

CvP - Werkcollege 1

Exercise 1 Consider the following grammar G :

$$\begin{aligned}\langle assign \rangle &\rightarrow \langle id \rangle = \langle expr \rangle \\ \langle expr \rangle &\rightarrow \langle expr \rangle + \langle expr \rangle \mid \langle expr \rangle * \langle expr \rangle \mid (\langle expr \rangle) \mid \langle id \rangle \\ \langle id \rangle &\rightarrow A \mid B \mid C\end{aligned}$$

- (a) List all lexemes, tokens, terminals, and non-terminals used in G .
- (b) Give two distinct right-most derivations of “ $A = A + B * C$ ”, and conclude the ambiguity of this grammar.
- (c) Count the number of sentential forms in each of the derivations in (b).
- (d) Rewrite G to make it unambiguous. Argue (no full proof) why your grammar is unambiguous.

Exercise 2 Consider the following grammar G :

$$\begin{aligned}\langle assign \rangle &\rightarrow \langle id \rangle = \langle expr \rangle \\ \langle expr \rangle &\rightarrow \langle expr \rangle + \langle expr \rangle \\ &\quad \mid \langle id \rangle \\ \langle id \rangle &\rightarrow A \mid B \mid C \mid D\end{aligned}$$

- (a) Draw a parse tree of “ $A = B + C + D$ ”.
- (b) Prove grammar G is ambiguous.
- (c) Suppose that $+$ associative. Does this change the ambiguity of the expression “ $A = B + C + D$ ” in G ? Explain your answer.
- (d) Rewrite G into an unambiguous grammar where $+$ right-associative.
- (e) Starting from G , write an attribute grammar, where
 - types cannot be mixed in expressions, and
 - types of both sides of an assignment match.

Exercise 3 Translate the following rules from EBNF to BNF:

- (a) $\langle A \rangle \rightarrow \langle B \rangle [\langle C \rangle]$
- (b) $\langle id_list \rangle \rightarrow \langle id \rangle \{, \langle id \rangle\}$
- (c) $\langle T \rangle \rightarrow \langle T \rangle (+ \mid - \mid 9) \langle F \rangle$

Exercise 4 Consider the attribute grammar in Figure 1.

EXAMPLE 3.6	An Attribute Grammar for Simple Assignment Statements
	1. Syntax rule: $\langle assign \rangle \rightarrow \langle var \rangle = \langle expr \rangle$ Semantic rule: $\langle expr \rangle.expected_type \leftarrow \langle var \rangle.actual_type$
	2. Syntax rule: $\langle expr \rangle \rightarrow \langle var \rangle[2] + \langle var \rangle[3]$ Semantic rule: $\langle expr \rangle.actual_type \leftarrow$ <div style="margin-left: 100px;"> if ($\langle var \rangle[2].actual_type = int$) and ($\langle var \rangle[3].actual_type = int$) then int else real end if </div>
	Predicate: $\langle expr \rangle.actual_type == \langle expr \rangle.expected_type$
	3. Syntax rule: $\langle expr \rangle \rightarrow \langle var \rangle$ Semantic rule: $\langle expr \rangle.actual_type \leftarrow \langle var \rangle.actual_type$ Predicate: $\langle expr \rangle.actual_type == \langle expr \rangle.expected_type$
	4. Syntax rule: $\langle var \rangle \rightarrow A \mid B \mid C$ Semantic rule: $\langle var \rangle.actual_type \leftarrow \text{look-up}(\langle var \rangle.string)$
	The look-up function looks up a given variable name in the symbol table and returns the variable's type.

Figure 1: Attribute grammar.

- (a) In this grammar, point out a synthesized attribute, an inherited attribute, and an intrinsic attribute.
- (b) Draw the parse tree of the expression " $A = A + B$ " and decorate the parse tree (add the flow of attributes to the tree).
- (c) Consider the expression " $A = A + B$ ". What would happen if the actual type of A is int and the actual type of B is real?