

# Trabajo Práctico 1

## Análisis de Lenguajes de Programación

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## Ejercicio 1

### Gramática Abstracta

$$\begin{aligned} \textit{intexp} &::= \textit{nat} \mid \textit{var} \mid -_u \textit{intexp} \\ &\mid \textit{intexp} + \textit{intexp} \\ &\mid \textit{intexp} -_b \textit{intexp} \\ &\mid \textit{intexp} \times \textit{intexp} \\ &\mid \textit{intexp} \div \textit{intexp} \\ &\mid \textit{var} = \textit{intexp} \\ &\mid \textit{intexp}, \textit{intexp} \\ \textit{boolexp} &::= \mathbf{true} \mid \mathbf{false} \\ &\mid \textit{intexp} == \textit{intexp} \\ &\mid \textit{intexp} \neq \textit{intexp} \\ &\mid \textit{intexp} < \textit{intexp} \\ &\mid \textit{intexp} > \textit{intexp} \\ &\mid \textit{boolexp} \wedge \textit{boolexp} \\ &\mid \textit{boolexp} \vee \textit{boolexp} \\ &\mid \neg \textit{boolexp} \\ \textit{comm} &::= \mathbf{skip} \\ &\mid \textit{var} = \textit{intexp} \\ &\mid \textit{comm}; \textit{comm} \\ &\mid \mathbf{if} \textit{boolexp} \mathbf{then} \textit{comm} \mathbf{else} \textit{comm} \\ &\mid \mathbf{repeat} \textit{comm} \mathbf{until} \textit{boolexp} \end{aligned}$$

### Gramática Concreta

Estan mal algunos simbolos

```

digit ::= '0' | '1' | ... | '9'
letter ::= 'a' | ... | 'Z'
nat ::= digit | digit nat
var ::= letter | letter var
intexp ::= nat
          | var
          | '-' intexp
          | intexp '+' intexp
          | intexp '-' intexp
          | intexp '*' intexp
          | intexp '/' intexp
          | '(' intexp ')'
          | var '=' intexp
          | intexp ',' intexp
boolexp ::= 'true' | 'false'
          | intexp '==' intexp
          | intexp '!=' intexp
          | intexp '!' intexp
          | intexp '!' intexp
          | boolexp '&&' boolexp
          | boolexp '——' boolexp
          | '!' boolexp
          | '(' boolexp ')'
comm ::= skip
          | var '=' intexp
          | comm ';' comm
          | 'if' boolexp '{' comm '}'
          | 'if' boolexp '{' comm '}' 'else' '{' comm
          | 'repeat' comm 'until' boolexp 'end'

```

## Ejercicio 6

Habría que enunciar las nuevas reglas (ejercicio 4) y cambiar el nombre aca, y enunciar las reglas de la clausura transitiva.

A:	$\frac{\langle \text{skip}; \text{repeat } x = x - y \text{ until } x == 0, [[\sigma x:1] y:1] \rangle \rightsquigarrow \langle \text{repeat } x = x - y \text{ until } x == 0, [[\sigma x:1] y:1] \rangle}{\langle \text{skip}; \text{repeat } x = x - y \text{ until } x == 0, [[\sigma x:1] y:1] \rangle \rightsquigarrow^* \langle \text{repeat } x = x - y \text{ until } x == 0, [[\sigma x:1] y:1] \rangle} \text{SEQ}_1 \text{ T}_1$	Análisis de Lenguajes de Programación
B:	$\frac{\langle \text{repeat } x = x - y \text{ until } x == 0, [[\sigma x:1] y:1] \rangle \rightsquigarrow \langle x = x - y; \text{if } x == 0 \text{ then skip else repeat } x = x - y \text{ until } x == 0, [[\sigma x:1] y:1] \rangle}{\langle \text{repeat } x = x - y \text{ until } x == 0, [[\sigma x:1] y:1] \rangle \rightsquigarrow^* \langle x = x - y; \text{if } x == 0 \text{ then skip else repeat } x = x - y \text{ until } x == 0, [[\sigma x:1] y:1] \rangle} \text{REPEAT T}_1$	
C:	$\frac{\frac{\langle x, [[\sigma x:1] y:1] \rangle \Downarrow_{exp} \langle \mathbf{1}, [[\sigma x:1] y:1] \rangle}{\langle x - y, [[\sigma x:1] y:1] \rangle \Downarrow_{exp} \langle \mathbf{0}, [[\sigma x:1] y:1] \rangle} \text{VAR} \quad \frac{\langle y, [[\sigma x:1] y:1] \rangle \Downarrow_{exp} \langle \mathbf{1}, [[\sigma x:1] y:1] \rangle}{\langle x = x - y, [[\sigma x:1] y:1] \rangle \rightsquigarrow \langle \text{skip}, [[\sigma x:0] y:1] \rangle} \text{MINUS} \quad \text{ASS}$	
D:	$\frac{\langle x = x - y; \text{if } x == 0 \text{ then skip else repeat } x = x - y \text{ until } x == 0, [[\sigma x:1] y:1] \rangle \rightsquigarrow \langle \text{skip}; \text{if } x == 0 \text{ then skip else repeat } x = x - y \text{ until } x == 0, [[\sigma x:1] y:1] \rangle}{\langle x = x - y; \text{if } x == 0 \text{ then skip else repeat } x = x - y \text{ until } x == 0, [[\sigma x:1] y:1] \rangle \rightsquigarrow^* \langle \text{skip}; \text{if } x == 0 \text{ then skip else repeat } x = x - y \text{ until } x == 0, [[\sigma x:1] y:1] \rangle} \text{TP}_1$ $\frac{\frac{\langle 1, [[\sigma x:2] y:2] \rangle \Downarrow_{exp} \langle \mathbf{1}, [[\sigma x:2] y:2] \rangle}{\langle y = 1, [[\sigma x:2] y:2] \rangle \Downarrow_{exp} \langle \mathbf{1}, [[\sigma x:2] y:1] \rangle} \text{NVAL} \quad \text{NASS} \quad \frac{\langle x = y = 1, [[\sigma x:2] y:2] \rangle \rightsquigarrow \langle \text{skip}, [[\sigma x:1] y:1] \rangle}{\langle x = y = 1; \text{repeat } x = x - y \text{ until } x == 0, [[\sigma x:2] y:2] \rangle \rightsquigarrow \langle \text{skip}; \text{repeat } x = x - y \text{ until } x == 0, [[\sigma x:1] y:1] \rangle} \text{SEQ}_2 \text{ T}_1$	E:
F:	$\frac{\langle \text{skip}; \text{if } x == 0 \text{ then skip else repeat } x = x - y \text{ until } x == 0, [[\sigma x:0] y:1] \rangle \rightsquigarrow \langle \text{if } x == 0 \text{ then skip else repeat } x = x - y \text{ until } x == 0, [[\sigma x:0] y:1] \rangle}{\langle \text{skip}; \text{if } x == 0 \text{ then skip else repeat } x = x - y \text{ until } x == 0, [[\sigma x:0] y:1] \rangle \rightsquigarrow^* \langle \text{if } x == 0 \text{ then skip else repeat } x = x - y \text{ until } x == 0, [[\sigma x:0] y:1] \rangle} \text{TP}_1$	Cipullo, Sullivan

$$\begin{array}{c}
 \frac{\langle x, [[\sigma|x:0]]y:1 \rangle \Downarrow_{exp} \langle \mathbf{0}, [[\sigma|x:0]]y:1 \rangle}{\langle x == 0, [[\sigma|x:0]]y:1 \rangle \Downarrow_{exp} \langle \mathbf{true}, [[\sigma|x:0]]y:1 \rangle} \text{VAR} \quad \frac{\langle 0, [[\sigma|x:0]]y:1 \rangle \Downarrow_{exp} \langle \mathbf{0}, [[\sigma|x:0]]y:1 \rangle}{\langle x == 0 \text{ then skip else repeat } x = x - y \text{ until } x == 0, [[\sigma|x:0]]y:1 \rangle} \text{EQ} \quad \frac{\text{IF}_1}{\langle \text{if } x == 0 \text{ then skip else repeat } x = x - y \text{ until } x == 0, [[\sigma|x:0]]y:1 \rangle \rightsquigarrow^* \langle \mathbf{skip}, [[\sigma|x:0]]y:1 \rangle} \quad \frac{\text{T}_1}{\langle \text{if } x == 0 \text{ then skip else repeat } x = x - y \text{ until } x == 0, [[\sigma|x:0]]y:1 \rangle \rightsquigarrow^* \langle \mathbf{skip}, [[\sigma|x:0]]y:1 \rangle}
 \end{array}$$

**DEM:**

$$\begin{array}{c}
 D \quad A \quad \frac{\langle x == y = 1; \text{repeat } x = x - y \text{ until } x == 0, [[\sigma|x:2]]y:2 \rangle \rightsquigarrow^* \langle \mathbf{repeat } x = x - y \text{ until } x == 0, [[\sigma|x:1]]y:1 \rangle}{\langle x == y = 1; \text{repeat } x = x - y \text{ until } x == 0, [[\sigma|x:2]]y:2 \rangle \rightsquigarrow^* \langle x = x - y; \text{if } x == 0 \text{ then skip else repeat } x = x - y \text{ until } x == 0, [[\sigma|x:1]]y:1 \rangle} \text{T}_2 \quad B \quad \frac{\langle x == y = 1; \text{repeat } x = x - y \text{ until } x == 0, [[\sigma|x:2]]y:2 \rangle \rightsquigarrow^* \langle \mathbf{skip}; \text{if } x == 0 \text{ then skip else repeat } x = x - y \text{ until } x == 0, [[\sigma|x:0]]y:1 \rangle}{\langle x == y = 1; \text{repeat } x = x - y \text{ until } x == 0, [[\sigma|x:2]]y:2 \rangle \rightsquigarrow^* \langle \text{if } x == 0 \text{ then skip else repeat } x = x - y \text{ until } x == 0, [[\sigma|x:0]]y:1 \rangle} \text{T}_2 \quad \frac{\langle x == y = 1; \text{repeat } x = x - y \text{ until } x == 0, [[\sigma|x:2]]y:2 \rangle \rightsquigarrow^* \langle \mathbf{skip}, [[\sigma|x:0]]y:1 \rangle}{\langle x == y = 1; \text{repeat } x = x - y \text{ until } x == 0, [[\sigma|x:2]]y:2 \rangle \rightsquigarrow^* \langle \mathbf{skip}, [[\sigma|x:0]]y:1 \rangle}
 \end{array}$$

## Ejercicio 2

```

1  — Expresiones , aritmeticas y booleanas
2  data Exp a where
3  — Expresiones enteras
4  Const :: Int -> Exp Int
5  Var :: Variable -> Exp Int
6  UMinus :: Exp Int -> Exp Int
7  Plus :: Exp Int -> Exp Int -> Exp Int
8  Minus :: Exp Int -> Exp Int -> Exp Int
9  Times :: Exp Int -> Exp Int -> Exp Int
10 Div :: Exp Int -> Exp Int -> Exp Int
11 EAssgn :: Variable -> Exp Int -> Exp Int
12 ESeq :: Exp Int -> Exp Int -> Exp Int
13
14 — Expresiones booleanas
15 BTrue :: Exp Bool
16 BFalse :: Exp Bool
17 Lt :: Exp Int -> Exp Int -> Exp Bool
18 Gt :: Exp Int -> Exp Int -> Exp Bool
19 And :: Exp Bool -> Exp Bool -> Exp Bool
20 Or :: Exp Bool -> Exp Bool -> Exp Bool
21 Not :: Exp Bool -> Exp Bool
22 Eq :: Exp Int -> Exp Int -> Exp Bool
23 NEq :: Exp Int -> Exp Int -> Exp Bool

```

## Ejercicio 4

### Semántica Operacional Big-Step para Expresiones

$$\begin{array}{c}
\frac{}{\langle nv, \sigma \rangle \Downarrow_{exp} \langle \mathbf{nv}, \sigma \rangle} \text{ NVAL} \qquad \frac{}{\langle x, \sigma \rangle \Downarrow_{exp} \langle \sigma x, \sigma \rangle} \text{ VAR} \\
\\
\frac{\langle e, \sigma \rangle \Downarrow_{exp} \langle n, \sigma' \rangle}{\langle -_u e, \sigma \rangle \Downarrow_{exp} \langle -n, \sigma' \rangle} \text{ UMINUS} \qquad \frac{\langle e_0, \sigma \rangle \Downarrow_{exp} \langle n_0, \sigma' \rangle \quad \langle e_1, \sigma' \rangle \Downarrow_{exp} \langle n_1, \sigma'' \rangle}{\langle e_0 + e_1, \sigma \rangle \Downarrow_{exp} \langle n_0 + n_1, \sigma'' \rangle} \text{ PLUS} \\
\\
\frac{\langle e_0, \sigma \rangle \Downarrow_{exp} \langle n_0, \sigma' \rangle \quad \langle e_1, \sigma' \rangle \Downarrow_{exp} \langle n_1, \sigma'' \rangle}{\langle e_0 - e_1, \sigma \rangle \Downarrow_{exp} \langle n_0 - n_1, \sigma'' \rangle} \text{ BMINUS} \\
\\
\frac{\langle e_0, \sigma \rangle \Downarrow_{exp} \langle n_0, \sigma' \rangle \quad \langle e_1, \sigma' \rangle \Downarrow_{exp} \langle n_1, \sigma'' \rangle}{\langle e_0 * e_1, \sigma \rangle \Downarrow_{exp} \langle n_0 * n_1, \sigma'' \rangle} \text{ MULT} \\
\\
\frac{\langle e_0, \sigma \rangle \Downarrow_{exp} \langle n_0, \sigma' \rangle \quad \langle e_1, \sigma' \rangle \Downarrow_{exp} \langle n_1, \sigma'' \rangle \quad n_1 \neq 0}{\langle e_0 \div e_1, \sigma \rangle \Downarrow_{exp} \langle n_0 \div n_1, \sigma'' \rangle} \text{ DIV} \\
\\
\frac{\langle e_0, \sigma \rangle \Downarrow_{exp} \langle n_0, \sigma' \rangle \quad \langle e_1, \sigma' \rangle \Downarrow_{exp} \langle n_1, \sigma'' \rangle}{\langle e_0 > e_1, \sigma \rangle \Downarrow_{exp} \langle n_0 > n_1, \sigma'' \rangle} \text{ GT} \\
\\
\frac{\langle e_0, \sigma \rangle \Downarrow_{exp} \langle n_0, \sigma' \rangle \quad \langle e_1, \sigma' \rangle \Downarrow_{exp} \langle n_1, \sigma'' \rangle}{\langle e_0 < e_1, \sigma \rangle \Downarrow_{exp} \langle n_0 < n_1, \sigma'' \rangle} \text{ LT}
\end{array}$$

$$\begin{array}{c}
\frac{\langle e_0, \sigma \rangle \Downarrow_{exp} \langle n_0, \sigma' \rangle \quad \langle e_1, \sigma' \rangle \Downarrow_{exp} \langle n_1, \sigma'' \rangle}{\langle e_0 \neq e_1, \sigma \rangle \Downarrow_{exp} \langle n_0 \neq n_1, \sigma'' \rangle} \text{ NOTEQ} \\
\\
\frac{\langle e_0, \sigma \rangle \Downarrow_{exp} \langle n_0, \sigma' \rangle \quad \langle e_1, \sigma' \rangle \Downarrow_{exp} \langle n_1, \sigma'' \rangle}{\langle e_0 == e_1, \sigma \rangle \Downarrow_{exp} \langle n_0 = n_1, \sigma'' \rangle} \text{ EQ} \quad \frac{}{\langle bv, \sigma \rangle \Downarrow_{exp} \langle bv, \sigma \rangle} \text{ BVAL} \\
\\
\frac{\langle p, \sigma \rangle \Downarrow_{exp} \langle b, \sigma' \rangle}{\langle \neg p, \sigma \rangle \Downarrow_{exp} \langle \neg b, \sigma' \rangle} \text{ NOT} \quad \frac{\langle p_0, \sigma \rangle \Downarrow_{exp} \langle b_0, \sigma' \rangle \quad \langle p_1, \sigma' \rangle \Downarrow_{exp} \langle b_1, \sigma'' \rangle}{\langle e_0 \vee e_1, \sigma \rangle \Downarrow_{exp} \langle b_0 \vee b_1, \sigma'' \rangle} \text{ OR} \\
\\
\frac{\langle p_0, \sigma \rangle \Downarrow_{exp} \langle b_0, \sigma' \rangle \quad \langle p_1, \sigma' \rangle \Downarrow_{exp} \langle b_1, \sigma'' \rangle}{\langle e_0 \wedge e_1, \sigma \rangle \Downarrow_{exp} \langle b_0 \wedge b_1, \sigma'' \rangle} \text{ AND}
\end{array}$$

Llamamos IASS a la asignación como expresión y ISEQ a la secuencialización de expresiones con el operador ,

$$\begin{array}{c}
\frac{\langle e, \sigma \rangle \Downarrow_{exp} \langle n, \sigma' \rangle}{\langle v = e, \sigma \rangle \Downarrow_{exp} \langle n, [\sigma'|v : n] \rangle} \text{ IASS} \\
\\
\frac{\langle e_0, \sigma \rangle \Downarrow_{exp} \langle n_0, \sigma' \rangle \quad \langle e_1, \sigma' \rangle \Downarrow_{exp} \langle n_1, \sigma'' \rangle}{\langle e_0, e_1, \sigma \rangle \Downarrow_{exp} \langle n_1, \sigma'' \rangle} \text{ ISEQ}
\end{array}$$

## Ejercicio 5

Determinismo de la relación de evaluación en un paso  $\rightsquigarrow$

## Ejercicio 10

Nueva gramática abstracta

for (intexp;boolexp;intexp) comm

Nueva semántica operacional para comandos

$$\begin{array}{c}
\frac{\langle e, \sigma \rangle \Downarrow_{exp} \langle n, \sigma' \rangle}{\langle v = e, \sigma \rangle \rightsquigarrow \langle \mathbf{skip}, [\sigma'|v : n] \rangle} \text{ ASS} \\
\\
\frac{}{\langle \mathbf{skip}; c_1, \sigma \rangle \rightsquigarrow \langle c_1, \sigma \rangle} \text{ SEQ1} \\
\\
\frac{\langle c_0, \sigma \rangle \rightsquigarrow \langle c'_0, \sigma' \rangle}{\langle \mathbf{skip}; c_1, \sigma \rangle \rightsquigarrow \langle c'_0; c_1, \sigma' \rangle} \text{ SEQ2}
\end{array}$$

$$\begin{array}{c}
\frac{\langle b, \sigma \rangle \Downarrow_{exp} \langle \mathbf{true}, \sigma' \rangle}{\langle \mathbf{if } b \mathbf{ then } c_0 \mathbf{ else } c_1, \sigma \rangle \rightsquigarrow \langle c_0, \sigma' \rangle} \text{IF}_1 \qquad \frac{\langle b, \sigma \rangle \Downarrow_{exp} \langle \mathbf{false}, \sigma' \rangle}{\langle \mathbf{if } b \mathbf{ then } c_0 \mathbf{ else } c_1, \sigma \rangle \rightsquigarrow \langle c_1, \sigma' \rangle} \text{IF}_2 \\
\\
\frac{}{\langle \mathbf{repeat } c \mathbf{ until } b, \sigma \rangle \rightsquigarrow \langle c; \mathbf{if } b \mathbf{ then skip else repeat } b \mathbf{ until } c, \sigma \rangle} \text{REPEAT}_1 \\
\\
\frac{\langle e, \sigma \rangle \Downarrow_{exp} \langle n, \sigma' \rangle}{\langle \mathbf{repeat } (c; e) \mathbf{ until } b, \sigma \rangle \rightsquigarrow} \text{REPEAT}_2 \\
\\
\frac{\langle e_1, \sigma \rangle \Downarrow_{exp} \langle n_1, \sigma' \rangle}{\langle \mathbf{for } (e_1; e_2; e_3) \ c, \sigma \rangle \rightsquigarrow \langle \mathbf{repeat } (c; e_3) \mathbf{ until } e_2, \sigma' \rangle} \text{FOR}
\end{array}$$