

Algorithms Tutorial 1

- ① (a) $f_1(n) \in \Theta(n!)$
 (b) $f_2(n) \in \Theta(n^{40})$ -
 (c) $f_3(n) \in \Theta(1)$ -
 (d) $f_4(n) = n^{3/2} + n^{1/2} \log n \in \Theta(n^{3/2})$ -
 (e) $f_5(n) \in \Theta(3^n)$
 (f) $f_6(n) = 2^{n+3} = 2^n \cdot 2^3 = 4 \cdot 2^n \in \Theta(2^n)$
 (g) $f_7(n) \in \Theta(n)$ -

Arranging them lowest to highest

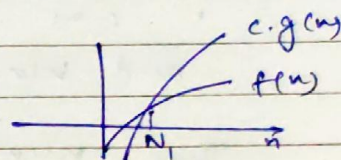
$$(c) < (g) < (b) = (d) < (f) < (e) < (a)$$

~~$$f_3(n) < f_7(n) < f_2(n) = f_4(n) < f_6(n) < f_5(n) < f_1(n)$$~~

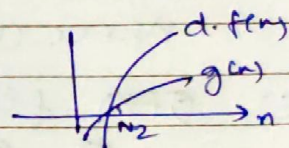
$$\Rightarrow f_3(n) < f_7(n) < f_2(n) = f_4(n) < f_6(n) < f_5(n) < f_1(n)$$

~~2. 1000~~

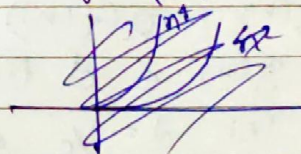
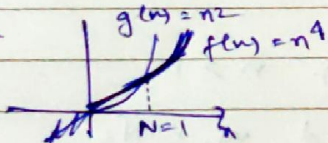
- ② (a) for $f(n) = O(g(n))$,
 $f(n) \leq c \cdot g(n)$ for ~~some~~ all $n \geq N_1$



- for $g(n) = O(f(n))$,
 $g(n) \leq d \cdot f(n)$ for all $n > N_2$



This is not true, since if $f(n) = n^2$ and $g(n) = n^4$,
 then



for $n = N_1$, ~~g(n)~~ $f(n) \leq 1 \cdot g(n)$ ($c=1$)

$$\text{Hence, } f(n) = O(g(n)) \Rightarrow n^2 = O(n^4)$$

But the opposite is not true, since for any d and any n , $g(n) = n^4$ will ~~never~~ sometime exceed $f(n) = n^2$

If we prove by contradiction, then, we say,

$$g(n) = O(f(n)) \Rightarrow n^4 = O(n^2)$$

$$\Rightarrow n^4 < d \cdot n^2 \text{ for some all } n \geq N_0$$

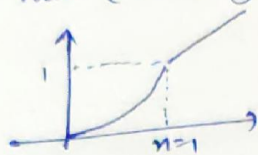
$$\Rightarrow n^2 < d$$

which is not true, since $n \geq N_0$
 Hence (a) is false

$$(b) \min(f(n), g(n)) = o(f(n) + g(n))$$

$$\text{If } f(n) = n, g(n) = n^2,$$

then $\min(f(n), g(n))$ is shown as



$$\text{But } o(f(n) + g(n)) = o(n + n^2) \neq o(n^2)$$

$$\min(f(n), g(n)) \leq c \cdot (f(n) + g(n)) \text{ for all } n \geq N$$

Now, ~~if~~ since $f(n), g(n) \geq 0$ for all n

$$\therefore f(n) + g(n) \geq f(n)$$

$$f(n) + g(n) \geq g(n)$$

$$\therefore c = 1$$

and hence $c \cdot (f(n) + g(n)) \geq \min(f(n), g(n))$ ~~for all~~ for all n .

$$\text{Thus, } \min(f(n), g(n)) = o(f(n) + g(n))$$

Hence (b) is true

$$(c) \text{ ~~if~~ } f(n) = \Omega(g(n))$$

$$\Rightarrow \text{~~if~~ } f(n) \geq c \cdot g(n) \quad n \geq N_1$$

$$\& g(n) = O(f(n))$$

$$\Rightarrow g(n) \leq d \cdot f(n) \quad n \geq N_2$$

$$\text{let } d = 1/c \text{ and } N_1 = N_2$$

then Both equations are true,

Hence (c) is True