# Games: Adversarial Search minimax and expectimax

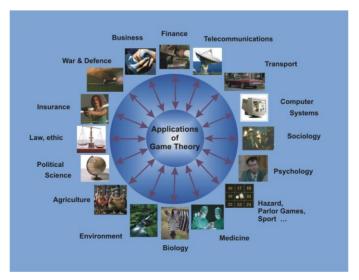
- Game Theory
- Optimal Decisions in Games
  - minimax decisions
  - $-\alpha$ - $\beta$  pruning
  - Expectimax
- Resource limits and approximate evaluation
- Games of imperfect information



Computer Science Department

1

# Game Theory in Artificial Intelligence by Pier Paolo



2

## Game Theory: Focus Here

- Two players
  - Max-min
  - Taking turns, fully observable
- Moves: ActionPosition: stateZero sum:
  - good for one player, bad for another
  - No win-win outcome.



Computer Science Department

3

## Types of Games



#### General Games

- Agents have independent utilities (values on outcomes)
- Cooperation, indifference, competition, and more are all possible
  - » We don't make AI to act in isolation, it should a) work around people and b) help people
  - » That means that every AI agent needs to solve a game



#### Zero-Sum Games

- Agents have opposite utilities (values on outcomes)
- Can reduce payoff to a single value that one player maximizes and the other minimizes
- Adversarial, pure competition



Computer Science Department

## Types of games

	deterministic	chance
perfect information	chess, checkers, go, <u>othello</u>	backgammon monopoly
imperfect information	battleships,	bridge, poker, scrabble

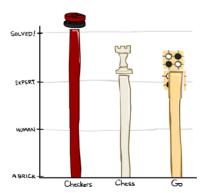


Computer Science Department

5

#### Zero-Sum Game Games ©

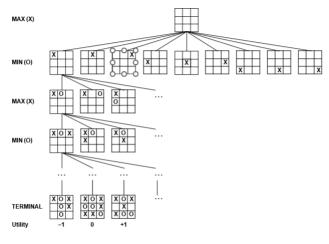
- Checkers: 1950: First computer player. 1994: First computer champion: Chinook ended 40-year-reign of human champion Marion Tinsley using complete 8-piece endgame. 2007: Checkers solved!
- Chess: 1997: Deep Blue defeats human champion Gary Kasparov in a six-game match. Deep Blue examined 200M positions per second, used very sophisticated evaluation and undisclosed methods for extending some lines of search up to 40 ply. Current programs are even better, if less historic
- Go :2016: Alpha GO defeats human champion. Uses Monte Carlo Tree Search, learned evaluation function.





Computer Science Department

## Game tree (2-player, deterministic, turns)





Computer Science Department

7

7

#### **Deterministic Games with Terminal Utilities**

- Many possible formalizations, one is:
  - States: S (start at s<sub>0</sub>)
  - Players: P={1...N} (usually take turns)
  - Actions: A (may depend on player / state)
  - Transition Function:  $SxA \rightarrow S$
  - Terminal Test: S →  $\{t,f\}$
  - Terminal Utilities: SxP → R

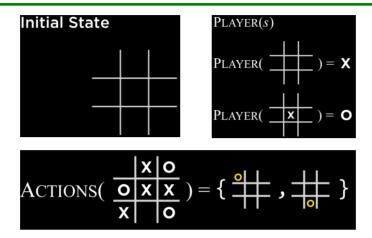


Solution for a player is a policy: S → A



Computer Science Department

#### Tic Tac Toe



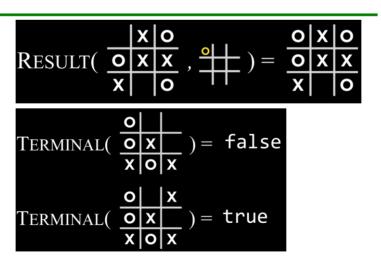
CAL POLY

Computer Science Department

9

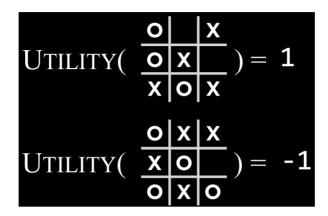
9

## Tic Tac Toe





Computer Science Department



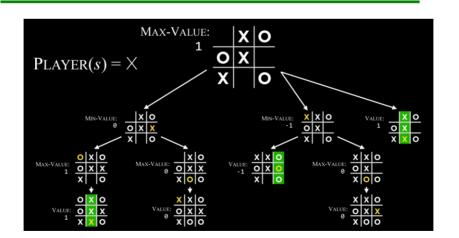


Computer Science Department

11

11

## Players: Max, Min





Computer Science Department

12

#### New: Cost -> Utility!

- no longer minimizing cost!
- agent now wants to maximize its score/utility!



Computer Science Department

13

13

## Maximizing utility not minimizing cost

- Utilities are functions from outcomes (states of the world) to real numbers that describe an agent's preferences
- Where do utilities come from?
  - In a game, may be simple (+1/-1)
  - Utilities summarize the agent's goals
  - Theorem: any "rational" preferences can be summarized as a utility function
- We hard-wire utilities and let behaviors emerge
  - Why don't we let agents pick utilities?
  - Why don't we prescribe behaviors?



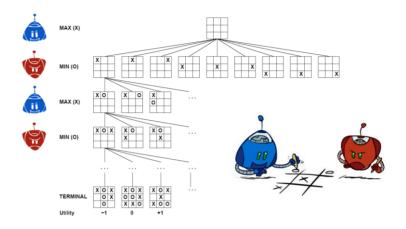






Computer Science Department

#### Tic-Tac-Toe Game Tree





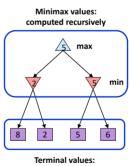
Computer Science Department

15

15

## Adversarial Search (Minimax)

- Deterministic, zero-sum games:
  - Tic-tac-toe, chess, checkers
  - One player maximizes result
  - The other minimizes result
- Minimax search:
  - A state-space search tree
  - Players alternate turns
  - Compute each node's minimax value: the best achievable utility against a rational (optimal) adversary



Terminal values: part of the game



Computer Science Department

16

## Minimax Implementation

CAL POLY

Computer Science Department

17

17

## Minimax Implementation (Dispatch)

```
def value(state):
    if the state is a terminal state: return the state's utility
    if the next agent is MAX: return max-value(state)
    if the next agent is MIN: return min-value(state)

def max-value(state):
    initialize v = -∞
    for each successor of state:
    v = max(v, value(successor))
    return v

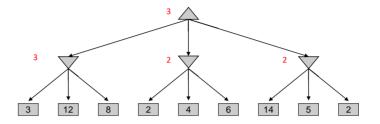
def min-value(state):
    initialize v = +∞
    for each successor of state:
    v = min(v, value(successor))
    return v
```

CAL POLY

Computer Science Department

18

## Minimax Example: Trace Code



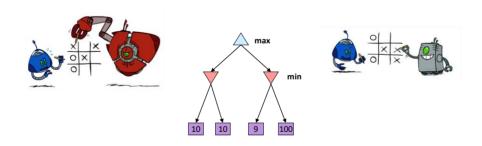


Computer Science Department

19

19

## **Minimax Properties**



Optimal against a perfect player. Otherwise?

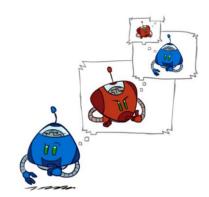


Computer Science Department

20

## Minimax Efficiency

- How efficient is minimax?
  - Just like (exhaustive) DFS
  - Time: O(b<sup>m</sup>)
  - Space: O(bm)
- Example: For chess, b ≈ 35, m ≈ 100
  - Exact solution is completely infeasible
  - But, do we need to explore the whole tree?



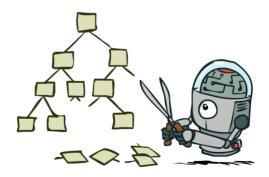


Computer Science Department

21

21

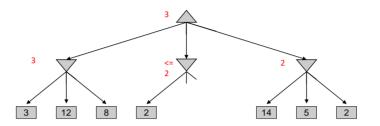
## **Game Tree Pruning**





Computer Science Department

## Minimax Example





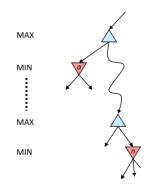
Computer Science Department

23

23

## Alpha-Beta Pruning

- General configuration (MIN version)
  - We're computing the MIN-VALUE at some node n
  - We're looping over n's children
  - n's estimate of the childrens' min is dropping
  - Who cares about n's value? MAX
  - Let a be the best value that MAX can get at any choice point along the current path from the root
  - If n becomes worse than a, MAX will avoid it, so we can stop considering n's other children (it's already bad enough that it won't be played)
- MAX version is symmetric



CAL POLY

Computer Science Department

24

## Alpha-Beta Implementation

α: MAX's best option on path to rootβ: MIN's best option on path to root

```
\label{eq:defmax-value} \begin{split} &\text{def max-value(state, } \alpha, \beta)\text{:} \\ &\text{initialize } v = -\infty \\ &\text{for each successor of state:} \\ &v = \text{max}(v, \text{value(successor, } \alpha, \beta)) \\ &\text{if } v \geq \beta \text{ return } v \\ &\alpha = \text{max}(\alpha, v) \\ &\text{return } v \end{split}
```

```
def min-value(state , \alpha, \beta):
    initialize v = +\infty
    for each successor of state:
        v = min(v, value(successor, \alpha, \beta))
        if v \le \alpha return v
        \beta = min(\beta, v)
    return v
```



Computer Science Department

25

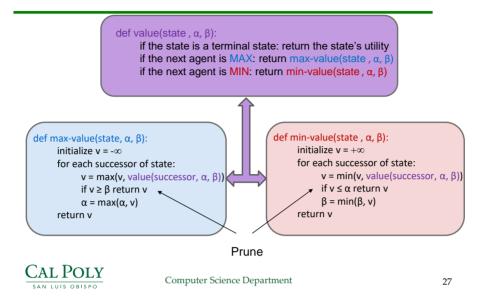
25

## Minimax Implementation (Dispatch)



Computer Science Department

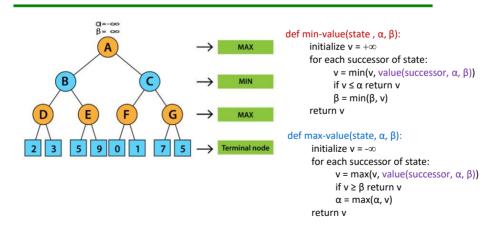
26



27

#### α-β Pruning: Detailed Example -1

https://people.cs.pitt.edu/~litman/courses/cs2710/lectures/pruningReview.pdf

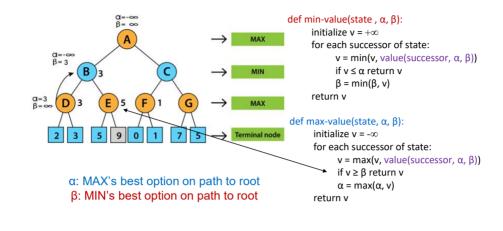




Computer Science Department

28

## $\alpha$ - $\beta$ Pruning: Detailed Example - 2



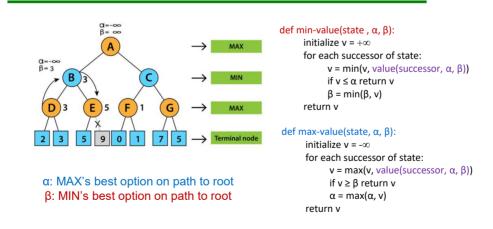
CAL POLY

Computer Science Department

29

29

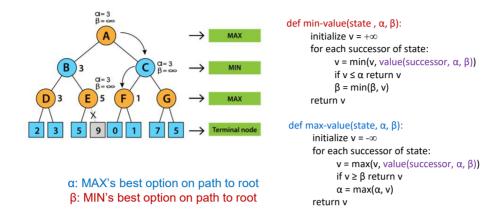
## $\alpha$ - $\beta$ Pruning: Detailed Example - 3



CAL POLY

Computer Science Department

## $\alpha$ - $\beta$ Pruning: Detailed Example - 4



CAL POLY

Computer Science Department

31

31

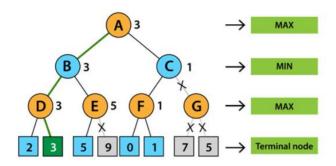
## $\alpha$ - $\beta$ Pruning: Detailed Example - 5

```
def min-value(state, \alpha, \beta):
           A
                                                                 initialize v = +\infty
                                                                 for each successor of state:
                                                                        v = min(v, value(successor, \alpha, \beta))
B
                                                                       if v ≤ α return v
                                                                        \beta = \min(\beta, v)
              F
                                                                 return v
                                                            def max-value(state, \alpha, \beta):
                                                                  initialize v = -\infty
                                                                  for each successor of state:
                                                                        v = max(v, value(successor, \alpha, \beta))
                                                                         if v \ge \beta return v
α: MAX's best option on path to root
                                                                         \alpha = \max(\alpha, v)
β: MIN's best option on path to root
                                                                  return v
```

CAL POLY

Computer Science Department

## $\alpha$ - $\beta$ Pruning: Detailed Example - 6





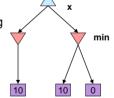
Computer Science Department

33

33

## Alpha-Beta Pruning Properties

- This pruning has no effect on minimax value computed for the root!
- Values of intermediate nodes might be wrong
  - Important: children of the root may have the wrong value
  - So the most naïve version won't let you do action selection
- Good child ordering improves effectiveness of pruning
- With "perfect ordering":
  - Time complexity drops to O(b<sup>m/2</sup>)
  - Doubles solvable depth!
  - Full search of, e.g. chess, is still hopeless...



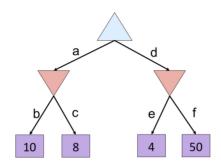
ma

This is a simple example of metareasoning (computing about what to compute)



Computer Science Department

## Alpha-Beta Quiz



 $\alpha$ : MAX's best option on path to root  $\beta$ : MIN's best option on path to root

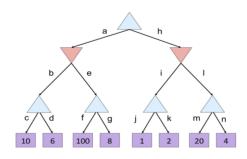


Computer Science Department

35

35

## Alpha-Beta Quiz 2



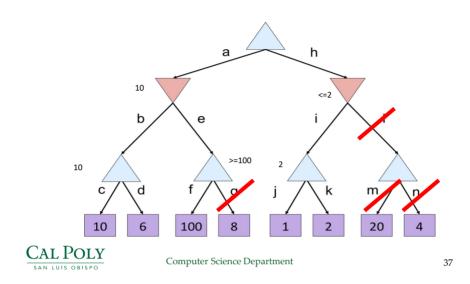
 $\alpha : MAX's$  best option on path to root  $\beta : MIN's$  best option on path to root



Computer Science Department

36

## Alpha-Beta Quiz 2



37

## **Resource Limits**

- Problem: In realistic games, cannot search to leaves!
- Solution: Depth-limited search
  - Instead, search only to a limited depth in the tree
  - Replace terminal utilities with an evaluation function for non-terminal positions
- Example:
  - Suppose we have 100 seconds, can explore 10K nodes / sec
  - So can check 1M nodes per move
  - $\alpha$ - $\beta$  reaches about depth 8 decent chess program
- Guarantee of optimal play is gone
- More plies makes a BIG difference
- Use iterative deepening for an anytime algorithm



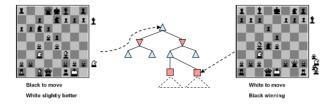
Computer Science Department

38

max min

#### **Evaluation Functions**

Evaluation functions score non-terminals in depth-limited search



- Ideal function: returns the actual minimax value of the position
- In practice: typically weighted linear sum of features:

$$Eval(s) = w_1 f_1(s) + w_2 f_2(s) + ... + w_n f_n(s)$$

• e.g.  $f_1(s) = \text{(num white queens - num black queens), etc.}$ 



Computer Science Department

39

39

## **Depth Matters**

- Evaluation functions are always imperfect
- The deeper in the tree the evaluation function is buried, the less the quality of the evaluation function matters
- An important example of the tradeoff between complexity of features and complexity of computation





[Demo: depth limited (L6D4, L6D5)]

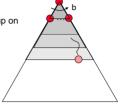


Computer Science Department

#### **Iterative Deepening**

Iterative deepening uses DFS as a subroutine:

- Do a DFS which only searches for paths of length 1 or less. (DFS gives up on any path of length 2)
- 2. If "1" failed, do a DFS which only searches paths of length 2 or less.
- 3. If "2" failed, do a DFS which only searches paths of length 3 or less.



Why do we want to do this for multiplayer games?

Note: wrongness of eval functions matters less and less the deeper the search goes!



Computer Science Department

41

41

## Summary

- Game playing is best modeled as a search problem
- Game trees represent alternate computer/opponent moves
- Minimax is a procedure which chooses moves by assuming that the opponent will always choose the move which is best for them
  - Basically, it avoids all worst-case outcomes for Max, to find the best.
  - If the opponent makes an error, Minimax will take optimal advantage of that error and make the best possible play that exploits the error.
- Cutting off search
  - In general, it is infeasible to search the entire game tree.
  - In practice, Cutoff-Test decides when to stop searching further.
- Static heuristic evaluation function
  - Estimate quality of a given board configuration for the Max player.
  - Called when search is cut off, to determine value of position found.



Computer Science Department

## Other Game Types





Computer Science Department

43

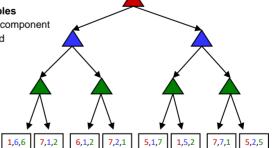
43

## Multi-Agent

- What if the game is not zero-sum, or has multiple players?
- Generalization of minimax:
  - Node values are also utility tuples Each player maximizes its own component Can give rise to cooperation and

competition dynamically...

Terminals have utility tuples





Computer Science Department

44

#### **Uncertain Outcomes**





Computer Science Department

45

45

#### What are Probabilities?

- Objectivist / frequentist answer:
  - Averages over repeated experiments
  - E.g. empirically estimating P(rain) from historical observation
  - Assertion about how future experiments will go (in the limit)
  - Makes one think of inherently random events, like rolling dice
- Subjectivist / Bayesian answer:
  - Degrees of belief about unobserved variables
  - E.g. an agent's belief that it's raining, given the temperature
  - E.g. player<sub>1</sub>'s belief that player<sub>2</sub> will take an action, given the state
  - Often learn probabilities from past experiences (more later)
  - New evidence updates beliefs (more later)



Computer Science Department

#### Probability: Basic Concepts

- Uncertainty can be represented as set of events and the likelihood, or probability, of each of them happening.
  - Event can be thought of as a possible state of a world or aspect of the world
- Axioms: Let S be the <u>sample space (all possible outcomes)</u>
  - $0 \le P(x) \le 1$ : every probability value must range between 0 and 1
  - $-\sum_{x\in S} P(x) = 1$  the sum of the probabilities over the sample space = 1
  - Random Variable: F: S → Possible Values in Real Numbers
  - Expected Value:  $E(\mathbf{F}) = \sum_{values\ that\ F\ takes\ on} F(x) * P(F(x))$



Computer Science Department

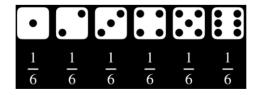
47

47

## Example: Rolling a single die

- Fair Die = Probability of each outcome is 1/6
- Can think of the value as a very simple random variable
- Expected value of rolling a single fair die

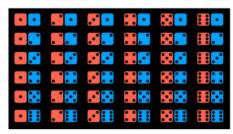
$$= 1 \bullet \frac{1}{6} + 2 \bullet \frac{1}{6} + 3 \bullet \frac{1}{6} + 4 \bullet \frac{1}{6} + 5 \bullet \frac{1}{6} + 6 \bullet \frac{1}{6} = \frac{21}{6} = 3\frac{1}{2}$$



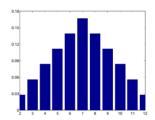


Computer Science Department

#### Example: Rolling a pair of dice: red and blue



- Probability of each event, e.g.  $(2_R, 3_B) = 1/36$
- Random variable: Sum of values of the pair
- Probability distribution of Sum





Computer Science Department

## Reminder: Expectations

The expected value of a function of a random variable the average, weighted by the probability distribution over outcomes



min

Example: How long to get to the airport?





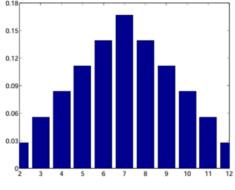




Computer Science Department

## Quiz: What is the expected value of the sum from rolling two dice

- Make a rough guess ( to check logic as go along )
  - remember E(X) is a weighted (by probability) average
- ???? Write the expression
- Do the calculation 0.18





Computer Science Department

51

51

## Summary: Probabilities

- A random variable represents an event whose outcome is unknown
- A probability distribution is an assignment of weights to outcomes
- Example: Traffic on freeway
  - Random variable: T = whether there's traffic
  - Outcomes: T in {none, light, heavy}
  - Distribution: P(T=none) = 0.25, P(T=light) = 0.50, P(T=heavy) = 0.25
- Some laws of probability (more later):
  - Probabilities are always non-negative
  - Probabilities over all possible outcomes sum to one
- As we get more evidence, probabilities may change:
  - P(T=heavy) = 0.25, P(T=heavy | Hour=8am) = 0.60
  - We'll talk about methods for reasoning and updating probabilities later



Computer Science Department

## The Dangers of Optimism and Pessimism

## Dangerous Optimism Assuming chance when the world is adversarial



## Dangerous Pessimism Assuming the worst case when it's not likely



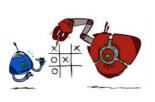


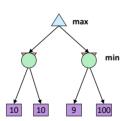
Computer Science Department

53

53

## Worst-Case vs. Average Case







- Idea: Uncertain outcomes controlled by adversary vs chance!
- Can be a combination of the two!!



Computer Science Department

#### **Expectimax Search**

- Why wouldn't we know what the result of an action will be?
  - Explicit randomness: rolling dice
  - Unpredictable opponents: humans are not perfect
  - Actions can fail: when moving a robot, wheels might slip
- Values should now reflect average-case (expectimax) outcomes, not worst-case (minimax) outcomes
- Expectimax search: compute the average score under optimal play
  - Max nodes as in minimax search
  - Chance nodes are like min nodes but the outcome is uncertain
  - Calculate their expected utilities
- Later, we'll learn how to formalize the underlying uncertain-result problems as Markov Decision Processes

chance | 10 | 10 | 9 | 10

[Demo: min vs exp (L7D1,2)]



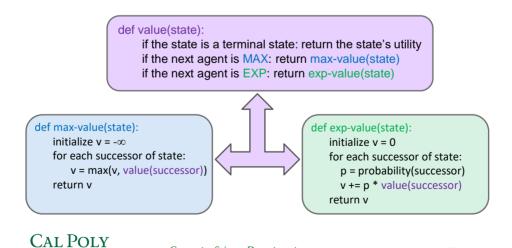
Computer Science Department

55

56

55

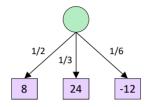
## **Expectimax Pseudocode**



Computer Science Department

## **Expectimax Pseudocode**

# def exp-value(state): initialize v = 0 for each successor of state: p = probability(successor) v += p \* value(successor) return v



v = (1/2)(8) + (1/3)(24) + (1/6)(-12) = 10



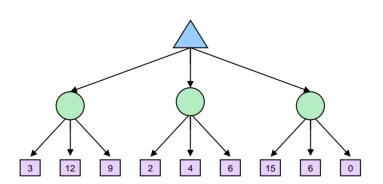
Computer Science Department

57

57

## Expectimax Example:

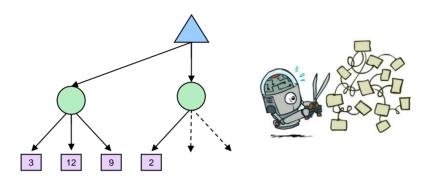
## Again -a postorder traversal





Computer Science Department

## **Expectimax Pruning?**



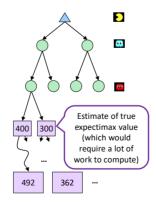


Computer Science Department

59

59

## **Depth-Limited Expectimax**





Computer Science Department

60

#### **Quiz: Informed Probabilities**

- Let's say you know that your opponent is actually running a depth 2 minimax, using the result 80% of the time, and moving randomly otherwise
- Question: What tree search should you use?
- Answer: Expectimax!
  - To figure out EACH chance node's probabilities, you have to run a simulation of your opponent
  - This kind of thing gets very slow very quickly
  - Even worse if you have to simulate your opponent simulating you...

This is basically how you would model a human, except for their utility: their utility might be the same as yours (i.e. you try to help them, but they are depth 2 and noisy), or they might have a slightly different utility (like another person navigating in the office)



Computer Science Department

61