

# Trabajo Práctico 1

## Análisis de Lenguajes de Programación

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## Ejercicio 1

### Gramática Abstracta

$$\begin{aligned} \text{intexp} &::= \text{nat} \mid \text{var} \mid -_u \text{intexp} \\ &\mid \text{intexp} + \text{intexp} \\ &\mid \text{intexp} -_b \text{intexp} \\ &\mid \text{intexp} \times \text{intexp} \\ &\mid \text{intexp} \div \text{intexp} \\ &\mid \text{var} = \text{intexp} \\ &\mid \text{intexp}, \text{intexp} \\ \text{boolexp} &::= \mathbf{true} \mid \mathbf{false} \\ &\mid \text{intexp} == \text{intexp} \\ &\mid \text{intexp} \neq \text{intexp} \\ &\mid \text{intexp} < \text{intexp} \\ &\mid \text{intexp} > \text{intexp} \\ &\mid \text{boolexp} \wedge \text{boolexp} \\ &\mid \text{boolexp} \vee \text{boolexp} \\ &\mid \neg \text{boolexp} \\ \text{comm} &::= \mathbf{skip} \\ &\mid \text{var} = \text{intexp} \\ &\mid \text{comm}; \text{comm} \\ &\mid \mathbf{if} \text{ boolexp} \mathbf{then} \text{ comm} \mathbf{else} \text{ comm} \\ &\mid \mathbf{repeat} \text{ comm} \mathbf{until} \text{ boolexp} \end{aligned}$$

### Gramática Concreta

Estan mal algunos simbolos

```

digit ::= '0' | '1' | ... | '9'
letter ::= 'a' | ... | 'Z'
nat ::= digit | digit nat
var ::= letter | letter var
intexp ::= nat
        | var
        | '-' intexp
        | intexp '+' intexp
        | intexp '-' intexp
        | intexp '*' intexp
        | intexp '/' intexp
        | '(' intexp ')'
        | var '=' intexp
        | intexp ';' intexp
boolexp ::= 'true' | 'false'
        | intexp '==' intexp
        | intexp '!=' intexp
        | intexp '!' intexp
        | intexp '!' intexp
        | boolexp '&&' boolexp
        | boolexp '——' boolexp
        | '!' boolexp
        | '(' boolexp ')'
comm ::= skip
        | var '=' intexp
        | comm ';' comm
        | 'if' boolexp '{' comm '}'
        | 'if' boolexp '{' comm '}' 'else' '{' comm
        | 'repeat' comm 'until' boolexp 'end'

```

## Ejercicio 6

Habría que enunciar las nuevas reglas (ejercicio 4) y cambiar el nombre aca, y enunciar las reglas de la clausura transitiva.



$$\begin{array}{c}
 \frac{\langle x, [[\sigma|x:0]|y:1] \rangle \Downarrow_{exp} \langle \mathbf{0}, [[\sigma|x:0]|y:1] \rangle}{\langle x == 0, [[\sigma|x:0]|y:1] \rangle \Downarrow_{exp} \langle \mathbf{true}, [[\sigma|x:0]|y:1] \rangle} \text{VAR} \quad \frac{\langle 0, [[\sigma|x:0]|y:1] \rangle \Downarrow_{exp} \langle \mathbf{0}, [[\sigma|x:0]|y:1] \rangle}{\langle x == 0 \text{ then skip else repeat } x = x - y \text{ until } x == 0, [[\sigma|x:0]|y:1] \rangle \rightsquigarrow^* \langle \mathbf{skip}, [[\sigma|x:0]|y:1] \rangle} \text{EQ} \\
 \frac{\langle x == 0 \text{ then skip else repeat } x = x - y \text{ until } x == 0, [[\sigma|x:0]|y:1] \rangle \rightsquigarrow^* \langle \mathbf{skip}, [[\sigma|x:0]|y:1] \rangle}{\langle x == 0 \text{ then skip else repeat } x = x - y \text{ until } x == 0, [[\sigma|x:0]|y:1] \rangle \rightsquigarrow^* \langle \mathbf{skip}, [[\sigma|x:0]|y:1] \rangle} \text{IF}_1 \quad \text{T}_1
 \end{array}$$

**DEM:**

$$\begin{array}{c}
 D \quad A \quad \frac{\langle x == y = 1; \text{repeat } x = x - y \text{ until } x == 0, [[\sigma|x:2]|y:2] \rangle \rightsquigarrow^* \langle \mathbf{repeat } x = x - y \text{ until } x == 0, [[\sigma|x:1]|y:1] \rangle}{\langle x == y = 1; \text{repeat } x = x - y \text{ until } x == 0, [[\sigma|x:2]|y:2] \rangle \rightsquigarrow^* \langle x = x - y; \text{if } x == 0 \text{ then skip else repeat } x = x - y \text{ until } x == 0, [[\sigma|x:1]|y:1] \rangle} \text{T}_2 \quad B \\
 \frac{\langle x == y = 1; \text{repeat } x = x - y \text{ until } x == 0, [[\sigma|x:2]|y:2] \rangle \rightsquigarrow^* \langle \mathbf{skip}; \text{if } x == 0 \text{ then skip else repeat } x = x - y \text{ until } x == 0, [[\sigma|x:0]|y:1] \rangle}{\langle x == y = 1; \text{repeat } x = x - y \text{ until } x == 0, [[\sigma|x:2]|y:2] \rangle \rightsquigarrow^* \langle \text{if } x == 0 \text{ then skip else repeat } x = x - y \text{ until } x == 0, [[\sigma|x:0]|y:1] \rangle} \text{T}_2 \\
 \frac{\langle x == y = 1; \text{repeat } x = x - y \text{ until } x == 0, [[\sigma|x:2]|y:2] \rangle \rightsquigarrow^* \langle \mathbf{skip}, [[\sigma|x:2]|y:2] \rangle \rightsquigarrow^* \langle \mathbf{skip}, [[\sigma|x:0]|y:1] \rangle}{\langle x == y = 1; \text{repeat } x = x - y \text{ until } x == 0, [[\sigma|x:2]|y:2] \rangle \rightsquigarrow^* \langle \mathbf{skip}, [[\sigma|x:0]|y:1] \rangle}
 \end{array}$$

## Ejercicio 2

```

1  — Expresiones, aritmeticas y booleanas
2  data Exp a where
3  — Expresiones enteras
4  Const :: Int -> Exp Int
5  Var :: Variable -> Exp Int
6  UMinus :: Exp Int -> Exp Int
7  Plus :: Exp Int -> Exp Int -> Exp Int
8  Minus :: Exp Int -> Exp Int -> Exp Int
9  Times :: Exp Int -> Exp Int -> Exp Int
10 Div :: Exp Int -> Exp Int -> Exp Int
11 EAssgn :: Variable -> Exp Int -> Exp Int
12 ESeq :: Exp Int -> Exp Int -> Exp Int
13
14 — Expresiones booleanas
15 BTrue :: Exp Bool
16 BFalse :: Exp Bool
17 Lt :: Exp Int -> Exp Int -> Exp Bool
18 Gt :: Exp Int -> Exp Int -> Exp Bool
19 And :: Exp Bool -> Exp Bool -> Exp Bool
20 Or :: Exp Bool -> Exp Bool -> Exp Bool
21 Not :: Exp Bool -> Exp Bool
22 Eq :: Exp Int -> Exp Int -> Exp Bool
23 NEq :: Exp Int -> Exp Int -> Exp Bool

```

## Ejercicio 4

### Semántica Operacional Big-Step para Expresiones

$$\begin{array}{c}
\frac{}{\langle nv, \sigma \rangle \Downarrow_{exp} \langle \mathbf{nv}, \sigma \rangle} \text{ NVAL} \qquad \frac{}{\langle x, \sigma \rangle \Downarrow_{exp} \langle \sigma x, \sigma \rangle} \text{ VAR} \\
\\
\frac{\langle e, \sigma \rangle \Downarrow_{exp} \langle n, \sigma' \rangle}{\langle -_u e, \sigma \rangle \Downarrow_{exp} \langle -n, \sigma' \rangle} \text{ UMINUS} \qquad \frac{\langle e_0, \sigma \rangle \Downarrow_{exp} \langle n_0, \sigma' \rangle \quad \langle e_1, \sigma' \rangle \Downarrow_{exp} \langle n_1, \sigma'' \rangle}{\langle e_0 + e_1, \sigma \rangle \Downarrow_{exp} \langle n_0 + n_1, \sigma'' \rangle} \text{ PLUS} \\
\\
\frac{\langle e_0, \sigma \rangle \Downarrow_{exp} \langle n_0, \sigma' \rangle \quad \langle e_1, \sigma' \rangle \Downarrow_{exp} \langle n_1, \sigma'' \rangle}{\langle e_0 - e_1, \sigma \rangle \Downarrow_{exp} \langle n_0 - n_1, \sigma'' \rangle} \text{ BMINUS} \\
\\
\frac{\langle e_0, \sigma \rangle \Downarrow_{exp} \langle n_0, \sigma' \rangle \quad \langle e_1, \sigma' \rangle \Downarrow_{exp} \langle n_1, \sigma'' \rangle}{\langle e_0 * e_1, \sigma \rangle \Downarrow_{exp} \langle n_0 * n_1, \sigma'' \rangle} \text{ MULT} \\
\\
\frac{\langle e_0, \sigma \rangle \Downarrow_{exp} \langle n_0, \sigma' \rangle \quad \langle e_1, \sigma' \rangle \Downarrow_{exp} \langle n_1, \sigma'' \rangle \quad n_1 \neq 0}{\langle e_0 \div e_1, \sigma \rangle \Downarrow_{exp} \langle n_0 \div n_1, \sigma'' \rangle} \text{ DIV} \\
\\
\frac{\langle e_0, \sigma \rangle \Downarrow_{exp} \langle n_0, \sigma' \rangle \quad \langle e_1, \sigma' \rangle \Downarrow_{exp} \langle n_1, \sigma'' \rangle}{\langle e_0 > e_1, \sigma \rangle \Downarrow_{exp} \langle n_0 > n_1, \sigma'' \rangle} \text{ GT} \\
\\
\frac{\langle e_0, \sigma \rangle \Downarrow_{exp} \langle n_0, \sigma' \rangle \quad \langle e_1, \sigma' \rangle \Downarrow_{exp} \langle n_1, \sigma'' \rangle}{\langle e_0 < e_1, \sigma \rangle \Downarrow_{exp} \langle n_0 < n_1, \sigma'' \rangle} \text{ LT}
\end{array}$$

$$\begin{array}{c}
\frac{\langle e_0, \sigma \rangle \Downarrow_{exp} \langle n_0, \sigma' \rangle \quad \langle e_1, \sigma' \rangle \Downarrow_{exp} \langle n_1, \sigma'' \rangle}{\langle e_0 \neq e_1, \sigma \rangle \Downarrow_{exp} \langle n_0 \neq n_1, \sigma'' \rangle} \text{ NOTEQ} \\
\\
\frac{\langle e_0, \sigma \rangle \Downarrow_{exp} \langle n_0, \sigma' \rangle \quad \langle e_1, \sigma' \rangle \Downarrow_{exp} \langle n_1, \sigma'' \rangle}{\langle e_0 == e_1, \sigma \rangle \Downarrow_{exp} \langle n_0 = n_1, \sigma'' \rangle} \text{ EQ} \quad \frac{}{\langle bv, \sigma \rangle \Downarrow_{exp} \langle bv, \sigma \rangle} \text{ BVAL} \\
\\
\frac{\langle p, \sigma \rangle \Downarrow_{exp} \langle b, \sigma' \rangle}{\langle \neg p, \sigma \rangle \Downarrow_{exp} \langle \neg b, \sigma' \rangle} \text{ NOT} \quad \frac{\langle p_0, \sigma \rangle \Downarrow_{exp} \langle b_0, \sigma' \rangle \quad \langle p_1, \sigma' \rangle \Downarrow_{exp} \langle b_1, \sigma'' \rangle}{\langle e_0 \vee e_1, \sigma \rangle \Downarrow_{exp} \langle b_0 \vee b_1, \sigma'' \rangle} \text{ OR} \\
\\
\frac{\langle p_0, \sigma \rangle \Downarrow_{exp} \langle b_0, \sigma' \rangle \quad \langle p_1, \sigma' \rangle \Downarrow_{exp} \langle b_1, \sigma'' \rangle}{\langle e_0 \wedge e_1, \sigma \rangle \Downarrow_{exp} \langle b_0 \wedge b_1, \sigma'' \rangle} \text{ AND}
\end{array}$$

Llamamos IASS a la asignación como expresión y ISEQ a la secuencialización de expresiones con el operador ,

$$\begin{array}{c}
\frac{\langle e, \sigma \rangle \Downarrow_{exp} \langle n, \sigma' \rangle}{\langle v = e, \sigma \rangle \Downarrow_{exp} \langle n, [\sigma'|v : n] \rangle} \text{ IASS} \\
\\
\frac{\langle e_0, \sigma \rangle \Downarrow_{exp} \langle n_0, \sigma' \rangle \quad \langle e_1, \sigma' \rangle \Downarrow_{exp} \langle n_1, \sigma'' \rangle}{\langle e_0, e_1, \sigma \rangle \Downarrow_{exp} \langle n_1, \sigma'' \rangle} \text{ ISEQ}
\end{array}$$

## Ejercicio 5

Determinismo de la relación de evaluación en un paso  $\rightsquigarrow$

## Ejercicio 10

Nueva gramática abstracta

for (intexp;boolexp;intexp) comm

Nueva semántica operacional para comandos

$$\begin{array}{c}
\frac{\langle e, \sigma \rangle \Downarrow_{exp} \langle n, \sigma' \rangle}{\langle v = e, \sigma \rangle \rightsquigarrow \langle \mathbf{skip}, [\sigma'|v : n] \rangle} \text{ ASS} \\
\\
\frac{}{\langle \mathbf{skip}; c_1, \sigma \rangle \rightsquigarrow \langle c_1, \sigma \rangle} \text{ SEQ1} \\
\\
\frac{\langle c_0, \sigma \rangle \rightsquigarrow \langle c'_0, \sigma' \rangle}{\langle \mathbf{skip}; c_1, \sigma \rangle \rightsquigarrow \langle c'_0; c_1, \sigma' \rangle} \text{ SEQ2}
\end{array}$$

$$\begin{array}{c}
\frac{\langle b, \sigma \rangle \Downarrow_{exp} \langle \mathbf{true}, \sigma' \rangle}{\langle \mathbf{if } b \mathbf{ then } c_0 \mathbf{ else } c_1, \sigma \rangle \rightsquigarrow \langle c_0, \sigma' \rangle} \text{IF}_1 \qquad \frac{\langle b, \sigma \rangle \Downarrow_{exp} \langle \mathbf{false}, \sigma' \rangle}{\langle \mathbf{if } b \mathbf{ then } c_0 \mathbf{ else } c_1, \sigma \rangle \rightsquigarrow \langle c_1, \sigma' \rangle} \text{IF}_2 \\
\\
\frac{}{\langle \mathbf{repeat } c \mathbf{ until } b, \sigma \rangle \rightsquigarrow \langle c; \mathbf{if } b \mathbf{ then skip else repeat } b \mathbf{ until } c, \sigma \rangle} \text{REPEAT}_1 \\
\\
\frac{\langle e, \sigma \rangle \Downarrow_{exp} \langle n, \sigma' \rangle}{\langle \mathbf{repeat } (c; e) \mathbf{ until } b, \sigma \rangle \rightsquigarrow} \text{REPEAT}_2 \\
\\
\frac{\langle e_1, \sigma \rangle \Downarrow_{exp} \langle n_1, \sigma' \rangle}{\langle \mathbf{for } (e_1; e_2; e_3) \ c, \sigma \rangle \rightsquigarrow \langle \mathbf{repeat } (c; e_3) \mathbf{ until } e_2, \sigma' \rangle} \text{FOR}
\end{array}$$