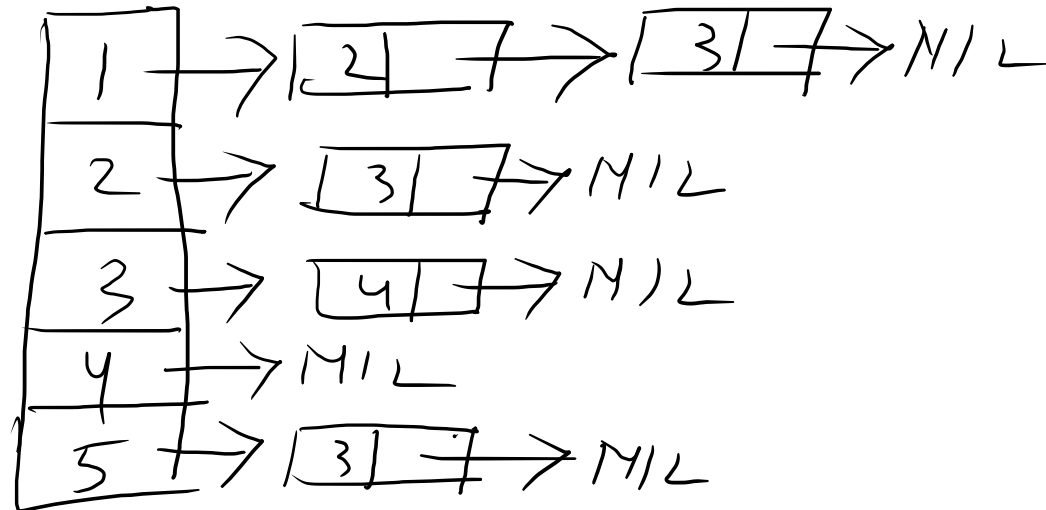
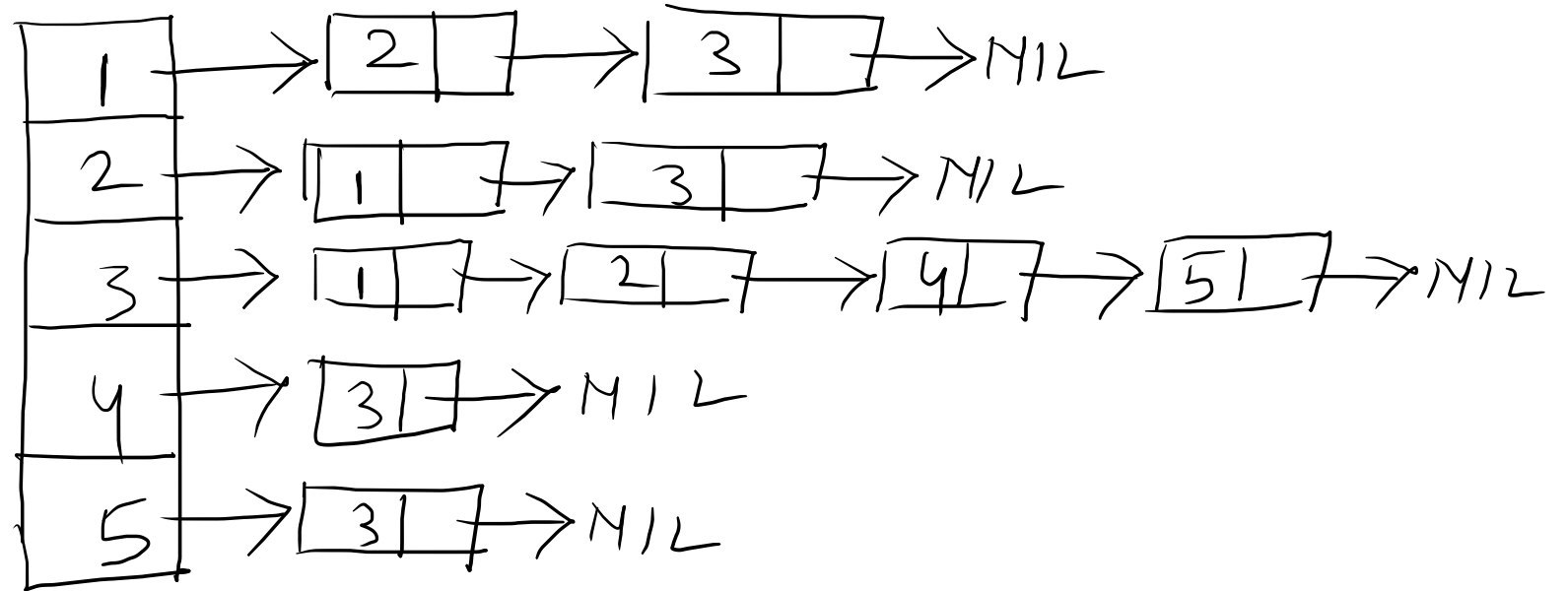
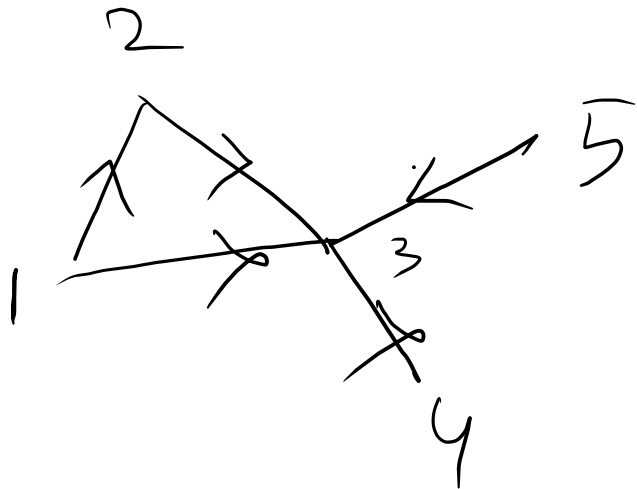
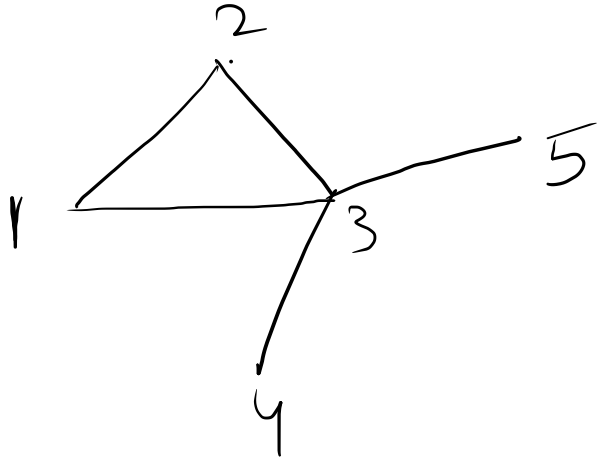


Representations of graphs

Adjacency - List



Storage
 $O(|V| + |E|)$

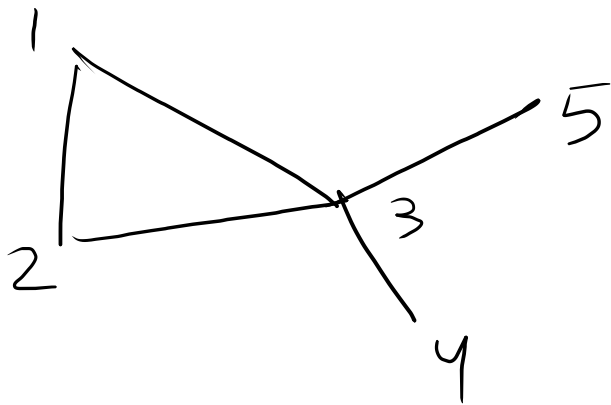
Adjacency matrix

v_1, v_2, \dots, v_n

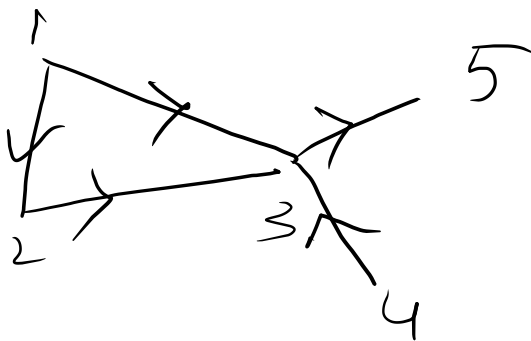
storage: $O(|V|^2)$

$A_{n \times n}$

$$a_{ij} = \begin{cases} 1 & \text{if } v_i \text{ is connected with } v_j \\ 0 & \end{cases}$$



	1	2	3	4	5
1	0	1	1	0	0
2	1	0	1	0	0
3	1	1	0	1	1
4	0	0	1	0	0
5	0	0	1	0	0

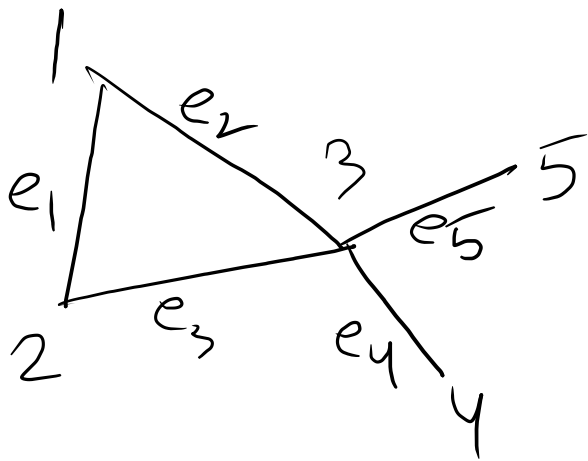


	1	2	3	4	5
1	0	1	1	0	0
2	0	0	1	0	0
3	0	0	0	0	1
4	0	0	1	0	0
5	0	0	0	0	0

Incidence matrix

$A_{n \times m}$

$$a_{ij} = \begin{cases} 1 & e_i \text{ is incident on } j\text{-th vertex} \\ 0 & \text{otherwise} \end{cases}$$



	e_1	e_2	e_3	e_4	e_5
1	1	1	0	0	0
2	1	0	1	0	0
3	0	1	1	1	1
4	0	0	0	1	0
5	0	0	0	0	1

Storage:

$$O(|V| |E|)$$

Graph searching algorithms

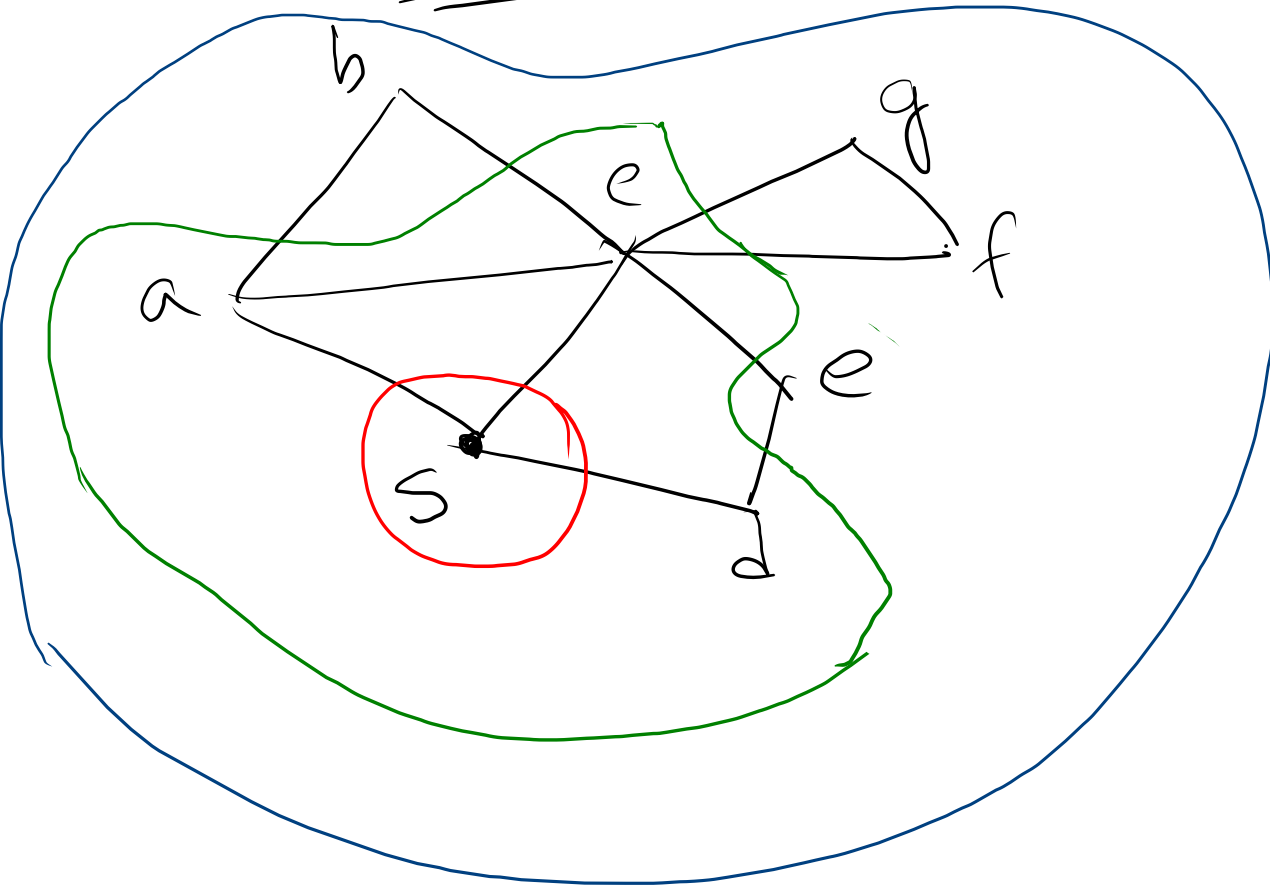
- systematic searching of each edge or vertices of the graph.
- For directed or undirected graph.
- Applications.
 - maze search.
 - connected component.

Two popular algorithms.

i) BFS

ii) DFS

BFS



White: This vertex is not yet discovered.

Gray: It is discovered, but not all its neighbours are discovered.

Black: Already discovered, and all its neighbours are discovered.

BFS (G, s)

for each $u \in G.V$
 $u.color = white$
 $u.dist = \infty$
 $u.pred = NIL$

$s.color = gray$

$s.dist = 0$

$Q = new\ queue$

$Q.enqueue(s)$

while Q is not empty.

$u = Q.dequeue()$

 for $v \in u.adj$

 if $v.color = white$

$v.color = gray$

$v.dist = u.dist + 1$

$v.pred = u$

$Q.enqueue(v)$

$u.color = black$.

Running time: $O(|V| + |E|)$

