## CSCI 2021: Binary Floating Point Numbers

Chris Kauffman

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### Logistics

### Reading Bryant/O'Hallaron

- ► Ch 2.4-5: Floats, Wed/Fri
- ► 2021 Quick Guide to GDB
- Next week: Ch 3.1-7: Assembly Intro

#### Goals this Week

- Discuss Bitwise ops from Integer Rep Slides
- ► Floating Point layout
- gdb introduction

#### P2: Released Wed, due 8 days

- 1. Bit shift operations for Battery (50%)
- 2. Puzzlebox via debugger (50% + Makeup Credit)

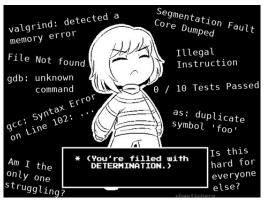
#### Feedback Survey

- Open on Canvas
- Anonymous: be honest!
- Due Wed 5-Oct for 1 EP 83% response rate so far

#### Labs/HW

- ► Lab05: Bit operations
- ► HW05: Bits, Floats, GDB

## Don't Give Up, Stay Determined!



- ▶ If Project 1 / Exam 1 went awesome, count yourself lucky
- If things did not go well, Don't Give Up
- Spend some time contemplating why things didn't go well, talk to course staff about it, learn from any mistakes
- There is a LOT of semester left and plenty of time to recover from a bad start

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#### Parts of a Fractional Number

The meaning of the "decimal point" is as follows:

$$123.406_{10} = 1 \times 10^{2} + 2 \times 10^{1} + 3 \times 10^{0} + 123 = 100 + 20 + 3$$
$$4 \times 10^{-1} + 0 \times 10^{-2} + 6 \times 10^{-3} \quad 0.406 = \frac{4}{10} + \frac{6}{1000}$$
$$= 123.406_{10}$$

Changing to base 2 induces a "binary point" with similar meaning:

$$110.101_2 = 1 \times 2^2 + 1 \times 2^1 + 0 \times 2^0 + 6 = 4 + 2$$
$$1 \times 2^{-1} + 0 \times 2^{-2} + 1 \times 2^{-3} \qquad 0.625 = \frac{1}{2} + \frac{1}{8}$$
$$= 6.625_{10}$$

One *could* represent fractional numbers with a **fixed point** e.g.

- ▶ 32 bit fractional number with
- ▶ 10 bits left of Binary Point (integer part)
- ▶ 22 bits right of Binary Point (fractional part)

**BUT** most applications require a more flexible scheme

#### Scientific Notation for Numbers

"Scientific" or "Engineering" notation for numbers with a fractional part is

Standard	Scientific	printf("%.4e",x);
123.456	$1.23456 \times 10^{2}$	1.2346e+02
50.01	$5.001 \times 10^{1}$	5.0010e+01
3.14159	$3.14159 \times 10^{0}$	3.1416e+00
0.54321	$5.4321 \times 10^{-1}$	5.4321e-01
0.00789	$7.89 \times 10^{-3}$	7.8900e-03

- Always includes one non-zero digit left of decimal place
- Has some significant digits after the decimal place
- Multiplies by a power of 10 to get actual number

### Binary Floating Point Layout Uses Scientific Convention

- Some bits for integer/fractional part
- Some bits for exponent part
- ► All in base 2: 1's and 0's, powers of 2

### Conversion Example

Below steps convert a decimal number to a fractional binary number equivalent then adjusts to scientific representation.

```
float fl = -248.75;

7 6 5 4 3 2 1 0 -1 -2

-248.75 = -(128+64+32+16+8+0+0+0).(1/2+1/4)

= -11111000.11 *2^0

76543210 12

= -1111100.011 *2^1

6543210 123

= -111110.0011 *2^2

543210 1234
```

. . .

MANTISSA EXPONENT = -1.111100011 \* 2^7 0 123456789

 $Mantissa \equiv Significand \equiv Fractional Part$ 

## Principle and Practice of Binary Floating Point Numbers

- ▶ In early computing, computer manufacturers used similar principles for floating point numbers but varied specifics
- ► Example of Early float data/hardware
  - ▶ Univac: 36 bits, 1-bit sign, 8-bit exponent, 27-bit significand¹
  - ▶ IBM: 32 bits, 1-bit sign, 7-bit exponent, 24-bit significand<sup>2</sup>
- Manufacturers implemented circuits with different rounding behavior, with/without infinity, and other inconsistencies
- Troublesome for reliability: code produced different results on different machines
- ► This was resolved with the adoption of the IEEE 754 Floating Point Standard which specifies
  - ▶ Bit layout of 32-bit float and 64-bit double
  - Rounding behavior, special values like Infinity
- ► Turing Award to William Kahan for his work on the standard

<sup>&</sup>lt;sup>1</sup>Floating Point Arithmetic

<sup>&</sup>lt;sup>2</sup>IBM Hexadecimal Floats

### IEEE 754 Format: The Standard for Floating Point

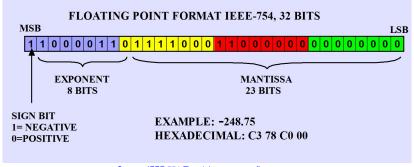
float	double	Property
32	64	Total bits
1	1	Bits for sign (1 neg / 0 pos)
8	11	Bits for Exponent multiplier (power of 2)
23	52	Bits for Fractional part or mantissa
7.22	15.95	Decimal digits of accuracy <sup>3</sup>

- Most commonly implemented format for floating point numbers in hardware to do arithmetic: processor has physical circuits to add/mult/etc. for this bit layout of floats
- Numbers/Bit Patterns divided into three categories

Category	Description	Exponent
Normalized	most common like 1.0 and -9.56e37	mixed $0/1$
Denormalized	very close to zero and 0.0	all 0's
Special	${\sf extreme/error\ values\ like\ Inf\ and\ NaN}$	all 1's

<sup>&</sup>lt;sup>3</sup>Wikipedia: IEEE 754

## Example float Layout of -248.75: float\_examples.c



Source: IEEE-754 Tutorial, www.puntoflotante.net

```
Color: 8-bit blocks, Negative: highest bit, leading 1
```

```
Exponent: high 8 bits, 2<sup>7</sup> encoded with bias of -127

1.111100011...

1000_0110 - 0111_1111

= 128+4+2 - 127

= 134 - 127

= 7

implied leading 1
not in binary layout
```

### Normalized Floating Point: General Case

- ➤ A "normalized" floating point number is in the standard range for float/double, bit layout follows previous slide
- ► Example: -248.75 = -1.111100011 \* 2^7

### Exponent is in **Bias Form** (not Two's Complement)

- Unsigned positive integer minus constant bias number
- ► **Consequence**: exponent of 0 is not bitstring of 0's
- ➤ Consequence: tiny exponents like -125 close to bitstring of 0's; this makes resulting number close to 0
- ▶ 8-bit exponent 1000 0110 = 128+4+2 = 134 so exponent value is 134 127 = 7

### Integer and Mantissa Parts

- ► The leading 1 before the binary point is **implied** so does not show up in the bit string
- ► Remaining fractional/mantissa portion shows up in the low-order bits

## Fixed Bit Standards for Floating Point

#### **IEEE Standard Layouts**

Kind	Sign	Exponent			Mantissa
	Bit	Bits	Bias	Exp Range	Bits
float	31 (1)	30-23 (8 bits)	-127	-126 to +127	22-0 (23 bits)
double	63 (1)	62-52 (11 bits)	-1023	-1022 to $+1023$	51-0 (52 bits)

Standard allows hardware to be created that is as efficient as possible to do calculation on these numbers

#### Consequences of Fixed Bits

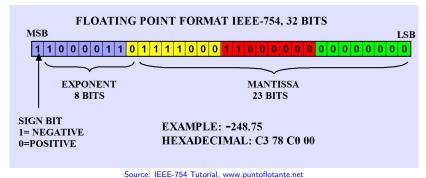
- ➤ Since a fixed # of bit is used, **some numbers cannot be exactly represented**, happens in any numbering system:
- ▶ Base 10 and Base 2 cannot represent  $\frac{1}{3}$  in finite digits
- ▶ Base 2 cannot represent  $\frac{1}{10}$  in finite digits float f = 0.1;

```
float f = 0.1;
printf("0.1 = %.20e\n",f);
0.1 = 1.00000001490116119385e-01
```

Try show\_float.c to see this in action

### Exercise: Quick Checks

- 1. What distinct parts are represented by bits in a floating point number (according to IEEE)
- 2. What is the "bias" of the exponent for 32-bit floats
- 3. Represent 7.125 in binary using "binary point" notation
- 4. Lay out 7.125 in IEEE-754 format
- 5. What does the number 1.0 look like as a float?



The diagram above may help in recalling IEEE 754 layout

## Special Cases: See float\_examples.c

#### Denormalized values: Exponent bits all 0

- ► Fractional/Mantissa portion evaluates *without* implied leading one, still an unsigned integer though
- **Exponent** is Bias + 1:  $2^{-126}$  for float
- ► Result: very small numbers close to zero, smaller than any other representation, degrade uniformly to 0
- Zero: bit string of all 0s, optional leading 1 (negative zero);

#### Special Values

- ▶ **Infinity**: exponent bits all 1, fraction all 0, sign bit indicates  $+\infty$  or  $-\infty$
- Infinity results from overflow/underflow or certain ops like float x = 1.0 / 0.0;
- #include <math.h> gets macro INFINITY and -INFINITY
- ▶ NaN: not a number, exponent bits all 1, fraction has some 1s
- ► Errors in floating point like 0.0 / 0.0

#### Other Float Notes



DURING A COMPETITION, I TOLD THE PROGRAMMERS ON OUR TEAM THAT  $\sigma^{\sigma}-\pi$  was a standard test of Flating-Point Handlers -- It would come out to 20 unless they had rounding errors.





Source: XKCD #217

### Approximations and Roundings

- Approximate  $\frac{2}{3}$  with 4 digits, usually 0.6667 with standard rounding in base 10
- Similarly, some numbers cannot be exactly represented with fixed number of bits: <sup>1</sup>/<sub>10</sub> approximated
- ► IEEE 754 specifies various rounding modes to approximate numbers

### Clever Engineering

- IEEE 754 allows floating point numbers to sort using signed integer sorting routines
- Bit patterns for float follows are ordered nearly the same as bit patterns for signed int
- Integer comparisons are usually fewer clock cycles than floating comparisons

#### Sidebar: The Weird and Wonderful Union

- Bitwise operations like & are not valid for float/double
- ► Can use pointers/casting to get around this OR...
- Use a union: somewhat unique construct to C
- Defined like a struct with several fields
- ▶ BUT fields occupy the same memory location (!?!)
- ► Allows one to treat a byte position as multiple different types, ex: int / float / char[]
- Memory size of the union is the max of its fields

```
// union.c
typedef union { // shared memory
               // an float
 float fl;
 int in; // a int
 char ch[4]; // char array
} flint t;
               // 4 bytes total
int main(){
 flint t flint;
 flint.in = 0xC378C000;
 printf("%.4f\n", flint.fl);
 printf("%08x %d\n",flint.in,flint.in);
 for(int i=0; i<4; i++){
    unsigned char c = flint.ch[i];
   printf("%d: %02x '%c'\n",i,c,c);
```

## Floating Point Operation Efficiencies

- ► Floating Point Operations per Second, **FLOPS** is a major measure for numerical code/hardware efficiency
- Often used to benchmark and evaluate scientific computer resources, (e.g. top super computers in the world)
- ► Tricky to evaluate because of
  - ► A single FLOP (add/sub/mul/div) may take 3 clock cycles to finish: latency 3
  - Another FLOP can start before the first one finishes: pipelined
  - Enough FLOPs lined up can get average 1 FLOP per cycle
  - ► FP Instructions may automatically operate on multiple FPs stored in memory to feed pipeline: **vectorized ops**
  - Generally referred to as superscalar
  - Processors schedule things out of order too
- All of this makes micro-evaluation error-prone and pointless
- Run a real application like an N-body simulation and compute

$$FLOPS = \frac{number of floating ops done}{time taken in seconds}$$

Top 5 Super Computers Worldwide, June 2022

				- I	
Rank	System	#Cores	Rmax (TFlop/s)	Rpeak (TFlop/s)	Power (kW)
1	Frontier, <i>USA / Oak Ridge</i> Cray EX235a, AMD EPYC 2GHz (x86-64)	8,730,112	1,102.00	1,685.65	21,100
2	Fugaku, <i>Japan / Fujitsu</i> Fujitsu A64FX 2.2GHz (Arm)	7,630,848	442,010.0	537,212.0	29,899
3	LUMI <i>Finland / EuroHPC</i> Cray EX235a, AMD EPYC 2GHz (x86-64)	1,110,144	151.90	214.35	2,942
4	Summit <i>United States</i> IBM POWER9 22C 3.07GHz (Power)	2,414,592	148,600.0	200,794.9	10,096
5	Sierra <i>United States</i> IBM POWER9 22C 3.1GHz (Power)	1,572,480	94,640.0	125,712.0	7,438

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3	Sierra <i>United States</i> IBM POWER9 22C 3.1GHz (Power)	1,572,480	94,640.0	125,712.0	7,438
4	Sunway TaihuLight <i>China</i> Sunway SW26010 (custom RISC)	10,649,600	93,014.6	125,435.9	15,371
5	Perlmutter, <i>United States</i> AMD EPYC 2.45GHz, Cray (x86-64)	706,304	64,590.0	89,794.5	2,528

Rank	System	#Cores	Rmax (TFlop/s)	Rpeak (TFlop/s)	Power (kW)
1	Fugaku, <i>Japan / Fujitsu</i> Fujitsu A64FX 2.2GhZ (Arm)	7,299,072	415,530.0	513,854.7	28,335
2	Summit <i>United States</i> IBM POWER9 22C 3.07GHz (Power)	2,397,824	143,500.0	200,794.9	10,096
3	Sierra <i>United States</i> IBM POWER9 22C 3.1GHz (Power)	1,572,480	94,640.0	125,712.0	7,438
4	Sunway TaihuLight <i>China</i> Sunway SW26010 (custom RISC)	10,649,600	93,014.6	125,435.9	15,371
5	Selene <i>USA</i> , <i>NVIDIA/AMD</i> AMD EPYC 7742 64C 2.25GHz (x86-64)	555,520	63,460.0	79,215.0	2,646

Rank	System	#Cores	Rmax (TFlop/s)	Rpeak (TFlop/s)	Power (kW)
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2	Summit <i>United States</i> IBM POWER9 22C 3.07GHz (Power)	2,397,824	143,500.0	200,794.9	10,096
3	Sierra <i>United States</i> IBM POWER9 22C 3.1GHz (Power)	1,572,480	94,640.0	125,712.0	7,438
4	Sunway TaihuLight <i>China</i> Sunway SW26010 (custom RISC)	10,649,600	93,014.6	125,435.9	15,371
5	Tianhe-2A <i>China</i> Intel Xeon 2.2GHz (x86-64)	4,981,760	61,444.5	100,678.7	18,482

	_		Rmax	Rpeak	Power
Rank	System	#Cores	(TFlop/s)	(TFlop/s)	(kW)
1	Summit <i>United States</i> IBM POWER9 22C 3.07GHz	2,397,824	143,500.0	200,794.9	9,783
2	Sierra <i>United States</i> IBM POWER9 22C 3.1GHz,	1,572,480	94,640.0	125,712.0	7,438
3	Sunway TaihuLight <i>China</i> Sunway MPP	10,649,600	93,014.6	125,435.9	15,371
4	Tianhe-2A <i>China</i> Xeon 2.2GHz	4,981,760	61,444.5	100,678.7	18,482
5	Frontera, <i>United States</i> Dell 6420, Xeons 2.7GHz	448,448	23,516.4	38,745.9	??

https://www.top500.org/list/2019/11/

Rank	System	#Cores	Rmax (TFlop/s)	Rpeak (TFlop/s)	Power (kW)
1	Summit <i>United States</i> IBM POWER9 22C 3.07GHz	2,397,824	143,500.0	200,794.9	9,783
2	Sierra <i>United States</i> IBM POWER9 22C 3.1GHz,	1,572,480	94,640.0	125,712.0	7,438
3	Sunway TaihuLight <i>China</i> Sunway MPP	10,649,600	93,014.6	125,435.9	15,371
4	Tianhe-2A <i>China</i> TH-IVB-FEP Cluster	4,981,760	61,444.5	100,678.7	18,482
5	Piz Daint <i>Switzerland</i> Cray XC50, Xeon E5-2690v3	387,872	21,230.0	27,154.3	2,384

https://www.top500.org/list/2018/11/

Rank	System	#Cores	Rmax (TFlop/s)	Rpeak (TFlop/s)	Power (kW)
1	Sunway TaihuLight <i>China</i> Sunway MPP	10,649,600	93,014.6	125,435.9	15,371
2	Tianhe-2 (MilkyWay-2) <i>China</i> TH-IVB-FEP Cluster	3,120,000	33,862.7	54,902.4	17,808
3	Piz Daint <i>Switzerland</i> Cray XC50	361,760	19,590.0	25,326.3	2,272
4	Gyoukou <i>Japan</i> ZettaScaler-2.2 HPC system	19,860,000	19,135.8	28,192.0	1,350
5	Titan <i>USA</i> Cray XK7	560,640	17,590.0	27,112.5	8,209

https://www.top500.org/lists/2017/11/