

## Network-Flow Problem

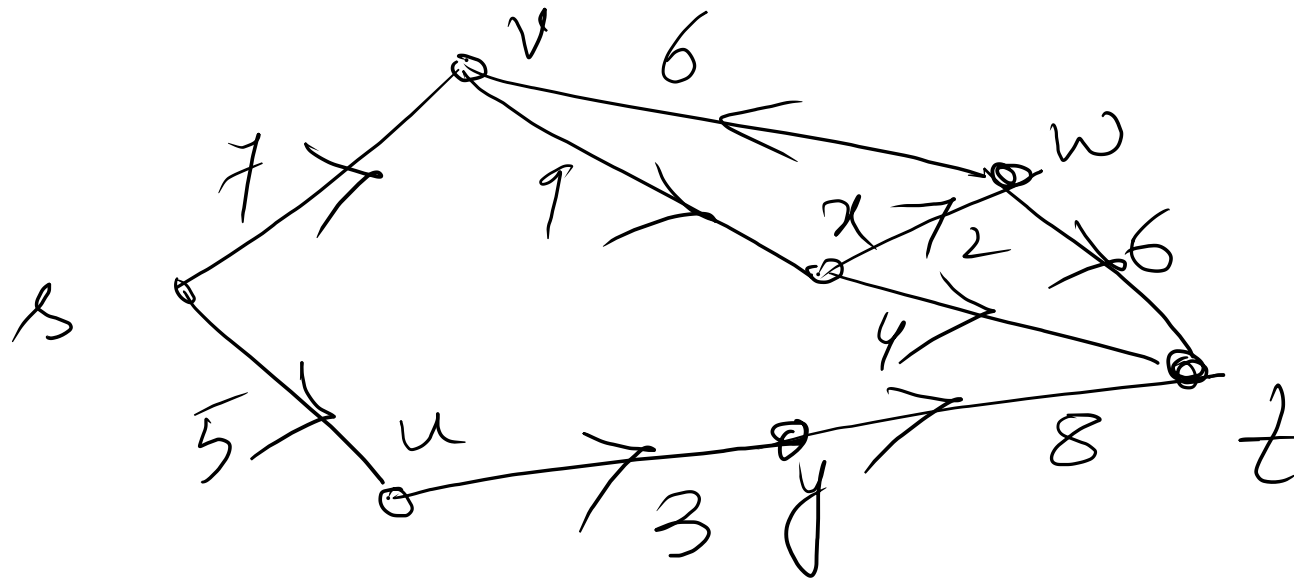
Flow network : Input:  $G = (V, E)$ ,  $s, t$ ,  $c: E \rightarrow \mathbb{R}^+$

A directed graph  $G$

$s$  : source

$t$  : sink

For each edge  $e \in E$   $c(e) \geq 0$  called the capacity of the edge.



## Maximum flow problem

An  $s-t$  flow  $f$  that satisfies the following two conditions / constraints:

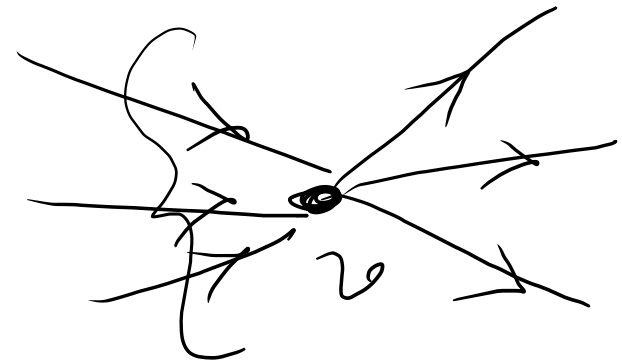
① capacity constraint:

for each edge  $e \in E$ ,  $0 \leq f(e) \leq c(e)$

② flow conservation constraint

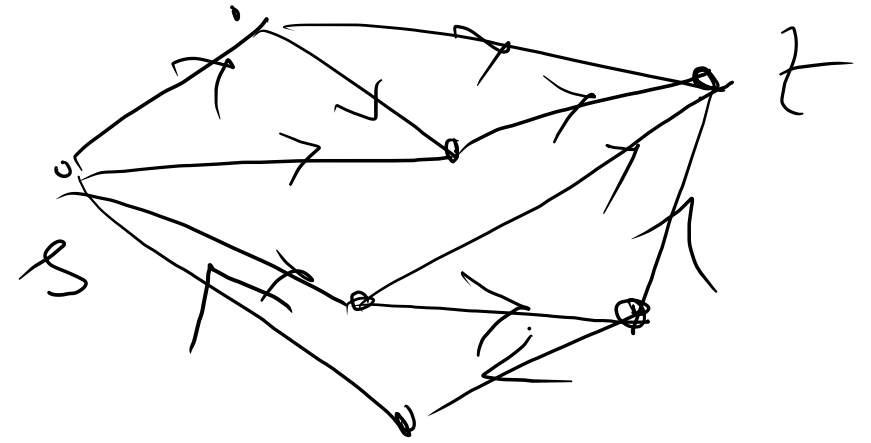
for each vertex  $v \in V \setminus \{s, t\}$

$$\sum_{\substack{e \text{ is incident} \\ \text{towards } v}} f(e) = \sum_{\substack{e \text{ is incident} \\ \text{outward } v}} f(e)$$

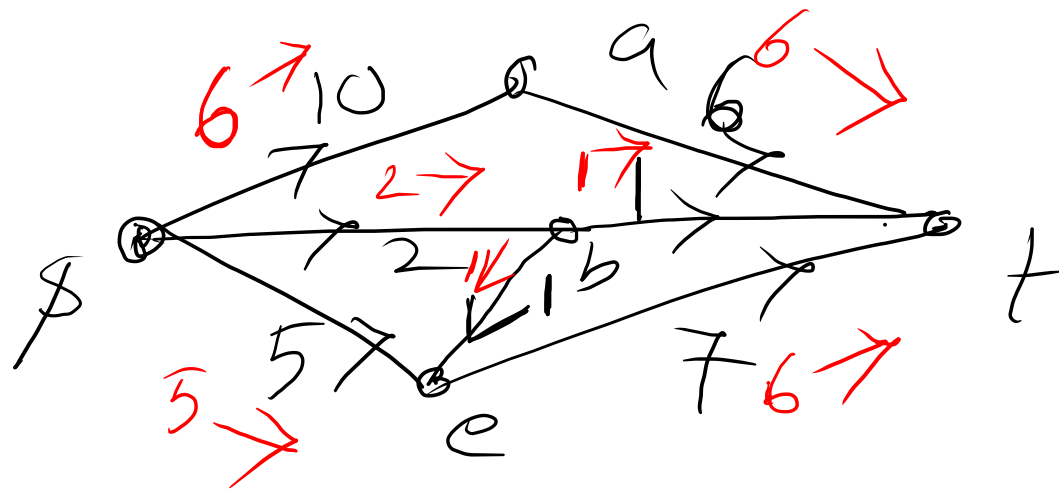


The value of the flow

$$\text{val}(f) = \sum_{e \text{ is going out of } s} f(e) - \sum_{e \text{ is going inside } s} f(e)$$



Objective: Find a flow of maximum value.



## An algorithm

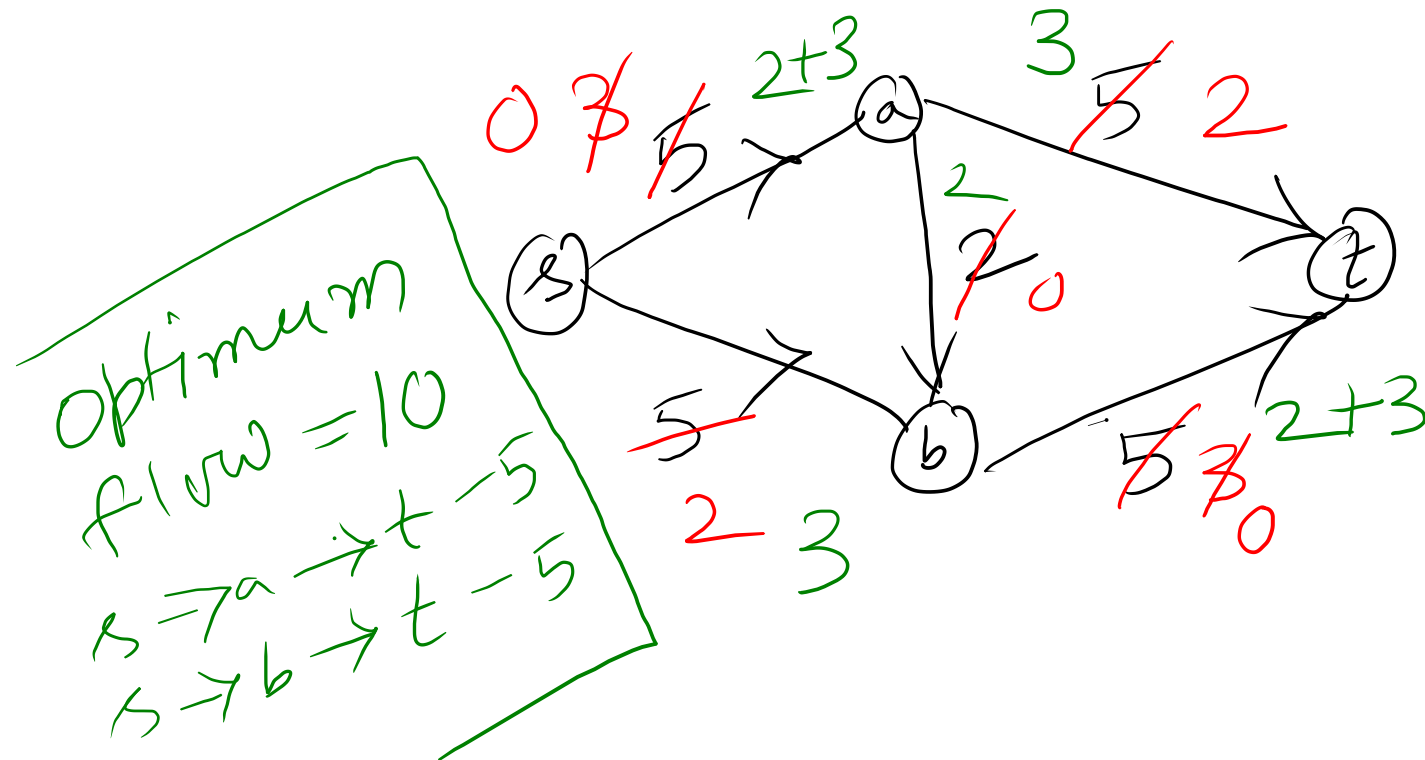
- for each edge flow is 0 i.e.  $f(e) = 0$
- if there is a  $s \rightarrow t$  path where  $f(e) < c(e)$
- augment flow along the path
- Repeat until no  $s \rightarrow t$  path exist.

$s \rightarrow a \rightarrow b \rightarrow t \rightarrow 2$

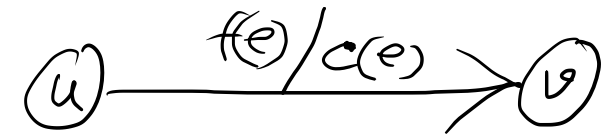
$s \rightarrow a \rightarrow t - 3$

$s \rightarrow b \rightarrow t - 3$

Total flow = 8

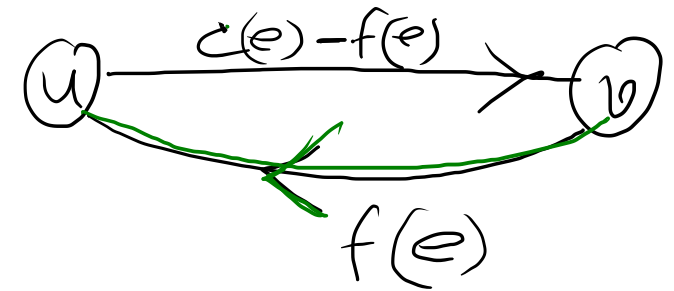


## Residual network



Residual capacity

$$c_f(e) = \begin{cases} c(e) - f(e) & \text{if } e \in E \\ f(e^{\text{reverse}}) & \text{if } e^{\text{reverse}} \in E \end{cases}$$



Residual network:  $G_f(V, E_f)$ ,  $s, t$ ,  $c_f$  is residual capacity function.

$$E_f = \{ e : f(e) < c(e) \} \cup \{ e : f(e^{\text{reverse}}) > 0 \}$$

## Ford-Fulkerson algorithm

Ford-Fulkerson( $G$ )

for each edge  $e \in E$   
 $f(e) = 0$

$G_f \leftarrow$  residual network of  $G$  w.r.t.  $f$

while there is an  $s \rightarrow t$  path  $P$  in  $G_f$

$f \leftarrow \text{augment}(G, P, c)$

update  $G_f$

return  $f$ .

$\text{augment}(G, P, c)$

$\delta \leftarrow$  bottleneck capacity  
in  $P$

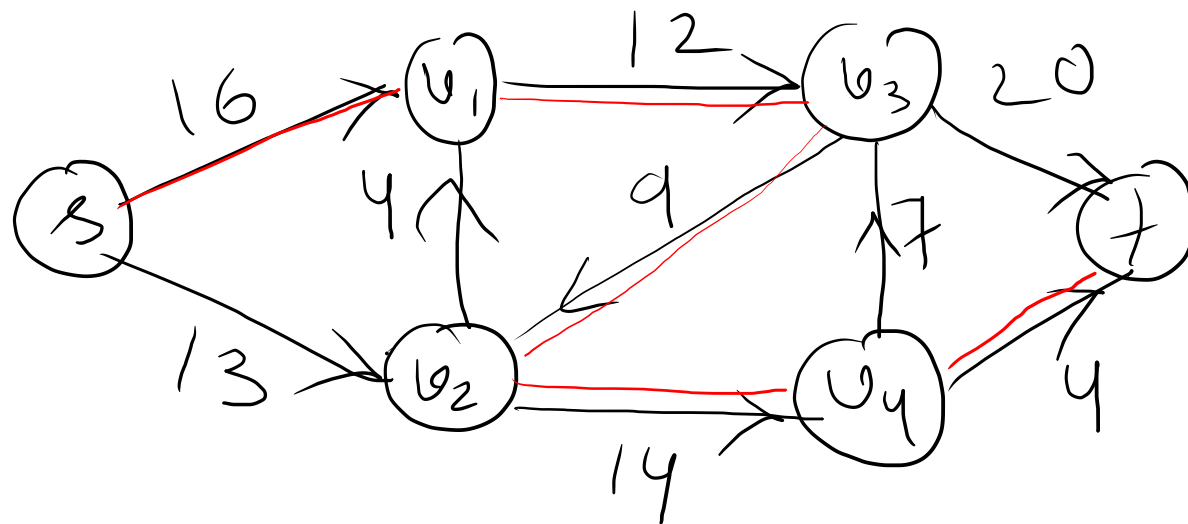
for each edge  $e$  in  $P$   
if  $e \in E$

$f(e) \leftarrow f(e) + \delta$

else

$f(e^{\text{reverse}}) =$   
 $f(e^{\text{reverse}}) - \delta$

Return  $f$ .



$s \rightarrow v_1 \rightarrow v_3 \rightarrow v_2 \rightarrow v_4 \rightarrow t$

