

CMSC216: Assembly Basics and x86-64

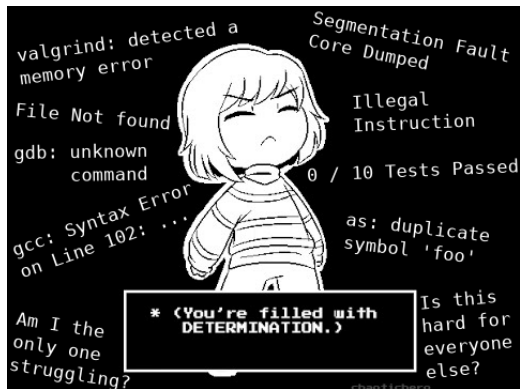
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Logistics

Announcements

Don't Give Up, Stay Determined!



- ▶ If Project 1 / Exam 1 went awesome, count yourself lucky
- ▶ If things did not go well, **Don't Give Up**
- ▶ Spend some time contemplating **why** things didn't go well, talk to course staff about it, learn from mistakes
- ▶ There is a LOT of semester left and time to recover from a bad start

Wrapping Up Integer Affairs

Pre-exam one slides contain info on byte-ordering (little / big endian) and bit-wise operations to be discussed here prior to assembly.

The **Many** Assembly Languages

- ▶ Most **microprocessors** are created to understand a **binary machine language**
- ▶ Machine Language provides means to manipulate internal memory, perform arithmetic, etc.
- ▶ The Machine Language of one processor is **not understood** by other processors

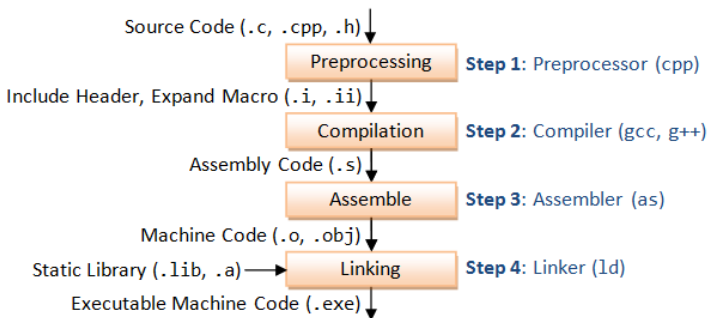
MOS Technology 6502

- ▶ 8-bit operations, limited addressable memory, **1 general purpose register**, powered notable gaming systems in the 1980s
- ▶ Apple IIe, Atari 2600, Commodore
- ▶ Nintendo Entertainment System / Famicom

IBM Cell Microprocessor

- ▶ Developed in early 2000s, 64-bit, many cores (execution elements), many registers (32 on the PPE), large addressable space, fast multimedia performance, is a **pain** to program
- ▶ Playstation 3 and Blue Gene Supercomputer

Assemblers and Compilers



- ▶ **Compiler**: chain of tools that translate high level languages to lower ones, may perform optimizations
- ▶ **Assembler**: translates text description of the machine code to binary, formats for execution by processor, late compiler stage
- ▶ **Consequence**: The compiler can **generate assembly code**
- ▶ Generated assembly is a pain to read but is often quite fast
- ▶ **Consequence**: A compiler on an Intel chip can generate assembly code for a different processor, **cross compiling**

Our focus: The x86-64 Assembly Language

- ▶ x86-64 Targets Intel/AMD chips with 64-bit word size
Reminder: 64-bit “word size” \approx size of pointers/addresses
- ▶ Lineage of x86 family
 - ▶ 1970s: 16-bit systems like Intel 8086
 - ▶ 1990s: IA32 (Intel 32-bit systems like 80386 and 80486)
 - ▶ 2000s: x86-64 (64-bit extension by AMD)
- ▶ x86-64 is backwards compatibility, consequently much cruft
 - ▶ Can run compiled code from the 70's / 80's on modern processors without much trouble BUT means 50-year-old instructions must be preserved
 - ▶ x86-64 is not the assembly language you would design from scratch today, it's the assembly you have to code against
 - ▶ RISC-V is a new assembly language that is “clean” as it has no history to support (and few CPUs run it)
- ▶ **Warning:** Lots of information available on the web for Intel assembly programming **BUT** some of it is dated, IA32 info which may not work on 64-bit systems

x86-64 Assembly Language Syntax(es)

- ▶ Different assemblers understand different syntaxes for the same assembly language
- ▶ GCC use the GNU Assembler (GAS, command 'as file.s')
- ▶ GAS and Textbook favor AT&T syntax so **we will too**
- ▶ NASM assembler favors Intel, may see this online

AT&T Syntax (Our Focus)

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- ▶ Use of % to indicate registers
- ▶ Use of q/l/w/b to indicate 64 / 32 / 16 / 8-bit operands

Intel Syntax

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- ▶ Register names are bare
- ▶ Use of QWORD etc. to indicate operand size

Generating Assembly from C Code

- ▶ `gcc -S file.c` will stop compilation at assembly generation
- ▶ Leaves assembly code in `file.s`
 - ▶ `file.s` and `file.S` conventionally assembly code though sometimes `file.asm` is used
- ▶ By default, compiler generates code that is often difficult for humans to interpret, may include re-arrangements, “conservative” compatibility assembly, etc. increasing size of assembly considerably
- ▶ `gcc -Og file.c`: optimize for debugging, generally makes it easier to read generated assembly, aligns somewhat more closely to C code

Example of Generating Assembly from C

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```
gcc -Og -S mstore.c
```

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Every Programming Language

Look for the following as it should almost always be there

- ☐ Comments
- ☐ Statements/Expressions
- ☐ Variable Types
- ☐ Assignment
- ☐ Basic Input/Output
- ☐ Function Declarations
- ☐ Conditionals (if-else)
- ☐ Iteration (loops)
- ☐ Aggregate data (arrays, structs, objects, etc)
- ☐ Library System

Exercise: Examine col_simple_asm.s

Take a simple sample problem to demonstrate assembly:

Computes Collatz Sequence starting at $n=10$:

*if n is ODD $n=n*3+1$; else $n=n/2$.*

*Return the number of steps to converge to 1 as the **return code** from `main()`*

The following codes solve this problem

| Code | Notes |
|------------------|---|
| col_simple_asm.s | Hand-coded assembly for obvious algorithm Straight-forward reading |
| col_unsigned.c | Unsigned C version Generated assembly is reasonably readable |
| col_signed.c | Signed C version Generated assembly is ... interesting |

- ▶ Kauffman will Compile/Run code
- ▶ Students should study the code and predict what lines do
- ▶ Illustrate tricks associated with gdb and assembly

Exercise: col_simple_asm.s

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Answers: x86-64 Assembly Basics for AT&T Syntax

- ▶ *Comments* are one-liners starting with #
- ▶ *Statements*: each line does ONE thing, frequently text representation of an assembly instruction
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- ▶ Assembler directives and labels are also possible:
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- ▶ *Variables*: mainly **registers**, also memory ref'd by registers maybe some named global locations
- ▶ *Assignment*: instructions like movX that put bits into registers and memory
- ▶ *Conditionals/Iteration*: assembly instructions that jump to code locations
- ▶ *Functions*: code locations that are **labeled** and global
- ▶ *Aggregate data*: none, use the stack/multiple registers
- ▶ *Library System*: link to other code

So what *are* these Registers?

- ▶ Memory locations directly wired to the CPU
- ▶ Usually very fast to access, faster than **main memory**
- ▶ Most instructions involve registers, access or change reg val

Example: Adding Together Integers

- ▶ Ensure registers have desired values in them
- ▶ Issue an addX instruction involving the two registers
- ▶ Result will be stored in a register
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- ▶ Note instruction and register names indicate whether 32-bit int or 64-bit long are being added

x86-64 “General Purpose” Registers

Many “general purpose” registers have special purposes and conventions associated such as

- ▶ Return Value:
%rax / %eax / %ax
- ▶ Function Args 1 to 6:
%rdi, %rsi, %rdx,
%rcx, %r8, %r9
- ▶ Stack Pointer (top of stack): %rsp
- ▶ Old Code Base Pointer:
%rbp, **historically start of current stack frame** but is not used that way in modern codes

Note: There are also Special Registers like %rip and %eflags which we will discuss later.

| 64-bit | 32-bit | 16-bit | 8-bit | Notes |
|---------------------|--------|----------------------------|-------|------------|
| %rax | %eax | %ax | %al | Return Val |
| %rbx | %ebx | %bx | %bl | |
| %rcx | %ecx | %cx | %cl | Arg 4 |
| %rdx | %edx | %dx | %dl | Arg 3 |
| %rsi | %esi | %si | %sil | Arg 2 |
| %rdi | %edi | %di | %dil | Arg 1 |
| %rsp | %esp | %sp | %spl | Stack Ptr |
| %rbp | %ebp | %bp | %bpl | Base Ptr? |
| %r8 | %r8d | %r8w | %r8b | Arg 5 |
| %r9 | %r9d | %r9w | %r9b | Arg 6 |
| %r10 | %r10d | %r10w | %r10b | |
| %r11 | %r11d | %r11w | %r11b | |
| %r12 | %r12d | %r12w | %r12b | |
| %r13 | %r13d | %r13w | %r13b | |
| %r14 | %r14d | %r14w | %r14b | |
| %r15 | %r15d | %r15w | %r15b | |
| Caller Save: | | Restore after calling func | | |
| Callee Save: | | Restore before returning | | |

Register Naming Conventions

- ▶ AT&T syntax identifies registers with prefix %
- ▶ Naming convention is a historical artifact
- ▶ Originally 16-bit architectures in x86 had
 - ▶ General registers `ax,bx,cx,dx`,
 - ▶ Special Registers `si,di,sp,bp`
- ▶ *Extended* to 32-bit: `eax,ebx,...,esi,edi,...`
- ▶ Grew again to 64-bit: `rax,rbx,...,rsi,rdi,...`
- ▶ Added Eight 64-bit regs `r8,r9,...,r14,r15` with 32-bit portion `r8d,r9d,...`, 16-bit `r8w,r9w...`, etc.
- ▶ Instructions must match registers sizes:
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- ▶ When hand-coding assembly, easy to mess this up, assembler will error out

Hello World in x86-64 Assembly : Not that Easy

- ▶ Non-trivial in assembly because **output is involved**
 - ▶ Try writing `helloworld.c` without `printf()`
- ▶ Output is the business of the **operating system**, always a request to the almighty OS to put something somewhere
 - ▶ **Library call**: `printf("hello");` mangles some bits but eventually results with a ...
 - ▶ **System call**: Unix system call directly implemented in the OS **kernel**, puts bytes into files / onto screen as in
`write(1, buf, 5); // file 1 is screen output`

This gives us several options for hello world in assembly:

1. `hello_printf64.s`: via calling `printf()` which means the C standard library must be (painfully) linked
2. `hello64.s` via direct system `write()` call which means no external libraries are needed: OS knows how to write to files/screen. Use the 64-bit Linux calling convention.
3. `hello32.s` via direct system call using the older 32 bit Linux calling convention which “traps” to the operating system.

(Optional): The OS Privilege: System Calls

- ▶ Most interactions with the outside world happen via Operating System Calls (or just “system calls”)
- ▶ User programs indicate what service they want performed by the OS via making system calls
- ▶ System Calls differ for each language/OS combination
 - ▶ x86-64 Linux: set `%rax` to system call number, set other args in registers, issue `syscall`
 - ▶ IA32 Linux: set `%eax` to system call number, set other args in registers, issue an **interrupt**
 - ▶ C Code on Unix: make system calls via `write()`, `read()` and others (studied in CSCI 4061)
 - ▶ Tables of Linux System Call Numbers
 - ▶ 64-bit (335 calls)
 - ▶ 32-bit (190 calls)
 - ▶ Mac OS X: very similar to the above (it's a Unix)
 - ▶ Windows: use OS wrapper functions
- ▶ OS executes **privileged** code that can manipulate any part of memory, touch internal data structures corresponding to files, do other fun stuff discussed in OS courses

Basic Instruction Classes

- ▶ **Remember:** Goal is to understand assembly as a *target* for higher languages, not become expert “assemblists”
- ▶ Means we won't hit all 4,834 pages of the [Intel x86-64 Manual](#)
- ▶ [Brown University's x64 Cheat Sheet](#) has a good overview
- ▶ [x86 Assembly Guide from Yale](#) is also good but is limited to 32-bit coverage

| Kind | Assembly Instructions |
|---------------------------|-------------------------|
| <i>Fundamentals</i> | |
| - Memory Movement | mov |
| - Stack manipulation | push,pop |
| - Addressing modes | (%eax),12(%eax,%ebx)... |
| <i>Arithmetic/Logic</i> | |
| - Arithmetic | add,sub,mul,div,lea |
| - Bitwise Logical | and,or,xor,not |
| - Bitwise Shifts | sal,sar,shr |
| <i>Control Flow</i> | |
| - Compare / Test | cmp,test |
| - Set on result | set |
| - Jumps (Un)Conditional | jmp,je,jne,jl,jg,... |
| - Conditional Movement | cmovc,cmovg,... |
| <i>Procedure Calls</i> | |
| - Stack manipulation | push,pop |
| - Call/Return | call,ret |
| - System Calls | syscall |
| <i>Floating Point Ops</i> | |
| - FP Reg Movement | vmov |
| - Conversions | vcvts |
| - Arithmetic | vadd,vsub,vmul,vdiv |
| - Extras | vmins,vmaxs,sqrts |

Data Movement: `movX` instruction

`movX SOURCE, DEST` # move/copy source value to dest

Overview

- ▶ Moves data...
 - ▶ Reg to Reg
 - ▶ Mem to Reg
 - ▶ Reg to Mem
 - ▶ Imm to ...
- ▶ Reg: register
- ▶ Mem: main memory
- ▶ Imm: “immediate” value (constant) specified like
 - ▶ `$21` : decimal
 - ▶ `$0x2f9a` : hexadecimal
 - ▶ **NOT** `1234` (mem adder)
- ▶ More info on operands next

Examples

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Note variations

- ▶ `movq` for 64-bit (8-byte)
- ▶ `movl` for 32-bit (4-byte)
- ▶ `movw` for 16-bit (2-byte)
- ▶ `movb` for 8-bit (1-byte)

Operands and Addressing Modes

In many instructions like `movX`, operands can have a variety of forms called **addressing modes**, may include constants and memory addresses

| Style | Address Mode | C-like | Notes |
|--------------------------------------|-----------------------------------|---------------------------------------|--|
| \$21 \$0xD2 | immediate | 21 | value of constant like 21 or 0xD2 = 210 |
| %rax (%rax) 8(%rax) 4(%rdx) | register indirect displaced | rax *rax *(rax+2) rdx->field | to/from register contents reg holds memory address, deref base plus constant offset, often used for struct field derefs |
| (%rax,%rbx) | indexed | *(rax+rbx) char_arr[rbx] | base plus offset in given reg actual value of rbx is used, NOT multiplied by sizeof() |
| (%rax,%rbx,4) (%rax,%rbx,8) | scaled index | rax[rbx] rax[rbx] | like array access with sizeof(..)=4 "" with sizeof(..)=8 |
| 1024 | absolute | ... | Absolute address #1024 Rarely used |

Exercise: Show movX Instruction Execution

Code `movX_exercise.s`

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Registers/Memory

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Lookup...

May need to look up addressing conventions for things like...

```
movX %y,%x    # reg y to reg x
movX $5,(%x)  # 5 to address in %x
```

Answers Part 1/2: movX Instruction Execution

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#WARNING!: On 64-bit systems, ALWAYS use a 64-bit reg name like `%rdx` and `movq` to copy memory addresses; using smaller name like `%edx` will miss half the memory addressing leading to major memory problems

Answers Part 2/2: movX Instruction Execution

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`gdb` Assembly: Examining Memory

`gdb` commands `print` and `x` allow one to print/examine memory of interest. Try on `movX_exercises.s`

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Many of these tricks are needed to debug assembly.

Register Size and Data Movement

- ▶ `%rax` is 64-bit register, `%eax` is its lower 32 bits
- ▶ Data movement involving small registers **may NOT overwrite** higher bits in extended register
- ▶ Moving data to low 32-bit regs automatically zeros high 32-bits
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- ▶ Moving data to other small regs DOES NOT ALTER high bits
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- ▶ Gives rise to two other families of movement instructions for moving little registers (X) to big (Y) registers, see `movz_examples.s`
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Exercise: movX differences in Main Memory

| Instr | # bytes |
|-------|---------|
| movb | 1 byte |
| movw | 2 bytes |
| movl | 4 bytes |
| movq | 8 bytes |

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Show the result of each of the following copies to main memory in sequence.

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Answers: movX to Main Memory 1/2

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Answers: movX to Main Memory 2/2

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addX : A Quintessential ALU Instruction

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- ▶ Addition represents most 2-operand ALU instructions well
- ▶ Second operand A is modified by first operand B, No change to B
- ▶ Variety of register, memory, constant combinations honored
- ▶ addX has variants for each register size: addq, addl, addw, addb

Optional Exercise: Addition

Show the results of the following addX/movX ops at each of the specified positions

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Answers: Addition

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The Other ALU Instructions

- ▶ Most ALU instructions follow the same pattern as `addX`: two operands, second gets changed.
- ▶ Some one operand instructions as well.

| Instruction | Name | Effect | Notes |
|-------------------------|------------|------------------|---------------------------------------|
| <code>addX B, A</code> | Add | $A = A + B$ | Two Operand Instructions |
| <code>subX B, A</code> | Subtract | $A = A - B$ | |
| <code>imulX B, A</code> | Multiply | $A = A * B$ | <i>Has a limited 3-arg variant</i> |
| <code>andX B, A</code> | And | $A = A \& B$ | |
| <code>orX B, A</code> | Or | $A = A B$ | |
| <code>xorX B, A</code> | Xor | $A = A \wedge B$ | |
| <code>salX B, A</code> | Shift Left | $A = A \ll B$ | B is constant or <code>%cl</code> reg |
| <code>shlX B, A</code> | | $A = A \ll B$ | |
| <code>sarX B, A</code> | | $A = A \gg B$ | Arithmetic: Sign carry |
| <code>shrX B, A</code> | | $A = A \gg B$ | Logical: Zero carry |
| <code>incX A</code> | Increment | $A = A + 1$ | One Operand Instructions |
| <code>decX A</code> | Decrement | $A = A - 1$ | |
| <code>negX A</code> | Negate | $A = -A$ | |
| <code>notX A</code> | Complement | $A = \sim A$ | |

leax: Load Effective Address

- ▶ Memory addresses must often be loaded into registers
- ▶ Often done with a leax, usually leaq in 64-bit platforms
- ▶ Sort of like “address-of” op & in C but a bit more general

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Compiler sometimes uses leax for multiplication as it is usually faster than imulx but less readable.

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Division: It's a Pain (1/2)

- ▶ `idivX` operation has some special rules
- ▶ Dividend must be in the `rax` / `eax` / `ax` register
- ▶ Sign extend to `rdx` / `edx` / `dx` register with `cqto`
- ▶ `idivX` takes one **register** argument which is the divisor
- ▶ At completion
 - ▶ `rax` / `eax` / `ax` holds quotient (integer part)
 - ▶ `rdx` / `edx` / `dx` holds the remainder (leftover)

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Compiler avoids division whenever possible: compile `col_unsigned.c` and `col_signed.c` to see some tricks.

Division: It's a Pain (2/2)

- ▶ When performing division on 8-bit or 16-bit quantities, use instructions to sign extend small reg to all rax register

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