

Course Code: MT-3001	Course Name: Graph Theory
Instructor Names :	Dr. Nazish Kanwal, Mr. M Abdul Basit Khan
Student Roll No:	Section No:

Instructions:

1. Answer all questions on answer script . Credit will be awarded for correct content and clarity of presentation.
2. There are 3 questions and 2 pages.

Time: 60 minutes.

Max Points : 30

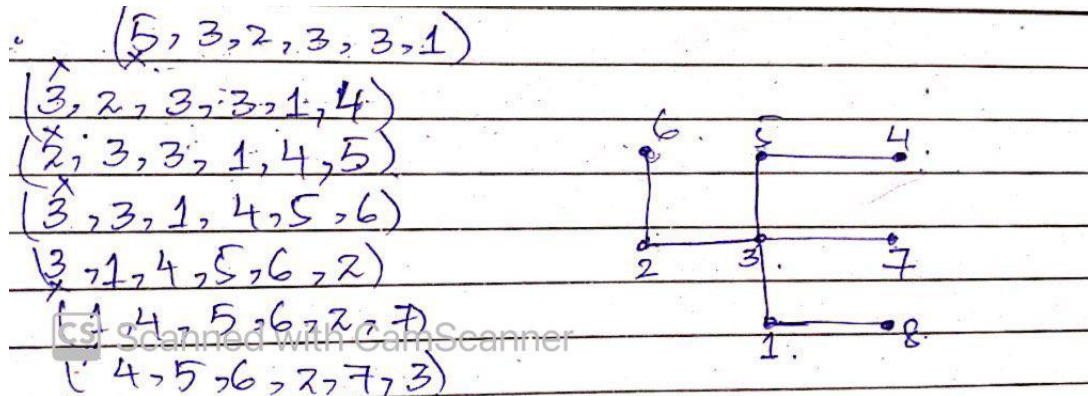
Question 1: CLO 2 12 points

Attempt each part.

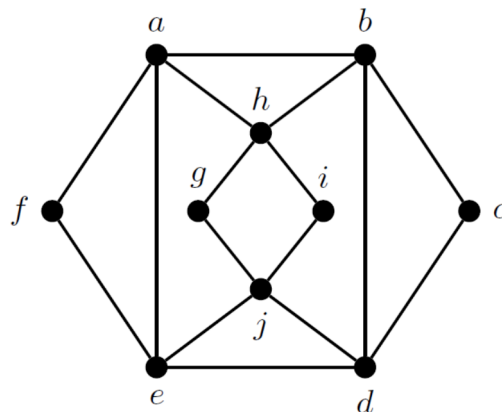
- (a) 4 points Draw the labelled tree corresponds to the given prüfer sequence.

(5, 3, 2, 3, 3, 1)

Solution:

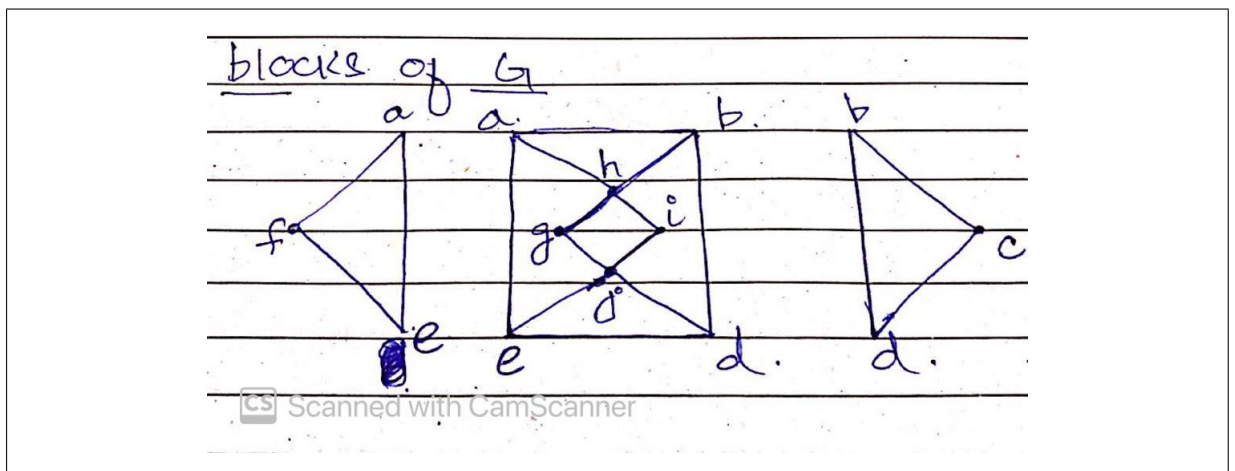


- (b) Consider the following graph G , and answer each subparts.



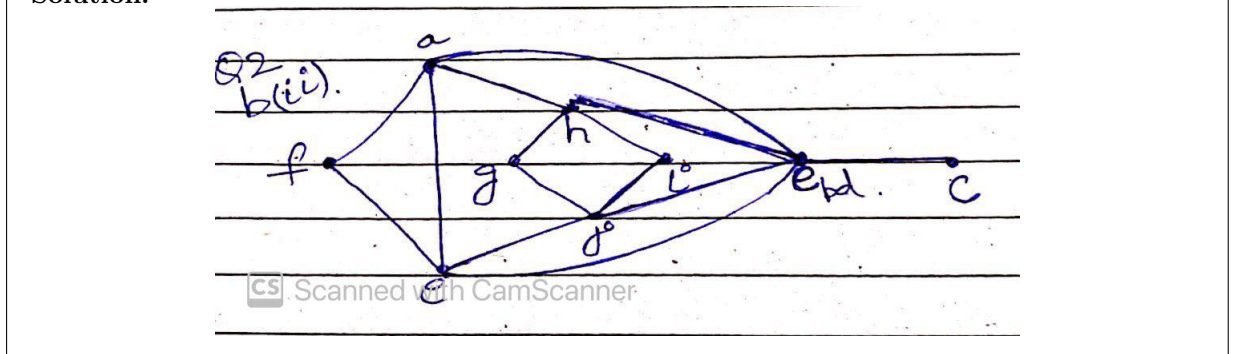
- i. 2 points Determine the blocks of G .

Solution:



- ii. 1 point Find G/e , the graph obtained by contracting the edge $e = bd$.

Solution:



- iii. 2 points Find all the edge-disjoint paths from $h - e$.

Solution: hae (or $hafe$), $hije$ (or $hgje$ or $hgjde$), $hbde$ (or $hbcde$ or $habcde$).

- iv. 2 points Find all internally disjoint paths from $h - j$.

Solution: hij , hgj , $haej$ (or $hafej$), $hbdj$ (or $hbdcj$).

- v. 1 point Find a vertex-separating set for vertex h and j .

Solution: $\{a, b, i, g\}$.

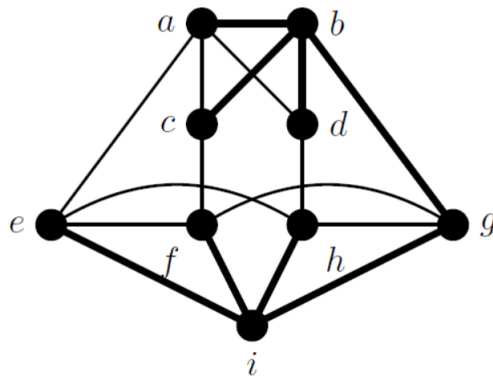
Question 2: CLO 2 11 points

- (a) 4 points Prove that, an edge e is a bridge of G if and only if there exist vertices x and y such that e is on every $x - y$ path.

Solution: Suppose that an edge $e = ab$ is a bridge of G , then $G - e$ is disconnected. This implies $G - e$ must have at least two components. Let x and y be vertices in different components of $G - e$. Since G is connected so $x - y$ path exists in G .

Conversely, let e be an arbitrary edge and there exist vertices x and y in G such that e is on every $x - y$ path, then none of these path exists in $G - e$, this implies e disconnect G . Hence e is a bridge for G .

- (b) Consider the following graph G .



- i. 1 point State Whitney's Theorem.

Solution: For any graph G , $\kappa(G) \leq \kappa'(G) \leq \delta(G)$.

- ii. 4 points For the graph G , prove the Whitney's theorem.

Solution: We need to prove that G satisfy the inequality: $\kappa(G) \leq \kappa'(G) \leq \delta(G)$. Since $\delta(G) = 3$ and vertex-cut set of G is $\{a, b, f\}$ and $\{a, b, h\}$, so $\kappa(G) = 3$ and edge-cut set of G is $\{ac, bc, cf\}$ and $\{ad, bd, hd\}$ so $\kappa'(G) = 3$. Hence

$$\kappa(G) \leq \kappa'(G) \leq \delta(G) \Rightarrow 3 \leq 3 \leq 3.$$

- iii. 1 point Find vertex-cut set for G .

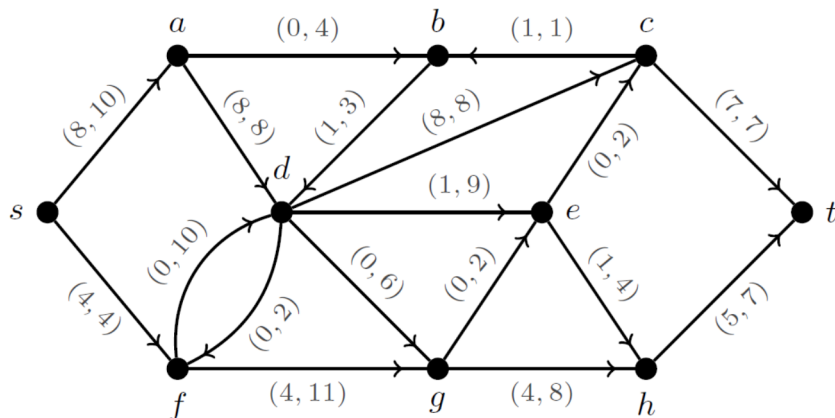
Solution: $\{a, b, f\}$ and $\{a, b, h\}$.

- iv. 1 point Find edge-cut set for G .

Solution: $\{ac, bc, cf\}$ and $\{ad, bd, hd\}$.

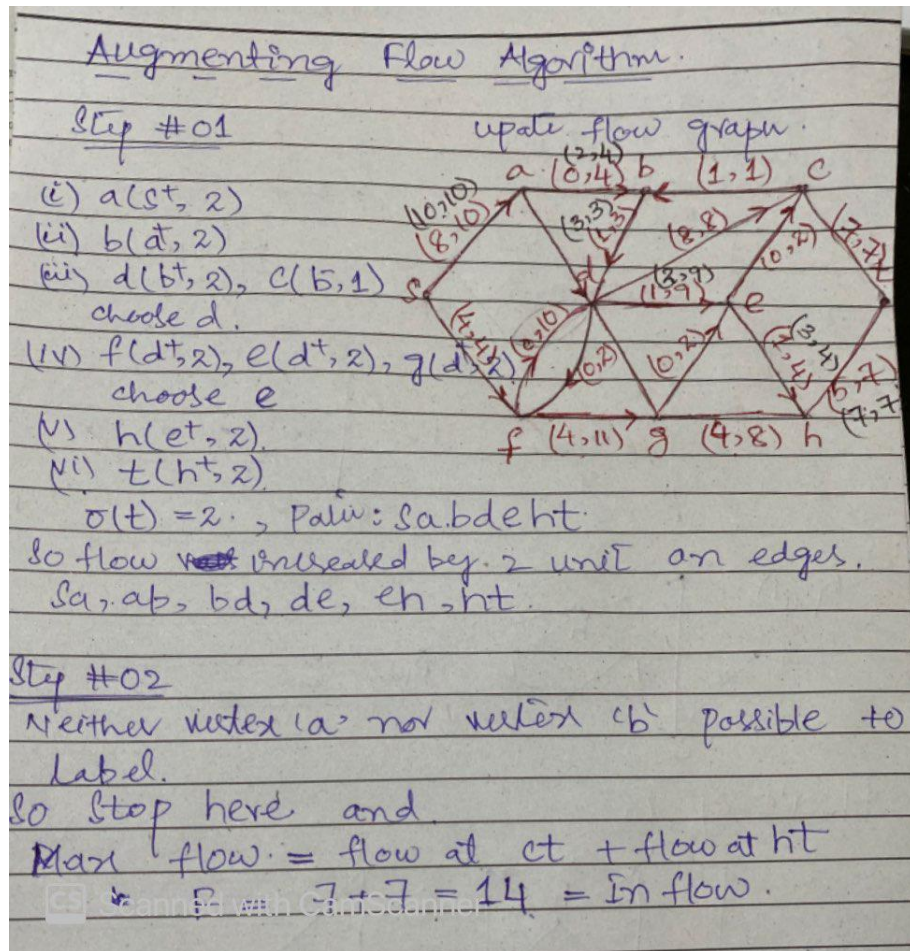
Question 3:CLO 2 7 points

- (a) For the problem below:



- i. 4 points Use Augmenting Flow Algorithm to maximize the flow.

Solution:



- ii. 1 point Use Min-Cut method to find a minimum cut.

Solution: In the last step of Augmenting flow algorithm, it is not possible to label any vertex. So, $P = \{s\}$ and $\bar{P} = \{a, b, c, d, e, f, g, h, t\}$ and the edges involve in minimum cut are sa and sb . Therefore, capacity of minimum cut is $C = c(sa) + c(sb) = 10 + 4 = 14$, and equal to value of maximum flow.

- (b) 2 points The algorithms given below are written for connected graphs as input. Determine the output of each of the algorithms below if the input was a disconnected graph.
- Kruskal's Algorithm.
 - Prim's Algorithm
 - Depth-First Search
 - Breadth-First Search

Solution: In above mentioned algorithms the output of the graphs are disconnected graphs of atleast two components in which each component is itself a tree. OR output is the forest.