# Merge Interval

Leetcode Problem 56

#### Problem Description

Given an array of intervals where intervals[i] = [starti, endi], merge all overlapping intervals, and return an array of the non-overlapping intervals that cover all the intervals in the input.

#### Constraints:

- 1 <= intervals.length <= 104</li>
- intervals[i].length == 2
- 0 <= starti <= endi <= 104

#### Example:

#### Input:

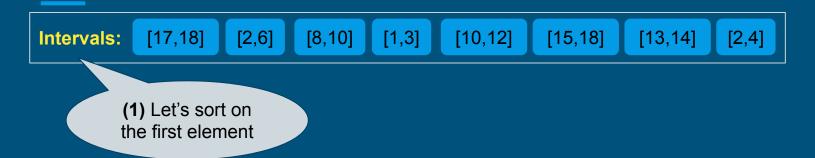
• intervals = [[1,3],[2,6],[8,10],[15,18]]

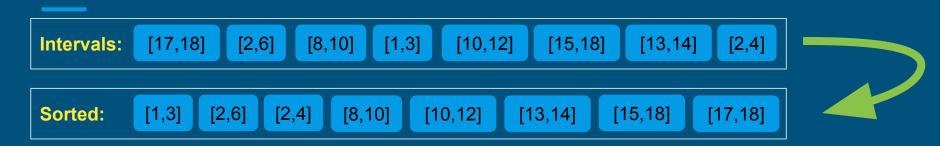
#### Output:

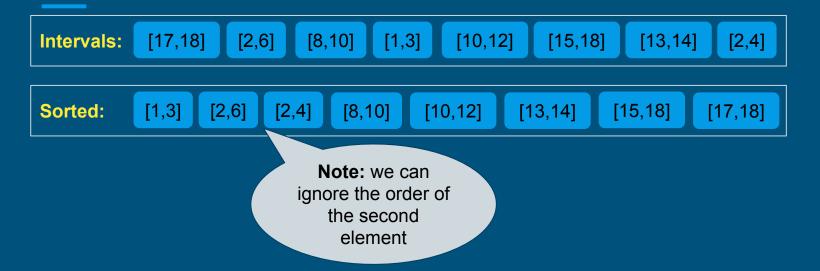
[[1,6],[8,10],[15,18]]

- 1. Sort the array on the first element
- 2. Add the first element to the result
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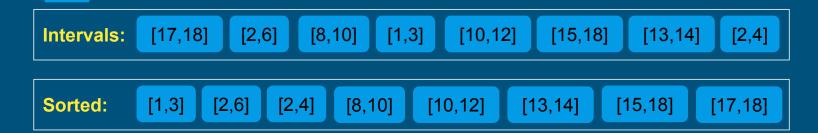
Intervals: [17,18] [2,6] [8,10] [1,3] [10,12] [15,18] [13,14] [2,4]



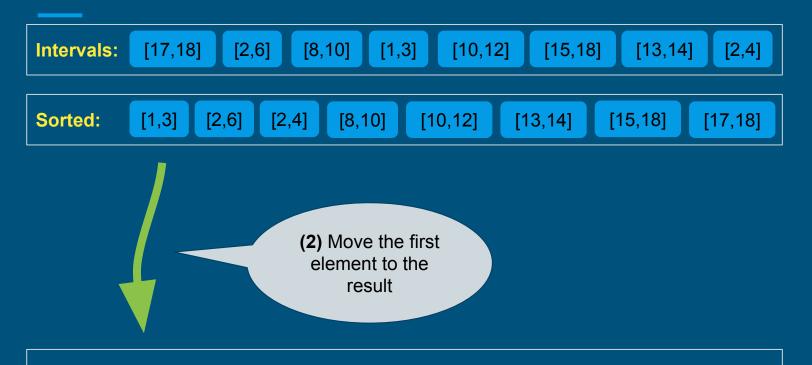


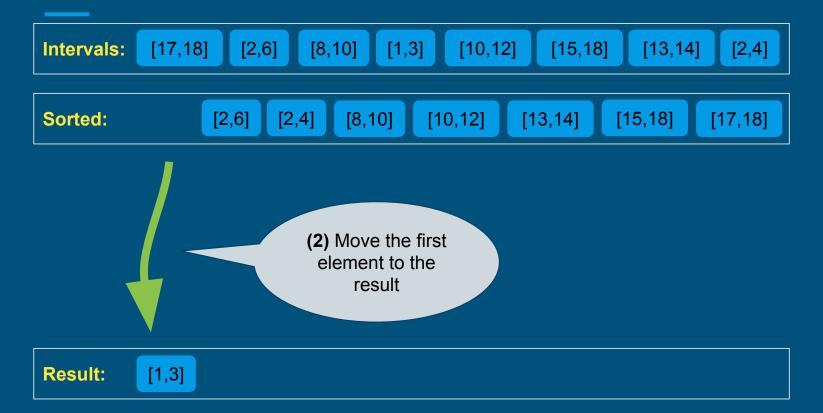


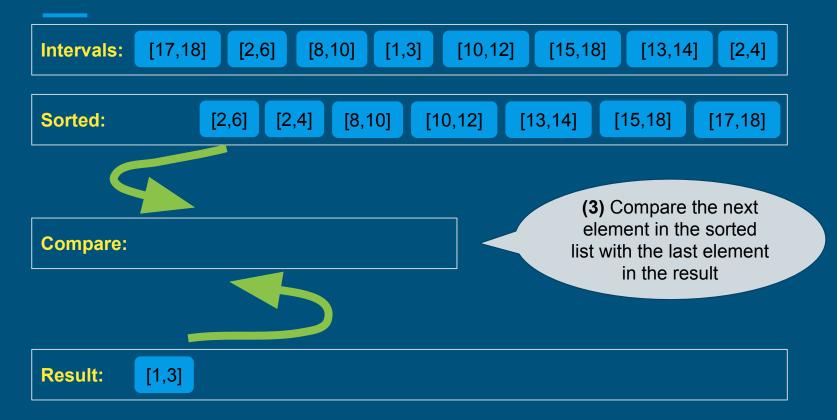
**Result:** 

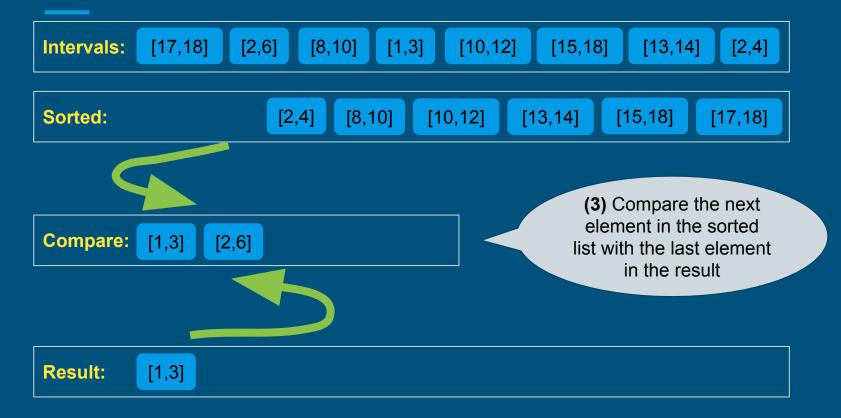


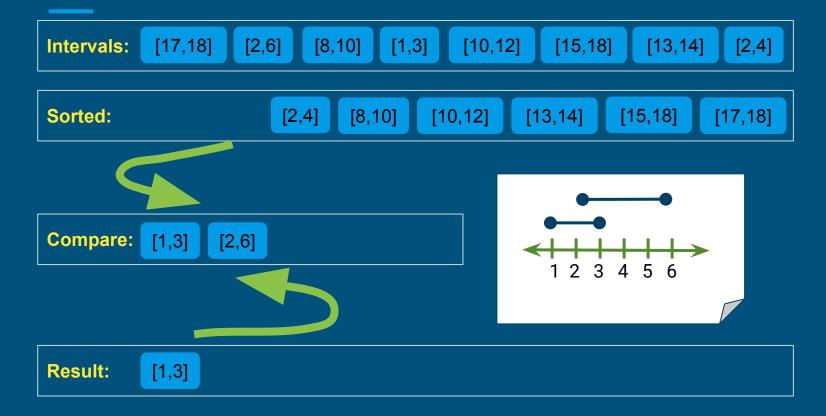
**Result:** 

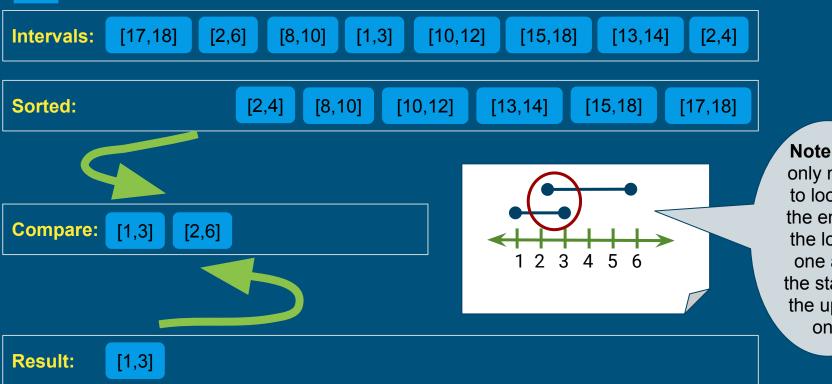




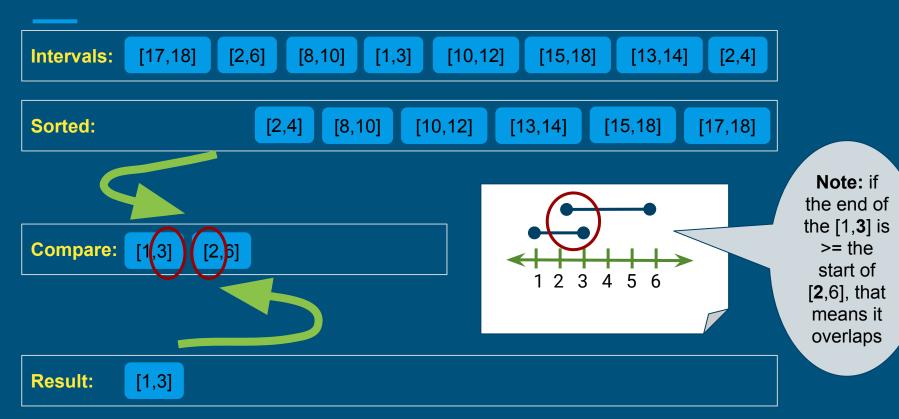


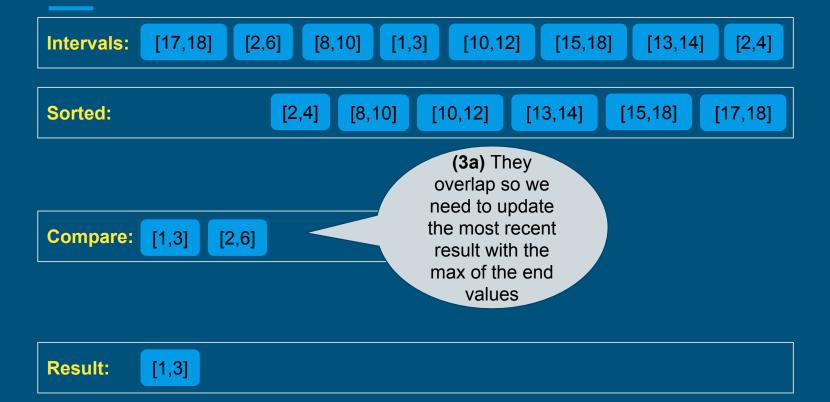


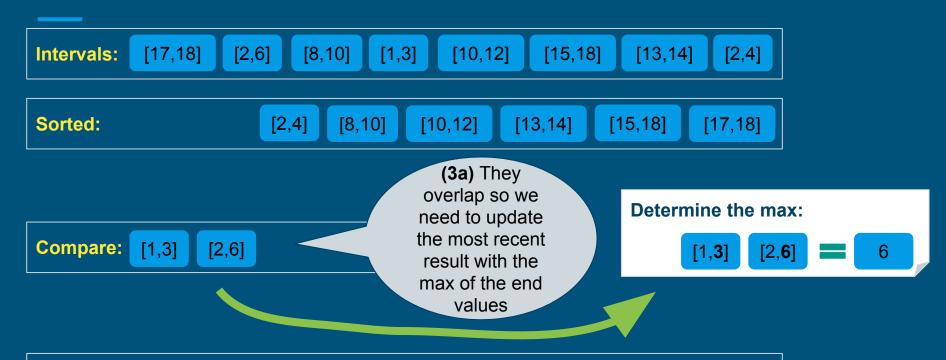




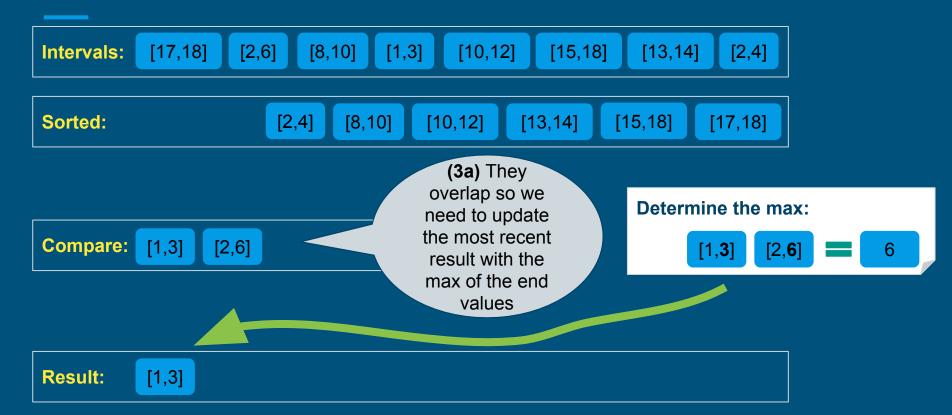
Note: we only need to look at the end of the lower one and the start of the upper one

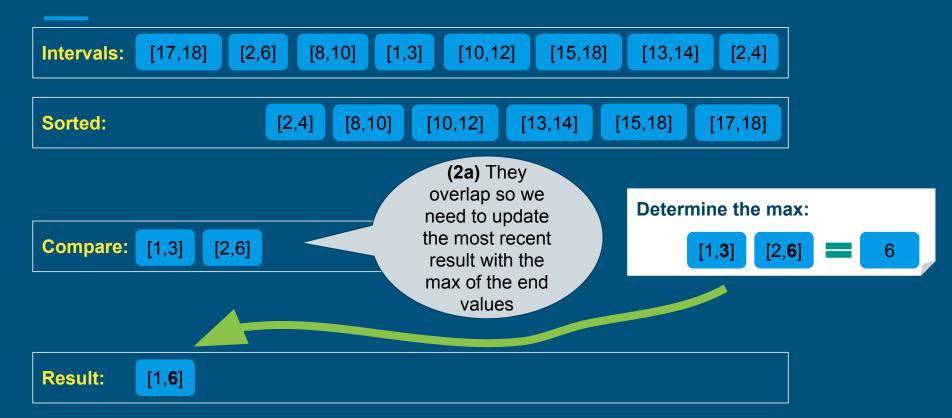


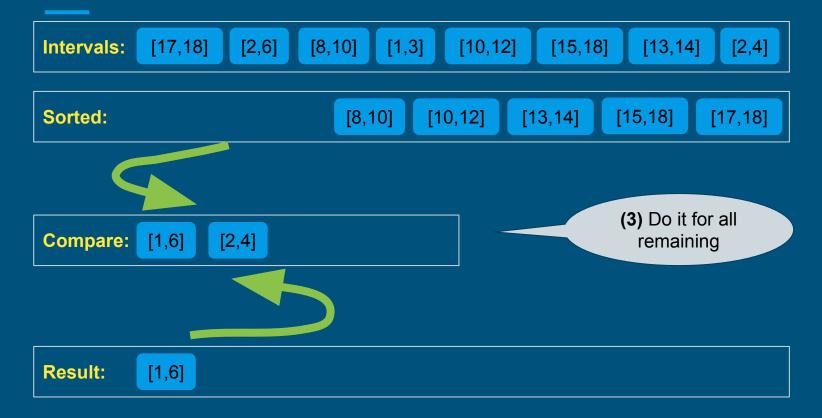




Result: [1,3]

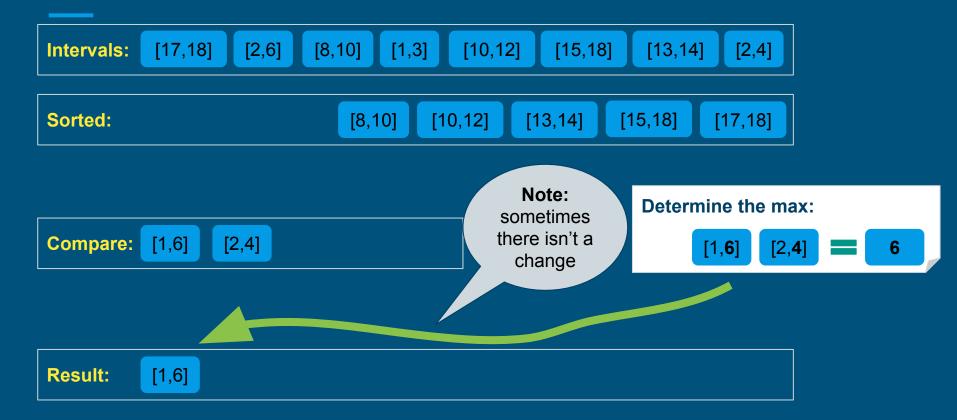








Result: [1,6]

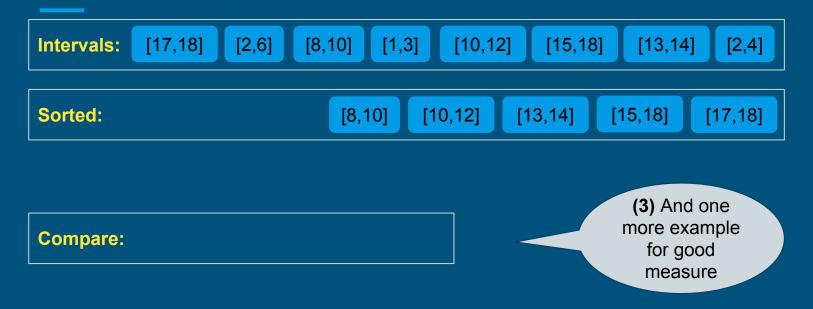


Intervals: [17,18] [2,6] [8,10] [1,3] [10,12] [15,18] [13,14] [2,4]

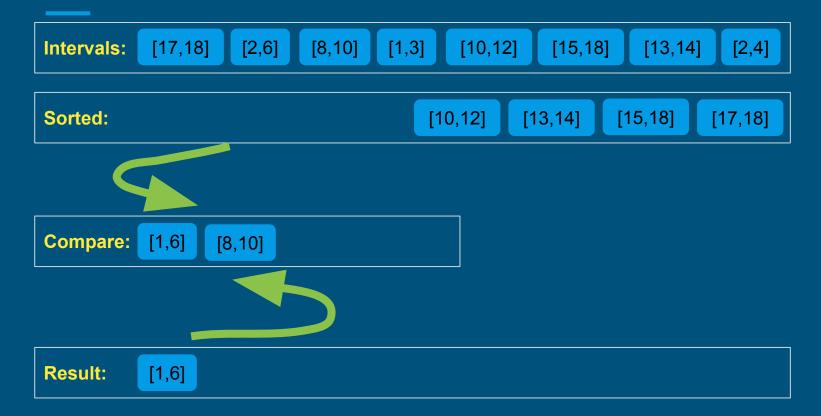
Sorted: [8,10] [10,12] [13,14] [15,18] [17,18]

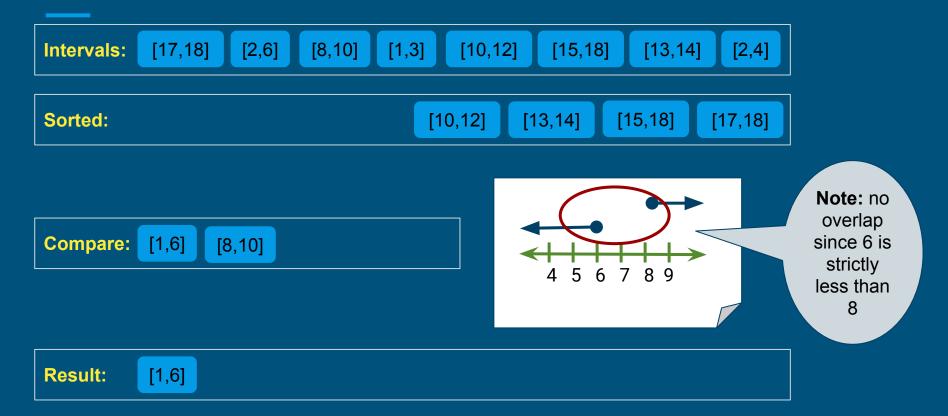
Compare:

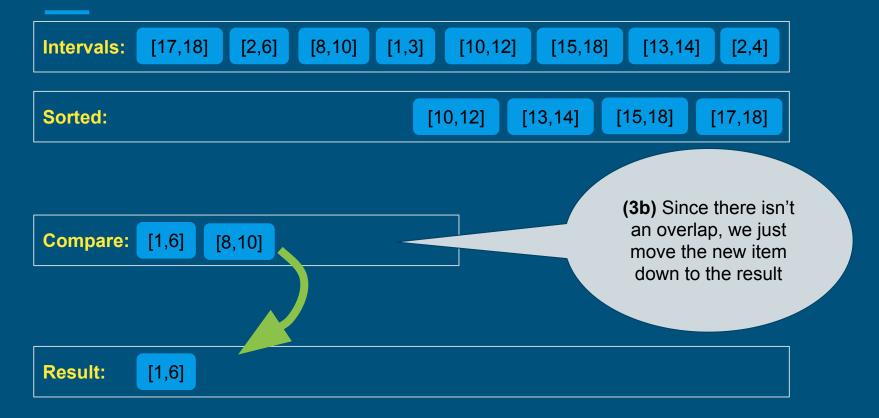
Result: [1,6]

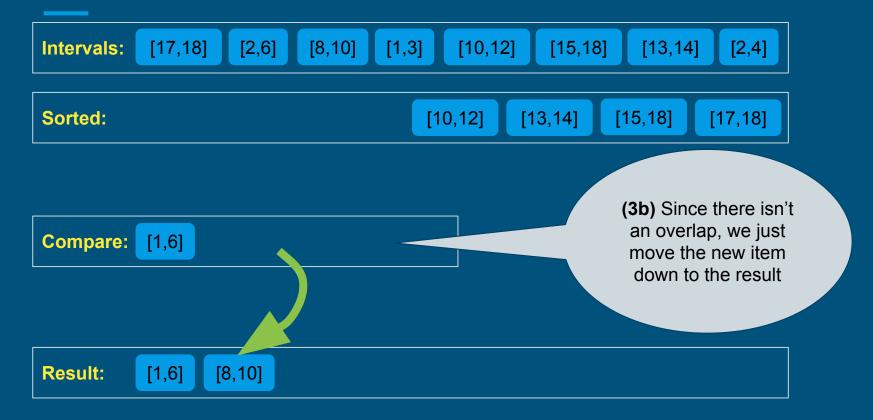


Result: [1,6]







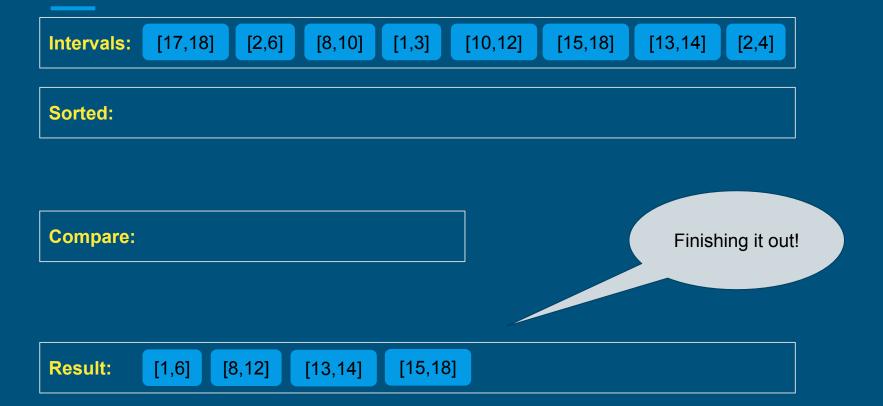


Intervals: [17,18] [2,6] [8,10] [1,3] [10,12] [15,18] [13,14] [2,4]

Sorted: [10,12] [13,14] [15,18] [17,18]

**Compare:** 

**Result:** [1,6] [8,10]



Intervals: [17,18] [2,6] [8,10] [1,3] [10,12] [15,18] [13,14] [2,4]

Sorted:

Compare:

**Result:** [1,6] [8,12] [13,14] [15,18]

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# Let's code it!

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Note: This can depend, but I'm using Java's Arrays.sort(...) in Java 8 which is duval-pivot quicksort

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O(1)

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