

## Depth First Search:

Goal: to find a path from  $s$  to every other reachable vertex, visiting each vertex at most once.

- Mark  $v$ .
- For each unmarked adjacent vertex,
  - set  $edgeTo[w] = v$ .
  - $dfs(w)$ .

Demo @ 8:30 in lecture.

DFS = general term of any graph search that goes deep before backtracking.

\* HasPathTo: if  $marked[v] == True$ !

Runtime =  $\Theta(V+E)$

Each edge is used exactly once as well.

- Each vertex  $x$  is visited once, each visit costs constant time.

Space =  $\Theta(V)$

- Call stack depth is at most  $V$ .

There are a lot of graph traversals too!

- DFS preorder: order of DFS calls

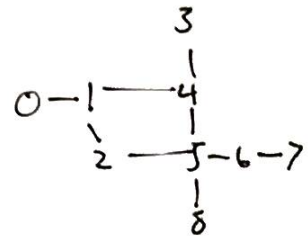
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- DFS postorder: order of returns from DFS

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- Level order: increasing distance from  $s$ .

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If we do level order and a vertex has 2 edges, use the shortest one.

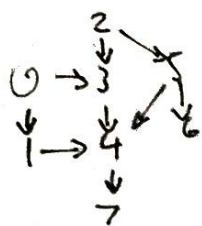
## Topological Sort

- DFS postorder from every vertex with indegree 0, not clearing markings in between traversals.

→ Demo @ 2:45

- Record DFS post order in a list.

- Topological ordering is the reverse of this list.



→ DFS post orders

$[7, 4, 1, 3, 0, 6, 5, 2]$

→ Topological ordering

$[2, 5, 6, 0, 3, 1, 4, 7]$

When you topological sort, all arrows point to the right.

topological sort

- for finding an ordering of vertices consistent w/ directed edges.

Efficiency  $\rightarrow \Theta(V+E)$  time,  $\Theta(V)$  space.

### Breadth First Search

- Level-order  $\rightarrow$  order visited by breadth first search.

1) Create a queue w/ starting vertex  $s$ . Mark it.  
- this is the fringe.

2) Repeat until queue is empty:

a) Remove vertex  $v$  from queue.

b) Add any unmarked vertices adjacent to  $v$  into the queue and mark them.

Demo @ 39:15

### Breadth First Path

Goal: Find a path from  $s$  to every reachable vertex.

1) Initialize the fringe  $\rightarrow$  a queue w/ starting vertex  $s$  and mark it.

2) Repeat until ~~the~~ queue is empty:

a) Remove vertex  $v$ .

b) For each unmarked neighbor of  $v$ : ~~add~~  
mark, add to queue, set  $edgeTo = v$ .

Invariant:

fringe queue always has:

$\geq 0$  vertices of distance  $k$  from  $s$ .

$\geq 0$  vertices of distance  $k+1$  from  $s$ .

- Useful for finding shortest path

$\Theta(V+E)$

$\Theta(V+E)$

space:

$\Theta(V)$

Breadth first is not good for GMaps because we want to minimize drive time.