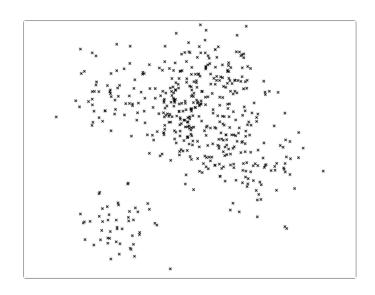
## Introduction to Data Science

APPROXIMATE
K-NEAREST NEIGHBORS
BRIAN D'ALESSANDRO

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# APPROXIMATE NEAREST NEIGHBORS

As mentioned before, searching for a nearest neighbor takes O(n) time, which can be prohibitive for large n. If we are willing to make some tradeoffs, we can design a nearest neighbor search in O(log(n)) time.

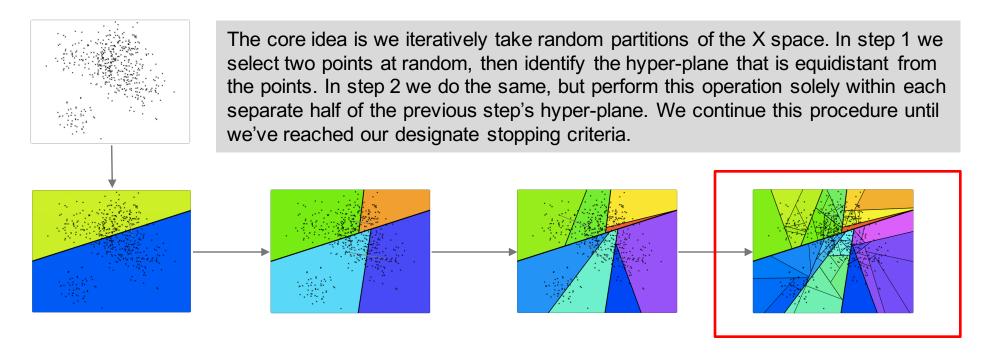


Instead of searching through all points, we will partition the *X* space to create neighborhoods, and search only within those neighborhoods.

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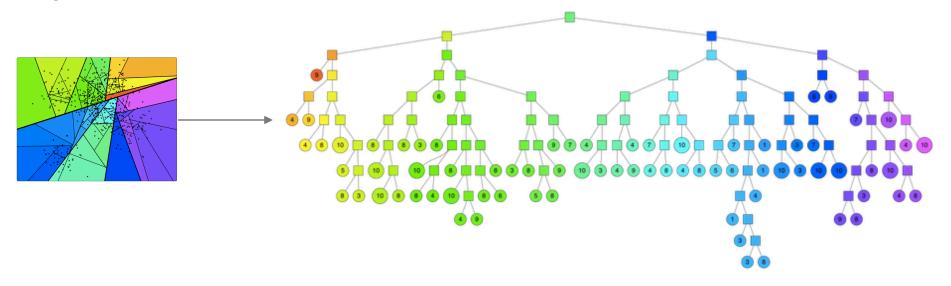
#### **CREATING THE NEIGHBORHOODS**

There are multiple approaches, so we'll use the ANNOY \* package as an example



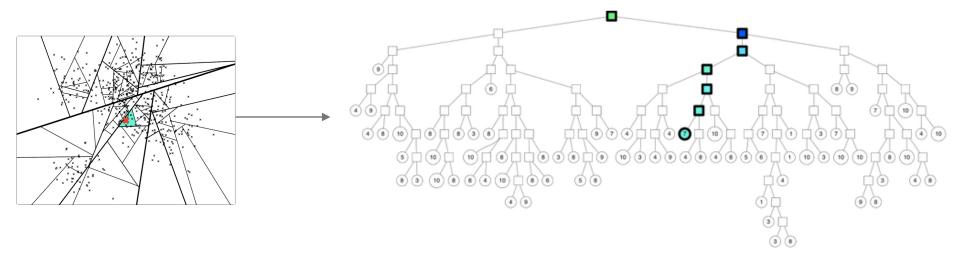
### **SEARCHING FOR NEIGHBORS**

After K recursive splits, our partitioned data can be represented with a binary tree data structure. Finding a set of neighbors amounts to taking an input x, traversing the tree to see which node it lands in, and taking the neighbors from that node.



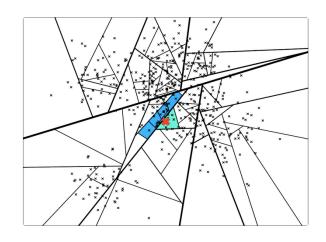
### **SEARCHING FOR NEIGHBORS**

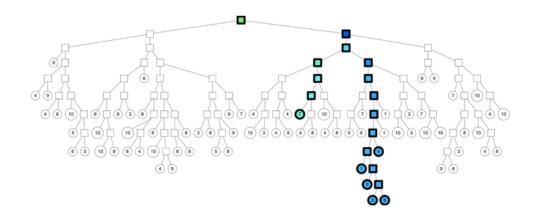
After K recursive splits, our partitioned data can be represented with a binary tree data structure. Finding a set of neighbors amounts to taking an input x, traversing the tree to see which node it lands in, and taking the neighbors from that node.



## **THE TRADEOFFS**

**Problem:** You are limited to just the neighbors in the terminal node of the tree

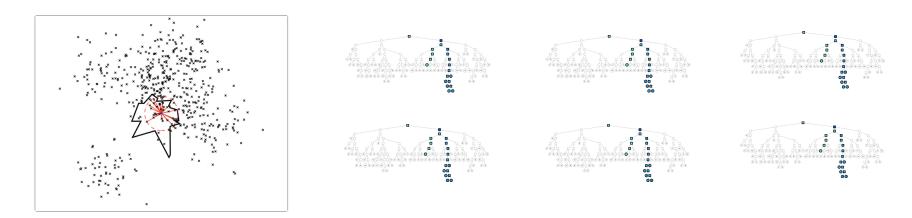




**Solution:** You don't have to use a single terminal node. Nearby splits (regions) can also be considered)

## **THE TRADEOFFS**

**Problem:** You may miss true nearest neighbors that are on the boundaries of the partitions



**Solution:** You can build a forest using multiple trees. Run the full algorithm D times to get D separate neighborhoods for a point. When scoring, the candidate set becomes the union of all D neighborhoods.