Logic Programming

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Outline

- Introduction to Prolog
- Knowledge representation
- Prolog's search strategy

Introduction to Prolog

- PROgramming in LOGic.
- Prolog is based on First-Order logic.
- Prolog is known to be a difficult language to master.
 - The system does not give too much help to the programmer to employ structured programming concepts.

Introduction to Prolog (cont')

- The special attention should be given to understanding variables in Prolog.
- The programmer should pay careful attention to the issue of backtracking.

Knowledge Representation

- How do we represent what we know?
 - Knowledge how procedural knowledge such as how to drive a car.
 - Knowledge that declarative knowledge such as knowing the speed limit for a car on a motorway.

Knowledge Representation (cont')

- Propositional calculus.
 - The propositional calculus is based on statements which have truth values (true or false).
 - An example:
 - If p stands for "fred is rich" and q for "fred is tall".
 - Connection symbol ∧, ∨, ⇔, ⇒, ¬

Knowledge Representation (cont')

- First order predicate calculus.
 - Permits the description of relations and the use of variables.
 - Variables: X, Y, Z
 - Constant: x, y, john, jean
 - Predicate: brother (x, y)
 - Formulae: brother(x, y) ∧brother(y, z)
 - Sentence: brother (john, jean)

We turn to Prolog

- "the capital of france is paris"
 - has capital(france, paris).
- "jim is tall"
 - jim(tall).
 - tall(jim).*
- There must be no space between the predicate name and "(". The whole thing also ends in a "."

Prolog Constants

- Constant is a string of characters starting with a lower case letter.
 - loves (jane, jim).
 - jane and jim are constant.
- No predicate may be variable
 - X(jane, jim).

Examples of statement in Prolog form

- bill likes ice-cream.
- bill is tall
- jane hits jimmy with the cricket bat
- john travels to london by train

Multiple Clauses

- Examples
 - bill only eats chocolate, bananas or cheese.
 - the square root of 16 is 4 or -4.
 - wales, ireland and scotland are all countries.

Rules

The format is:

```
divisible_by_two:- even.
```

- This is a non-unit clause.
- The head is divisible by two.
- The body is even.
- divisible by two is true if even is true

Rules (continue)

- No more than one goal is allowed in the head.
- We cannot translate rich(fred) ⇒
 happy(fred) ∧ powerful(fred)
 directly into the Prolog version
 happy(fred), powerful(fred): rich(fred).

Rules (continue)

- A number is divisible by two if it is even.
- We can express this with the help of the logical variable. Here is the improved rule:
- divisible by two(X):- even(X).

Semantics of x Municipal 2 200 no healer rule

divisible by two (X):- even (X).

- If we can find a value of X that satisfies the goal even (X) then we have also found a number that satisfies the goal divisible by two(X).

The Logical Variable

- If an object is referred to by a name starting with a capital letter then the object has the status of a *logical* variable.
- The same variables in one rule refer to the same object.

The Logical Variable (continue)

- The logical variable cannot be overwritten with a new value.
- For example, in Pascal:
 - X := 1; X := 2;
- In Prolog:
 - $X=1 \land X=2$.
 - cannot be true as X cannot be both '2' and '1' simultaneously. An attempt to make a logical variable take a new value will fail.

Rules and Conjunction

- A man is happy if he is rich and famous
- might translate to:

```
happy(Person):-
  man(Person),
  rich(Person),
  famous(Person).
```

Rules and Disjunctions

- Someone is happy if they are healthy, wealthy or wise.
- · translates to:

```
happy(Person):- healthy(Person).
happy(Person):- wealthy(Person).
happy(Person):- wise(Person).
```

Both Disjunctions and Conjunctions

```
happy (Person) :- healthy (Person), woman (Person).
happy (Person): -wealthy (Person), woman (Person).
happy (Person): -wise (Person), woman (Person).
```

Prolog's Search Strategy

- Prolog use depth-first search.
- Prolog enables the programmer to implement other search methods quite easily.
- Prolog is an interactive system.

Queries and Disjunctions

- A query is a goal which is submitted to Prolog in order to determine whether this goal is true or false.
- Prolog normally expects queries it prints the prompt:

Queries and Disjunctions (continue)

 Perhaps we would like to determine whether or not

```
woman (jane)
```

We can type this

```
?- woman (jane).
```

A Simple Conjunction

Program Database

A Simple Conjunction (continue)

- From the knowledge base, if we would like to find "Is Jean happy?".
- Prolog try to match happy(jean) with the knowledge base.

A Simple Conjunction (continue)

- So two sub goals are:
 - woman(jean)
 - wealthy(jean)
- There is a possible match but we cannot unify
 - wealthy(fred) **With** wealthy(jean)

Unification

- Unification is a two way matching process
 - book(waverley, X) and book(Y, scott)
 - X/scott, Y/waverley
- Examples
 - X=fred. succeeds
 - jane=fred fail
 - Y=fred, X=Y succeeds
 - X=happy(jim) succeeds

Recursion

- Prolog depends heavily upon it
- Examples
 - One of my ancestors is one of my parents or one of their ancestors.
 - A string of characters is a single character or a single character followed by a string of characters

Recursion (continue)

 Examples X (Thurshord Y • ancestor(X, Y) :- parent(X, Y). ancestor(X,Y) :- parent(X,Z), \rangle V ancestor (Z, Y). • Program Database ancestor (X,Y):- parent (X,Y). ancestor(X,Y) :- parent(X,Z), ancestor (Z, Y). parent (bob, bill). parent (bill, jane).

parent (jane, jim). parent (bob, babe).

Lists

- Lists can be regarded as special Prolog structures.
- Lists can be used to represent an ordered sequence of Prolog terms.
- Examples
 - [ice_cream, coffee, chocolate]
 - [a,b,c,c,d,e]

- List destruction
 - [X|Y] = [f,r,e,d].
 - will result in X=f.
 - The first element of the list is known as the HEAD of the list.
 - Y = [r, e, d].
 - The list formed by deleting the head is the TAIL of the list.

- List construction
 - X = [r, e, d].
 - Result Wanted=[b|X].
 - will result in Result_Wanted=[b,r,e,d].

- Bigger Chunks:
 - Suppose you want to stick the elements a,
 and c onto the front of the list x to make
 a new list y.
 - This can be done with Y=[a,b,c|X]

- Bigger Chunks:
 - Conversely, suppose you want to take three elements off the front of a list x.
 - This can be done with X=[A,B,C|Y].
 - Y is available for use as the remaining list.

Some Possible Matches

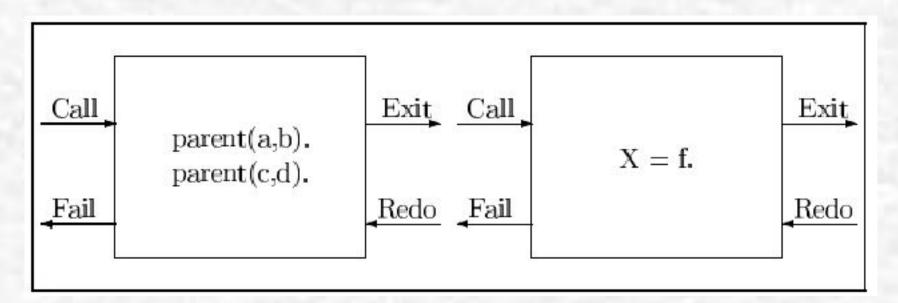
```
    [b,a,d]=[d,a,b] fails
    [X]=[b,a,d] fails
    [X|Y]=[he,is,a,cat] succeeds
    [X,Y|Z]=[a,b,c,d] succeeds
    [X|Y]=[] fails
    [X|Y]=[[a,[b,c]],d] succeeds
    [X|Y]=[a] succeeds
```

A recursive program using list

- write/1 is a build in predicate, the argument is Prolog term.
- n1/0 is new line command

The Box Model of Execution

• We can think in term of *procedure* rather than *predicate*.



The Box Model of Execution

(continue)

- The control flows into the box through the Call port.
- Then, Prolog seek a clause with a head that unify with the goal.
- Seek solutions to all the subgoals in the body of the successful clause.
- If the unification fails for all clauses then control would pass out of the *Fail* port.

The Box Model of Execution

(continue)

- Control reaches the *Exit* port if the procedure succeeds.
- The Redo port can only be reached if the procedure call has been successful and some subsequent goal has failed.

The Flow of Control

Call: parent(X,Y)Exit: parent(a,b)

Call: a=f

Fail: a=f

Redo: parent(X,Y)

Exit: parent(c,d)

Call: c=f

Fail: c=f

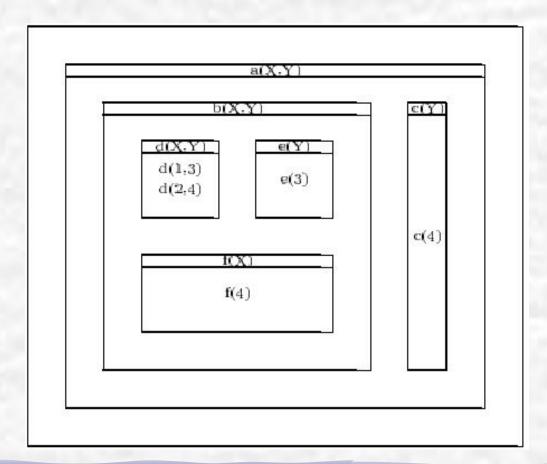
Now backtracking

Now backtracking

Redo: parent(X,Y)Fail: parent(X,Y)

The Flow of Control (continue)

```
Program Database
a(X,Y):-
           b(X,Y),
           c(Y).
b(X,Y):-
           d(X,Y),
           e(Y).
b(X,Y):-
           f(X).
c(4).
d(1,3).
d(2,4).
e(3).
f(4).
```



The Flow of Control (continue)

```
Call:
          a(X,Y)
                               Call:
                                         b(X,Y)
                                                              Call:
                                                                        d(X,Y)
                                                             Exit:
                                                                        d(1,3)
                                                              Call:
                                                                        e(3)
                                                             Exit:
                                                                        e(3)
                               Exit:
                                         b(1,3)
                               Call:
                                         c(3)
                               Fail:
                                         c(3)
                                       Now backtracking
                               Redo:
                                         b(X,Y)
                                                              Redo:
                                                                        e(3)
                                                             Fail:
                                                                        e(3)
                                                              Redo:
                                                                        d(X,Y)
                                                             Exit:
                                                                        d(2,4)
                                                              Call:
                                                                        e(4)
                                                             Fail:
                                                                        e(4)
                                       Now backtracking
                                                              Call:
                                                                        f(X)
                                                             Exit:
                                                                        f(4)
                               Exit:
                                         b(4,Y)
                               Call:
                                         c(Y)
                               Exit:
                                         c(4)
Exit:
          a(4,4)
```

Practical Matter

- Exit Prolog by the following commands
 - ?- halt.
 - ^D
 - end of file.
- Loading file(s)
 - ?- consult (filename) .
 - ?- consult('test.pl').
 - ?- consult('test').
 - ?- consult(test).

Practical Matter (continue)

- Undeclared predicates could be handled by
 - :- dynamic predicate/arity.
 - :- dynamic foo/1.
 - :- dynamic two hand/1.
 - ?- unknown(X,Y).
 - ?- unknown(X, trace).
 - ?- unknown(X, fail).

Practical Matter (continue)

- Singleton variable
 - A singleton variable occurs in a clause if a variable is mentioned once and once only.
 - mammal(Animal): dog(Aminal).
 - Warning: singleton variable Animal in procedure mammal/1
 - Warning: singleton variable Aminal in procedure mammal/1

Practical Matter (continue)

- Singleton variable
 - You can use the anonymous variable.
 - The *anonymous* variables can be written by the string starting with the underscore (_).
 - member(X, [X|]).

Programming Techniques and Lists Processing

The reversibility of Prolog program

```
Program Database square(1,1).
square(2,4).
square(3,9).
square(4,16).
square(5,25).
square(6,36).
```

We can ask all the following queries.

```
?- square(2,X).
?- square(X,5).
?- square(X,Y).
?- square(2,3).
```

- Evaluation in Prolog
 - Pascal: Y:=2+1;
 - Prolog: Y=2+1.
 - This command only binds the variable Y with the clause 2+1 and gives the value of Y as 2+1
 - Prolog: Y is 2+1.
 - After the above command, the value of Y is 3.

- successor(X,Y):- Y is X + 1.
 - ?- successor(3,X).
 - ?- successor(X, 4).
 - ?- successor(X,Y).
 - ?- successor(3,5).

- Calling patterns
 - For any given predicate with arity greater than 0, each argument may be intended to have one of three calling patterns:
 - Input |indicated with a +
 - Output |indicated with a -
 - Indeterminate |indicated with a ? (+ or -)

- For example, successor/2 above requires a calling pattern of
 - 1st argument must be +
 - 2nd argument can be + or and is therefore ?
 - We write this as
 - :- mode successor(+,?).

Lists Processing

Program Patterns

Test for Existence

```
    list_existence_test(Info,[Head|Tail]):-
        element_has_property(Info,Head).
    list_existence_test(Info,[Head|Tail]):-
        list_existence_test(Info,Tail).
```

Example

```
nested_list([Head|Tail]):- sublist(Head).
nested_list([Head|Tail]):- nested_list(Tail).
sublist([]).
```

Example

- member (Element, [Element | Tail]).
- member (Element, [Head|Tail]): -member (Element, Tail).

Test all elements

```
• test_all_have_property(Info,[]).
```

```
    test_all_have_property(Info,[Head|Tail]):-
        element_has_property(Info,Head),
        test_all_have_property(Info,Tail).
```

Example

```
all_digits([]).
```

```
    all_digits([Head|Tail]):-
        member(Head,[0,1,2,3,4,5,6,7,8,9]),
        all_digits(Tail).
```

Return a Result — Having processed one element

```
    return_after_event(Info,[H|T],Result):-
        property(Info,H),
        result(Info,H,T,Result).
```

```
return_after_event(Info,[H|T],Ans):-
return_after_event(Info,T,Ans).
```

Example

```
    everything_after_a([Head|Tail],Ans):-
    everything_after_a(Tail,Ans).
```

Return a Result — Having processed all elements

```
process_all(Info,[],[]).
process_all(Info,[H1|T1],[H2|T2]):-
process_one(Info,H1,H2),
```

process all(Info,T1,T2).

Example

```
triple([],[]).
triple([H1|T1],[H2|T2]):-
H2 is 3*H1,
triple(T1,T2).
```

Control and Negation

- Some useful predicate for control
 - true/0
 - father (jim, fred) .
 - father(jim, fred) :- true.
 - fail/0
 - lives_forever(X) :- fail.
 - repeat/0
 - test :- repeat, write (hello) red fail. block norman

Control and Negation (continue)

- call/1
 - call(write(hello)).

1+ 60 7

Negation

- \+/1
- Equivalent to close for assumption hatoroxizons ev. 2 onulas
 - man(jim).
 - man(fred).
 - woman $(X) :- \ + \ man(X)$.

Some General Program Schemata

• Generate – Test

- Finite generator
- Infinite generator

Generate - Test

```
integer_with_two_digit_square(X) :- int(X),
     test square(X).

test square(X) :- Y is X*X, Y >= 10, Y < 100.</pre>
```

Finite generator

```
int(1). int(2). int(3). int(4). int(5).
```

Infinite Generator

```
int(1).
int(N) :- int(N1), N is N1+1.
```

Test - Process

```
    test_process(Info,X,Y):-
        test(Info,X),
        process(Info,X,Y).
```

Example

```
• parity(X, Y) :- odd(X), Y = odd.
```

• parity(X,Y) :- $\+$ (odd(X)), Y=even.

Failure-Driven Loop

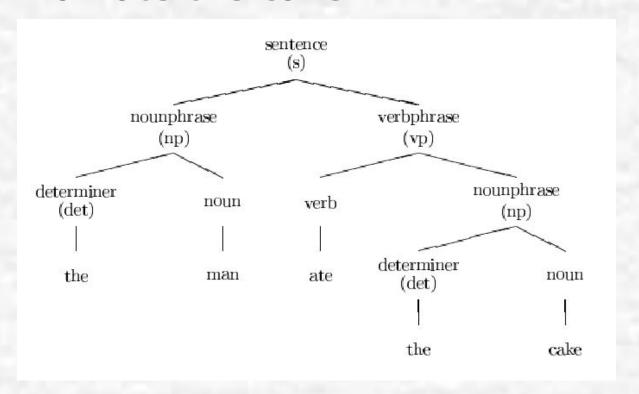
- failure_driven_loop(Info): generate(Info, Term),
 fail.
- failure_driven_loop(Info).
- An example
 - int(1). int(2). int(3). int(4). int(5).
 - print int :- int(X), write(X), nl, fail.
 - print_int.

Parsing in Prolog

- Simple English Syntax
 - Unit: sentence Constructed from noun_phrase followed by a verb_phrase
 - Unit: noun phrase Constructed from: proper noun or determiner followed by a noun
 - Unit: verb phrase Constructed from: verb or verb followed by noun_phrase
 - Unit: determiner Examples: a, the
 - Unit: noun Examples: man, cake
 - Unit verb: Examples: ate

Simple English Syntax

• The man ate the cake.



First Attempt at Parsing

- sentence(S):-append(NP, VP, S),
 noun_phrase(NP), verb_phrase(VP).
- noun_phrase(NP):-append(Det, Noun, NP),
 determiner(Det), noun(Noun).
- verb_phrase(VP):-append(Verb,NP,VP),
 verb(Verb), noun phrase(NP).
- determiner([a]). determiner([the]).
- noun([man]). noun([cake]).
- verb([ate]).

A Second Approach

- sentence(S,S0):-noun_phrase(S,S1),
 verb_phrase(S1,S0).
- noun_phrase(NP,NP0): determiner(NP,NP1), noun(NP1,NP0).
- verb_phrase(VP, VP0):-verb(VP, VP1),
 noun phrase(VP1, VP0).
- determiner([a|Rest], Rest).
- determiner([the|R],R).
- noun([man|R],R). noun([cake|R],R).
- verb([ate|R],R).

Prolog Grammar Rules

- sentence --> noun_phrase, verb_phrase.
- noun phrase --> determiner, noun.
- verb_phrase --> verb, noun_phrase.
- determiner --> [a].
- determiner --> [the].
- noun --> [man].
- noun --> [cake].
- verb --> [ate].
- noun phrase --> determiner, adjectives, noun.
- adjectives --> adjective.
- adjectives --> adjective, adjectives.
- adjective --> [young].

Prolog Grammar Rules

- Prolog enables you to write in:
 - sentence --> noun phrase, verb phrase.
- and Prolog turns this into:
 - sentence(S,S0):-noun_phrase(S,S1),
 verb phrase(S1,S0).
- And
 - adjective --> [young].
- into
 - adjective(A, A0):-'C'(A, young, A0).
 - 'C'([H|T],H,T).

Prolog Grammar Rules

- How to Extract a Parse Tree
 - sentence([[np,NP],[vp,VP]]) --> noun_phrase(NP),
 verb phrase(VP).
 - noun_phrase([[det,Det],[noun,Noun]]) --> determiner(Det), noun(Noun).
 - determiner(the) --> [the].
- So what structure is returned from solving the goal:
 - sentence(Structure,[the,man,ate,a,cake],[])
- The result is:
 - [[np,[[det,the],[noun,man]]],[vp,[...

Special Control Predicate

- Cut (!/0)
 - On backtracking, all attempts to redo a subgoal to the left of the cut results in the subgoal immediately failing.
 - On backtracking, into the predicate once the call had exited: if one of the clauses defining the predicate had previously contained a cut that had been executed then no other clauses for that predicate may be used to resatisfy the goal being redone.

Special Control Predicate (continue)

Commit

```
• skin(X,Y) :- thai(X), Y=yellow. skin(X,Y) :- asia(X), Y=yellow.
```

Special Control Predicate (continue)

Make determination

```
sum(1,1).
sum(N,Ans): NewN is N-1,
sum(NewN,Ans1),
Ans is Ans1+N.
sum(1,1):-!.
sum(N,Ans): NewN is N-1,
sum(NewN,Ans1),
Ans is Ans1+N.
```

Special Control Predicate (continue)

Fail Goal Now

```
• \+ (Goal):- call(Goal),!,fail. 
\+ (Goal).
```

```
    woman(X):- man(X),!, fail.
    woman(X).
```

Changing the Program

- assert (C): Assert clause C
- asserta (C): Assert C as first clause
- assertz(C): Assert C as last clause
- retract(C): Erase the first clause of form C
- abolish (Name, Arity): Abolish the procedure named F with arity N

Changing the Program (continue)

- assert (man (bill)).
- asserta (man (bob)).
- assertz (man (jim)).
- retract (man (bob)).
- abolish (man, 1).

Input/Output

Input/Output and File

- see/1: Take input from the named file
- seen/0: Close the current input stream and take input from user
- seeing/1: Returns name of current input stream

Input/Output and File (continue)

- test:-read(X),\+(X = -1),
 double(X,Y),write(Y),nl,test.
- go:-see(in), test, seen.

Input/Output and File (continue)

- tell/1: Send output to the named file
- told/0: Close the current output stream and send output to user
- telling/1: Returns name of current output stream

Input/Output and File (continue)

- test:-read(X),\+(X = -1),
 double(X,Y),write(Y),nl,test.
- go:-tell(out), see(in), test, seen, told.

Prolog Examples

- move(state(middle,onbox,middle,hasnot),grasp, state(middle,onbox,middle,has)).
- move(state(P,onfloor,P,H),climb,state(P,onbox,P,H)).
- move(state(P1,onfloor,P1,H),push(P1,P2), state(P2,onfloor,P2,H)).
- move(state(P1,onfloor,B,H),walk(P1,P2), state(P2,onfloor,B,H)).

Prolog Examples (continue)

- depth_first(State, State, []).
- depth_first(StartState, GoalState, Ans): findsuccessor(StartState, Successor, Operator),
 depth_first(Successor, GoalState, Ans1),
 Ans=[Operator|Ans1].
- findsuccessor (OldState, NewState, Operator) : move (OldState, Operator, NewState).