



National Textile University

Department of Computer Science

Subject

Operating System

Submitted to:

Sir Nasir Mehmood

Submitted by:

Kashmir Jamshaid

Registration Number

23-NTU-CS-1167

Lab No.

07

Semester

5th

Task_01 :

Input

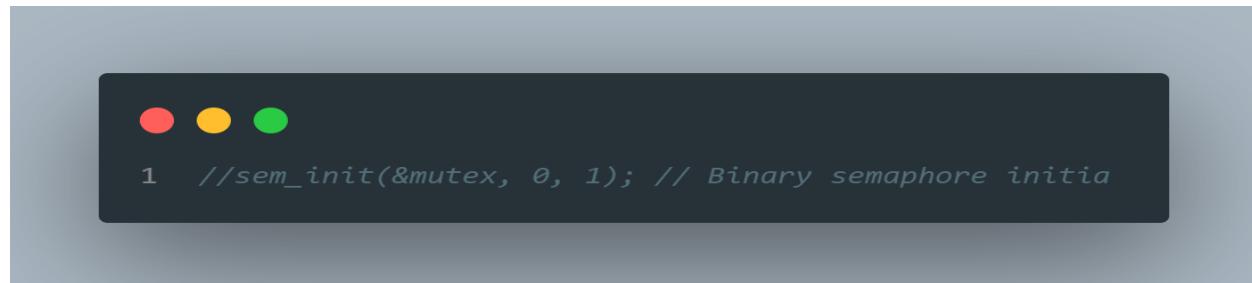


```
1 #include <stdio.h>
2 #include <pthread.h>
3 #include <semaphore.h>
4 #include <unistd.h>
5 sem_t mutex; // Binary semaphore
6 int counter = 0;
7 void* thread_function(void* arg) {
8     int id = *(int*)arg;
9     for (int i = 0; i < 5; i++) {
10         printf("Thread %d: Waiting...\n", id);
11         sem_wait(&mutex); // Acquire
12
13     // Critical section
14     counter++;
15     printf("Thread %d: In critical section | Counter = %d\n", id,
16     counter);
17     sleep(1);
18     sem_post(&mutex); // Release
19     sleep(1);
20 }
21 return NULL;
22 }
23 int main() {
24     sem_init(&mutex, 0, 1); // Binary semaphore initialized to 1
25     pthread_t t1, t2;
26     int id1 = 1, id2 = 2;
27     pthread_create(&t1, NULL, thread_function, &id1);
28     pthread_create(&t2, NULL, thread_function, &id2);
29     pthread_join(t1, NULL);
30     pthread_join(t2, NULL);
31     printf("Final Counter Value: %d\n", counter);
32     sem_destroy(&mutex);
33     return 0;
34 }
```

Out put :

```
> TERMINAL
● kashmirkj@Kashmirpc:~/Operating_System/Lab_07$ gcc Task1.c -o Task1 -lpthread
● kashmirkj@Kashmirpc:~/Operating_System/Lab_07$ ./Task1
Thread 1: Waiting...
Thread 1: In critical section | Counter = 1
Thread 2: Waiting...
Thread 2: In critical section | Counter = 2
Thread 1: Waiting...
Thread 1: In critical section | Counter = 3
Thread 2: Waiting...
Thread 2: In critical section | Counter = 4
Thread 1: Waiting...
Thread 1: In critical section | Counter = 5
Thread 2: Waiting...
Thread 2: In critical section | Counter = 6
Thread 1: Waiting...
Thread 1: In critical section | Counter = 7
Thread 2: Waiting...
Thread 2: In critical section | Counter = 8
Thread 1: Waiting...
Thread 1: In critical section | Counter = 9
Thread 2: Waiting...
Thread 2: In critical section | Counter = 10
Final Counter Value: 10
○ kashmirkj@Kashmirpc:~/Operating_System/Lab_07$
```

Code:



```
1 //sem_init(&mutex, 0, 1); // Binary semaphore initia
```

Output:

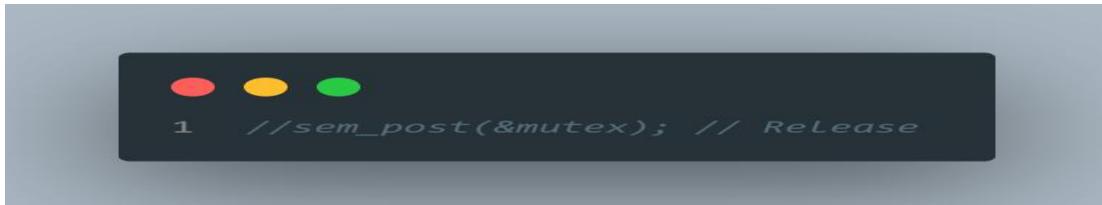
```
Final Counter Value: 10
● kashmirkj@Kashmirpc:~/Operating_System/Lab_07$ gcc Task1.c -o Task1 -lpthread
○ kashmirkj@Kashmirpc:~/Operating_System/Lab_07$ ./Task1
Thread 1: Waiting...
Thread 2: Waiting...

^Z
[1]+  Stopped                  ./Task1
○ kashmirkj@Kashmirpc:~/Operating_System/Lab_07$
```

Description:

- **Both threads will block immediately**
- **DEADLOCK occurs** - Neither thread can ever enter the critical section
- **Program hangs forever**

Code :

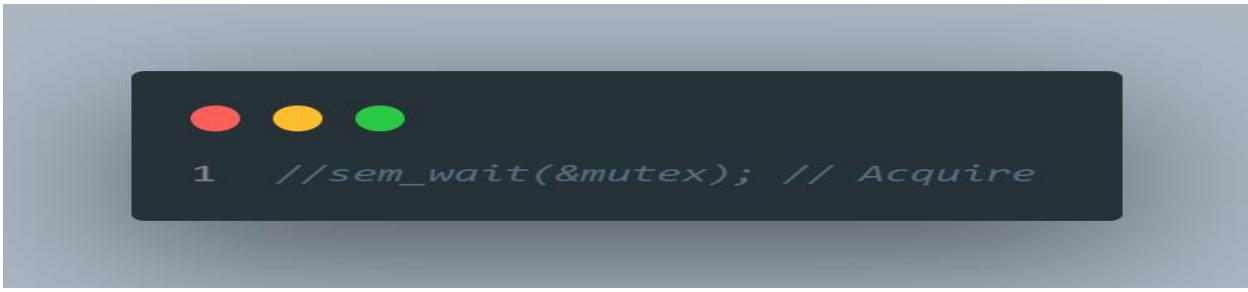


Output:

```
> ▾ TERMINAL
  • kashmirkj@Kashmirpc:~/Operating_System/Lab_07$ gcc Task1.c -o Task1 -lpthread
  ◉ kashmirkj@Kashmirpc:~/Operating_System/Lab_07$ ./Task1
    Thread 1: Waiting...
    Thread 1: In critical section | Counter = 1
    Thread 2: Waiting...
    Thread 1: Waiting...
    Thread 1: Waiting...
    ^Z
  [2]+  Stopped                  ./Task1
  ◉ kashmirkj@Kashmirpc:~/Operating_System/Lab_07$
```

- **First thread enters once, then deadlocks:** Whichever thread gets the semaphore first will enter the critical section, increment counter to 1, but never release the semaphore.
- **Second thread and all subsequent attempts block forever:** The first thread's second iteration and the other thread will all get stuck waiting at sem_wait() since the semaphore is never released.
- **Program hangs**

Code:



Output:

```
● kashmirkj@Kashmirpc:~/Operating_System/Lab_07$ gcc Task1.c -o Task1 -lpthread
● kashmirkj@Kashmirpc:~/Operating_System/Lab_07$ ./Task1
Thread 1: Waiting...
Thread 1: In critical section | Counter = 1
Thread 2: Waiting...
Thread 2: In critical section | Counter = 2
Thread 1: Waiting...
Thread 1: In critical section | Counter = 3
Thread 2: Waiting...
Thread 2: In critical section | Counter = 4
Thread 2: Waiting...
Thread 2: In critical section | Counter = 5
Thread 1: Waiting...
Thread 1: In critical section | Counter = 6
Thread 2: Waiting...
Thread 2: In critical section | Counter = 7
Thread 1: Waiting...
Thread 1: In critical section | Counter = 8
Thread 2: Waiting...
Thread 2: In critical section | Counter = 9
Thread 1: Waiting...
Thread 1: In critical section | Counter = 10
Final Counter Value: 10
○ kashmirkj@Kashmirpc:~/Operating_System/Lab_07$
```

Description:

- **No mutual exclusion:** Both threads can enter the critical section simultaneously, causing potential race conditions on the counter variable.
- **Unpredictable counter increments:** Due to race conditions, the final counter value may be less than 10 (could be anywhere from 2-10, though likely close to 10 in practice).

Task-2

Input:



```
1 #include <stdio.h>
2 #include <pthread.h>
3 #include <semaphore.h>
4 #include <unistd.h>
5 sem_t mutex; // Binary semaphore
6 int counter = 0;
7 void* thread_function(void* arg){
8     int id = *(int*)arg;
9     for (int i = 0; i < 5; i++) {
10         printf("Thread %d: Waiting...\n", id);
11         sem_wait(&mutex); // Acquire
12
13     // Critical section
14     counter++;
15     printf("Thread %d: In critical section | Counter = %d\n", id,
16     counter);
17     sleep(1);
18     sem_post(&mutex); // Release
19     sleep(1);
20 }
21     return NULL;
22 }
23 void* thread_function2(void* arg){
24     int id = *(int*)arg;
25     for (int i = 0; i < 5; i++) {
26         printf("Thread %d: Waiting...\n", id);
27         sem_wait(&mutex); // Acquire
28
29     // Critical section
30     counter--;
31     printf("Thread %d: In critical section | Counter = %d\n", id,
32     counter);
33     sleep(1);
34     sem_post(&mutex); // Release
35     sleep(1);
36 }
37     return NULL;
38 }
39 int main() {
40     sem_init(&mutex, 0, 1); // Binary semaphore initialized to 1
41     pthread_t t1, t2;
42     int id1 = 1, id2 = 2;
43     pthread_create(&t1, NULL, thread_function, &id1);
44     pthread_create(&t2, NULL, thread_function2, &id2);
45     pthread_join(t1, NULL);
46     pthread_join(t2, NULL);
47     printf("Final Counter Value: %d\n", counter);
48     sem_destroy(&mutex);
49     return 0;
50 }
```

Output:

```
TERMINAL
● kashmirkj@Kashmirpc:~/Operating_System/Lab_07$ gcc Task2.c -o Task2 -lpthread
● kashmirkj@Kashmirpc:~/Operating_System/Lab_07$ ./Task2
Thread 1: Waiting...
Thread 1: In critical section | Counter = 1
Thread 2: Waiting...
Thread 2: In critical section | Counter = 0
Thread 1: Waiting...
Thread 1: In critical section | Counter = 1
Thread 2: Waiting...
Thread 2: In critical section | Counter = 0
Thread 1: Waiting...
Thread 1: In critical section | Counter = 1
Thread 2: Waiting...
Thread 2: In critical section | Counter = 0
Thread 1: Waiting...
Thread 1: In critical section | Counter = 1
Thread 2: Waiting...
Thread 2: In critical section | Counter = 0
Thread 1: Waiting...
Thread 1: In critical section | Counter = 1
Thread 2: Waiting...
Thread 2: In critical section | Counter = 0
Thread 1: Waiting...
Thread 1: In critical section | Counter = 1
Thread 2: Waiting...
Thread 2: In critical section | Counter = 0
Final Counter Value: 0
```

Code:



```
TERMINAL
● kashmirkj@Kashmirpc:~/Operating_System/Lab_07$ gcc Task2.c -o Task2 -lpthread
● kashmirkj@Kashmirpc:~/Operating_System/Lab_07$ ./Task2
Thread 1: Waiting...
Thread 2: Waiting...
Thread 1: In critical section | Counter = 1
Thread 2: Waiting...
Thread 2: In critical section | Counter = 0
Thread 1: Waiting...
Thread 1: In critical section | Counter = 1
Thread 2: Waiting...
Thread 2: In critical section | Counter = 0
Thread 1: Waiting...
Thread 1: In critical section | Counter = 1
Thread 2: Waiting...
Thread 2: In critical section | Counter = 0
Thread 1: Waiting...
Thread 1: In critical section | Counter = 1
Thread 2: Waiting...
Thread 2: In critical section | Counter = 0
Thread 1: Waiting...
Thread 1: In critical section | Counter = 1
Thread 2: Waiting...
Thread 2: In critical section | Counter = 0
Thread 1: Waiting...
Thread 1: In critical section | Counter = 1
Thread 2: Waiting...
Thread 2: In critical section | Counter = 0
Final Counter Value: 0
```

Output:

```
TERMINAL
● kashmirkj@Kashmirpc:~/Operating_System/Lab_07$ gcc Task2.c -o Task2 -lpthread
● kashmirkj@Kashmirpc:~/Operating_System/Lab_07$ ./Task2
Thread 1: Waiting...
Thread 2: Waiting...
^Z
[3]+  Stopped                  ./Task2
● kashmirkj@Kashmirpc:~/Operating_System/Lab_07$
```

2. Commit the `//sem_post(&mutex);` in 2nd function i.e is in Counter –

Code:



Output

A screenshot of a terminal window titled 'TERMINAL'. It shows the execution of a C program named Task2.c. The output indicates two threads: Thread 1 and Thread 2. Thread 1 enters a critical section and increments the counter to 1. Thread 2 then enters its critical section and sees the counter as 0. Both threads then wait again. The user then stops the process with ^Z and kills the thread with ^C. The terminal window has tabs for 'bash - Lab_07' and 'bash Lab_07'.

Commit the `//sem_post(&mutex);` in First function i.e is in Counter ++

Code:



Output:

The terminal window shows the following session:

```
kashmirkj@Kashmirpc:~/Operating_System/Lab_07$ gcc Task2.c -o Task2 -lpthread
kashmirkj@Kashmirpc:~/Operating_System/Lab_07$ ./Task2
Thread 2: Waiting...
Thread 2: In critical section | Counter = -1
Thread 1: Waiting...
Thread 1: In critical section | Counter = 0
Thread 2: Waiting...
Thread 1: Waiting...
^Z
[5]+ Stopped ./.Task2
kashmirkj@Kashmirpc:~/Operating_System/Lab_07$
```

Description:

- Thread 1 acquires semaphore on first iteration but never releases it
- Thread 2 gets stuck waiting after Thread 1's first entry, deadlock occurs
- Final counter = 1 (only Thread 1's first increment completes), program hangs

Commit the `//sem_wait(&mutex);` Second function i.e is in Counter--

Code:



Output:

```
> ▾ TERMINAL
  ● kashmirkj@Kashmirpc:~/Operating_System/Lab_07$ gcc Task2.c -o Task2 -lpthread
  ● kashmirkj@Kashmirpc:~/Operating_System/Lab_07$ ./Task2
    Thread 1: Waiting...
    Thread 1: In critical section | Counter = 1
    Thread 2: Waiting...
    Thread 2: In critical section | Counter = 0
    Thread 1: Waiting...
    Thread 1: In critical section | Counter = 1
    Thread 2: Waiting...
    Thread 2: In critical section | Counter = 0
    Thread 1: Waiting...
    Thread 1: In critical section | Counter = 1
    Thread 2: Waiting...
    Thread 2: In critical section | Counter = 0
    Thread 1: Waiting...
    Thread 1: In critical section | Counter = 1
    Thread 2: Waiting...
    Thread 2: In critical section | Counter = 0
    Thread 1: Waiting...
    Thread 1: In critical section | Counter = 1
    Thread 2: Waiting...
    Thread 2: In critical section | Counter = 0
    Final Counter Value: 0
  ○ kashmirkj@Kashmirpc:~/Operating_System/Lab_07$
```

Commit the `//sem_wait(&mutex);` First function i.e is in Counter ++

Code:



Output:

```

● kashmirkj@Kashmirpc:~/Operating_System/Lab_07$ gcc Task2.c -o Task2 -lpthread
● kashmirkj@Kashmirpc:~/Operating_System/Lab_07$ ./Task2
Thread 2: Waiting...
Thread 2: In critical section | Counter = -1
Thread 1: Waiting...
Thread 1: In critical section | Counter = 0
Thread 2: Waiting...
Thread 2: In critical section | Counter = -1
Thread 1: Waiting...
Thread 1: In critical section | Counter = 0
Thread 2: Waiting...
Thread 2: In critical section | Counter = 1
Thread 2: Waiting...
Thread 2: In critical section | Counter = 0
Thread 2: Waiting...
Thread 2: In critical section | Counter = -1
Thread 1: Waiting...
Thread 1: In critical section | Counter = 0
Thread 2: Waiting...
Thread 2: In critical section | Counter = -1
Thread 1: Waiting...
Thread 1: In critical section | Counter = 0
Final Counter Value: 0
● kashmirkj@Kashmirpc:~/Operating_System/Lab_07$ 

```

Description:

- Thread 1 enters critical section without acquiring semaphore, while Thread 2 properly waits
- Race condition: Thread 1 can increment while Thread 2 is decrementing (both in critical section simultaneously)
- Final counter unpredictable (likely close to 0 but not guaranteed), Thread 1's increments may be lost due to concurrent access

Task _3 Comparison between Binary Semaphore Mutex?

#	Difference	Binary Semaphore	Mutex	Why it matters in practice
1	Who can release it?	Any thread or process can call sem_post()	Only the thread that locked it can unlock	If the wrong thread unlocks a mutex → crash or deadlock. This is the biggest difference.
2	Ownership	No owner at all	Has a strict owner (thread ID stored)	Mutex prevents accidental unlocks by other threads → safer for protecting data.

#	Difference	Binary Semaphore	Mutex	Why it matters in practice
3	Priority Inversion Protection	Usually none (high-priority threads can be blocked)	Automatic priority inheritance (Linux futex does it)	With semaphores, high-priority threads can starve; mutexes reduce this problem.
4	Intended Use	Signaling / notifications (e.g., “work is done”)	Mutual exclusion (protect counters, lists, shared memory)	Using a semaphore for shared data protection can cause bugs.
5	Behavior if owner thread dies	Stays locked forever → program hangs	Can be robust →	