

CS/EE120A: Logic Design

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Logic Design?

Basic organization of the circuitry of a digital computer.

- ✓ Two-valued logic system (binary code): 1/0, on/off
- ✓ Logic gates made of integrated circuits for calculations
- ✓ Three basic kinds of logic gates: AND, OR, NOT

Computer



A programmable usually electronic device that can
store, retrieve, and process
data

Digital Computers



Analog Computers



The Antikythera mechanism, dating between 150 and 100 BC, was an early analog computer: calculate astronomical positions



The Norden bombsight was a highly sophisticated optical/mechanical analog computer used by the United States Army Air Force during World War II to aid the pilot of a bomber aircraft in dropping bombs accurately.

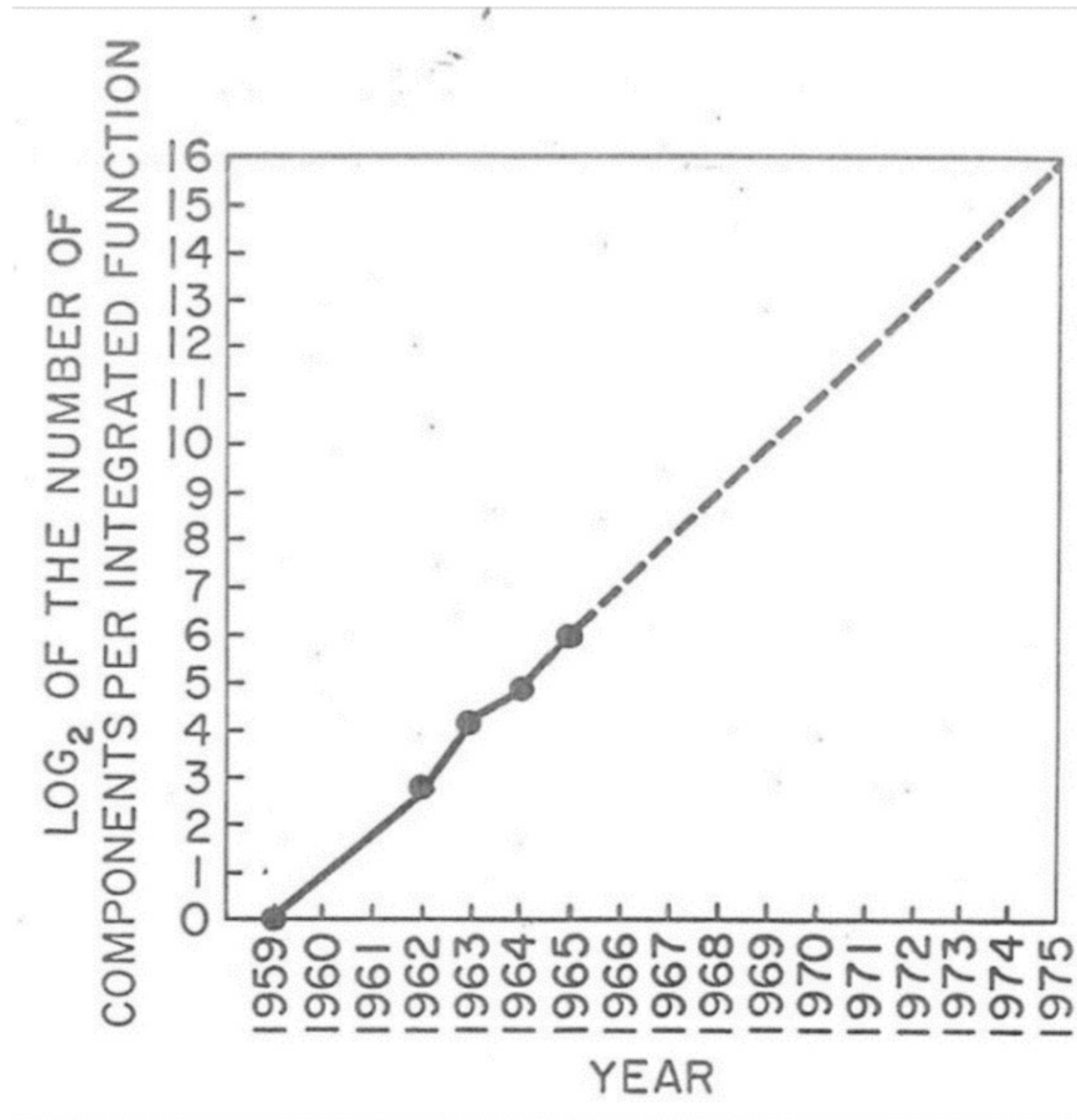


MONIAC (Monetary National Income Analogue Computer), was created in 1949 by the New Zealand economist Bill Phillips to model the national economic processes of the UK.

Why are digital computers more popular now?

- Please identify how many of the following statements explains why digital computers are now more popular than analog computers.
 - ① The cost of building systems with the same functionality is lower by using digital computers.
 - ② Digital computers can express more values than analog computers in the same range.
 - ③ Digital signals are less fragile to defective components.
 - ④ Digital data are easier to store.
- A. 0
B. 1
C. 2
D. 3
E. 4

Moore's Law

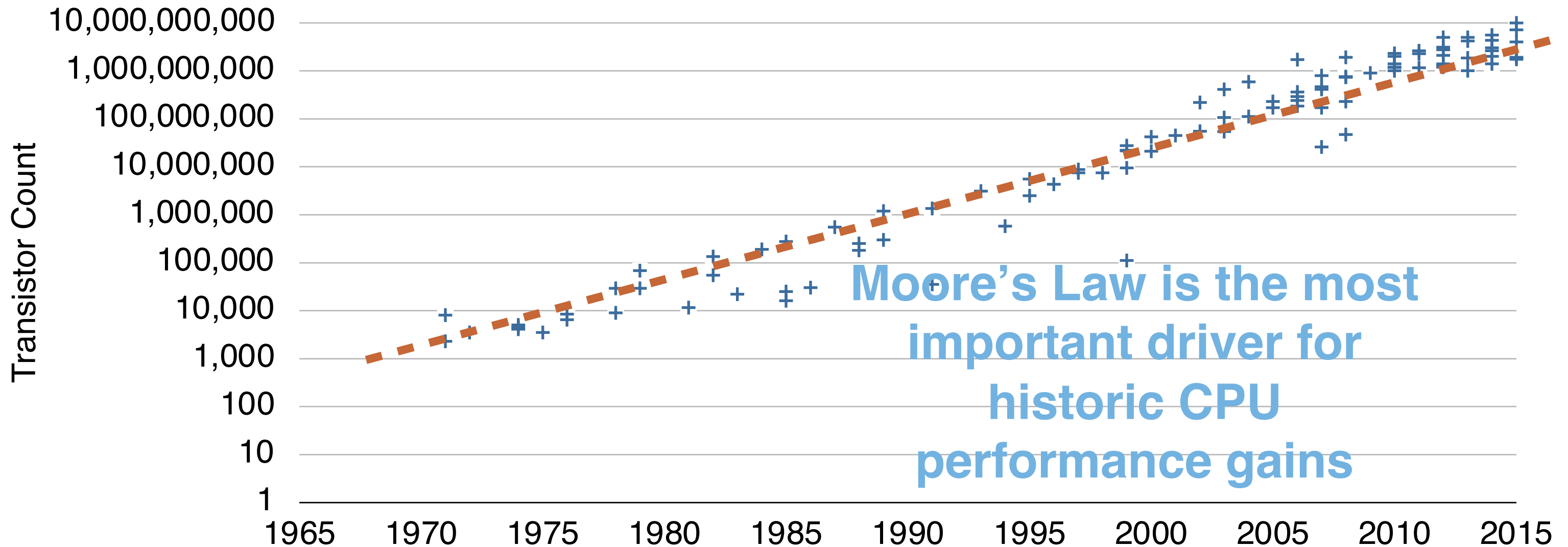


Gordon Moore in 1965

Moore, G. E. (1965), 'Cramming more components onto integrated circuits', Electronics 38 (8) .


Moore's Law⁽¹⁾

- The number of transistors we can build in a fixed area of silicon doubles every 12 ~ 24 months.



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A. 0

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C. 2

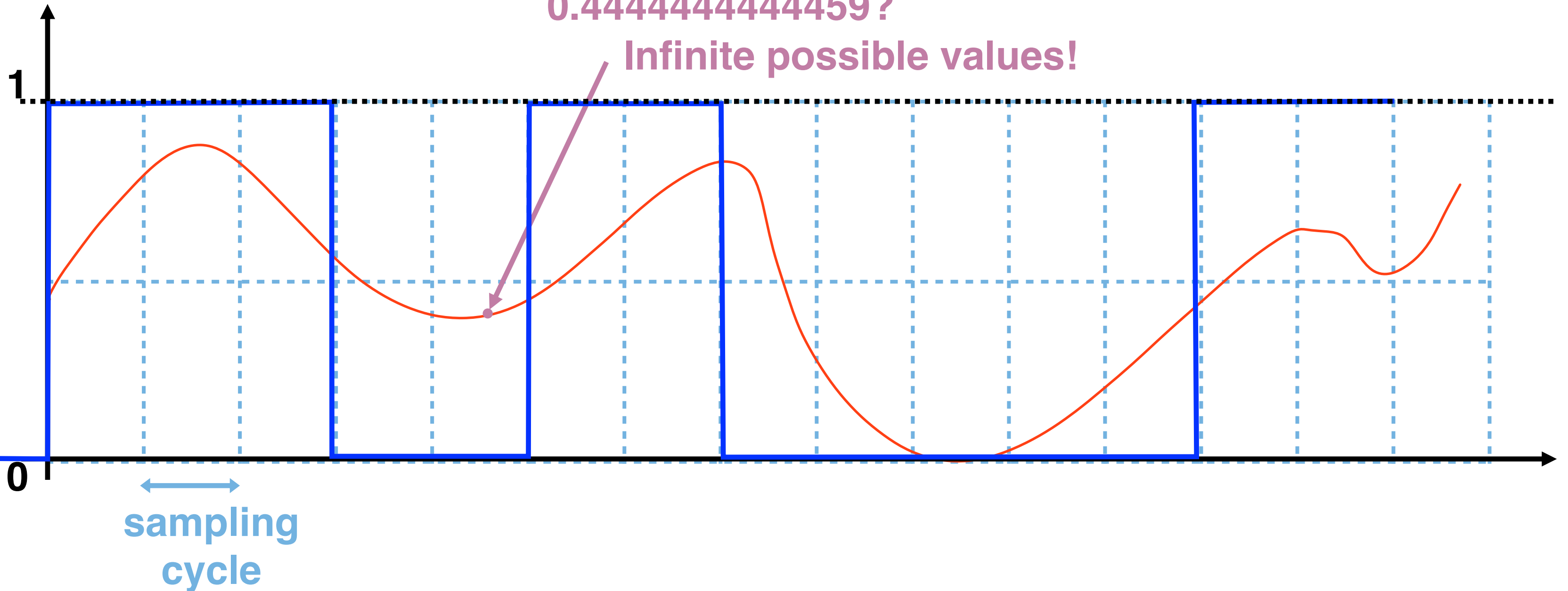
D. 3

E. 4

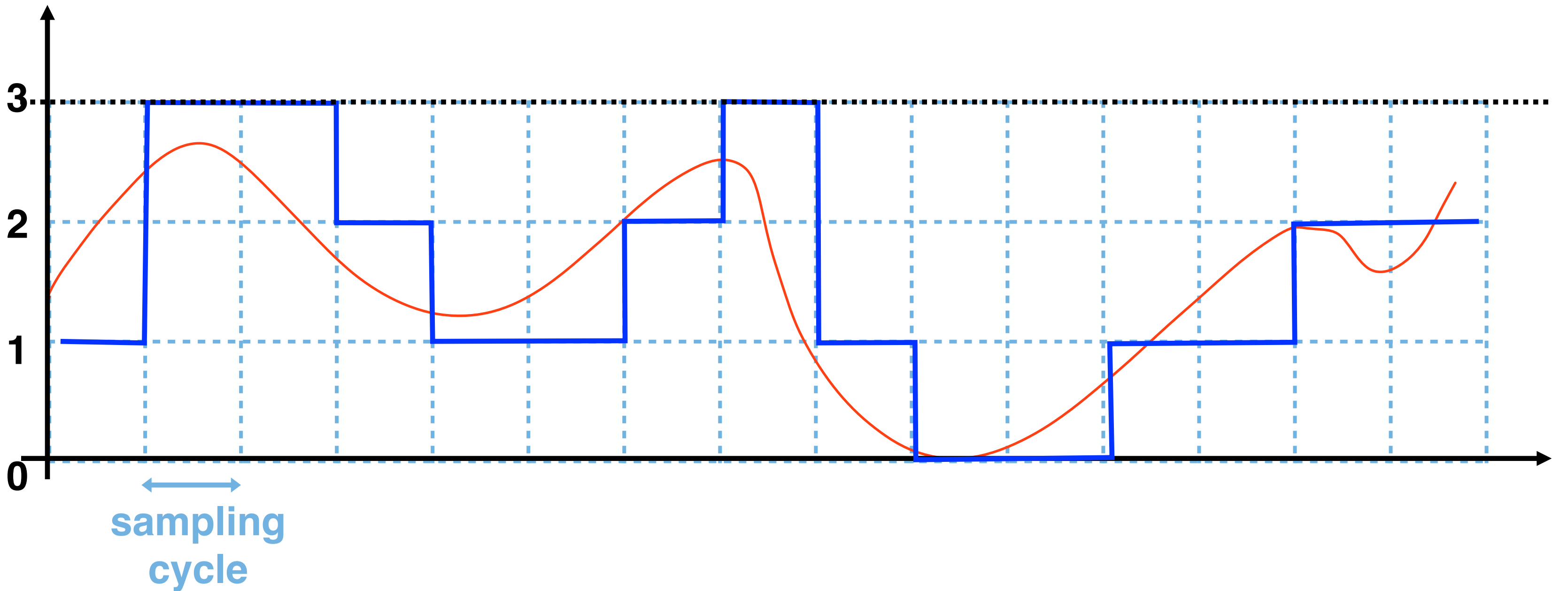
Analog v.s. digital signals

0.5? 0.4? 0.45?
0.445? 0.4445? or
0.44444444444459?

Infinite possible values!



Analog v.s. digital signals



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③ Digital signals are less fragile to defective components.

④ Digital data are easier to store.

A. 0

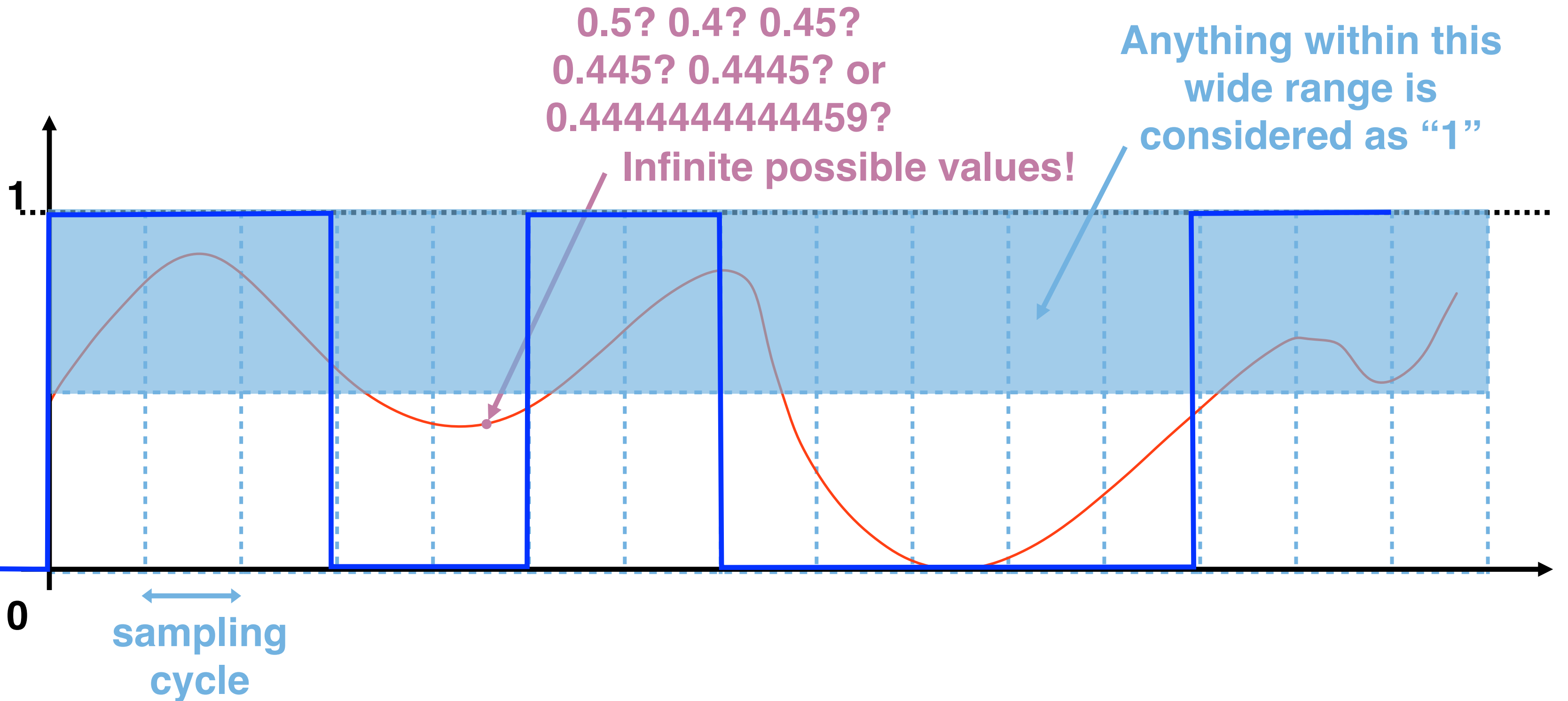
B. 1

C. 2

D. 3

E. 4

Analog v.s. digital signals



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④ Digital data are easier to store.

A. 0

B. 1

C. 2

D. 3

E. 4

Analog data storage



VS

Digital data storage



Samples per second

- CD Audio = 44,100
- DVD Audio = 192,000

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C. 2

D. 3

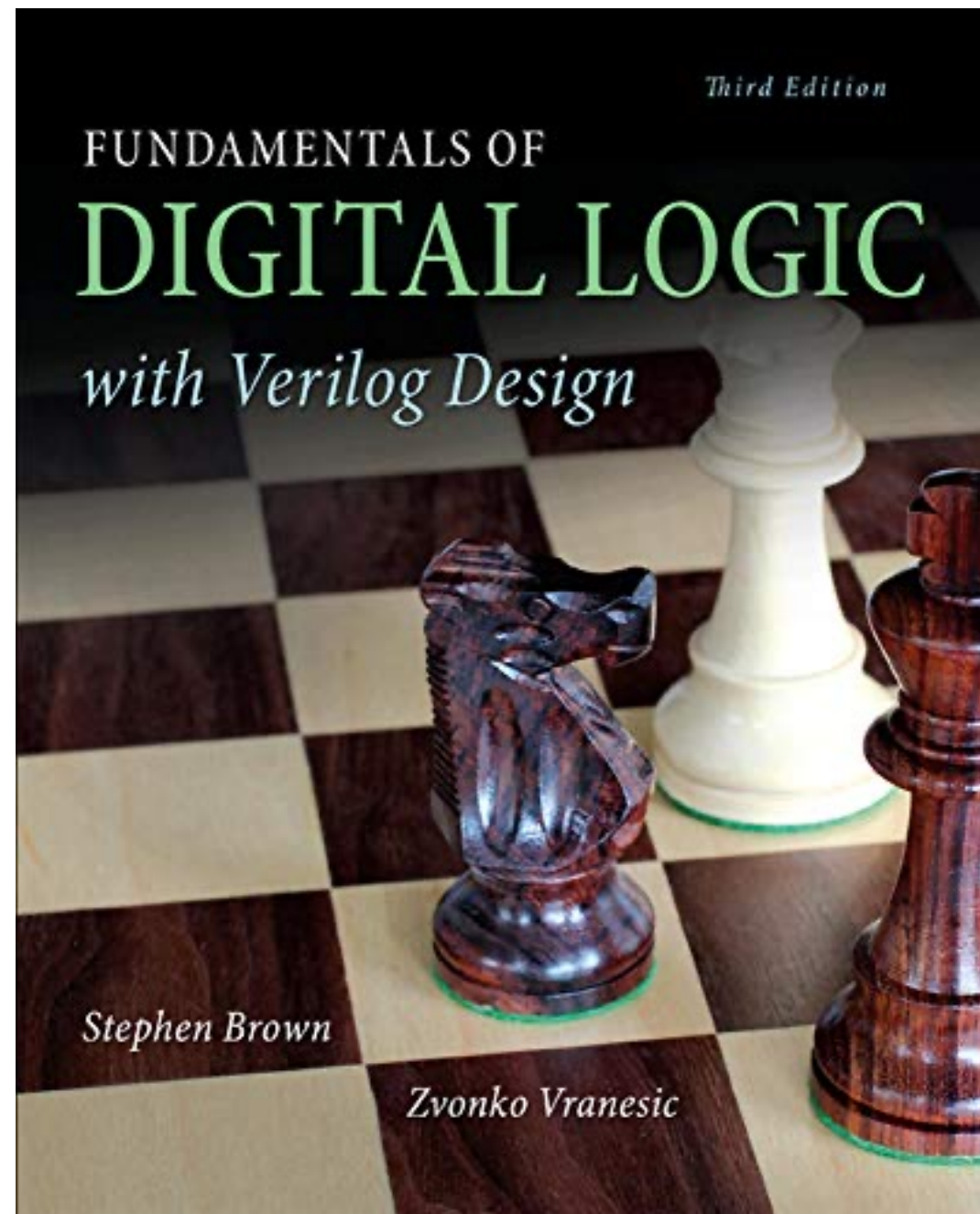
E. 4

Topics of this quarter

- Combinational Logic
 - Logic gates
 - Boolean algebra
 - K-map
- Sequential Logic
 - Finite state machines
 - Clock
 - Flip-flops
- Datapath Components
 - Adder, mux, multipliers, comparator, encoder, decoder
 - Registers, shifters, counter
- HLSM (High-Level State Machine)
- RTL (register transfer level) Design
- Verilog programming in Labs

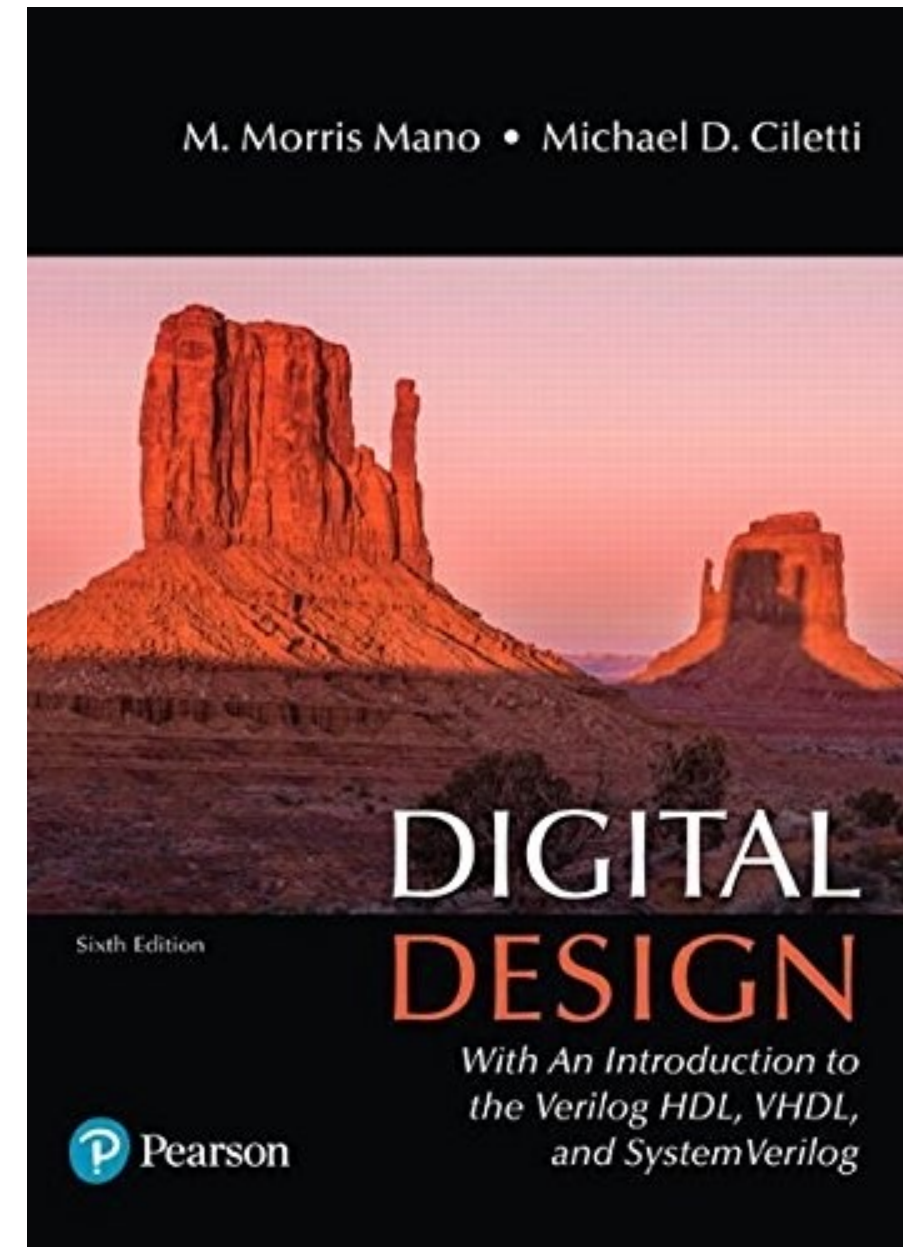
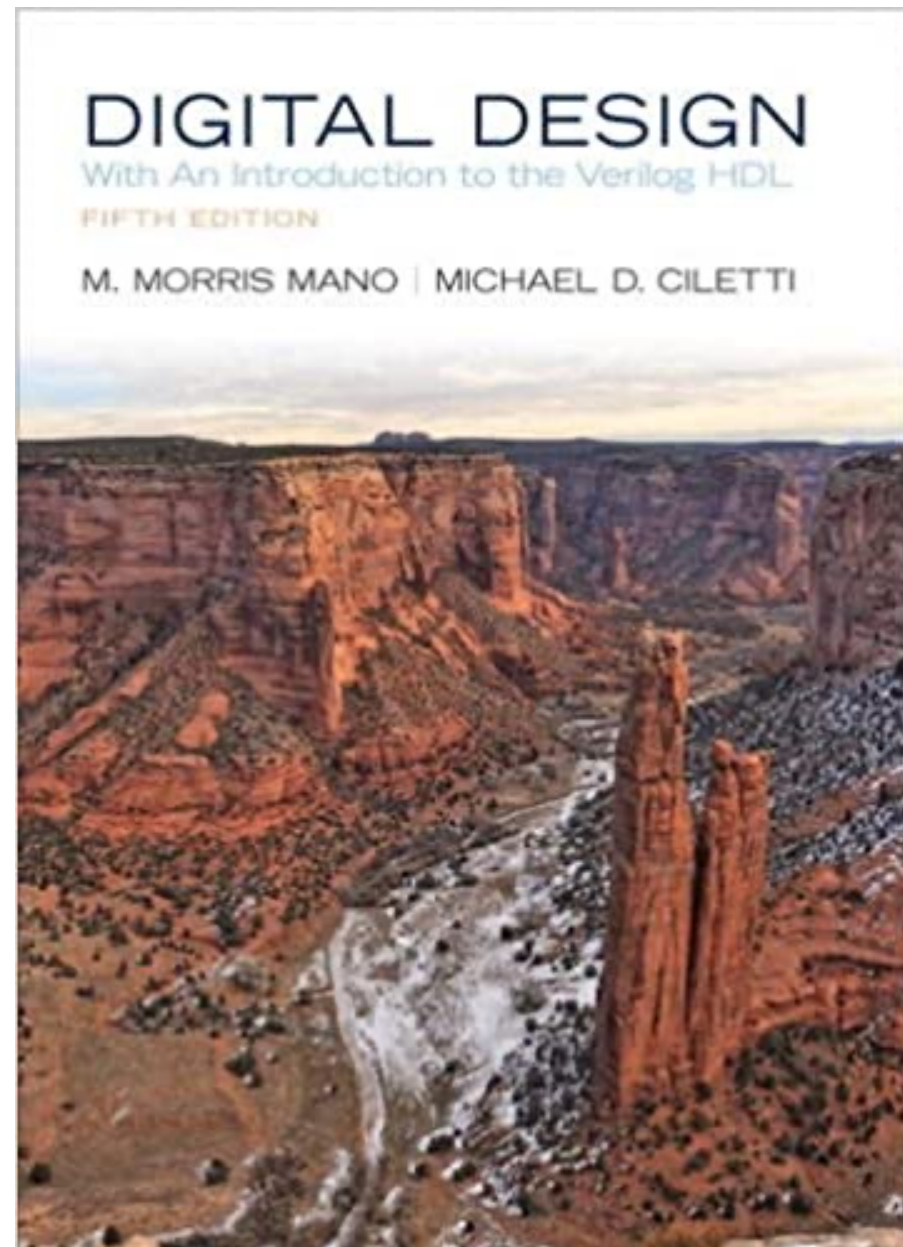
Textbook references

S. Brown and Z. Vranesic, Fundamentals of Digital Logic with Verilog Design, McGraw Hill, Third Edition, ISBN 978-0-07-338054-4.



Textbook references

M. Morris Mano and Michael D. Ciletti, Digital Design with an Introduction to the Verilog HDL, PEARSON, Fifth Edition, ISBN-13: 978-0-13-277420-8.



Lab

- We will have 5-6 labs
 - Using Verilog
 - Using simulation tools to verify and evaluate your design
- 3 students per lab group
- One lab report submission per group per assignment

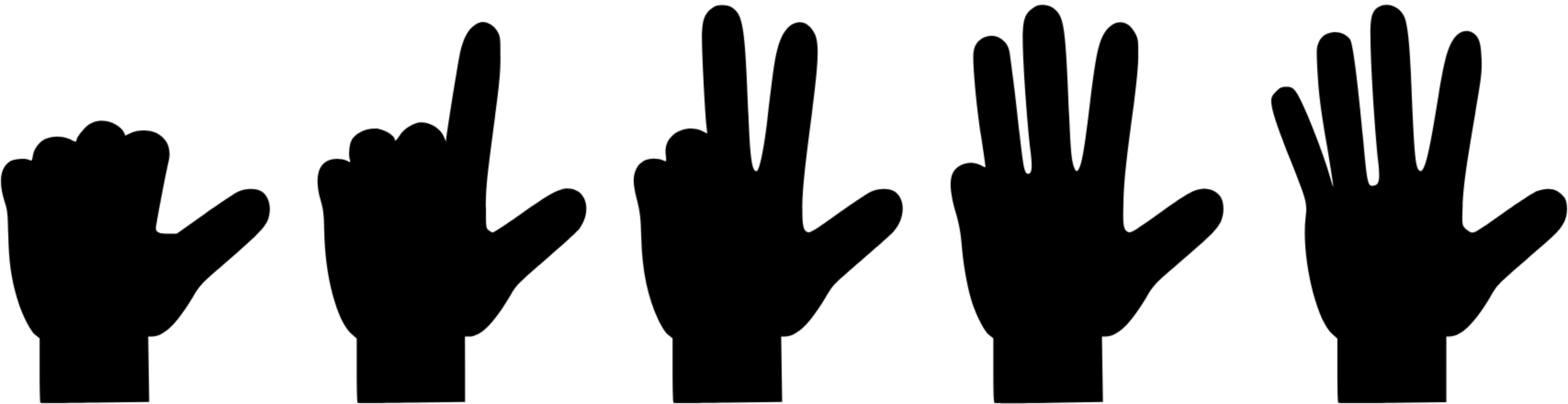
Logistics

- Lecture: MWF 1:00 – 1:50 pm
- Instructor: Jia Chen
 - Email: jia@ucr.edu
 - Office location: Bourns Hall A 149
 - Office hours: Tuesday 3:00-5:00 pm in person or by appointment.
- All the materials/announcements can be found on Canvas
- Discussion: Piazza (see signup info. in Syllabus)

Grading

- Homework: 20%
- Labs: 30%
- Midterm: 20%
- Final exam: 30%

10-based number systems is the human-nature



10-based number system is popular since thousands of years ago

1: |

10: ∩

100: ρ

1000: ⌋

10000: ⌋

⌋⌋ ∩ ∩ | | | | = ?

100000: ⌋

1000000: ⌋



10-based number system is popular since thousands of years ago

1: |

10: ∩

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1000: ⌋

10000: ⌋

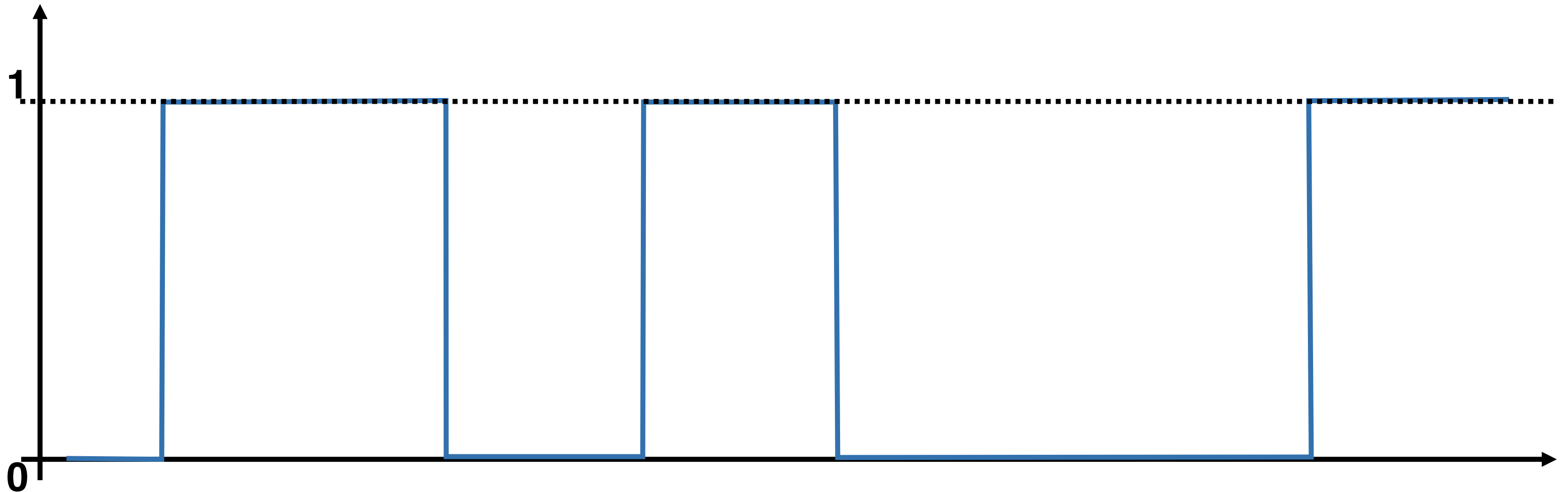
⌋⌋ ∩ ∩ | | | | = 2024

100000: 🦎

1000000: 🦎



But digital circuits only have 0s and 1s...



Connection to CS 061

Machine Organization and Assembly Language Programming

Survey

Have you taken CS 061?

Survey

Are you familiar with unsigned and signed numbers?

Review: unsigned binary numbers

The basic idea of a number system

- Each position represents a quantity; symbol in position means how many of that quantity

$$\begin{array}{r} 10^2 \quad 10^1 \quad 10^0 \\ \times \\ \underline{3} \quad + \quad \underline{2} \quad + \quad \underline{1} \quad = 300 \\ + 20 \\ + 1 \\ \hline = 321 \end{array}$$

- Decimal (base 10)

- Ten symbols: 0, 1, 2, 3, 4, 5, 6, 7, 8, 9
- More than 9: next position
- Each position is incremented by power of 10

- Binary (base 2)

- Two symbols: 0, 1
- More than 1: next position
- Each position is incremented by power of 2

$$\begin{array}{r} 2^3 \quad 2^2 \quad 2^1 \quad 2^0 \\ \times \\ \underline{1} \quad + \quad \underline{0} \quad + \quad \underline{0} \quad + \quad \underline{1} \quad = 1 \quad \times \quad 2^3 \\ + 1 \quad \times \quad 2^0 \\ = 1 \quad \times \quad 8 \\ + 1 \quad \times \quad 1 \\ \hline = 9 \end{array}$$

Converting from decimal to binary

2		321	
2		160 1
2		80 0
2		40 0
2		20 0
2		10 0
2		5 0
2		2 1
		1 0
		

$$321 = (101000001)_2$$

Other frequently used number systems

- Octal — base of 8

- 8 symbols: 0, 1, 2, 3, 4, 5, 6, 7

- More than 7: next position

$$321 = (101000001)_2$$

- Each position is incremented by power of 8

$$321 = (101\ 000\ 001)_2$$

- Easy conversion from binary — merge 3-digit into one

$$= (5\ 0\ 1)_8$$

- Hexadecimal — base of 16

- 16 symbols: 0, 1, 2, 3, 4, 5, 6, 7, 8, 9, A, B, C, D, E, F

- More than 15: next position

$$321 = (1\ 0100\ 0001)_2$$

- Each position is incremented by power of 16

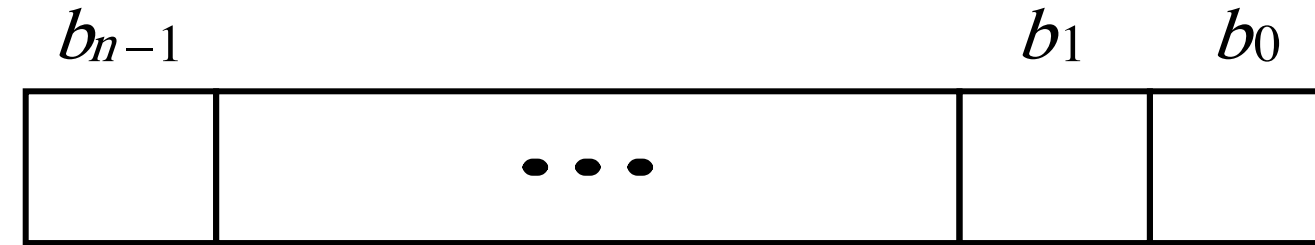
$$= (1\ 4\ 1)_{16}$$

- Easy conversion from binary — merge 4-digit into one

Review: signed binary numbers

Unsigned numbers vs signed numbers

(a) Unsigned number

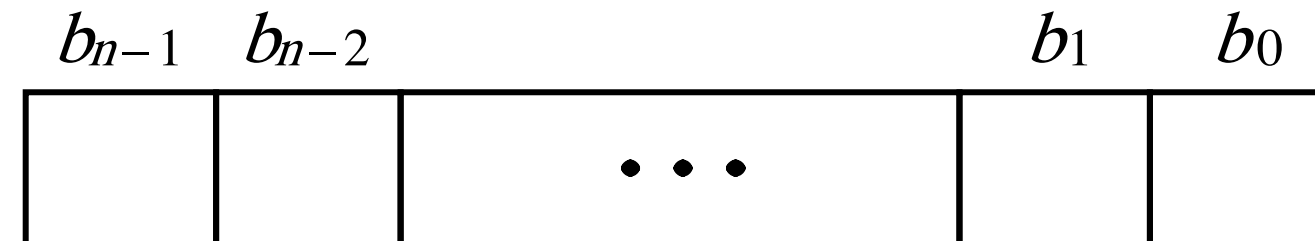


Magnitude

MSB

(most significant bit)

(b) Signed number



Sign

0 denotes +

1 denotes -



MSB

Magnitude

Three types of signed number representations

- Sign-and-magnitude
- 1's complement
- 2's complement

Positive Numbers

The three representations are the same.

Ex: signed number $0110_2 = 4 + 2 = 6_{10}$



“0” means positive

Negative Numbers

Three representations are different.

- Sign-and-magnitude
- 1's complement
- 2's complement

Sign-and-Magnitude

- ❑ Straightforward:

- Sign bit = 0 positive; sign-bit = 1 negative

- The rest bits indicate the magnitude (like unsigned number)

- For example

$$+5 = 0101$$

$$-5 = 1101$$

- ❑ Easy for us, but not suited for efficient operations in computers.

1's Complement

- ❑ Positive number: same as sign-and-magnitude positive number.
- ❑ Negative number: **complement each bit** of the corresponding positive number, including the sign bit.
- ❑ For example

$$\begin{array}{rcl} +5 = & 0101 & \\ & \downarrow \downarrow \downarrow \downarrow & \text{complement each bit} \\ -5 = & 1010 & \end{array}$$

2's Complement

- ❑ Positive number: same as sign-and-magnitude and 1's complement positive number.
- ❑ Negative number: complement each bit of the corresponding positive number, including the sign bit, and then add it with 1.
- ❑ For example

$$\begin{array}{r} +5 = 0101 \\ \quad \downarrow \downarrow \downarrow \downarrow \text{complement each bit} \\ \quad 1010 \\ + \quad \quad \quad 1 \text{ add it with 1} \\ \hline -5 = 1011 \end{array}$$

2's Complement

- ❑ A simple rule to find the 2's complement of a negative number:
 - *Examine the bits of the corresponding positive number from the right to the left, **copy** all bits that are **0s** and the **first** bit that is **1**, then **complement the rest of the bits**.*

Using the rule: +5 = 0101

-5 = 1011

 copy

Using the rule: +2 = 0010

-2 = 1110


 complement

2's Complement Addition

$$\begin{array}{r} (+5) \\ + (+2) \\ \hline (+7) \end{array} \quad \begin{array}{r} 0101 \\ + 0010 \\ \hline 0111 \end{array}$$

$$\begin{array}{r} (-5) \\ + (+2) \\ \hline (-3) \end{array} \quad \begin{array}{r} 1011 \\ + 0010 \\ \hline 1101 \end{array}$$

$$\begin{array}{r} (+5) \\ + (-2) \\ \hline (+3) \end{array} \quad \begin{array}{r} 0101 \\ + 1110 \\ \hline 10011 \end{array}$$


ignore

$$\begin{array}{r} (-5) \\ + (-2) \\ \hline (-7) \end{array} \quad \begin{array}{r} 1011 \\ + 1110 \\ \hline 11001 \end{array}$$


ignore

- ❑ If there is a **carry-out** from the sign-bit position, it is simply **ignored**.
- ❑ The 2's complement notation is highly **suitable for addition**.

2's Complement Subtraction

□ Easy way: find the 2's complement of the subtrahend, and add

$$\begin{array}{r} (+5) \\ - (+2) \\ \hline (+3) \end{array} \quad \begin{array}{r} 0101 \\ - 0010 \\ \hline \end{array} \Rightarrow \begin{array}{r} 0101 \\ + 1110 \\ \hline 10011 \end{array}$$

↑
ignore

$$\begin{array}{r} (-5) \\ - (+2) \\ \hline (-7) \end{array} \quad \begin{array}{r} 1011 \\ - 0010 \\ \hline \end{array} \Rightarrow \begin{array}{r} 1011 \\ + 1110 \\ \hline 11001 \end{array}$$

↑
ignore

$$\begin{array}{r} (+5) \\ - (-2) \\ \hline (+7) \end{array} \quad \begin{array}{r} 0101 \\ - 1110 \\ \hline \end{array} \Rightarrow \begin{array}{r} 0101 \\ + 0010 \\ \hline 0111 \end{array}$$

$$\begin{array}{r} (-5) \\ - (-2) \\ \hline (-3) \end{array} \quad \begin{array}{r} 1011 \\ - 1110 \\ \hline \end{array} \Rightarrow \begin{array}{r} 1011 \\ + 0010 \\ \hline 1101 \end{array}$$

Arithmetic Overflow

❖ For n bits, if the result is **not** in the range of -2^{n-1} to $2^{n-1} - 1$, arithmetic **overflow** has occurred.

$$\begin{array}{r} (+7) \quad 0111 \\ + (+2) \quad +0010 \\ \hline (+9) \quad 1001 \end{array}$$

Overflow $c_4 = 0$
 $c_3 = 1$

$$\begin{array}{r} (-7) \quad 1001 \\ + (+2) \quad +0010 \\ \hline (-5) \quad 1011 \end{array}$$

$c_4 = 0$
 $c_3 = 0$ **No overflow**

$$\begin{array}{r} (+7) \quad 0111 \\ + (-2) \quad +1110 \\ \hline (+5) \quad 10101 \end{array}$$

No overflow $c_4 = 1$
 $c_3 = 1$

$$\begin{array}{r} (-7) \quad 1001 \\ + (-2) \quad +1110 \\ \hline (-9) \quad 10111 \end{array}$$

$c_4 = 1$
 $c_3 = 0$ **Overflow**

- ❖ When the numbers have **opposite** signs, there is **no overflow**.
- ❖ When the numbers have the **same** sign, **overflow might** occur.
- ❖ When **carry-outs** from MSB (most significant bit) and sign-bit have **different** values, there is **overflow**.

$$\text{overflow} = c_{n-1} \oplus c_n$$

$$\begin{array}{r}
 (+7) \quad 0 \ 1 \ 1 \ 1 \\
 + (+2) \quad + 0 \ 0 \ 1 \ 0 \\
 \hline
 (+9) \quad 1 \ 0 \ 0 \ 1 \\
 c_4 = 0 \\
 c_3 = 1
 \end{array}$$

$$\begin{array}{r}
 (-7) \quad 1 \ 0 \ 0 \ 1 \\
 + (+2) \quad + 0 \ 0 \ 1 \ 0 \\
 \hline
 (-5) \quad 1 \ 0 \ 1 \ 1 \\
 c_4 = 0 \\
 c_3 = 0
 \end{array}$$

$$\begin{array}{r}
 (+7) \quad 0 \ 1 \ 1 \ 1 \\
 + (-2) \quad + 1 \ 1 \ 1 \ 0 \\
 \hline
 (+5) \quad 1 \ 0 \ 1 \ 0 \ 1 \\
 c_4 = 1 \\
 c_3 = 1
 \end{array}$$

$$\begin{array}{r}
 (-7) \quad 1 \ 0 \ 0 \ 1 \\
 + (-2) \quad + 1 \ 1 \ 1 \ 0 \\
 \hline
 (-9) \quad 1 \ 0 \ 1 \ 1 \ 1 \\
 c_4 = 1 \\
 c_3 = 0
 \end{array}$$

- Carry-outs from MSB and sign-bit have **different** values -> **overflow**

$$\text{overflow} = C_{n-1} \oplus C_n$$

- Alternative way to check overflow: if both summands have the same sign but the resulting sum has a different sign.

$$\text{EX: } X = x_3x_2x_1x_0, Y = y_3y_2y_1y_0, S = X + Y$$

$$\text{overflow} = x_3y_3\bar{s}_3 + \bar{x}_3\bar{y}_3s_3$$