AMS 553.361 - Introduction to Optimization

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Introduction

Math 110.304 is one of the semi-important courses that is required/recommended for the engineering-based majors at Johns Hopkins University.

These notes are being live-TeXed, through I edot for Typos and add diagrams requiring the TikZ package separately. I am using Texpad on Mac OS X.

I would like to thank Zev Chonoles from The University of Chicago and Max Wang from Harvard University for providing me with the inspiration to start live-TeXing my notes. They also provided me the starting template for this, which can be found on their personal websites.

Please email any corrections or suggestions to ksriniv4@jhu.edu.

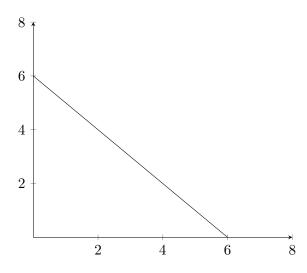
Lecture 1 (2018-08-30)

Introduction

Example 1. Up to 6 units of two nutrients can be added to solution and we require the number of units of nutrient 2 has to be at least (natural) logarithm of # units of nutrient one. Goal: choose x_1 = number units of nutrient 1, choose x_2 = number units of nutrient 2. To maximize expected height of plant $1 + x_1^2(x_2 - 1)^3 e^{-x_1 - x_2}$

Maximise
$$1 + x_1^2(x_2 - 1)^3 e^{-x_1 - x_2}$$

 $x_1 + x_2 \le 6$
 $x_2 > \log x_1$
 $x_1 \ge 0$
 $x_2 \ge 0$



generic optimization problem: say $\subseteq \mathbb{R}$

$$f: \underbrace{\mathbb{S}}_{\text{feasible region}} \to \underbrace{\mathbb{R}}_{\text{objective function}}$$

$$x^* \begin{bmatrix} 2 \\ 4 \end{bmatrix}$$
 optimal

however optimal objective function value is 1.2677

Definition 1. We say x^* is an optimal solution if

- $x^* \in \mathbb{S}$
- For any $y \in \mathbb{S}$

Example 2. Find min $\log x$ s.t. $-\infty \le x \le 7 \to \text{unbounded}$, has no solution.

Example 3. Find min $\log x$ s.t. $1 < x \le 7 \to \text{bounded}$, but is also no solution. (as we can go 1.000001)

Example 4. Find min $\log x$ s.t. x > 1, $x \le 0.5 \rightarrow$ infeasible!!!

Example 5. minimize $3 + (x - 1)^2$ s.t. $1 \le x \le 3 \to$ feasible as optimal solution $x^* = 2$ is an interior part of feasible region.

$$f(x) = 3 + (x - 1)^{2}$$
$$f'(x) = 0$$
$$f''(x) > 0$$

Example 6. Minimize $3 + (x - 2)^2$ s.t. $x \ge 10 \to \text{optimal solution is } x^* = 10$. but $f'(x) \ne 0$. But is not an interior point - it is a boundary point of feasible region.

Lecture 2 (2018-08-31)

Definition 1. $\forall x \in \mathbb{R}^n$ Euclidian length of x is $||x|| = \left(\sum_{i=1}^n x_i^2\right)^{\frac{1}{2}}$

 $\forall x, y \in \mathbb{R}^n$ Euclidian distance from x to y is ||x - y||

Definition 2. $\forall S \subseteq \mathbb{R}^n$ point $x \in S$ is an interior point of S if J a neighborhood of x which is a subset of S point $S \in \mathbb{R}^n$ is a boundary point of S if every neighborhood of x contains a point in S and a point not in S

- Set $S \subseteq \mathbb{R}^n$ is open if every point in S is an interior point. example: open ball.
- Set $S \subseteq \mathbb{R}^n$ is <u>closed</u> if S contains all boundary points of S.

Note: $\forall S \subseteq \mathbb{R}^n$, S is open **iff** S^c is closed.

In (P) min f(x) s.t. $x \in S$, suppose $x^* \in S$:

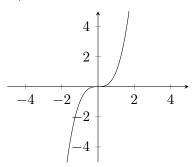
- x^* is a global maximizer if $\forall y \in S, f(x^*) \leq f(y)$
- x^* is a strict global minimizer if $\forall y \in S$, s.t. $y \neq x^*$, $f(x^*) < f(y)$
- x^* is a <u>local minimizer</u> if \exists any neighborhood N of x s.t. $\forall y \in N \cap S, f(x^*) \leq f(y)$
- x^* is a <u>strict</u> local minimizer if \exists any neighborhood N of x s.t. $\forall y \in (S \cap N) \underbrace{\setminus x^*}_{\text{besides}}$, $f(x^*) < f(y)$

Note: If $S \subseteq \mathbb{R}^1$, x^* is interior of S, f is suitably differentiable at x^* .

If x^* is a local minimizer $\Rightarrow f'(x^*) = 0$

 x^* is a local maximizer $\Rightarrow f'(x^*) = 0$

BUT $f'(x^*) = 0 \implies x^* \text{ local min/max}$



BUT

$$f'(x^*) = 0\&f''(x^*) > 0 \Longrightarrow x^*$$
 strict local max $f'(x^*) = 0\&f''(x^*) < 0 \Longrightarrow x^*$ strict local min

If $S \subseteq \mathbb{R}^n$, x^* interior part of S, f strictly differentiable, x^* local min/max $\Longrightarrow \nabla f(x^*) = \vec{0}$

$$\nabla f(x^*) = \vec{0}$$
 & [?] $\Rightarrow x^*$ strict local min/max

Lecture 3 (2018-09-05)

Linear Programming

Example 1. Diet Problem: You will pick levels of four ingredients for chicken feed.

• $x_1 = \text{units of ingredient } 1$

• x_3 = units of ingredient 3

• x_2 = units of ingredient 2

• x_4 = units of ingredient 4

All are real numbers, so fractions allowed. Given minimum levels of 3 nutrients $\begin{bmatrix} 6.2\\11.9\\10 \end{bmatrix} = \begin{bmatrix} \text{nutrient 1}\\ \text{nutrient 2}\\ \text{nutrient 3} \end{bmatrix}$

Given how many units of nutrient per unit of ingredient

nutrients \ ingredients	1	2	3	4
1 (protein)	1.2	2.6	0	9.2
2 (carbs)	3.9	1	.8	2
3 (cholesterol)	6	0	4	3.1

Problem: Find yand of ingredients that meet nutritional requirements cheaply as possible.

Minimum Cost of the ingredients $\begin{bmatrix} 6.\\2\\1.\\3. \end{bmatrix}$

minimize
$$6.2x_1 + 2x_2 + 1.6x_3 + 3.2x_4$$

$$\begin{cases}
1.2x_1 + 2.6x_2 + 0x_3 + 9.2x_4 \ge 6.2 \\
3.9x_1 + x_2 + 0.8x_3 + 2x_4 \ge 11.9 \\
3.9x_1 + 1x_2 + 0.8x_3 + 2x_4 \ge 11.9 \\
6x_1 + 0x_2 + 4x_3 + 3.1x_4 \ge 11.9
\end{cases}$$
 and $x_1, x_2, x_3, x_4 \ge 0$

What makes a ingredient list special?

- 1. proportionality
- 2. additivity
- 3. divisibility
 - fractions are allowed referring to x_1, x_2, x_3

(Linear Programming) in canonical form, an example

minimize
$$c_1x_1 + c_2x_2 + \dots + c_nx_n$$

s.t. $a_{11}x_1 + a_{12}x_2 + \dots + a_{1n}x_n \ge b_1$
 $a_{21}x_1 + a_{22}x_2 + \dots + a_{2n}x_n \ge b_2$
 \vdots
 $a_{m1}x_1 + a_{m2}x_2 + \dots + a_{mn}x_n \ge b_m$
where $x_1, x_2, \dots, x_n \ge 0$

Example 2. Transportation problem: There are 3 electricity generation plants α, β, γ and two cities u, v. α, β, γ respectively produce 65.2, 98.6, 32.5 units of electricity. u, v respectively use 86.2, 110.1 units of electricity

Question: How to supply transportation as cheaply as possible?

generator to city	cost variable	cost
$\alpha \to u$	x_1	31.7
$\alpha \to v$	x_2	28.6
$\beta \to u$	x_3	17.6
$\beta \to v$	x_4	37.4
$\gamma \to u$	x_5	22.8
$\gamma \to v$	x_6	29.7

min $31.7x_1 + 28.6x_2 + 17.6x_3 + 374x_4 + 22.8x_5 + 29.7x_6$

$$x_1 + x_2 \le 65.2 \qquad \alpha$$

$$x_3 + x_4 \le 98.6 \qquad \beta$$

$$x_5 + x_6 \le 32.5 \qquad \gamma$$
electricity produced at α, β, γ

$$x_1 + x_3 + x_5 \ge 86.2$$

$$v \qquad x_2 + x_4 + x_6 \ge 110.1$$
Need for electricity for u, v

(LP) in standard form minimize $c^T x$ s.t. $Ax = b, x \ge \vec{0}$ Here

$$A = \begin{bmatrix} 1 & 1 & 0 & 0 & 0 & 0 \\ 0 & 0 & 1 & 1 & 0 & 0 \\ 0 & 0 & 0 & 0 & 1 & 1 \\ 1 & 0 & 1 & 0 & 1 & 0 \\ 0 & 1 & 0 & 1 & 0 & 1 \end{bmatrix} \qquad b = \begin{bmatrix} 65.2 \\ 98.6 \\ 32.5 \\ 86.2 \\ 110.1 \end{bmatrix}$$

$$c^{T} = \begin{bmatrix} 31.7 & 28.6 & 17.6 & 37.4 & 22.8 & 29.7 \end{bmatrix}$$

Converting one form of (LP) to another

$$\max \rightarrow \min$$
 $\min \rightarrow \max \implies \max 6x_1 + 3x_2 - 4x_3 \leftrightarrow \min -6x_1 - 3x_2 + 4x_3$

$$Ax \ge b \to \le$$

$$Ax \ge b \leftrightarrow \boxed{-A}x \le \boxed{-b}$$

Standard form \rightarrow canonical form

$$\begin{vmatrix}
\min c^T x \\
\text{s.t. } Ax = b \\
x \ge 0
\end{vmatrix}
\leftrightarrow
\begin{cases}
\min c^T x \\
\text{s.t. } Ax \ge b \\
Ax \le b \\
x \ge \vec{0}
\end{cases}
\leftrightarrow
\begin{cases}
\min c^T x \\
\text{s.t. } Ax \ge b \\
-Ax \ge -b \\
x \ge \vec{0}
\end{cases}
\leftrightarrow
\begin{cases}
\min c^T x \\
A \\
-A \\
x \ge \vec{0}
\end{cases}$$

Canoconical form \rightarrow Standard form