

Lecture 39: Reductions, Algorithmic Bounds, NP Completeness.

Best self-extracting bitstream is the program that you used to make the image.

Given a bitstream B , what's the shortest bitstream C_B that outputs B .

Kolmogorov Complexity:

The Java-Kolmogorov complexity $K_j(B)$ is the length of the shortest Java program that generates B . length in Bytes.

— There is an answer! It might be hard to find.

Fact 1: Kolmogorov Complexity is independent of language.

Fact 2: It's impossible to write a program that calculates Kolmogorov Complexity of any bitstream.

— Corollary: It's impossible to write the "perfect" compression algorithm

How hard is the independent set?

Does there exist a ^{independent} set of size k ? (Such that vertices don't touch)

— For each of the 2^N possible colorings:

— Check if # of colored vertices is equal to k : $O(N)$

— For each red vertex, check that neighbors are white $O(kN)$

— If both checks succeed, return true.

— If either check fails, go on to next coloring.

— RT: $O(kN \cdot 2^N) = O(N^2 \cdot 2^N)$ $k \leq N$

Reductions:

— Useful tool. Using one algorithm to solve another prob.

— X reduces to Y if instance of X can be transformed into an instance of Y that solves X .

3SAT: Given a boolean formula, does there exist a truth value for boolean vars that obey a set of 3 var disjunctive constraints?

Example: $(x_1 \vee x_2 \vee !x_3) \wedge (x_1 \vee !x_1 \vee x_1) \wedge (x_2 \vee x_3 \vee x_4)$

Solution: $x_1 = T, x_2 = T, x_3 = T, x_4 = F$

Independence-set cracks 3SAT.

Complexity Classes:

Decision problem = yes or no problem.

a problem is in complexity class P if:

- It is a decision problem
- An answer can be found in $O(N^k)$ for some k.

a problem is in class NP if:

- it is a decision problem
- A "yes" answer can be verified in $O(N^k)$ time for some k. We can verify a specific answer "yes" in $O(N^k)$.

a problem is NP-complete if:

- it is a member of NP
- it cracks every other problem in NP.

3SAT is NP complete. \rightarrow it cracks any problem for which a "yes" answer can be efficiently verified.

As NP-complete problems grow it gets easier to find more NP-complete problems.