3 Aries

a) Dirty Page Table:

PageID	recLSN
$P_{\scriptscriptstyle A}$	0

Transaction Table:

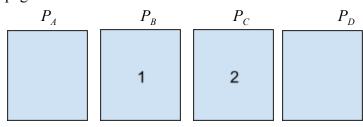
transID	status	lastLSN
T_I	running	0
T_2	running	1
T_3	running	2

Log:

LSN	transID	type	pageID	undoNextLSN	prevLSN
0	T_{I}	update	$P_{\scriptscriptstyle A}$		工
1	T_2	update	$P_{\scriptscriptstyle B}$		Т
2	T_3	update	P_{C}		Т

Where recLSN $\sim lastLSN \sim LSN$

pageLSN on disk:



b) Dirty Page Table:

PageID	recLSN
P_{C}	8
$P_{\scriptscriptstyle B}$	10

P_A	0
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Transaction Table:

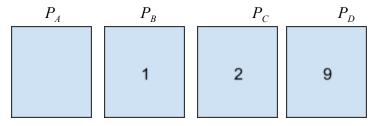
transID	status	lastLSN
T_2	running	10
T_3	running	11

Log:

<u> </u>				1	
LSN	transID	type	pageID	undoNextLSN	prevLSN
0	T_I	update	$P_{\scriptscriptstyle A}$		Т
1	T_2	update	$P_{\scriptscriptstyle B}$		Т
2	T_3	update	P_{C}		Т
3	begin checkpoint				
3	end checkpoint				
4	T_2	update	$P_{\scriptscriptstyle D}$		1
5	T_{I}	update	$P_{\scriptscriptstyle A}$		0
6	T_{I}	commit	1		5
7	T_{I}	end	1		6
8	T_3	update	P_{C}		2
9	T_2	update	P_{D}		4
10	T_2	update	P_{B}		9
11	T_3	update	$P_{\scriptscriptstyle A}$		8

Where recLSN \sim lastLSN \sim LSN

pageLSN on disk of each page:



c) All log records prior and including the final force-write W2(D) are written to stable storage (LSN 0-9).

d) Dirty Page Table:

PageID	recLSN
$P_{\scriptscriptstyle A}$	0
P_{C}	8
P_D	4

Transaction Table:

transID	status	lastLSN
T_2	aborted	9
T_3	aborted	8

Log:

Log.					
LSN	transID	type	pageID	undoNextLSN	prevLSN
0	T_I	update	$P_{\scriptscriptstyle A}$		Т
1	T_2	update	P_{B}		Т
2	T_3	update	P_{C}		Т
3	begin checkpoint				
3	end checkpoint				
4	T_2	update	P_D		1
5	T_{I}	update	$P_{\scriptscriptstyle A}$		0

6	T_I	commit	1	5
7	T_{I}	end	-	6
8	T_3	update	P_{C}	2
9	T_2	update	$P_{\scriptscriptstyle D}$	4

Where recLSN \sim lastLSN \sim LSN

e) LSN $0 \rightarrow \text{Redone}$

LSN 1 \rightarrow Does not need be redone because affected page is not in the dirty page table

LSN 2 \rightarrow Does not need to be redone because for P_{C} , recLSN = 8 > LSN = 2

LSN 4 \rightarrow Does not need to be redone because for P_D , pageLSN = 9 \geq LSN = 4

LSN $5 \rightarrow \text{Redone}$

LSN 8 → Redone

LSN 9 \rightarrow Does not need to be redone because for P_D , pageLSN = 9 \geq LSN = 9

f) System must undo T_2 and T_3 ...

Dirty Page Table:

PageID	recLSN
$P_{\scriptscriptstyle A}$	11
$P_{\scriptscriptstyle B}$	10
P_{C}	8
P_D	9

Transaction Table:

transID	status	lastLSN	

Log:

LSN	transID	type	pageID	undoNextLSN	prevLSN
0	T_I	update	$P_{\scriptscriptstyle A}$		T

1	T_2	update	$P_{\scriptscriptstyle B}$		Т
2	T_3	update	P_C		Т
3	begin checkpoint				
3	end checkpoint				
4	T_2	update	P_D		1
5	T_I	update	$P_{\scriptscriptstyle A}$		0
6	T_I	commit	-		5
7	T_I	end	-		6
8	T_3	update	P_C		2
9	T_2	update	P_D		4
10	T_2	CLR	Undo T ₂ LSN 9	4	9
11	T_3	CLR	Undo T ₃ LSN 8	2	8
12	T_2	CLR	Undo T ₂ LSN 4	1	10
13	T_3	CLR	Undo T ₃ LSN 2	Т	11
14	T_3	end	-		13
15	T_2	CLR	Undo T ₂ LSN 1	Т	12
16	T_2	end	-		15