Parallel Languages and Platforms

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Logistics

Today

- Finish Applications
- Parallel Languages and Platforms

Assignments

- ▶ Questions on A2?
- ► A3 Canceled: all students who take Final Exam or Submit A2 will get +10% for A3*

Schedule

Mon 12/13	Applications Parallel Languages
Wed 12/15	Review A2 Due Course Evals Due
Fri 12/17	A2 Late Deadline
Mon 12/20	Final Exam
	1:30-3:30pm

Lecture Location

* Originally was "take Final Exam" but added "submit A2"

Course Feedback

Official Student Rating of Teaching (SRTs)

- ▶ Official UMN Evals are done online this semester
- ► Available here: https://srt.umn.edu/blue
- ► EVALUATE YOUR LECTURE SECTION: 001
- **▶ Due** Wed 12/15/2021 by 9:45am
- ▶ Response Rate $\geq 80\% \rightarrow$ One Final Exam Question Revealed during lecture on 12/15

Menagerie of Parallel Languages and Platforms

Distributed Memory Only

Erlang, Map+Reduce / Hadoop, Job Schedulers

Shared Memory Only

Cilk, Clojure

Distributed + Shared

Unified Parallel C, Chapel

Device Concurrency / GPUs

CUDA / OpenCL, OpenACC

Erlang



- Developed for distributed computation, telephony systems
- Virtual machine which mirrors many OS functions
- Process spawn to create lightweight procs
- send/receive clauses to share information among processes
- Facilities to contact a remote Erlang VM and talk to its processes

Erlang Sample straight from Wikipedia

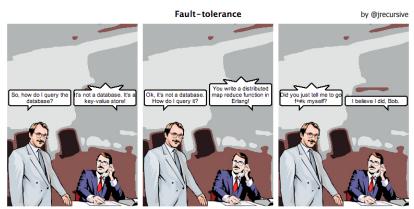
```
% Create a process on this machine and invoke the function
% web:start_server(Port,MaxConnections)
ServerProcess = spawn(web, start server, [Port, MaxConnections]),
% Create a remote process and invoke the function
% web:start server(Port, MaxConnections)
% on machine RemoteNode
RemoteProcess = spawn(RemoteNode, web, start server, [Port, MaxConnections]),
% Send a message to ServerProcess (asynchronously). The message
% consists of a tuple with the atom "pause" and the number "10".
ServerProcess ! {pause, 10},
% Receive messages sent to this process
receive
    a_message -> do_something;
    {data, DataContent} -> handle(DataContent);
    {hello, Text} -> io:format("Got hello message: ~s", [Text]);
    {goodbye, Text} -> io:format("Got goodbye message: ~s", [Text])
end.
```

Erlang's Nature and Target

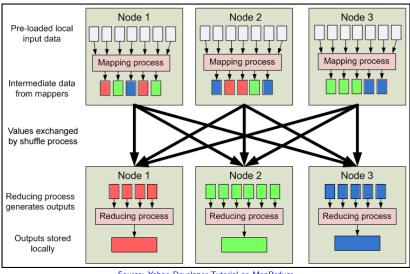
- Syntax and semantics are somewhat odd/archaic but can be "gotten used to"
- Targeted at client server architectures, computation distributed across many nodes
- Well known for robustness of the VM, fault-tolerance features to keep application going if participating nodes go down
- Not targeted at high-performance computation / scientific problems, more towards business, IT, web services

MapReduce (or more properly Map, Shuffle, Reduce)

- ► A style of programming, inspired by functional programming (def doub-sum (reduce + 0 (map double '(1 2 3 4 5))))
- ► Targeted at big data: large distributed stores of data
 - ► Map: Transform / filter data in some way
 - ▶ Shuffle: Move data with same properties to same node
 - ▶ Reduce: Combine results on individual nodes



Basic Architecture of MapReduce



Source: Yahoo Developer Tutorial on MapReduce

Shameless Wikipedia Example: Document Word Counts

Pseudocode

```
function map(String name,
             String document):
  // name: document name
  // document: document contents
  for each word w in document:
    emit(w, 1)
function reduce(String word,
                Iterator partialCounts):
  // word: a word to count
  // partialCounts: list of
                    partial counts
  sum = 0
  for each pc in partialCounts:
    sum += pc
  emit (word, sum)
```

- Goal: produce frequency of each word in a document
- Nodes are each fed the document
- During reduce() emit
 pairs like ("apple",1)
 and ("Dell",1)
- System automatically sends pairs with key apple to the same nodes (redistribute)
- Nodes run reduce() to count apple occurrences, may redistribute further

Variety of Languages for MapReduce Framework

```
public static class MapClass extends MapReduceBase
  implements Mapper < Long Writable, Text, Text, IntWritable >
  private final static
    IntWritable one = new IntWritable(1):
  private Text word = new Text():
  public void
  map (LongWritable key, Text value,
      OutputCollector<Text, IntWritable> output,
      Reporter reporter) throws IOException
    String line = value.toString():
    StringTokenizer itr = new StringTokenizer(line):
    while (itr.hasMoreTokens()) {
      word.set(itr.nextToken()):
      output.collect(word, one):
  }
public static class Reduce extends MapReduceBase
  implements Reducer < Text, IntWritable, Text, IntWritable >
  public void
  reduce(Text key, Iterator < IntWritable > values,
         OutputCollector<Text, IntWritable> output.
         Reporter reporter) throws IOException
    int sum = 0:
    while (values.hasNext()) {
      sum += values.next().get();
    output.collect(kev. new IntWritable(sum)):
```

Code for Word Frequencies in a Document

← Java (Source)

Pig Latin \downarrow (Source)

```
A = load './input.txt';
B = foreach A generate
    flatten(TOKENIZE((chararray)$0))
    as word;
C = group B by word;
D = foreach C generate COUNT(B), group;
store D into './wordcount';
```

MapReduce Framework Notes

- Primary contribution of implementations is distributing load across many machines efficiently
- When machines both store some data and participate in MapReduce, gain locality for speed
- Alternative to large database processing, may open up opportunities for parallelism to avoid read/write locks in traditional DBs
- Most frequently referenced implementation is Apache Hadoop



But other implementations exist, some proprietary

- All implementation implement a MapReduce server/scheduler to which jobs are submitted
 - > javac MRWordCount.java
 - > java MRWordCount &
 - > bin/hadoop job -list
 - 1 jobs currently running
 - JobId State job_0001

Job Schedulers

- Has long been the need for many parallel jobs to be run on individual systems/clusters
- Job schedulers offer frameworks for this: submit many programs to run, scheduler assigns resources
- slurm: a common scheduler used for instance on MSI systems
- Generic form of easy concurrency for variety of different programs, resources etc.

A sample slurm job submission script from MSI

```
#!/bin/bash -1
#SBATCH --time=8:00:00
#SBATCH --ntasks=8
#SBATCH --mem=10g
#SBATCH --tmp=10g
#SBATCH --mail-type=ALL
#SBATCH --mail-user=sample_email@umn.edu
cd ~/program_directory
module load intel
module load ompi/intel
mpirun -np 8 program_name < inputfile
> outputfile
```

Cilk and CilkPlus (supported in gcc/g++)

- ► Tutorial Here
- \triangleright Extensions to C/C++ which enable easy spawning of threads (strands) to run functions concurrently
- Primary Additions are keywords to easily spawn functions

```
// Run func concurrently - separate thread
int x = cilk spawn func(n);
// Wait for all running functions to finish
cilk sync;
// Compile with cilk features enabled
> gcc -fcilkplus cilk-fib.c
```

- - Examine: cilk-fib.c
 - Contrast with PThread startup
 - ► Alas, Cilk dropped from gcc

Clojure and Software Transactional Memory

► A lisp which runs on the Java Virtual Machine



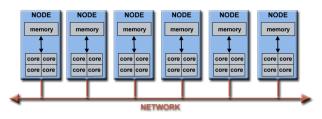
- Design Goal: allow shared memory parallelism
- ► Realization:
 - ► Each data element has well-defined local/shared semantics
 - Data is immutable by default
 - Provides atom data types for atomic alterations
 - Other alterations to shared resources occur in dosync blocks
 - Software Transactional Memory (STM) system: try changing a shared area, if it changes, try again with the new current value
- Runs as an executable JAR

```
> java -jar clojure-1.8.0.jar
Clojure 1.8.0
user=> (def x (atom 0))
#'user/x
user=> x
#object[clojure.lang.Atom 0x6c372fe6 {:status :ready, :val 0}]
user=> @x
0
```

Picalc in Clojure

```
;; Serial version with atomic updates
(defn calc-pi-atoms [iterations]
  (let [hits (atom 0)]
    (dotimes [i iterations]
      (let [x (rand) y (rand)]
        (if (<= (+ (* x x) (* y y)) 1)
          (swap! hits inc))))
    (double (* (/ @hits iterations) 4))))
;; Parallel version with atomic updates
(defn calc-pi-atoms [iterations nthreads]
  (let [hits (atom 0)]
    (doall (pmap ;; parallel map, force evaluation
      (fn [x]
        (dotimes [i (/ iterations nthreads)]
          (let [x (rand) y (rand)]
            (if (<= (+ (* x x) (* y y)) 1)
              (swap! hits inc)))))
      (range nthreads))) ;; map onto number of threads
    (double (* (/ @hits iterations) 4))))
```

Common HPC Parallel Platform



- Similar config to medusa cluster we used
- Cluster of machines, each with multiple cores
- Options to program:
 - Serial execution on each core/machine
 - Parallel shared memory execution on each machine
 - Parallel distributed memory execution on each core
 - ► **Mixed**: Distributed/Shared parallel execution



Mixing MPI and OpenMP

Unified Parallel C (requires special compiler)

- ► Extensions to the C language
- Aimed at BOTH shared memory and distributed memory
- ► Automatic THREADS and id MYTHREAD variables
- ► Thread is more generalized: might be same machine (shared) or different machine (distributed)
- Shared memory blocks
 shared int all_hits[THREADS];
 Access like int x = all_hits[1]; will work locally (shared)
 or via MPI-style message passing if remote (distributed)
- Compiler/runtime automatically sets up sharing
- Standard locks
- Automatic loop parallelization via upc_forall(..) with some affinity control: which thread executes which iteration
- Control over layout of shared blocks of memory: which thread gets what section
- Examine upc-shared-picalc.c