F) = V \Lambda V^{-1}$$
.

Then.

$$\frac{d}{dt}\Big(V^{-1}\big(y(t)-y_F\big)\Big) = \Lambda\left(V^{-1}\big(y(t)-y_F\big)\right) + O\Big(\big\|y(t)-y_F\big\|^2\Big),$$

i.e. each of the *n* components of the vector $w(t) = V^{-1}(y(t) - y_F)$ satisfies

$$\dot{w}_k(t) = \lambda_k w_k(t) + O(\|w\|^2).$$

Proof. We have

$$\frac{\frac{d}{dt}(y(t) - y_F)}{dt} = \frac{\dot{y}(t) - \frac{d}{dt}y_F}{y_F} = f(y(t)) + 0 \\
= \underbrace{f(y_F)}_{0} + \underbrace{\nabla f(y_F)}_{VAV^{-1}}(y(t) - y_F) + O(\|y(t) - y_F\|^2).$$

The claim follows by multiplying from the left with V^{-1} .

Discussion

The above theorem may be summarised as follows.

Any ODE $\dot{y} = f(y)$ behaves like the simple ODE $\dot{w} = \lambda w$ in a neighbourhood of its fixed points y_F .

Terminology: Linearisation

 $\dot{w} = \lambda w$ is often called the *linearisation* of $\dot{y} = f(y)$ because it is derived from the linear approximation

$$f(y) \approx f(y_F) + \nabla f(y_F)(y - y_F).$$

Discussion (continued)

The theorem further says that the parameters λ_k in the linearised ODEs are given by the eigenvalues of the Jacobian $\nabla f(y_F)$ and may therefore be complex even if $f: \mathbb{R}^n \to \mathbb{R}^n$ is real.

In fact, we will see on the next slide that the real and imaginary parts of λ provide us with some insight into the local behaviour of y(t).

Interpretation of eigenvalues of $\nabla f(y_F)$

The solution to the linarised ODE $\dot{w} = \lambda w$ is given by

$$w(t) = y_0 \exp(\lambda t)$$

= $y_0 \exp(\text{Re}(\lambda) t) \left(\cos(\text{Im}(\lambda) t) + i \sin(\text{Im}(\lambda) t)\right).$

We therefore conclude the following.

- ▶ y(t) converges to y_F if all eigenvalues of $\nabla f(y_F)$ have a negative real part $\text{Re}(\lambda) < 0$, and it diverges from y_F if at least one eigenvalue has a positive real part $\text{Re}(\lambda) > 0$.
- y(t) oscillates if $Im(\lambda) \neq 0$, and y(t) is monotonous if $Im(\lambda) = 0$.

The first item motivates the following definition.

Def: Attractive fixed point

A fixed point y_F of $\dot{y} = f(y)$ is said to be *attractive* if all eigenvalues of $\nabla f(y_F)$ have a negative real part.

Example

Recall from slide 5 the ODE $\ddot{x} = -x$, or equivalently

$$\dot{y} = \begin{pmatrix} \dot{y}_1 \\ \dot{y}_2 \end{pmatrix} = \begin{pmatrix} y_2 \\ -y_1 \end{pmatrix} = f(y)$$
 with $\begin{pmatrix} y_1 \\ y_2 \end{pmatrix} = \begin{pmatrix} x \\ \dot{x} \end{pmatrix}$.

Even though f(y) is real, the eigenvalues $\lambda = \pm i$ of

$$abla f = egin{pmatrix} 0 & 1 \ -1 & 0 \end{pmatrix}$$

are complex, and in fact they are purely imaginary which indicates that the solutions oscillate with constant amplitudes.

This of course agrees with the particular solution

$$y(t) = \begin{pmatrix} \cos(t) \\ -\sin(t) \end{pmatrix}$$

determined on slide 5.

Discussion

Our findings so far may be summarised as follows.

Euler's method and the trapezoidal rule applied to an ODE converging to an attractive fixed point suffer from a step size constraint which forces them to invest compute time in regions where we should not be investing compute time.

The mathematical literature has a special term for such ODEs.

Terminology: Stiffness

ODEs for which some¹ Runge-Kutta methods require finite time steps even in regions where there is virtually no change in the solution are called *stiff*.

¹Specifically, stiff ODEs are ODEs for which *explicit* Runge-Kutta methods require finite time steps. However, the notion of explicit and implicit Runge-Kutta methods will be introduced only later.

Discussion

Managing stiffness is perhaps the single most important source of complications when solving ODEs in practice.

In simple cases like the $\dot{y}=-y$ example discussed above, "managing stiffness" amounts to recognising that we should stop the Runge-Kutta solver once we reach a time T such that $f(\tilde{y}(T))\approx 0$ and set $\tilde{y}(t)=\tilde{y}(T)$ for all T>t.

However, this work-around breaks down once we move to more complicated ODEs like the one discussed on the next slide.

Example

Consider the problem of finding $y:[0,\infty)\to\mathbb{R}^2$ such that

$$\dot{y}_1 = -1000 \, y_1, \qquad \dot{y}_2 = -y_2.$$

Since $\lambda_1=-1000$ is very negative, the Euler / trapezoidal rule solution will remain bounded only if we use a very small time step $\Delta t < \frac{2}{1000}$.

This would not be a problem if the ODE consisted of only y_1 , because a large λ_1 also means that $y_1(t)$ converges to the steady state quickly and hence the minimal number of time steps

$$n \ge \frac{\text{[time scale of interest]}}{\text{[largest admissible time step]}} = \frac{O(1/1000)}{O(1/1000)} = O(1)$$

would still be manageable for the single ODE $\dot{y}_1 = -1000 \, y_1$.

However, once we add the relatively slowly converging $\dot{y}_2 = -y_2$ to the simulation, then the time scale of interest will change to O(1) and thus the lower bound on the number of time steps becomes

$$n \ge \frac{\text{[time scale of interest]}}{\text{[largest admissible time step]}} = \frac{O(1)}{O(1/1000)} = O(1000).$$

This may be more than we can handle.

Discussion

An ODE $\dot{y} = f(y)$ like the above where some components converge / oscillate very rapidly while others converge / oscillate very slowly is said to exhibit a *separation of time scales*.

Separation of time scales is the primary reason why stiffness is a problem: the rapidly converging components force us to take very small time steps while the slowly converging components prevent us from terminating the simulation once the former components converge.

It is sometimes possible to eliminate a separation-of-time-scales-induced loss in performance by splitting the overall ODE into one ODE for the rapidly oscillating components (e.g. y_1) and another ODE for the slowly oscillating components (y_2).

However, this approach is not always applicable, and when it is then it is usually quite cumbersome.

Discussion (continued)

We would therefore prefer to simulate ODEs exhibiting separation of time scales using just a single Runge-Kutta method which does not suffer from a time step constraint when applied to stiff ODEs.

We have seen that neither Euler's method nor the trapezoidal rule satisfy this criterion, so the question arises whether we can design special Runge-Kutta schemes which do.

The answer is yes, but in order to design such Runge-Kutta schemes we need the abstract concepts introduced on the next slide.

Def: Stability function

The stability function associated with a Runge-Kutta scheme is the function $R:\mathbb{C}\to\mathbb{C}$ such that the numerical time propagator for the equation $\dot{y}=\lambda y$ is given by

$$\tilde{\Phi}(y_0,t)=R(\lambda t)\,y_0.$$

Def: Stability domain

The stability domain associated with a Runge-Kutta method is the set

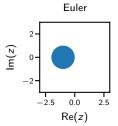
$$\Big\{z\in\mathbb{C}\ \big|\ |R(z)|<1\Big\}.$$

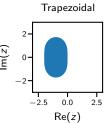
Example

We have derived on slide 63 two examples of stability functions, namely

- ightharpoonup R(z) = 1 + z for Euler's method, and
- ► $R(z) = 1 + z + \frac{z^2}{2}$ for the trapezoidal rule.

Plugging these functions into a plotting tool, we obtain the following stability domains.





The stability domain for Euler's method can also be determined using pen-and-paper analysis: the set

$$\{z \in \mathbb{C} \mid |R(z)| = |1+z| < 1\}$$

is simply a ball of radius 1 centred at z = -1.

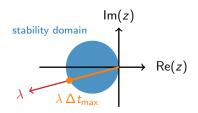
Determining step size constraints using stability domains

The purpose of the stability domain is to allow us to visually estimate the range of admissible time steps Δt such that the numerical solution

$$\tilde{y}(k\,\Delta t)=R(\lambda\,\Delta t)^ky_0$$

to $\dot{y}=\lambda y$ does not blow up: this range is proportional to the intersection between the stability domain $\{|R(z)|<1\}$ and the ray λ $[0,\infty)$.

Furthermore, the largest admissible time step $\Delta t_{\rm max}$ is proportional to the point of maximum modulus in this intersection.



Discussion

We have seen on slide 69 that the exact solution to $\dot{y}=\lambda y$ converges to 0 whenever $\text{Re}(\lambda)<0$. Combined with the insights from the previous slide, this means that a given Runge-Kutta scheme will suffer from a time step constraint unless its stability domain involves the entire left half plane $\{\text{Re}(z)<0\}$.

Devising a Runge-Kutta scheme with unbounded stability domain $\{|R(z)|<1\}$ will be easier once we have a formula for the stability function of a generic Runge-Kutta scheme.

Thm: Stability function formula

The stability function of a Runge-Kutta scheme with Butcher tableau

$$\left(egin{array}{c|c} \theta & W \\ \hline & w^T \end{array}
ight)$$
 is given by $R(z) = 1 + z w^T ig(I - z W ig)^{-1} \, \mathbf{1},$

where 1 denotes the vector of all ones.

Proof. The time propagator for the above Butcher tableau and the ODE $\dot{y}=\lambda y$ is given by

$$\tilde{\Phi}(y_0, t) = y_0 + \lambda t \, w^T \tilde{y}$$
 where $\tilde{y} = y_0 \, \mathbf{1} + \lambda t \, W \tilde{y}$.

Solving the second equation for the vector of auxiliary variables \tilde{y} , we obtain

$$\tilde{\mathbf{y}} = (\mathbf{I} - \lambda \mathbf{t} \, \mathbf{W})^{-1} \, \mathbf{y}_0 \, \mathbf{1},$$

and inserting this expression into the formula for $\tilde{\Phi}(y_0,t)$ yields

$$\tilde{\Phi}(y_0,t) = \left(1 + \lambda t \, w^T (I - \lambda t \, W)^{-1} \mathbf{1}\right) y_0 = R(\lambda t) \, y_0.$$

Example: Stability function for Euler's method

Butcher tableau:

$$\left(\begin{array}{c|c} 0 & \\ \hline & 1 \end{array}\right)$$

Stability function:

$$R(z) = 1 + z 1 (1 - 0)^{-1} 1 = 1 + z.$$

Example: stability function for the trapezoidal rule

Butcher tableau:

$$\begin{pmatrix}
0 & & \\
1 & 1 & \\
& & \frac{1}{2} & \frac{1}{2}
\end{pmatrix}$$

Stability function:

$$R(z) = 1 + z \begin{pmatrix} \frac{1}{2} & \frac{1}{2} \end{pmatrix} \begin{pmatrix} I - z \begin{pmatrix} 0 & 0 \\ 1 & 0 \end{pmatrix} \end{pmatrix}^{-1} \begin{pmatrix} 1 \\ 1 \end{pmatrix}$$

$$= 1 + z \begin{pmatrix} \frac{1}{2} & \frac{1}{2} \end{pmatrix} \begin{pmatrix} 1 & 0 \\ -z & 1 \end{pmatrix}^{-1} \begin{pmatrix} 1 \\ 1 \end{pmatrix}$$

$$= 1 + z \begin{pmatrix} \frac{1}{2} & \frac{1}{2} \end{pmatrix} \begin{pmatrix} 1 \\ 1 + z \end{pmatrix}$$

$$= 1 + z \begin{pmatrix} 1 + \frac{z}{2} \end{pmatrix}$$

$$= 1 + z + \frac{z^2}{3}.$$

Discussion

Now that we have a general formula for the stability function R(z), let us look into what this formula tells us about the possible shapes of the stability domain.

The following result will be useful for this endeavour.

Lemma

$$R(z) = 1 + zw^{T} (I - zW)^{-1} \mathbf{1}$$

is a rational function of z in general, and it is a polynomial in z if W is strictly lower triangular (i.e. $i \le j \implies W_{ii} = 0$).

Proof. $(I-zW)^{-1}\mathbf{1}$ can be evaluated using LU factorisation and back substitution. These algorithms involve only addition, multiplication and division; hence every entry of $(I-zW)^{-1}\mathbf{1}$ is a rational function of z, and thus R(z) is a rational function.

If W is strictly lower triangular, then I-zW is a triangular matrix and hence we can skip the LU factorisation. Furthermore, all diagonal entries of I-zW are 1 and thus back substitution does not involve any divisions. Every entry of $(I-zW)^{-1}\mathbf{1}$ is hence a polynomial in z, and thus R(z) is a polynomial in this case.

Stability domains for polynomial and rational stability functions

The above result is important because if R(z) is a polynomial, then we must necessarily have $\lim_{|z|\to\infty}|R(z)|=\infty$ and hence the stability domain $\big\{|R(z)|<1\big\}$ is bounded and the Runge-Kutta method suffers from a step size constraint.

By contrast, if R(z) is a rational, then $\lim_{|z|\to\infty}|R(z)|<\infty$ is possible and hence the stability domain may be unbounded and the Runge-Kutta method may be free from step size constraints.

Example

Consider the Runge-Kutta scheme

$$\begin{array}{c|c} 1 & 1 \\ \hline & 1 \end{array} \qquad \begin{array}{c|c} \tilde{y}_1(t) = \tilde{y}(0) + f(\tilde{y}_1(t)) \ t \\ \hline \tilde{y}(t) = \tilde{y}(0) + f(\tilde{y}_1(t)) \ t \end{array}$$

Since the expressions for $\tilde{y}(t)$ and $\tilde{y}_1(t)$ are exactly the same, we conclude that the one-step equation for this scheme may also be written as

$$\tilde{y}(t) = \tilde{y}(0) + f(\tilde{y}(t)) t.$$

Example (continued)

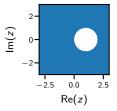
Inserting f(y) = y, we obtain

$$\tilde{y}(t) = \tilde{y}(0) + \tilde{y}(t) t \qquad \iff \qquad \tilde{y}(t) = \frac{\tilde{y}(0)}{1-t};$$

hence the stability function for this scheme is $R(z) = \frac{1}{1-z}$, and the stability domain is

$${z \mid |R(z)| < 1} = {z \mid \frac{1}{|1-z|} < 1} = {z \mid |1-z| > 1},$$

which is the complement of the ball of radius 1 around z = 1.



This stability domain includes the entire left half plane. The above Runge-Kutta scheme therefore does not suffer from a time step constraint near attractive fixed points.

Remark

The one-step equation of the Runge-Kutta scheme introduced above,

$$\tilde{y}(t) = \tilde{y}(0) + f(\tilde{y}(t))t,$$
 (3)

differs from the one step equation of Euler's method from slide 23,

$$\tilde{y}(t) = \tilde{y}(0) + f(\tilde{y}(0)) t,$$

only in the argument $\tilde{y}(t) / \tilde{y}(0)$ which is passed to f(y).

Furthermore, (3) is an *implicit* equation for $\tilde{y}(t)$, i.e. it does not provide a formula for evaluating $\tilde{y}(t)$ but instead defines $\tilde{y}(t)$ implicitly as the solution to the above equation. This motivates the following terminology.

Def: Implicit Euler's method

Approximate the solution to $\dot{y} = f(y)$ using

$$\tilde{y}(0) = y(0), \qquad \tilde{y}(t_k) = \tilde{y}(t_{k-1}) + f(\tilde{y}(t_k))(t_k - t_{k-1})$$

where $(t_k)_k$ denotes a given temporal mesh.

Discussion

More generally, we have observed on slide 46 that Runge-Kutta schemes cannot be evaluated using simple forward substitution whenever W is not strictly lower triangular. The explicit / implicit terminology for Euler's method therefore generalises as follows.

Terminology: Explicit and implicit Runge-Kutta schemes

Consider a Runge-Kutta scheme with Butcher tableau

$$\left(\begin{array}{c|c} \theta & W \\ \hline & w^T \end{array}\right)$$
.

This scheme is called *explicit* if W is strictly lower triangular, and *implicit* otherwise.

Numerical demonstration

See implicit_euler_step(), stability_example().

Convergence of implicit Euler method

We have now seen that the implicit Euler method allows us to take large time steps without risking spurious blow-up of the numerical solution. This is nice, but of course we also want that the numerical solution converges to the exact solution as the time step goes to zero, i.e. we want to have a proof that the implicit Euler method converges in the limit $\max_k |t_k-t_{k-1}| \to 0$.

Using the bound

$$\|\tilde{y}_n - y_n\| \le \sum_{k=1}^n \exp(L(t_n - t_k)) \|\tilde{\Phi}_k(\tilde{y}_{k-1}) - \Phi_k(\tilde{y}_{k-1})\|$$

from slide 27, we can once again reduce the problem of showing convergence to that of showing that the local consistency errors satisfy

$$\lim_{t\to 0} \|\tilde{\Phi}(y_0,t) - \Phi_k(y_0,t)\| = 0.$$

Lemma: Consistency of implicit Euler time propagator

The implicit Euler time propagator $\tilde{\Phi}(y_0,t)$ satisfies

$$\|\tilde{\Phi}(y_0,t)-\Phi(y_0,t)\|=O(t^2) \qquad \text{for } t\to 0.$$

 ${\it Proof.}$ We obtain by repeatedly differentiating the one-step equation for the implicit Euler method that

$$\begin{split} \tilde{y}(t) &= y_0 + f(\tilde{y}(t)) t, \\ \dot{\tilde{y}}(t) &= f'(\tilde{y}(t)) \dot{\tilde{y}}(t) t + f(\tilde{y}(t)), \\ \ddot{\tilde{y}}(t) &= f''(\tilde{y}(t)) (\dot{\tilde{y}}(t))^2 t + f'(\tilde{y}(t)) \ddot{\tilde{y}}(t) t + 2 f'(\tilde{y}(t)) \dot{\tilde{y}}(t). \end{split}$$

Setting t = 0, we conclude

$$\tilde{y}(0) = y_0 = y(0), \quad \dot{\tilde{y}}(0) = f(y_0) = \dot{y}(0), \quad \ddot{\tilde{y}}(0) = 2 f'(y_0) \dot{\tilde{y}}(0) = 2 \ddot{y}(0).$$

We observe that the zeroth and first derivatives of $\tilde{y}(t)$ and y(t) match, but the second derivatives do not. Hence the implicit Euler is second-order consistent like the explicit Euler method.

Corollary: Convergence of the implicit Euler rule

Denote by y(t) the solution to $\dot{y}=f(y)$ and by $\tilde{y}(t)$ the numerical approximation obtained using the implicit Euler rule with the equispaced temporal mesh $(t_k=\frac{k}{n}\,T)_{k=0}^n$. Then,

$$\|\tilde{y}(T) - y(T)\| = O\left(\exp(LT)\frac{T^2}{n}\right)$$
 for both $n, T \to \infty$,

assuming f(y) is Lipschitz continuous with Lipschitz constant L.

Proof. Analogous to the one on slide 30.

Numerical demonstration

See convergence().

Discussion

We have now seen how we can avoid time step constraints for stiff ODEs by switching from the explicit to the implicit Euler method.

The same is also possible for higher-order schemes. For example, we can transform the (explicit) trapezoidal rule scheme from slide 34 into an implicit scheme by modifying it as follows.

Def: Implicit trapezoidal rule

Approximate the solution to $\dot{y} = f(y)$ using $\tilde{y}(0) = y(0)$ and

$$\widetilde{y}(t_k) = \widetilde{y}(t_{k-1}) + \left(f(\widetilde{y}(t_{k-1})) + f(\widetilde{y}(t_k))\right) \frac{t_k - t_{k-1}}{2},$$

where $(t_k)_k$ denotes a given temporal mesh.

Discussion

We by now have quite a list of properties associated with a given Runge-Kutta scheme, namely

- a stepping function,
- ▶ the Butcher tableau,
- the order of consistency and the order of convergence,
- ▶ the stability function and stability domain.

In the following, I will derive each of these properties for the implicit trapezoidal rule.

Implementation of the implicit trapezoidal rule

See implicit_trapezoidal_step().

Butcher tableau for the implicit trapezoidal rule

The Butcher tableau for the implicit trapezoidal rule is given by

$$\begin{array}{c|c}
0 & \tilde{y}_1(t_{k-1}) = \tilde{y}(t_{k-1}) \\
\hline
1 & \frac{1}{2} & \frac{1}{2} \\
\hline
 & \frac{1}{2} & \frac{1}{2}
\end{array}$$

$$\begin{array}{c|c}
\tilde{y}_2(t_k) = \tilde{y}(t_{k-1}) + f(\tilde{y}_1) \frac{t_k - t_{k-1}}{2} + f(\tilde{y}_2) \frac{t_k - t_{k-1}}{2} \\
\hline
\tilde{y}(t_k) = \tilde{y}(t_{k-1}) + f(\tilde{y}_1) \frac{t_k - t_{k-1}}{2} + f(\tilde{y}_2) \frac{t_k - t_{k-1}}{2}
\end{array}$$

These equations are indeed equivalent to the equation

$$\tilde{y}(t_k) = \tilde{y}(t_{k-1}) + \left(f(\tilde{y}(t_{k-1})) + f(\tilde{y}(t_k))\right) \frac{t_k - t_{k-1}}{2}$$

given on slide 91 because the above equations for $\tilde{y}_2(t_k)$ and $\tilde{y}(t_k)$ are the same and hence $\tilde{y}(t_k) = \tilde{y}_2(t_k)$.

Lemma: Consistency of implicit trapezoidal rule time propagator The implicit trapezoidal rule time propagator $\tilde{\Phi}(y_0, t)$ satisfies

$$\|\tilde{\Phi}(y_0,t)-\Phi(y_0,t)\|=O(t^3) \qquad \text{for } t\to 0.$$

Proof. We obtain by repeatedly differentiating the one-step equation for the implicit trapezoidal rule that

$$\begin{split} \tilde{y}(t) &= y_0 + f(y_0) \frac{t}{2} + f(\tilde{y}(t)) \frac{t}{2}, \\ \dot{\tilde{y}}(t) &= f(y_0) \frac{1}{2} + f'(\tilde{y}(t)) \dot{\tilde{y}}(t) \frac{t}{2} + f(\tilde{y}(t)) \frac{1}{2}, \\ \ddot{\tilde{y}}(t) &= f''(\tilde{y}(t)) (\dot{\tilde{y}}(t))^2 \frac{t}{2} + f'(\tilde{y}(t)) \ddot{\tilde{y}}(t) \frac{t}{2} + 2f'(\tilde{y}(t)) \dot{\tilde{y}}(t) \frac{1}{2}. \end{split}$$

Setting t = 0, we conclude

$$\tilde{y}(0) = y_0 = y(0), \quad \dot{\tilde{y}}(0) = f(y_0) = \dot{y}(0), \quad \ddot{\tilde{y}}(0) = f'(y_0)\dot{\tilde{y}}(0) = \ddot{y}(0),$$

i.e. the zeroth to second derivatives of $\tilde{y}(t)$ and y(t) match and hence the implicit trapezoidal rule is third-order consistent like the explicit trapezoidal rule.

Corollary: Convergence of the implicit trapezoidal rule

Denote by y(t) the solution to $\dot{y}=f(y)$ and by $\tilde{y}(t)$ the numerical approximation obtained using the implicit trapezoidal rule with the equispaced temporal mesh $(t_k=\frac{k}{n}\,T)_{k=0}^n$. Then,

$$\|\tilde{y}(T) - y(T)\| = O\left(\exp(LT)\frac{T^3}{n^2}\right)$$
 for both $n, T \to \infty$,

assuming f(y) is Lipschitz continuous with Lipschitz constant L.

Proof. Analogous to the one on slide 30.

Numerical demonstration

See convergence().

Stability function of the implicit trapezoidal rule

Inserting f(y) = y into the one-step equation of the implicit trapezoidal rule, we obtain

$$\tilde{y}(t) = y_0 + y_0 \frac{t}{2} + \tilde{y}(t) \frac{t}{2} \qquad \Longleftrightarrow \qquad \tilde{y}(t) = \frac{1 + \frac{t}{2}}{1 - \frac{t}{2}} y_0.$$

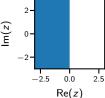
The stability function of the implicit trapezoidal rule is hence

$$R(z)=\tfrac{2+z}{2-z},$$

and the stability domain is

$$\left\{ \left| \frac{2+z}{2-z} \right| < 1 \right\} = \left\{ |2+z| < |2-z| \right\}.$$

This is the set of points which are closer to -2 than to +2, i.e. the left half plane $\{Re(z) < 0\}$.



Step size constraint for the implicit trapezoidal rule

We conclude from the above that the implicit trapezoidal rule incurs no step size constraint. This is demonstrated in stability_example().

Adaptive implicit time stepping

The implicit trapezoidal rule can be combined with a lower-order scheme to obtain an embedded method for step size selection. See embedded_implicit_ET_step(), adaptive_rk_example_2() and stepsize().

Explicit vs. implicit schemes

The key takeaway points of the last \sim 30 slides are the following.

- Explicit Runge-Kutta schemes applied to ODEs with attractive fixed points suffer from time step constraints which may significantly impact their performance.
- ▶ Implicit Runge-Kutta schemes can avoid time step constraints and the associated performance problems, but they require solving a potentially nonlinear equation in each time step.

Given these points, we conclude that implicit Runge-Kutta schemes *may* help when solving stiff ODEs, but it is not the case that they always outperform explicit Runge-Kutta methods.

In particular, if it is possible to isolate the stiff components and solving the remainder using an explicit scheme, then doing so is almost always the most efficient option.

Explicit vs. implicit schemes (continued)

If this is not possible, then it depends on the following criteria whether an explicit or implicit scheme is more efficient.

- ▶ How large is the separation of time scales? If the separation is small, then the increased step sizes enabled by an implicit scheme may not make up for the extra cost of solving equations in each time step.
- ▶ How expensive is it to solve the implicit equations? For example, if f(y) is a low-dimensional linear function, i.e. f(y) = Ay for some small matrix $A \in \mathbb{R}^{n \times n}$, then the implicit Euler scheme becomes

$$\tilde{y}(t) = y_0 + f(\tilde{y}(t)) t = y_0 + t A \tilde{y}(t)$$
 \iff $\tilde{y}(t) = (I - tA)^{-1} y_0$,

and this formula can be evaluated cheaply using an LU factorisation.

Summary

Theory of ODEs:

- $\dot{y} = f(y)$ has a unique (local) solution if f(y) is (locally) Lipschitz continuous.
- ▶ If f(y) has Lipschitz constant L, then y(t) is a Lipschitz continuous function of y(0), but the Lipschitz constant is $\exp(Lt)$.

Summary

Algorithm:

1:
$$\tilde{y}(0) = y(0)$$
,

2: **for**
$$k = 1, 2, ...$$
 do

3: **for**
$$i = 1, ..., s$$
 do

4:
$$\tilde{y}_{i}^{(k)} = \tilde{y}(t_{k-1}) + \sum_{i=1}^{i-1} f(\tilde{y}_{j}^{(k)}) W_{ij}(t_{k} - t_{k-1})$$

6:
$$\tilde{y}(t_k) = \tilde{y}(t_{k-1}) + \sum_{i=1}^{s} f(\tilde{y}_j^{(k)}) w_j(t_k - t_{k-1})$$

7: end for

$$\begin{array}{c|cccc} \theta_1 & w_{11} & \dots & w_{1s} \\ \vdots & \vdots & \ddots & \vdots \\ \theta_s & w_{s1} & \dots & w_{ss} \\ \hline & w_1 & \dots & w_s \end{array}$$

Summary

Convergence theory:

► Global error is weighted sum of local errors,

$$\|\tilde{y}_n - y_n\| \le \sum_{k=1}^n \exp(L(t_n - t_k)) \|\tilde{\Phi}_k(\tilde{y}_{k-1}) - \Phi_k(\tilde{y}_{k-1})\|.$$

▶ Local / consistency errors $\|\tilde{\Phi}_k(\tilde{y}_{k-1}) - \Phi_k(\tilde{y}_{k-1})\|$ can be estimated using Taylor's theorem.

Adaptive time stepping:

- ▶ Optimal efficiency is achieved if local errors are of approximately the same magnitude.
- ► Local errors can be estimated by comparing Runge-Kutta methods of different orders,

$$\|\tilde{\Phi}_{\mathbf{1}}(y_0,t) - \tilde{\Phi}_{\mathbf{2}}(y_0,t)\| = \|\tilde{\Phi}_{\mathbf{1}}(y_0,t) - \Phi(y_0,t)\| + O(t^{\rho_1+1}).$$

Appropriate step sizes can be determined using

$$\|\tilde{\Phi}(y_0,t)-\Phi(y_0,t)\|=O(t^p)\approx Ct^p.$$

Summary

Stability and implicit Runge-Kutta methods

- Explicit Runge-Kutta schemes applied to ODEs with attractive fixed points suffer from time step constraints which may significantly impact their performance.
- ▶ The time step constraint can be understood using linearisation, the stability function R(z) such that

$$\dot{y} = \lambda y \implies \tilde{y}(k \Delta t) = R(\lambda \Delta t)^k y_0,$$

and the stability domain $\{z \in \mathbb{C} \mid |R(z)| < 1\}$.

▶ Implicit Runge-Kutta schemes can avoid time step constraints and the associated performance problems, but they require solving a potentially nonlinear equation in each time step.