

Abstractions

- Abstraction: A concept in computer science where unnecessary details are hidden to focus on the essential features.
- Purpose: Simplifies complex systems by providing a general view, allowing us to concentrate on broader goals without distraction.
- Illustration Example:
- If we represent "birds" abstractly, we refer to general characteristics without specifying individual features.
- This helps us understand the concept of birds in general, rather than learning about one specific type.
- Everyday Computing:
- We use computers without understanding or building the hardware and operating systems from scratch.
- Developers rely on libraries instead of writing complex functionalities themselves, which speeds up development and makes work manageable.
- Benefits in Programming:
- Libraries: Offer ready-to-use functions while hiding implementation details, making development more efficient.
- Reusability: Well-designed abstractions can be used across multiple projects, saving time and effort.

Database Abstraction

Transactions:

- Provide a simple interface: **commit** for success or **abort** for failure.
- Ensure data moves from one consistent state to another, even with multiple concurrent operations.

Purpose:

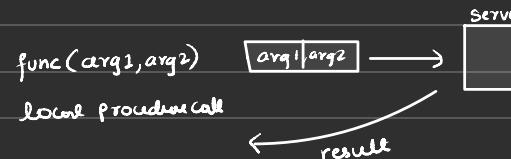
- Users don't have to worry about the complex issues of data conflicts or inconsistencies.
- Allows focusing on business logic rather than the technical details of concurrent data changes.

Distributed systems

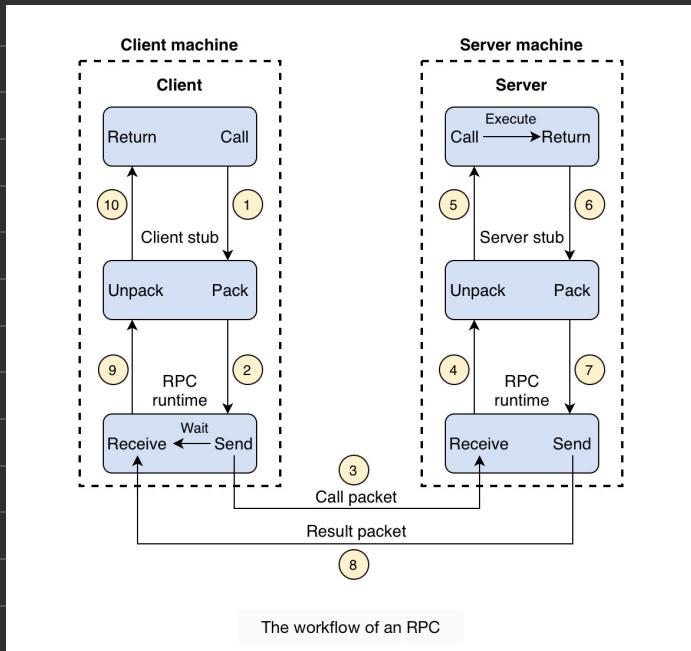
- Abstractions in Distributed SystemYs: Simplify the complexities of managing distributed environments.
- Purpose:
- Allow engineers to focus on application development instead of dealing with complex distributed system mechanics.
- Use in Industry:
- Popular in services like Amazon AWS, Google Cloud, and Microsoft Azure.
- These platforms provide different abstraction levels, hiding implementation details from users.
- Scalability:
- Essential for today's applications, which need to handle large numbers of users.
- Abstractions enable quick and efficient scaling by transitioning to distributed systems.

② RPC :

Remote procedure calls (RPCs) provide an abstraction of a local procedure call to the developers by hiding the complexities of packing and sending function arguments to the remote server, receiving the return values, and managing any network retries.



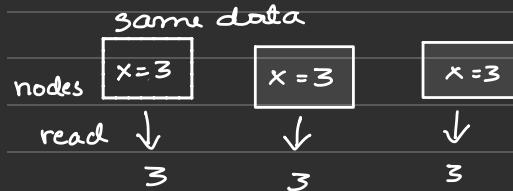
- RPC is called on a different machine



5 main parts

- client
- server
- client stub
- server stub
- RPC runtime

3 CONSISTENCY : strong
weak

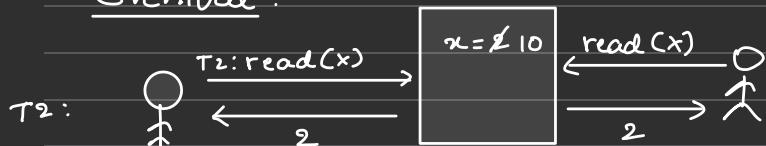


• ACID Consistency:

- Ensures that database rules are followed to maintain data integrity.
- Examples:
 - **Uniqueness Rule:** If a schema requires a unique value, the system ensures it remains unique across all actions.
 - **Foreign Key Rule:** If deleting a base row requires deleting associated rows, the system enforces this consistently.

- CAP Consistency:
 - Ensures all replicas in a distributed system have the same logical value.
 - Logical Guarantee: Values appear consistent to clients, even if physical replication takes time (e.g., due to network delays).
 - Client Access Control: Prevents clients from seeing different values across nodes, maintaining a coherent view of the data.

Eventual:



at T3: $x = 10$ is updated

- no strict ordering
- Write heavy: never converges
- all nodes have diff replicas

The domain name system is a highly available system that enables name lookups to a hundred million devices across the Internet. It uses an eventual consistency model and doesn't necessarily reflect the latest values

causally

Causal. — Dependent: related op
` Independent

Process:

$$\begin{aligned} x &= a \\ b &= x+5 \\ y &= b \end{aligned}$$

(P1) write(x) a
(P2)

since $b = x + 5$,

read(x) and write(y) are causally related

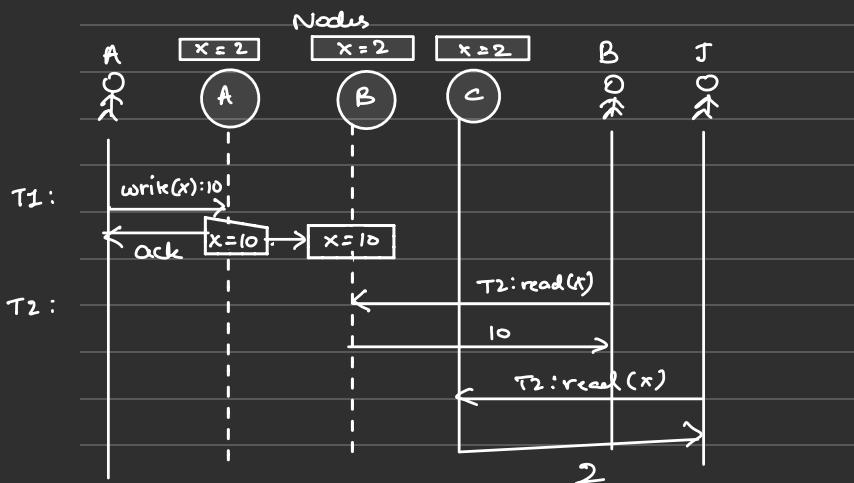
eg: commenting.

Sequential: stronger than causal

- preserves order
- same order as global clock X

Linearizability aka Strict consistency:

- read = latest value
- acknowledge.
- challenge in Distri system



- worse performance
- Strong assurance
- sear: reads time order

(4) FAILURE SPECTRUM :

easy → difficult

Fail Stop Crash Omission Temporal Byzantine



Fail stop:

- Halts permanently
- other nodes detect

Crash:

- node halts silently
- other nodes can't detect non working

Omission:

- node can't send/receive messages.
- send omission failure
or
receive omission failure

Temporal :

- too late to be useful
- bad algo / bad design
- loss of synch clocks

Byzantine :

- random behaviour

e.g.: transmit arbitrary msg
wrong results , stop midway