CSCI 2021: Binary Floating Point Numbers

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Logistics

Reading Bryant/O'Hallaron

- ► Ch 2.4-5: Floats, Wed/Fri
- ► 2021 Quick Guide to GDB
- Next week: Ch 3.1-7: Assembly Intro

Goals this Week

- Discuss Bitwise ops (Integer Slides)
- gdb introduction
- Floating Point layout

Feedback Survey

- Open on Canvas
- Anonymous: be honest!
- ▶ Due Wed 2/17 for 1 EP

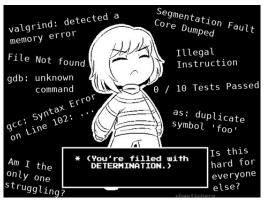
Labs/HW

- ► Lab05: Bit operations
- HW05: Bits, Floats, GDB

Project 2

Release Thursday

Don't Give Up, Stay Determined!



- ▶ If Project 1 / Exam 1 went awesome, count yourself lucky
- If things did not go well, Don't Give Up
- Spend some time contemplating why things didn't go well, talk to course staff about it, learn from any mistakes
- There is a LOT of semester left and plenty of time to recover from a bad start

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Parts of a Fractional Number

The meaning of the "decimal point" is as follows:

$$123.406_{10} = 1 \times 10^{2} + 2 \times 10^{1} + 3 \times 10^{0} + 123 = 100 + 20 + 3$$
$$4 \times 10^{-1} + 0 \times 10^{-2} + 6 \times 10^{-3} \quad 0.406 = \frac{4}{10} + \frac{6}{1000}$$
$$= 123.406_{10}$$

Changing to base 2 induces a "binary point" with similar meaning:

$$110.101_2 = 1 \times 2^2 + 1 \times 2^1 + 0 \times 2^0 + \qquad 6 = 4 + 2$$
$$1 \times 2^{-1} + 0 \times 2^{-2} + 1 \times 2^{-3} \qquad 0.625 = \frac{1}{2} + \frac{1}{8}$$
$$= 6.625_{10}$$

One could represent fractional numbers with a fixed point e.g.

- ▶ 32 bit fractional number with
- ▶ 10 bits left of Binary Point (integer part)
- ▶ 22 bits right of Binary Point (fractional part)

BUT most applications require a more flexible scheme

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Scientific Notation for Numbers

"Scientific" or "Engineering" notation for numbers with a fractional part is

Standard	Scientific	printf("%.4e",x);
123.456	1.23456×10^{2}	1.2346e+02
50.01	5.001×10^{1}	5.0010e+01
3.14159	3.14159×10^{0}	3.1416e+00
0.54321	5.4321×10^{-1}	5.4321e-01
0.00789	7.89×10^{-3}	7.8900e-03

- Always includes one non-zero digit left of decimal place
- Has some significant digits after the decimal place
- Multiplies by a power of 10 to get actual number

Binary Floating Point Layout Uses Scientific Convention

- Some bits for integer/fractional part
- Some bits for exponent part
- ► All in base 2: 1's and 0's, powers of 2

Conversion Example

Below steps convert a decimal number to a fractional binary number equivalent then adjusts to scientific representation.

```
float fl = -248.75;

7 6 5 4 3 2 1 0 -1 -2

-248.75 = -(128+64+32+16+8+0+0+0).(1/2+1/4)

= -11111000.11 *2^0

76543210 12

= -1111100.011 *2^1

6543210 123

= -111110.0011 *2^2

543210 1234
```

. . .

MANTISSA EXPONENT = -1.111100011 * 2^7 0 123456789

 $Mantissa \equiv Significand \equiv Fractional Part$

Principle and Practice of Binary Floating Point Numbers

- ▶ In early computing, computer manufacturers used similar principles for floating point numbers but varied specifics
- ► Example of Early float data/hardware
 - ▶ Univac: 36 bits, 1-bit sign, 8-bit exponent, 27-bit significand¹
 - ▶ IBM: 32 bits, 1-bit sign, 7-bit exponent, 24-bit significand²
- Manufacturers implemented circuits with different rounding behavior, with/without infinity, and other inconsistencies
- Troublesome for reliability: code produced different results on different machines
- ► This was resolved with the adoption of the IEEE 754 Floating Point Standard which specifies
 - ▶ Bit layout of 32-bit float and 64-bit double
 - Rounding behavior, special values like Infinity
- ► Turing Award to William Kahan for his work on the standard

¹Floating Point Arithmetic

²IBM Hexadecimal Floats

IEEE 754 Format: The Standard for Floating Point

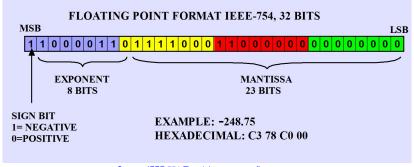
float	double	Property
32	64	Total bits
1	1	Bits for sign (1 neg / 0 pos)
8	11	Bits for Exponent multiplier (power of 2)
23	52	Bits for Fractional part or mantissa
7.22	15.95	Decimal digits of accuracy ³

- Most commonly implemented format for floating point numbers in hardware to do arithmetic: processor has physical circuits to add/mult/etc. for this bit layout of floats
- Numbers/Bit Patterns divided into three categories

Category	Description	Exponent
Normalized	most common like 1.0 and -9.56e37	mixed 0/1
Denormalized	very close to zero and 0.0	all 0's
Special	extreme/error values like ${\tt Inf}$ and ${\tt NaN}$	all 1's

³Wikipedia: IEEE 754

Example float Layout of -248.75: float_examples.c



Source: IEEE-754 Tutorial, www.puntoflotante.net

```
Color: 8-bit blocks, Negative: highest bit, leading 1
```

```
Exponent: high 8 bits, 2<sup>7</sup> encoded with bias of -127

1.111100011...

1000_0110 - 0111_1111

= 128+4+2 - 127

= 134 - 127

= 7

implied leading 1
not in binary layout
```

Normalized Floating Point: General Case

- ➤ A "normalized" floating point number is in the standard range for float/double, bit layout follows previous slide
- ► Example: -248.75 = -1.111100011 * 2^7

Exponent is in **Bias Form** (not Two's Complement)

- Unsigned positive integer minus constant bias number
- ► **Consequence**: exponent of 0 is not bitstring of 0's
- ➤ Consequence: tiny exponents like -125 close to bitstring of 0's; this makes resulting number close to 0
- ▶ 8-bit exponent 1000 0110 = 128+4+2 = 134 so exponent value is 134 127 = 7

Integer and Mantissa Parts

- ► The leading 1 before the binary point is **implied** so does not show up in the bit string
- ► Remaining fractional/mantissa portion shows up in the low-order bits

Fixed Bit Standards for Floating Point

IEEE Standard Layouts

Kind	Sign	Exponent			Mantissa
	Bit	Bits	Bias	Exp Range	Bits
float	31 (1)	30-23 (8 bits)	-127	-126 to +127	22-0 (23 bits)
double	63 (1)	62-52 (11 bits)	-1023	-1022 to $+1023$	51-0 (52 bits)

Standard allows hardware to be created that is as efficient as possible to do calculation on these numbers

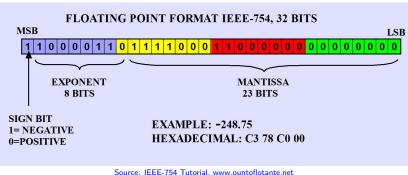
Consequences of Fixed Bits

- ➤ Since a fixed # of bit is used, **some numbers cannot be exactly represented**, happens in any numbering system:
- ▶ Base 10 and Base 2 cannot represent $\frac{1}{3}$ in finite digits
- Base 2 cannot represent 1/10 in finite digits
 float f = 0.1;
 printf("0.1 = %.20e\n",f);
 0.1 = 1.00000001490116119385e-01

Try show_float.c to see this in action

Exercise: Quick Checks

- 1. Represent 7.125 in binary using "binary point" notation
- 2. What distinct parts are represented by bits in a floating point number (according to IEEE)
- 3. What is the "bias" of the exponent for 32-bit floats
- 4. What does the number 1.0 look like as a float?



Source: IEEE-754 Tutorial, www.puntoflotante.net

The diagram above may help in recalling IEEE 754 layout

Answers: Quick Checks

- 1. Represent 7.125 in binary using a "binary point"
 - $ightharpoonup 7_{10} = 111_2$
 - $ightharpoonup 0.125_{10} = \frac{1}{8} = 2^{-3} = 0.001_2$
 - ightharpoonup 7.125₁₀ = 111.001₂
- 2. What distinct parts are represented by bits in a floating point number (according to IEEE 754)
 - Sign, Exponent, and Mantissa/Fractional Portion
- What is the "bias" of the exponent for 32-bit floats (according to IEEE 754)
 - ▶ Bias is -127 which is subtracted from the unsigned value of the 8 exponent bits to get the actual exponent
- 4. What does the number 1.0 look like as a float?
 - Positive: sign bit of 0
 - Exponent is 0, so sign bits total 127:
 - 8 4
 - Mantissa has implied leading 1 and all 0's so:
 000 0000 0000 0000 0000
 23 20 16 12 8 4

Special Cases: See float_examples.c

Denormalized values: Exponent bits all 0

- ► Fractional/Mantissa portion evaluates *without* implied leading one, still an unsigned integer though
- Exponent is Bias + 1: 2^{-126} for float
- ► Result: very small numbers close to zero, smaller than any other representation, degrade uniformly to 0
- Zero: bit string of all 0s, optional leading 1 (negative zero);

Special Values

- ▶ **Infinity**: exponent bits all 1, fraction all 0, sign bit indicates $+\infty$ or $-\infty$
- Infinity results from overflow/underflow or certain ops like float x = 1.0 / 0.0;
- #include <math.h> gets macro INFINITY and -INFINITY
- ▶ NaN: not a number, exponent bits all 1, fraction has some 1s
- ► Errors in floating point like 0.0 / 0.0

Other Float Notes



DURING A COMPETITION, I TOLD THE PROGRAMMERS ON OUR TEAM THAT e^{σ} - τ WAS A STANDARD TEST OF FLOATING-POINT HANDLERS -- IT WOULD COME OUT TO 20 UNLESS THEY HAD ROUNDING ERRORS.



THAT'S AWFUL HALF THEIR ALGORITHMS LOOKING FOR THE BUG BEFORE THEY FIGURED IT OUT.

Source: XKCD #217

Approximations and Roundings

- Approximate $\frac{2}{3}$ with 4 digits, usually 0.6667 with standard rounding in base 10
- Similarly, some numbers cannot be exactly represented with fixed number of bits: ¹/₁₀ approximated
- ► IEEE 754 specifies various rounding modes to approximate numbers

Clever Engineering

- IEEE 754 allows floating point numbers to sort using signed integer routines
- Bit patterns for float follows are ordered the same as bit patterns for signed int
- Integer comparisons are usually fewer clock cycles than floating comparisons

Sidebar: The Weird and Wonderful Union

- Bitwise operations like & are not valid for float/double
- Can use pointers/casting to get around this OR...
- Use a union: somewhat unique construct to C
- Defined like a struct with several fields
- ▶ BUT fields occupy the same memory location (!?!)
- Allows one to treat a byte position as multiple different types, ex: int / float / char[]
- Memory size of the union is the max of its fields

```
// union.c
typedef union { // shared memory
 float fl;
               // an int
 int in; // a float
 char ch[4]; // char array
} flint t;
               // 4 bytes total
int main(){
 flint t flint;
 flint.in = 0xC378C000;
 printf("%.4f\n", flint.fl);
 printf("%08x %d\n",flint.in,flint.in);
 for(int i=0; i<4; i++){
    unsigned char c = flint.ch[i];
   printf("%d: %02x '%c'\n",i,c,c);
```

Floating Point Operation Efficiencies

- ► Floating Point Operations per Second, **FLOPS** is a major measure for numerical code/hardware efficiency
- Often used to benchmark and evaluate scientific computer resources, (e.g. top super computers in the world)
- ► Tricky to evaluate because of
 - ► A single FLOP (add/sub/mul/div) may take 3 clock cycles to finish: latency 3
 - Another FLOP can start before the first one finishes: pipelined
 - Enough FLOPs lined up can get average 1 FLOP per cycle
 - ► FP Instructions may automatically operate on multiple FPs stored in memory to feed pipeline: **vectorized ops**
 - Generally referred to as superscalar
 - Processors schedule things out of order too
- ► All of this makes micro-evaluation error-prone and pointless
- Run a real application like an N-body simulation and compute

$$FLOPS = \frac{number of floating ops done}{time taken in seconds}$$

Top 5 Super Computers Worldwide, Nov 2017

Rank	System	#Cores	Rmax (TFlop/s)	Rpeak (TFlop/s)	Power (kW)
1	Sunway TaihuLight <i>China</i> Sunway MPP	10,649,600	93,014.6	125,435.9	15,371
2	Tianhe-2 (MilkyWay-2) <i>China</i> TH-IVB-FEP Cluster	3,120,000	33,862.7	54,902.4	17,808
3	Piz Daint <i>Switzerland</i> Cray XC50	361,760	19,590.0	25,326.3	2,272
4	Gyoukou <i>Japan</i> ZettaScaler-2.2 HPC system	19,860,000	19,135.8	28,192.0	1,350
5	Titan <i>USA</i> Cray XK7	560,640	17,590.0	27,112.5	8,209

https://www.top500.org/lists/2017/11/

Top 5 Super Computers Worldwide, Nov 2018

Rank	System	#Cores	Rmax (TFlop/s)	Rpeak (TFlop/s)	Power (kW)
1	Summit <i>United States</i> IBM POWER9 22C 3.07GHz	2,397,824	143,500.0	200,794.9	9,783
2	Sierra <i>United States</i> IBM POWER9 22C 3.1GHz,	1,572,480	94,640.0	125,712.0	7,438
3	Sunway TaihuLight <i>China</i> Sunway MPP	10,649,600	93,014.6	125,435.9	15,371
4	Tianhe-2A <i>China</i> TH-IVB-FEP Cluster	4,981,760	61,444.5	100,678.7	18,482
5	Piz Daint <i>Switzerland</i> Cray XC50, Xeon E5-2690v3	387,872	21,230.0	27,154.3	2,384

https://www.top500.org/list/2018/11/

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4	Tianhe-2A <i>China</i> Xeon 2.2GHz	4,981,760	61,444.5	100,678.7	18,482
5	Frontera, <i>United States</i> Dell 6420, Xeons 2.7GHz	448,448	23,516.4	38,745.9	??

https://www.top500.org/list/2019/11/

Top 5 Super Computers Worldwide, June 2020

Rank	System	#Cores	Rmax (TFlop/s)	Rpeak (TFlop/s)	Power (kW)
1	Fugaku, <i>Japan / Fujitsu</i> Fujitsu A64FX 2.2GhZ (Arm)	7,299,072	415,530.0	513,854.7	28,335
2	Summit <i>United States</i> IBM POWER9 22C 3.07GHz (Power)	2,397,824	143,500.0	200,794.9	10,096
3	Sierra <i>United States</i> IBM POWER9 22C 3.1GHz (Power)	1,572,480	94,640.0	125,712.0	7,438
4	Sunway TaihuLight <i>China</i> Sunway SW26010 (custom RISC)	10,649,600	93,014.6	125,435.9	15,371
5	Tianhe-2A <i>China</i> Intel Xeon 2.2GHz (x86-64)	4,981,760	61,444.5	100,678.7	18,482

https://www.top500.org/lists/top500/2020/06/

Top 5 Super Computers Worldwide, November 2020

Rank	System	#Cores	Rmax (TFlop/s)	Rpeak (TFlop/s)	Power (kW)
1	Fugaku, <i>Japan / Fujitsu</i> Fujitsu A64FX 2.2GhZ (Arm)	7,299,072	415,530.0	513,854.7	28,335
2	Summit <i>United States</i> IBM POWER9 22C 3.07GHz (Power)	2,397,824	143,500.0	200,794.9	10,096
3	Sierra <i>United States</i> IBM POWER9 22C 3.1GHz (Power)	1,572,480	94,640.0	125,712.0	7,438
4	Sunway TaihuLight <i>China</i> Sunway SW26010 (custom RISC)	10,649,600	93,014.6	125,435.9	15,371
5	Selene <i>USA</i> , <i>NVIDIA/AMD</i> AMD EPYC 7742 64C 2.25GHz (x86-64)	555,520	63,460.0	79,215.0	2,646

https://www.top500.org/lists/top500/2020/06/