

Introductory Lecture

Topics: basics, resources, stl, bitwise tricks

League of Programmers

ACA, IIT Kanpur

October 5, 2013

Outline

- 1 Aim and Clarifications
- 2 Common Problems
- 3 Parsing a problem
- 4 Standard Template Library
- 5 Using Bitwise
- 6 Problems

Aim

- Discussion camp not a lecture series. You need to show motivation.

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- There are handsome rewards - prestige, joy of learning new things, and yes lots of money!

Language specifications

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ACM-ICPC Official languages. Allowed in almost every contests
- But, Java is comparatively very slow, so sometimes an optimal algorithm might time out on the judge
- C has too restrictive and does not support stl/templates /classes
- Use Library functions and Data Structures instead of writing your own every time

Programming competitions

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- IOPC (IITK), Shaastra (IITM), Bitwise (IITKgp)

Websites for practice

- Compete against Indian coders in live contests: Codechef
- Short Programming Contests: Codeforces, Topcoder
- Problem set Archives: SPOJ, Project Euler, livearchive, acm.sgu.ru and many more

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Overflow

```
#include<stdio.h>
int main()
{
    int a, b;
    scanf("%d %d", &a, &b);
    printf("%d\n", a+b);
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- What if the given numbers are HUGE?
- Not all the input constraints are explicit
- Always think about the worst case scenario, edge cases, etc.

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- Comparing doubles
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double a, b;
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instead, do the following

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#define EPS 0.0000001
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(a-EPS<b and a+EPS>b)
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- Invalid memory reference
- Using too much memory than provided

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- $1\text{sec} \approx (1 - 2) * 10^8$ operations
- $x\text{MB} \approx x/4 * 10^6$ sized int arrays
- For example, 10 test cases with $N = 10000$ means $O(N^2)$ is required

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- 6 Coding the solution (the easiest part of all)
- 7 Debugging (TLE: time limit exceeded, WA : incorrect solution etc.)

Some optimizations

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 - operations on unsigned ints/long long are faster
 - bitwise operators and shift operators (& ^ | >> <<)
 - Using too much memory (> 10MB) slows down programmes

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- I/O: do NOT use cin/cout for large input output

Some Standard paradigms

- Sorting
- Searching
- Preprocessing
- Divide-and-Conquer
- Dynamic Programming
- Greedy Algorithms
- Graph
- Network Flow
- Backtracking
- Computational Geometry
- Pure maths
- Ad-hoc problems

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Standard template library

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- reverse
- swap

Stack

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- ③ C++ and Java have implementations of stack

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Query Operations

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- *IsEmpty(S)* : determine if the stack S is empty

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Query Operations

- $IsEmpty(S)$: determine if the stack S is empty
 - $Top(S)$: returns the element at the top of the stack
- Example: If S is a_1, a_2, \dots, a_n , then $Top(S)$ is a_1 .

Stack

Update Operations

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- *CreateEmptyStack(S)* : create an empty stack

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- $CreateEmptyStack(S)$: create an empty stack
- $Push(x, S)$: push x at the top of the stack

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x, a_1, a_2, \dots, a_n

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Example: If S is a_1, a_2, \dots, a_n , then after *Push(x, S)*, S becomes
 x, a_1, a_2, \dots, a_n
- *Pop(S)* : delete an element from the top of the stack S .
Example: If S is a_1, a_2, \dots, a_n , then after *Push(x, S)*, S becomes
 a_2, \dots, a_n

- 1 Try to write a code for the same and check it here:
[http://www2.cse.iitk.ac.in:
81/newonj/problem.php?problemID=42](http://www2.cse.iitk.ac.in:81/newonj/problem.php?problemID=42)
- 2 Then later on try this problem:
<http://codeforces.com/problemset/problem/344/D>

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- *Dequeue(Q)* : delete the element from the front of the queue.
Example: If Q is a_1, a_2, \dots, a_n , then after *Dequeue(Q)*, Q becomes
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- 2 Then later on try this problem:
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81/newonj/problem.php?problemID=47`
Hint: Its a Queue problem, now the (x, y) is a queue element.

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<http://www.spoj.com/problems/EZDIJKST/>
Hint: Its a a tutorial problem for the algorithm Dijkstra.

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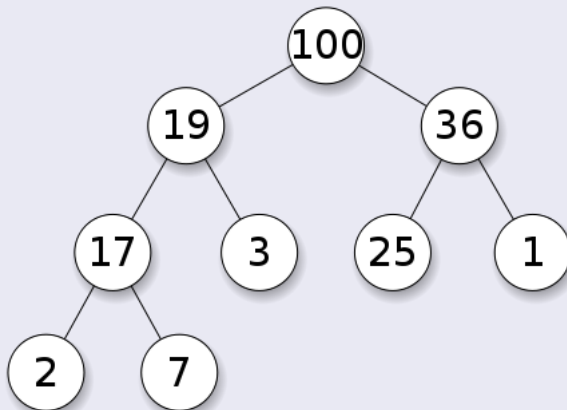
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- 5 May need rearrangement of some nodes

Heap



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- 4 Running time = tree height = $O(\log n)$

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Introduction to Bitwise Operators

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- 2 Our weapons:
 << (left shift); >> (right shift); & (bitwise and);
 | (bitwise or); ^ (bitwise xor); ~ (bitwise not)
- 3 Speed up the code by upto 100 times. **Caution: try to use bitwise operations on unsigned integers only**

Beauty of Bitwise

1 Example:

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- Generate all subsets of S in $2^{|S|}$ time
- Generate all subsets of $1..n$ changing one bit at a time
- Generate all subsets of S which have exactly t elements
- Count the number of elements of elements in a set S

Beauty of Bitwise

1 Example:

- Any subset of $0, 1, \dots, 31$ is a single int
- Do set union/intersection/complement in one operation
increment/decrement all elements by x in one operation

2 Even more:

- Find if $x \in S$.
- Generate all subsets of S in $2^{|S|}$ time
- Generate all subsets of $1..n$ changing one bit at a time
- Generate all subsets of S which have exactly t elements
- Count the number of elements of elements in a set S
- Remove smallest element from S

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3 Never multiply or divide or take remainder modulo power of 2

Outline

- 1 Aim and Clarifications
- 2 Common Problems
- 3 Parsing a problem
- 4 Standard Template Library
- 5 Using Bitwise
- 6 Problems**

Problems

Links:

- ① <http://spoj.pl/problems/WEIRDFN>
- ② <http://www.spoj.pl/problems/HOMO/>
- ③ <http://spoj.pl/problems/HISTOGRA>
- ④ <http://spoj.pl/problems/SUBSEQ>
- ⑤ <http://www.spoj.pl/problems/NGM2/>
- ⑥ <http://www.spoj.pl/problems/JOCHEF>
- ⑦ <http://www.spoj.pl/problems/SWTHIN/>
- ⑧ <http://www.spoj.pl/problems/LAZYPROG/>
- ⑨ <http://www.spoj.com/problems/NAKANJ/>
- ⑩ <http://www.spoj.com/problems/PPATH/>