CUDA Assignment 2

Kat Young

High Performance Distributed Systems Virginia Commonwealth University youngk6@vcu.edu

I. ORIGINAL DESIGN APPROACH

Generally speaking, I implemented KNN in CUDA in a similar algorithm to my implementation of KNN in MPI and sequential KNN. This assignment was different from the others, however, in that the functions we were able to use with the Arff Dataset were not able to be used in the kernel since they were host functions. Due to this, I had to convert the Arff Dataset into a 1D array, which would be able to be indexed in the kernel. I then got metrics that would be useful for indexing the various arrays that I would be using in the kernel such as instancecount, attributecount, and classcount which kept track of how many instances, attributes, and classes there were in the data respectively.

I had four arrays that were important in the kernel for computations: predictions (which was where the final predictions would be placed), Kdist (which was to hold the k nearest distances for each instance), Kclasses (which was to hold the corresponding classes for the k nearest distances in Kdist), and classVotes (which was to keep track of the votes for each class after all of the nearest distances were found). My KNN kernel function ended up taking dpredictions (results array), ddataset (dataset as a 1D array including both attributes and classes), k (nearest neighbors), instancecount, attributecount, classcount, dKdist, dKclasses, dclassVotes, inf (the value infinity for edge case use).

In the kernel, I began by getting the thread ID (tid). The first goal for the kernel was to get the k-nearest-neighbors. I accomplished this by filling up the positions in Kdist and Kclasses with the distances and classes of the first k instances in the dataset, keeping track of the largest distance in the array. As the algorithm continued to iterate through all of the other instances, if it found one of the instances had a distance from the instance being measured that was less than the largest distance in that instance's nearest neighbors array, it replaced the old distance and class with the class of the closer neighbor. The largest distance was recalculated after each time this replacement happened.

A problem occurred, for example, when tid was less than the k value. The distance from that instance to itself was zero, and therefore would never get replaced because there would always be distances greater than zero in the nearest neighbors. To handle this problem, when the first k elements of the nearest neighbors array were selected and tid; k, the distance associated with tid was infinity so that as a result that value would be replaced immediately when the algorithm

started replacing elements. When tid i, k, the algorithm was told to not look at elements where tid = j (j being the instance we are comparing the original instance with).

The result of this process was an array that contains the distances and an array that contains the classes of the k nearest neighbors for each instance. The next goal of the kernel was to count the votes and find the class with the most votes to be considered the prediction. Classes were voted on by going through the kclasses array for each instance (size k) and incrementing the index position in the class Values array. For example, if an instance's class votes looked like this: [1,0,4,5,1] for k=5, then the class Votes block allocated for that specific tid (assuming there are 8 classes) would look like this: [1,2,0,0,1,1,0,0]. Since index 1 has the highest value, the most votes went to classify the instance as 1, which would be placed in the predictions array in the spot corresponding to the tid as our prediction.

II. KERNEL CONFIGURATION, EXPERIMENTATION, AND RESULTS

To experiment with kernel configurations, I looked at changing the number of threads per block and changing my original one kernel to two kernels. To change the configuration to two kernels, all I did was move the class prediction to its own kernel. The first kernel put together the large array of nearest neighbors, and the second kernel essentially counted the votes. I decided to try this because kernel launch cost is negligible and it could potentially be faster due to fewer registers used in each one – more resources. Measuring improvements was somewhat hard because all of the times were in the same general range. I took the average of the first five runs to get the resulting numbers.

When looking at the first implementation I did with just one kernel, I tried two configurations of threads and blocks before I moved to trying a different number of kernels. I tried 256 threads per block as well as 1 thread per block. One thread per block performed better on the small dataset than 256 threads per block, while the reverse was true for the medium dataset.

When implementing two kernels, I tested first on 256 threads per block and saw improvements in the times. Then I tried other configurations of threads per block including 1 thread per block, 700 threads per block, 512 threads per block, and 32 threads per block. The results can be seen in the table below, but the best results were from two kernels with 256 threads per block. The differences in times however weren't very large for these datasets. It would potentially be helpful

to run these configurations on a massive dataset and see how each performs.

Run Times (ms)					
Method Sequential		small.arff	medium.arff		
		56 ms	9035 ms		
MPI	2 cores	266 ms	4823 ms		
	4 cores	288 ms	2610 ms		
	8 cores	318 ms	1634 ms		
	16 cores	344 ms	1277 ms		
CUDA	256 threads / block	179 ms	173 ms		
(one kernel)	1 thread / block	149 ms	200 ms		
CUDA (two kernels)	256 threads / block	117.2 ms	142 ms		
	1 thread / block	154 ms	193 ms		
	700 threads / block	153 ms	158 ms		
	512 threads / block	157 ms	156 ms		
	32 threads / block	156 ms	154 ms		

Fig. 1. Approximate run times

Running the code with the profiler, looking at the code written in one kernel, an average of 11.217 ms was spent in the KNN kernel. Running with two kernels, the first kernel cut its time down to 8.499 ms, and the second kernel runs in 5.7 us, which is equivalent to 0.0057 ms, which means the total time spent in the kernel was reduced by this change.

This was the profiler report from the initial configuration I had with 256 threads per block and one kernel.

Fig. 2. 256 Threads / Block and ONE kernel

```
==2560P== Profiling application: ./knn_cuda_xternels_256 datasets/medium.arff 5
==2560P== Profiling result:
Type Tine(s) Tine Calls Avg Mtn Max Name
670 activittes: 98.838 8.995ms 1 8.995ms 8.995ms 8.995ms 8.995ms 8.995ms 8.995ms 8.995ms 8.995ms 8.995ms 9.995ms 8.995ms 9.995ms 9.995ms
```

Fig. 3. 256 Threads / Block and TWO kernels

III. SPEEDUPS

Speedups were calculated using the following equation, where the numerator is the time for the sequential version to run, and the denominator is the time for the parallel version to run.

C 1	T_S	(1)
Speedup =	$\overline{T_P}$	(1)

Speedups						
Method		small.arff	medium.arff			
MPI	2 cores	0.21	1.87			
	4 cores	0.19	3.46			
	8 cores	0.18	5.53			
	16 cores	0.16	7.08			
CUDA (one kernel)	256 threads / block	0.31	52.23			
	1 thread / block	0.38	45.18			
CUDA (two kernels)	256 threads / block	0.48	63.62			
	1 thread / block	0.36	46.81			
	700 threads / block	0.37	57.18			
	512 threads / block	0.36	57.92			
	32 threads / block	0.36	58.67			

Fig. 4. Calculated speedups