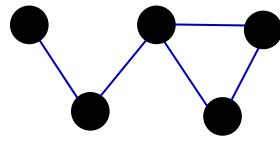
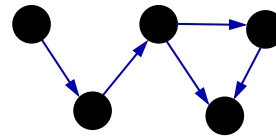


GRAPH THEORY DEFINITIONS



Undirected Graph

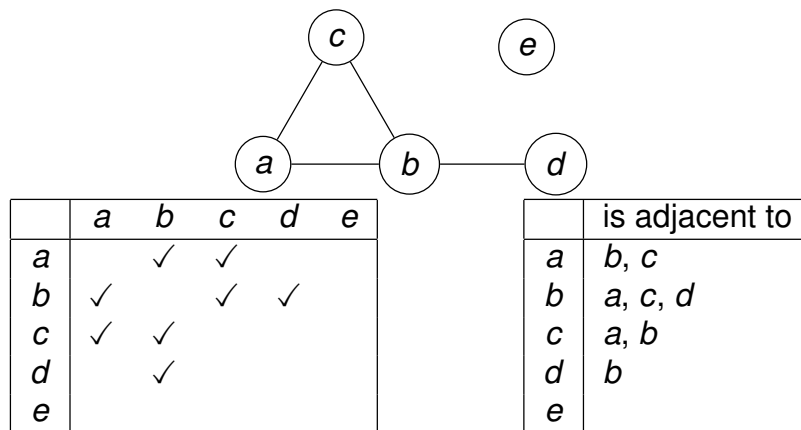


Directed Graph

- ▶ A *graph* $G = (V, E)$ consists of a set of *vertices* (or *nodes*) V and
- ▶ A set of *edges* E .
- ▶ Let $n = |V|$, the number of nodes, and $m = |E|$, the number of edges.
- ▶ In an *undirected* graph, each edge is a set of two vertices $\{u, v\}$ (so (u, v) and (v, u) are the same), and self-loops are *not allowed*. When it's clear from the context, we will use (u, v) for $\{u, v\}$.
- ▶ In a *directed* graph, each edge is an *ordered pair* of nodes (u, v) (so (u, v) is considered *different* from (v, u)); also, self-loops (edges of the form (u, u)) are allowed.

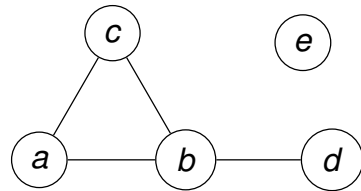
TERMINOLOGY: ADJACENT

Two vertices are *adjacent* iff there is an edge between them.



This brings us to two nice ways to store a graph. . .

STORING A GRAPH: ADJACENCY MATRIX

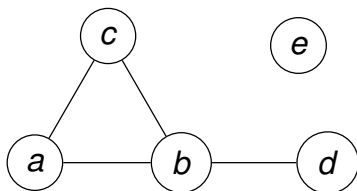


	a	b	c	d	e
a		✓	✓		
b	✓		✓	✓	
c	✓	✓			
d		✓			
e					

Adjacency matrix = store this in a 2D array.

- ▶ space: $\Theta(n^2)$
- ▶ “who are *adjacent* to v ?”: $\Theta(n)$ time
- ▶ “are v and w adjacent?”: $\Theta(1)$ time
- ▶ convenient for some other operations/queries
- ▶ This matrix is for an *undirected graph*. What would change for a *directed graph*?

STORING A GRAPH: ADJACENCY LISTS



	is adjacent to
a	b, c
b	a, c, d
c	a, b
d	b
e	

Adjacency lists = store *vertices* in a 1D *array* or *dictionary*.
Use a list or a set for each entry on the right.

- ▶ At entry $A[i]$, we store the *neighbours* of v_i
- ▶ If the graph is *directed*, we store only the *out-neighbours*.
- ▶ space: $\Theta(n + m)$
- ▶ “who are *adjacent* to v ?”: $\Theta(\deg(v))$ time, ie, length of adjacency list
- ▶ “are v and w adjacent?”: $\Theta(\deg(v))$ time if a *list*. Can we make this *faster*?
- ▶ *optimal* for graph searches
- ▶ How will the adjacency list look different for a *directed graph*?

OPERATIONS ON GRAPHS

Some standard operations on graphs are:

- ▶ **Add/Remove** a vertex/edge.
- ▶ **Edge Query**: given two vertices u, v , find out if the edge (u, v) (if the graph is directed) or the edge $\{u, v\}$ is in E .
- ▶ **Neighborhood**: given a vertex u in an *undirected graph*, get the set of vertices $\{v \mid \{u, v\} \in E\}$.
- ▶ **In-neighborhood, out-neighborhood**: given a vertex u in a directed graph, get the set of vertices $\{v \mid (v, u) \in E\}$ (in) or $\{v \mid (u, v) \in E\}$ (out).
- ▶ **Degree, in-degree, out-degree**: compute the *size* of the *neighborhood*, *in-neighborhood*, or *out-neighborhood*, respectively.