

COMS4032A

Applications of Algorithms

Assignment 5

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1 Part A

1.1 Test Input generation

We ran Graham Scan and Jarvis march on n randomly generated floating point data points using the java Random(). We ran multiple experiments where we varied n from 40 000 to 1 600 000 and increasing n by 40 000. For each n , we ran the 2 algorithms 20 times and recorded the average runtime. The graphs below are what resulted from these tests. Figure 1 shows the runtime for Graham scan and we see runtimes that resembles $O(n \log n)$. Figure 2a shows the runtime for Jarvis march as well as a curve representing $\frac{\text{jarvis_runtime}}{h}$ (h being the number of vertices in the convex hull). We expect the runtime for Jarvis march to be $O(nh)$. Figure 2b shows an almost $\log(n)$ relation between n and h and this is consistent to the graphs we obtained in 2a.

1.2 Graphs

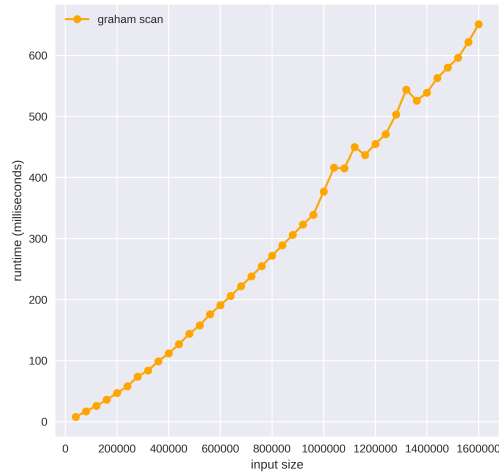


Figure 1: a) Graham Scan Runtime

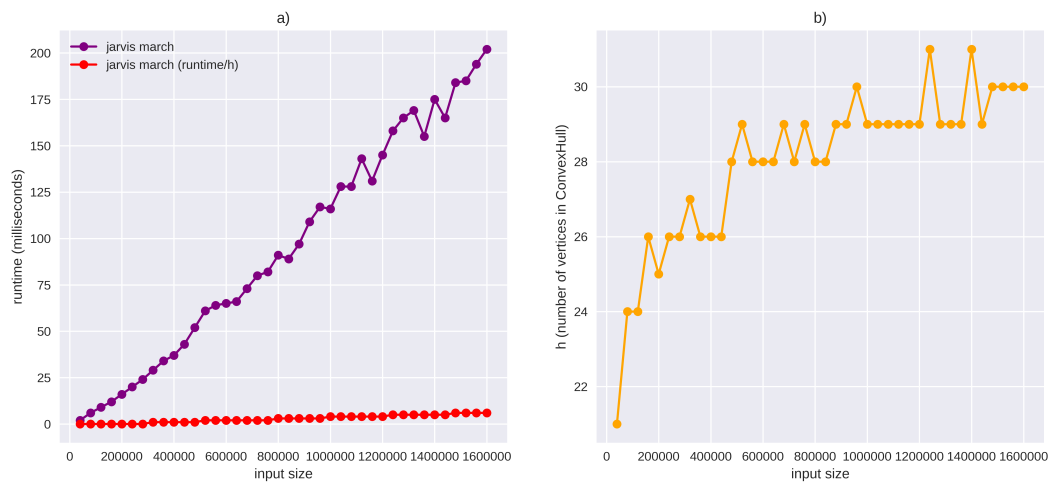


Figure 2: a) Jarvis March Runtime b) n vs h

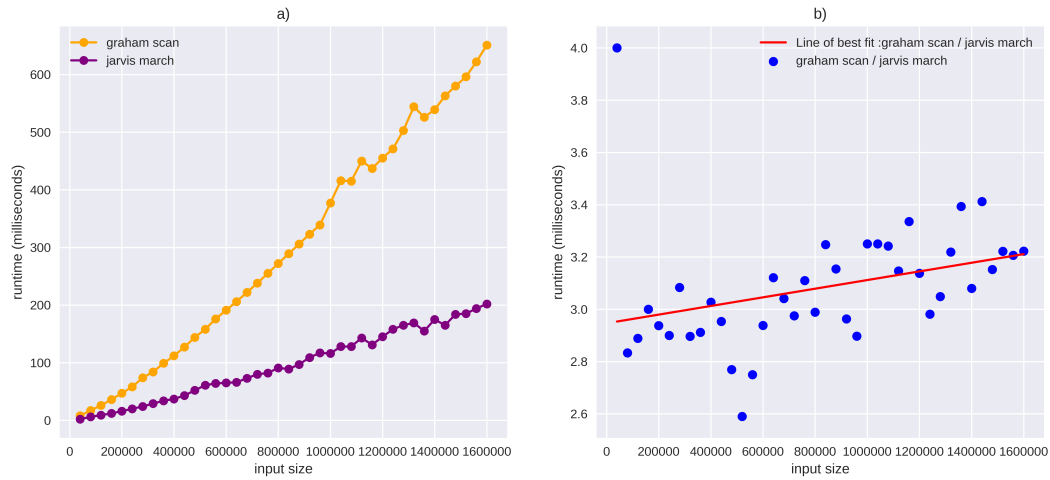


Figure 3: a) Graham Scan vs Jarvis March b) Jarvis March speedup over Graham Scan

2 Part B

Since the points in the polygon are already listed in a counterclockwise order, we can use Graham scan without the sorting part which would result in a complexity of $O(n)$.

Algorithm 1: CH_star_polygon(P, n)

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1:  $S = \text{empty\_stack}()$ 
2:  $S.\text{push}(P[0])$ 
3:  $S.\text{push}(P[1])$ 
4:  $S.\text{push}(P[2])$ 
5: for  $i = 3$  to  $n$  do
6:   while the angle formed by points  $S.\text{NextToTop}()$ ,  $S.\text{Top}()$ , and  $P[i]$  makes a non-left turn do
7:      $S.\text{pop}()$ 
8:   end while
9:    $S.\text{push}(P[i])$ 
10: end for
11: return  $S$ ;

```
