

CvP - Werkcollege 2

Exercise 1 Describe the meaning of

- (a) operational semantics,
- (b) denotational semantics,
- (c) axiomatic semantics.

Exercise 2 Consider the following EBNF grammar G

$$\begin{aligned} \langle stmt \rangle &\rightarrow \langle var \rangle = \langle expr \rangle \mid \langle stmt \rangle ; \langle stmt \rangle \mid \mathbf{do} \langle stmt \rangle \mathbf{until} \langle guard \rangle \mathbf{od} \\ \langle guard \rangle &\rightarrow \langle expr \rangle == \langle expr \rangle \\ \langle expr \rangle &\rightarrow \langle octal \rangle \mid \langle var \rangle \mid \langle expr \rangle (+ \mid *) \langle expr \rangle \\ \langle var \rangle &\rightarrow a \langle octal \rangle \\ \langle octal \rangle &\rightarrow 0 \mid (1 \mid 2 \mid 3 \mid 4 \mid 5 \mid 6 \mid 7) \{ (0 \mid 1 \mid 2 \mid 3 \mid 4 \mid 5 \mid 6 \mid 7) \} \end{aligned}$$

For every non-terminal x , let $T_G(x)$ denote the set of all parse trees generated by the grammar G with start symbol x . Let $\mathbb{N} = \{0, 1, \dots\}$ denote the set of all natural numbers. Let $\mathbb{B} = \{\mathbf{false}, \mathbf{true}\}$ denote the set of boolean values.

Represent the state as a function $s : L \rightarrow \mathbb{N}$, with $L \subseteq \mathbb{N}$ finite. Intuitively, a state s assigns the value $s(i)$ to the i -th memory cell, if $i \in \text{dom}(s)$, and leaves the value of the i -th memory cell undefined, otherwise. Let

$$S = \{s : L \rightarrow \mathbb{N} \mid L \subseteq \mathbb{N} \text{ finite}\} \cup \{\text{error}\}$$

be the set of all possible states (including some error state).

- (a) Define semantics $M_{\langle octal \rangle} : T_G(\langle octal \rangle) \rightarrow \mathbb{N}$ for octal numbers.
- (b) Define semantics $M_{\langle var \rangle} : T_G(\langle var \rangle) \rightarrow \mathbb{N}$ for variable names.
- (c) Define semantics $M_{\langle expr \rangle} : T_G(\langle expr \rangle) \rightarrow (S \rightarrow \mathbb{N})$ for expressions, and explain why $M_{\langle expr \rangle}$ is well-defined (although G is ambiguous).
- (d) Define semantics $M_{\langle guard \rangle} : T_G(\langle guard \rangle) \rightarrow (S \rightarrow \mathbb{B})$ for guards.
- (e) Define semantics $M_{\langle stmt \rangle} : T_G(\langle stmt \rangle) \rightarrow (S \rightarrow S)$ for statements, and explain why your definition is well-defined.
(Hint: use *substitution* $s[i \mapsto n] \in S$ and *composition* $f \circ g : S \rightarrow S$).

Exercise 3 Consider the statement S defined as $\{P\} x := 6y + 10; \{x > 16\}$.

- (a) Compute the weakest precondition P in S .
- (b) Give a different precondition of statement S .
- (c) Give a different postcondition of statement S .

Exercise 4 Prove or give a counter example:

- (a) $\{x < 5\} x := x + 4; \{x \leq 3\}$.
- (b) $\{y \leq 1\} x := (2y + 2)^2; \{x \leq 16\}$.