

## # INDEXING.

⇒ Database — RAM ✗ ⇒ Volatile.  
— HDD ✓ ⇒ Permanent Storage.

Students ⇒ HDD

id	name	batch-id	psp
1	X	3	80
2	Y	1	70
3	A	2	75.6
4	M	2	81.8
5	N	3	90.4
6	B	1	50.7

```
Select *  
from students  
where batch-id = 3 ;
```

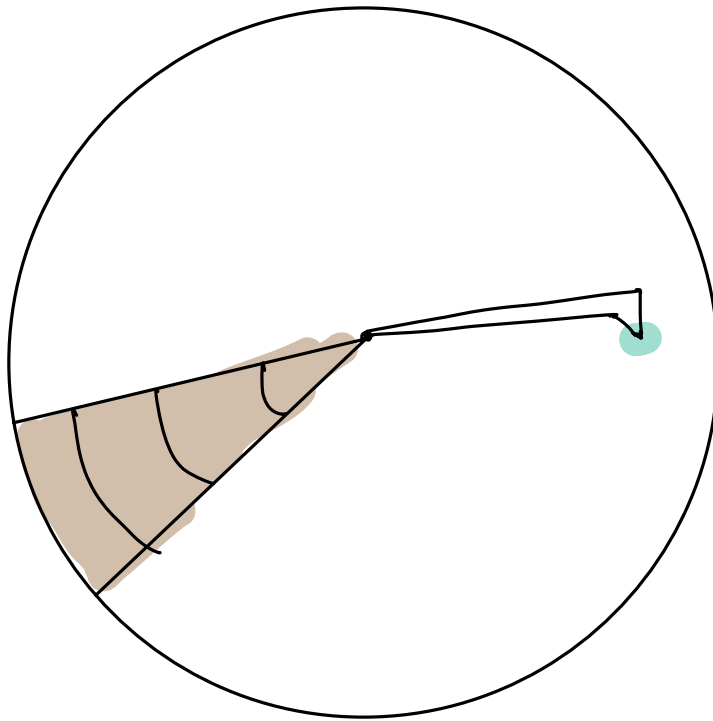
- CPU can't directly fetch the data from HDD.
- Data from HDD is brought to the RAM and then CPU will read the data from there.

Students  $\Rightarrow$  ADD

N rows

id	name	batch-id	psp
1	X	3	80
2	Y	1	90
3	A	2	75.6
4	M	2	81.8
5	N	3	90.4
6	B	1	50.4

→ B1  
→ B2  
→ B3



$\Rightarrow$  To get all the students from batch-id = 3, we'll have to go through all the blocks one by one & check in the block if there's a student with batch-id = 3 inside that block.

$\hookrightarrow$   $O(N)$

```

Select *
from students
where id = 4;

```

id → Primary Key  
 ↓  
 Sorted based on PK.  
 → Unique.

⇒ Iterate the table row by row & get the student with  $id = 4$ , we need not to go to the rows after  $id = 4$ .

Students

id	name	batch-id	PSP	
1	X	3	80	→ B1
2	Y	1	70	
3	A	2	75.6	→ B2
4	M	2	81.8	
5	N	3	90.4	→ B3
6	B	1	50.7	

```

Select *
from students
where id = 4;

```

Index table.

Block No	Address
B <sub>1</sub>	Ad <sub>1</sub>
B <sub>2</sub>	Ad <sub>2</sub>
B <sub>3</sub>	Ad <sub>3</sub>
⋮	⋮
⋮	⋮
⋮	⋮
⋮	⋮

PK

## Students

id	name	batch-id	psp
1	X	3	80
2	Y	1	70
3	A	2	75.6
4	M	2	81.8
5	N	3	90.4
6	B	1	50.7

AD6

7

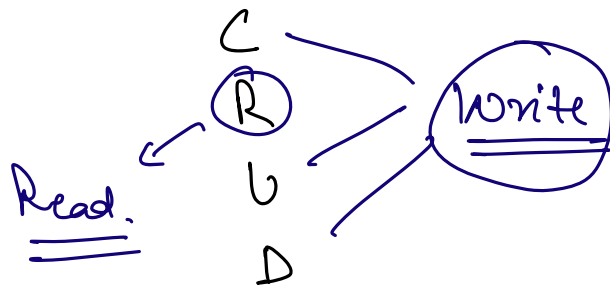
## Index table

1	AD1
2	AD2
3	AD3
4	AD4
.	.
6	AD6
.	.

Select \*

from students

where id = 6;



#

```
Select *  
from students  
where batch_id = 3 ;
```

Index

batch_id	
1	-
1	-
1	-
1	-
1	-

Indexing

- Makes our queries faster by reducing the no. of disk accesses.
- For any Create / Update or Delete operation on original table, Index table needs to be updated.
- Indexing makes writes (C|U|D) slower.

## Note

→ Create index if it is actually required.

→ Don't create index at the time of table creation, Create index based on access pattern.

type of queries we are getting

show indexes from film;

explain select \* from film where length = 105;

create index idx\_film\_length on film(length); -- creating the index

drop index idx\_film\_length on film; -- deleting the index.

explain select \* from film where film\_id = 105;

explain select \* from film where rental\_duration = 3;

Index table is also sorted.

index table

1	—
2	—
3	—
4	—
5	—
6	—

5	7	1	100	90	10	8	2
---	---	---	-----	----	----	---	---

Searching in a non sorted Array  $\Rightarrow$   $O(N)$ .

1	2	5	7	8	10	90	100
---	---	---	---	---	----	----	-----

②

$$N \rightarrow \frac{N}{2} \rightarrow \frac{N}{4} \rightarrow \dots \dots \dots \textcircled{1}$$

$\Rightarrow$  Binary Search. :  $O(\log N)$ .

Index table : B / B+ Trees.

$\hookrightarrow$   $O(\log N)$

## # Indexing on Strings.

Students

id	name	batch-id	psp
1	Sushant	3	80
2	Mukta	1	90
3	Murli	2	75.6
4	Abhigyan	2	81.8
5	Sai	3	90.4
6	Shlok	1	50.7

```
Select *  
from Student  
where name = 'Abhigyan';
```

⇒ Full Table Scan.

# Create index on name column.

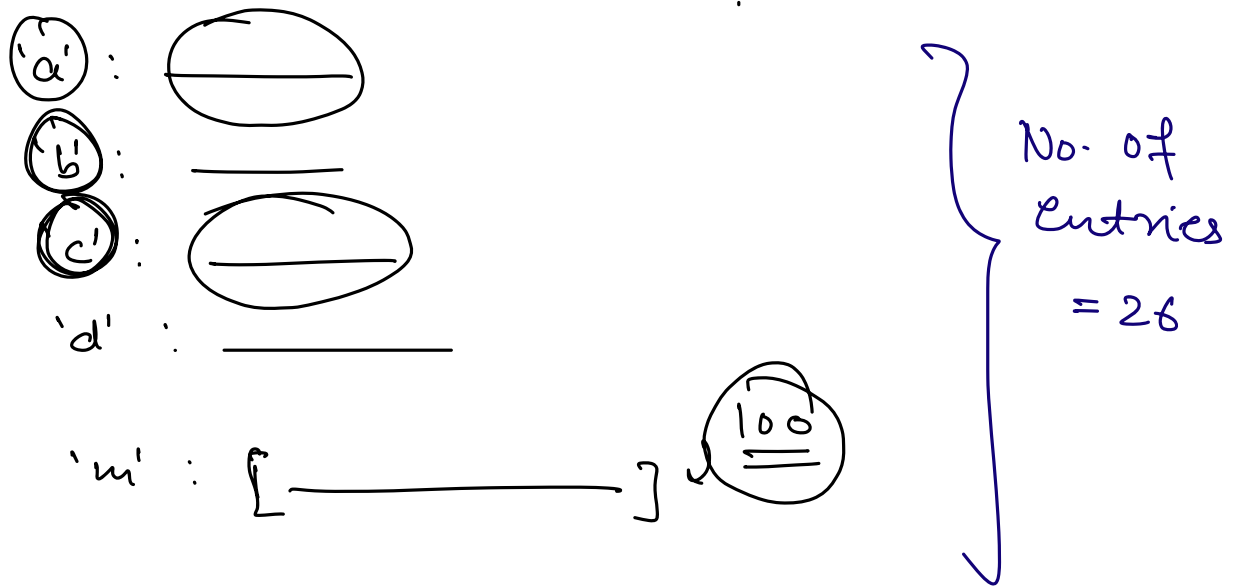
Sushant	—
Mukta	—
Murli	—
Deepak	—
Shlok	—
—	
—	
—	
—	
—	

→ Index on name column will create a very big table, which will consume a lot of space.

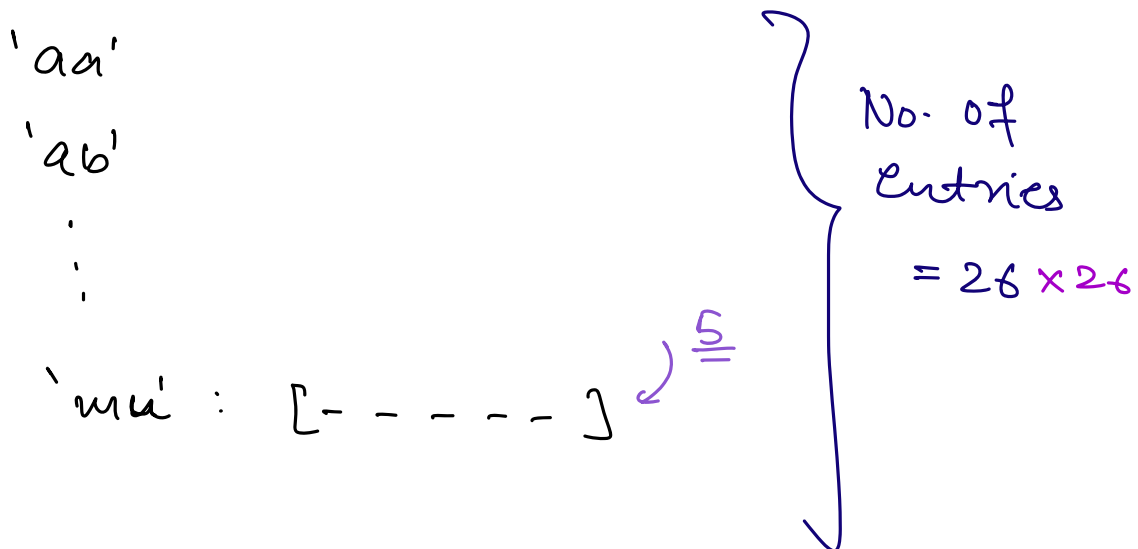


⇒ Index table is persisted permanently inside the HDP & there will be a copy of index table inside the RAM.

Index of first character only



Index of first 2 characters



```
select * from city;
```

```
explain select * from city where city = 'Goa';
```

```
create index idx_city on city(city(3));
```

```
drop index idx_city on city;
```

```
show indexes from city;
```

```
-- without index -> 600
```

```
-- with index on 1 character -> 13
```

```
-- with index on 2 characters -> 7
```

```
-- with index on 3 characters -> 4
```

```
-- with index on 4 characters -> 1
```

```
-- with index on complete city column -> 1
```