

# Computer OS and Programming Notes

Notes based on CS 2110 Textbook

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## 1.

## 2. Basic Storage Elements

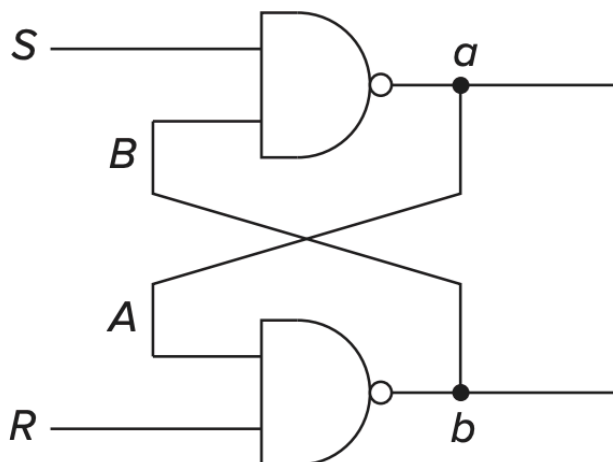
- The other kind of storage element are those that involve the storage of information and those that do not

### 2.1. The RS Latch

#### Definition 2.1.1

The **RS Latch** can store one bit of information, a 1 or a 0. Generally, two 2-input NAND gates are connected such that the output of each is connected to one of the inputs of the other. The other inputs are usually held to be zero.

Setting the latch to store a 1 is known as **setting** the latch, while setting the latch to store a 0 is referred to **resetting** the latch



## Definition 2.1.2

The **quiescent** (or quiet) state of a latch is the state when the latch is storing a value, either 0/1, and nothing attempts to change that value.

This happens when S and R are both equal to 1. So as long as the inputs S and R remain as 1, the state of the circuit will not change.

## Note 2.1.1

**Setting the latch to a 1 or 0**

The latch can be sent to 1 by momentarily setting S to 0, provided that we keep the value of R at 1. Similarly, we can set the patch to 0 by setting R to zero (known as clearing or resetting), provided we keep the value of S at 1.

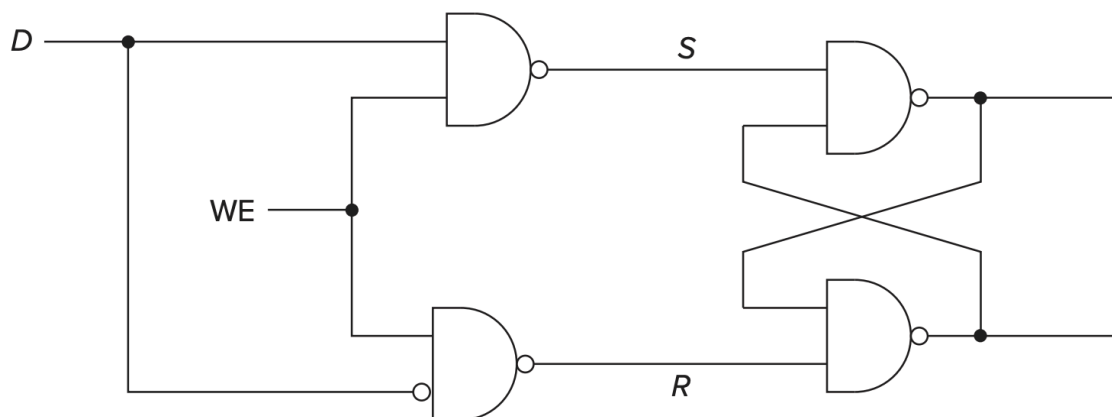
Logic behind setting to 1: If we set S to 0 for a brief period of time, this causes a and thus A to be equal to 1. Since R is 1 and A is 1, b must be 0. This causes B to be 0, which makes a equal to 1 again. Now when we return S to 1, a remains the same since B is also 0, and 1 0 input to a nand gate is enough to make sure that the NAND gate stays at 1.

**2.2. The Gated D Latch**

## Definition 2.2.1

The D latch helps control when a latch is set and when it is cleared. In the following figure, the latch is set to the value of D whenever WE is asserted. When WE is not asserted, the outputs S and R are both equal to 1.

When WE is momentarily set to 1, exactly one of the outputs S or R is set to 0 depending on the value of D. If D is set to 1, the S is set to 0, else



re 3.19 A gated D latch.

**3. The Concept of Memory**

**Memory** is made up of a (usually large) number of locations, each uniquely identifiable and each having the ability to store a value. We refer to the unique memory location as its *address*. We refer to the number of bits of information stored in each location as its *addressability*.

**3.1. Address Space**

## Definition 3.1.1

We refer to the total number of uniquely identifiable locations as the **memory's address space**.

For example, 2GB memory has two billion memory locations.

### 3.2. Addressability

## Definition 3.2.1

**Addressability**: the number of bits stored in each memory location.

### 3.3. $2^2$ by 3-Bit Memory Example

In this case, the memory has an address space of 4 locations and an addressability of three bits. Since it is  $2^2$  memory, it takes two bits to specify the address. We specify it using  $A[1:0]$ . Since its addressability is 3, that means in each location, it stores 3 bits worth of information/data.

## Note 3.3.1

When specifying a memory location in terms of  $A[\mathbf{high:low}]$ , we are starting from the rightmost spot as index of 0. This means we are looking at the sequence of  $h - l + 1$  bits such that high is the leftmost bit number, and low is the rightmost bit number in the sequence.

Access of memory first starts with *decoding the address bits*, using a decoder. We also have WE, which defines whether we are in write-enable mode or not.

The input of  $A[1:0]$  defines what the decoder has to select for the correct *word line*. From there, the decoder outputs a line of 1 which is anded across all three D-latches producing the output of that position.