MPI Basics

Chris Kauffman

Last Updated: Mon Sep 20 09:25:56 AM CDT 2021

Logistics

Reading: Grama Ch 6 + 4

- ► Ch 6: MPI basics
- Ch 4: Communication patterns

Assignment 1

- Today is the last day to submit late
- Questions?

Today

Begin discussion of MPI programming

Generic Send and Receive

Minimum required functionality to do distributed memory parallel computing:

```
send(void *sendbuf, int nelems, int dest)
receive(void *recvbuf, int nelems, int source)
```

Sample Use

```
1 P0 P1
2
3 a = 100; receive(&a, 1, 0)
4 send(&a, 1, 1); printf("%d\n", a);
5 a=0;
```

- Proc 0 sends a single integer to Proc 1
- Proc 0 then 0s that integer
- Proc 1 receives and prints the integer

More typical appearance

Will typically write this as a single program which every processor runs.

```
void exchange(){
  int a = 100:
  int my_proc = get_processor_number();
  if(my_proc == 0){
    send(&a, 1, 1);
    a=0;
  else if(my_proc == 1){
    receive(&a, 1, 0);
    printf("%d\n", a);
```

- Function to identify proc number
- Branching on proc number to take different actions

Flavors Send/Receive

Blocking Operations

Non-Blocking Operations

Buffered

Sending process returns after data has been copied into communication buffer Sending process returns after initiating DMA transfer to buffer. This operation may not be completed on return

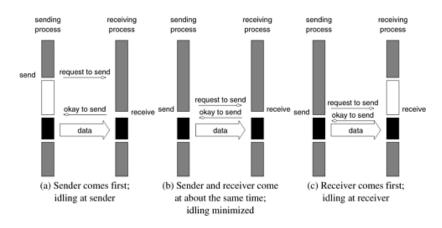
Non-Buffered

Sending process blocks until matching receive operation has been encountered

Send and Receive semantics assured by corresponding operation Programmer must explicitly ensure semantics by polling to verify completion

- Hardware/OS support for buffered communication tends to make things run faster
- Usually have function calls available to do send() (blocking) and send_nonblocking() but must have some hardware support for it

Blocked and Unbuffered



Blocking/Unbuffered: no extra buffer available to hold pending sends/receives so must wait, wait until message is sent to proceed Blocked processors are idle, do no work, which cuts into speedup

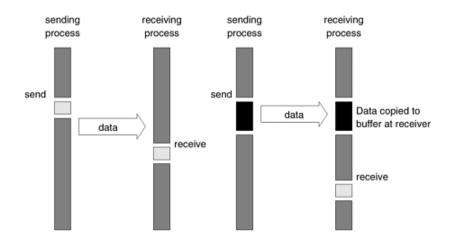
6

Ordering of Send Receive

```
1 P0 P1
2
3 send(&a, 1, 1); send(&a, 1, 0);
4 receive(&b, 1, 1); receive(&b, 1, 0);
```

Assuming send/receive blocked/unbuffered, what's wrong with the above code?

Blocking with Buffers



Hardware buffer support, sender and receiver have a memory minion

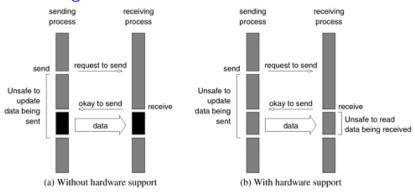
No buffer support: sender interrupts receiver

The Danger Continues

```
1 P0 P1
2
3 receive(&a, 1, 1); receive(&a, 1, 0);
4 send(&b, 1, 1); send(&b, 1, 0);
```

- receive() always blocks until message is obtained
- Does the above code work even in the buffered setting?

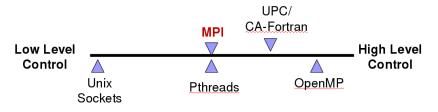
Non-blocking Communication



- ► Takes a bit more work on the programming side
- Must explicitly ensure that transaction completes with function calls
- isend(data,dest,status): send w/o waiting
- ireceive(data,dest,status): receive w/o waiting
- wait(status): wait until a message has been sent or received before moving one

MPI: Message Passing Interface

- ► Standardized library of functions for C/C++/Fortran
- Communicate between processors in a distributed memory machine
- Open source implementations: MPICH, Open MPI
- Proprietary: Intel, Platform, IBM, Platform, Cray
- Typically geared for particular architecture
- May exploit specifics of a particular machine



MPI In a Nutshell: 6 Essential Functions

```
// Initializes MPI.
int MPI_Init(int *argc, char ***argv) ;
// Terminates MPI.
int MPI_Finalize();
// Determines the number of processes.
int MPI_Comm_size(MPI_Comm comm, int *size);
// Determines the label of the calling process.
int MPI Comm rank(MPI Comm comm, int *rank);
// Sends a message.
int MPI Send(void *buf, int count, MPI Datatype datatype,
             int dest, int tag, MPI Comm comm);
// Receives a message.
int MPI Recv(void *buf, int count, MPI Datatype datatype,
             int source, int tag, MPI_Comm comm,
             MPI_Status *status);
```

MPI Setup: Hello World

```
#include <stdio.h>
#include <mpi.h>
int main (int argc, char *argv[]){
  int rank, size;
 MPI_Init (&argc, &argv);
                                         /* starts MPI */
 MPI_Comm_rank (MPI_COMM_WORLD, &rank); /* get current process id */
 MPI_Comm_size (MPI_COMM_WORLD, &size); /* get number of processes */
  int i;
  for(i=0; i<1; i++){
    printf( "Hello world from process %d of %d\n", rank, size );
 MPI_Finalize();
 return 0;
```

- Note the use of MPI_COMM_WORLD which is a predefined constant corresponding to all processors.
- Can also set up other communicators that correspond to subsets of processors

Compilation and Running

- Demo using openmpi implementation
- mpirun for interactive running
- mpirun -np 4
 progr sets number of
 "processors" to 4

```
lila [test-code]% mpicc -o hello hello.c
```

lila [test-code]% ./hello
Hello world from process 0 of 1

lila [test-code]% mpirun hello Hello world from process 0 of 4 Hello world from process 1 of 4 Hello world from process 2 of 4 Hello world from process 3 of 4

lila [test-code]% mpirun -np 2 hello Hello world from process 0 of 2 Hello world from process 1 of 2

lila [test-code]% mpirun -np 8 hello
Hello world from process 7 of 8
Hello world from process 0 of 8
Hello world from process 2 of 8
Hello world from process 3 of 8
Hello world from process 4 of 8
Hello world from process 6 of 8
Hello world from process 1 of 8
Hello world from process 5 of 8

MPI Implementations and OpenMPI Warnings

- Several Implementations of MPI:
 - OpenMPI and MPICH are free, open source, widely available
 - ► HPC Vendors like IBM and Cray provide their own tailored MPI versions
- Recent Versions of OpenMPI can complain a LOT about various items missing
- ► The CSE Labs machines with MPI are not all configured perfectly leading to additional errors
 - Example: --mca btl_base_warn_component_unused 0 to warn about missing HPC network components during mpirun
 - Example: --mca opal_warn_on_missing_libcuda 0 if not intending to use GPU libraries
- ► Exact nature of warnings/errors varies a lot, look at messages which often dictate how to disable them

Warning Suppression in OpenMPI

P 1: Hello world from process 1 of 4 (host: csel-plate02)

D 0. Halla ----ld f---- ------ 0 -f 4 (b--+, ---l -la+-00)

```
csel-plate02 [examples] mpirun -np 4 ./mpi hello
[[8230,1],1]: A high-performance Open MPI point-to-point messaging module
was unable to find any relevant network interfaces:
Module: OpenFabrics (openib)
  Host: csel-plate02
Another transport will be used instead, although this may result in
lower performance.
NOTE: You can disable this warning by setting the MCA parameter
btl base warn component unused to 0.
P 0: Hello world from process 0 of 4 (host: csel-plate02)
Hello from the root processor 0 of 4 (host: csel-plate02)
P 1: Hello world from process 1 of 4 (host: csel-plate02)
P 2: Hello world from process 2 of 4 (host: csel-plate02)
P 3: Hello world from process 3 of 4 (host: csel-plate02)
[csel-plate02:235926] 3 more processes have sent help message help-mpi-btl-base.txt / btl::
[csel-plate02:235926] Set MCA parameter "orte_base_help_aggregate" to 0 to see all help /
csel-plate02 [examples] mpirun --btl base warn component unused 0 -np 4 ./mpi hello
mpirun: Error: unknown option "--btl base warn component unused"
Type 'mpirun --help' for usage.
#
                                VVVVVVVVVVVVVVVVVVVVVVVVVVVVVVVVVV
csel-plate02 [examples]% mpirun --mca btl_base_warn_component_unused 0 -np 4 ./mpi_hello
P 0: Hello world from process 0 of 4 (host: csel-plate02)
Hello from the root processor 0 of 4 (host: csel-plate02)
```

MPI Hostfile

Default OpenMPI config is to use all processors on a single machine then start failing

./mpi_hello_plus

On some systems, like our lab machines, will can use a ${\tt hostfile.txt}$ which gives additional machines to harness - true distributed computation

```
csel-plate02 [examples]% mpirun -np 400 --hostfile hostfile-plate-ip.txt ./mpi_hello_plus P18: Hello world from process 18 of 400 (host: csel-plate02) P21: Hello world from process 21 of 400 (host: csel-plate02) P141:Hello world from process 141 of 400 (host: csel-plate03) P310:Hello world from process 310 of 400 (host: csel-plate04) P149:Hello world from process 149 of 400 (host: csel-plate03) ...
```

MPI Send and Recieve

```
int a[10], b[10];
int partner = 1;
....

// Send contents of a to partner proc with tag=1
MPI_Send(a, 10, MPI_INT, partner, 1, MPI_COMM_WORLD);

// Receive message into b from partner proc with tag=1,
// ignore status of receipt
MPI_Recv(b, 10, MPI_INT, partner, 1, MPI_COMM_WORLD, MPI_STATUS_IGNORE)
Analyze the program send receive test.c
```

Tags Make Messages Unique

```
int a[10], b[10], myrank;
MPI Status status;
MPI_Comm_rank(MPI_COMM_WORLD, &myrank);
if (mvrank == 0) {
  MPI_Send(a, 10, MPI_INT, 1, 1, MPI_COMM_WORLD);
  MPI Send(b, 10, MPI INT, 1, 2, MPI COMM WORLD);
}
else if (myrank == 1) {
  MPI_Recv(b, 10, MPI_INT, 0, 2, MPI_COMM_WORLD);
  MPI_Recv(a, 10, MPI_INT, 0, 1, MPI_COMM_WORLD);
```

- Tags must be honored on receive
- Above code may deadlock if not buffered due to the misordering of tags
- Mostly we will use tag=1 for simplicity

Issues with Untyped Data in MPI

- Type of buffer is always untyped (void* buf)
- ► To try to get at slightly better safety, MPI has standard datatypes

```
MPI_CHAR signed char
MPI_INT signed int
MPI_LONG signed long int
MPI_FLOAT float
MPI_DOUBLE double
MPI_BYTE Last two used for sending
MPI_PACKED structure arrays
```

Unsigned types also available

Exercise: Heat Transfer

- Discuss conversion of the following HW1 code to an MPI version
- ► How is data in H divided up?
- ► Is communication required?
- ▶ How would one arrange MPI_Send / MPI_Recv calls?

```
// Simulate the temperature changes for internal cells
for(t=0; t<max_time-1; t++){
  for(p=1; p<width-1; p++){
    double left_diff = H[t][p] - H[t][p-1];
    double right_diff = H[t][p] - H[t][p+1];
    double delta = -k*( left_diff + right_diff );
    H[t+1][p] = H[t][p] + delta;
}</pre>
```

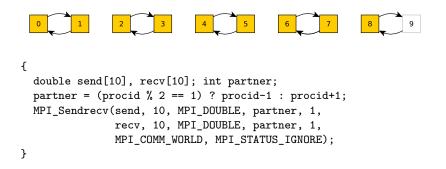
Some Patterns that occur in the problem

- ► Pair exchange of items: made easier with MPI_sendrecv
- ► Collecting final output for display: MPI_Gather

Exchange: Sendrecv for exchanging data between pairs

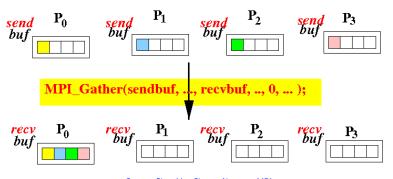
```
double send[10], recv[10]; int partner;
  if(procid % 2 == 1 ){ // odd procs send left, receive left
    partner = procid-1;
    MPI_Send(send, 10, MPI_DOUBLE, partner, 1, MPI_COMM_WORLD);
    MPI_Recv(recv, 10, MPI_DOUBLE, partner, 1, MPI_COMM_WORLD,
            MPI_STATUS_IGNORE);
  else{
                        // even procs receive right, send right
    partner = procid+1;
    MPI_Recv(recv, 10, MPI_DOUBLE, partner, 1, MPI_COMM_WORLD,
             MPI_STATUS_IGNORE);
    MPI_Send(send, 10, MPI_DOUBLE, partner, 1, MPI_COMM_WORLD);
{ // Sendrecv simplifies this pattern
  double send[10], recv[10]; int partner;
  partner = (procid % 2 == 1) ? procid-1 : procid+1;
 MPI_Sendrecv(send, 10, MPI_DOUBLE, partner, 1,
               recv, 10, MPI_DOUBLE, partner, 1,
               MPI COMM WORLD, MPI STATUS IGNORE);
```

Take Care: Pair exchange can hang



- ▶ With 9 processors, logic is broken
- Proc 8 will wait to communicate with a partner that doesn't exist
- Program never terminates

Gather



Source: Shun Yan Cheung Notes on MPI

- Every processor has computed columns
- One processor (usually procid 0) needs to gather all of the data
- Everyone calls MPI_Gather()

MPI_Gather Sample

Use of Gather

```
// Preamble for any code
MPI_Comm comm = MPI_COMM_WORLD;
int sendarray[100];
int procid, total_procs, *rbuf;
. . . ;
// Only proc 0 needs space for all
// data
if(procid == 0) {
  rbuf = malloc(total_procs*100*
                sizeof(int));
// Everyone calls gather
// proc 0 gets all data eventually
MPI_Gather(sendarray, 100, MPI_INT,
           rbuf, 100, MPI_INT,
           0, comm);
```

Equivalent Non-Gather Code

```
if(rank == 0){
  for(i=0; i<100; i++){
    rbuf[i] = sendarrav[i];
  for(i=1; i<total_procs; i++){</pre>
    int *rloc = &rbuf[i*100];
    MPI_Recv(rloc, 100,
             MPI INT, i,
             tag, MPI_COMM_WORLD,
             MPI STATUS IGNORE);
else{
  MPI_Send(sendarray, 100,
           MPI_INT, 0,
           tag, MPI_COMM_WORLD);
```