

- 3) Mathematically derive the average runtime complexity of non-random pivot. Since each element of the partitioning operations is compared with the pivot the recurrence relation for quick sort, where  $n$  is the array size, the partitioning process takes  $O(n)$  times.

$$T(n) = T(k) + T(n-k-1) + O(n)$$

In average case, the array is divided into two sub-arrays of almost equal size by the pivot.

So an average  $k \approx n/2$ , leading to the recurrence relation

$$T(n) = 2T(n/2) + O(n)$$

The Master theorem or the recursion tree approach can be used to solve this recurrence.

$$T(n) = aT(n/b) + O(n^d)$$

Now, we calculate the critical exponent.

$$\log_b a = \log_2 2 = 1$$

Since  $d = \log_b a$ , the time complexity is  $O(n^d \log n)$ , which simplifies to

$$T(n) = O(n \log n)$$

Thus, the average worst time complexity of quick sort is  $O(n \log n)$