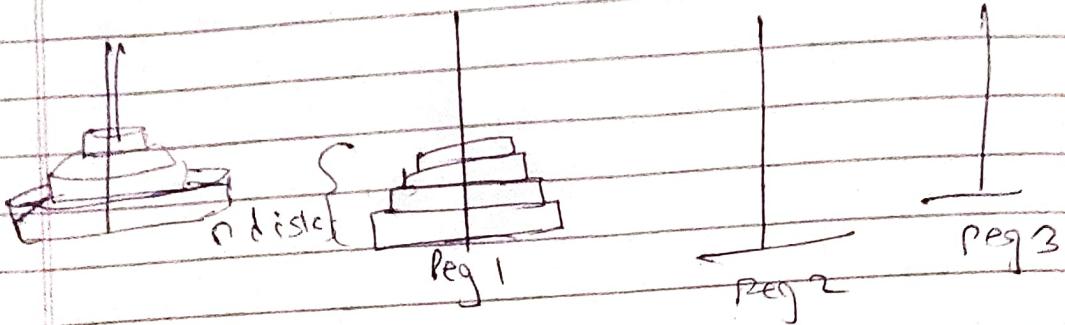


"Recurrence" relations

The Tower of Hanoi puzzle



A valid move step allows to move a disk from one peg to another ensuring that a disk is never placed on top of a smaller disk.
question: What is the minimum of steps required to move all n disks from peg 1 to peg 2?

- one disk ✓
- two disks

$$P_1 \rightarrow P_3 \quad P_1 \rightarrow P_2 \quad P_3 \rightarrow P_2$$

Let a_n be the answer for n disks

$$P_1 \xrightarrow{n} P_2$$

$$P_1 \xrightarrow{n+1} P_3$$

$$a_{n+1}$$

$$P_1 \rightarrow P_2$$

largest
disk

$$P_3 \xrightarrow{n-1} P_2$$

$$a_{n-1}$$

$$a_n = a_{n-1} + a_{n-1}$$

$$= 2a_{n-1} + 1$$

a_6

$$a_1 \ a_2 \ a_3 \ a_4 \ a_5 \ a_6 \\ | \quad 3 \quad 7 \quad 15 \quad 31 \quad 63 \\ a_1 = 1$$

$\boxed{n \geq 2}$

$\boxed{a_1 = 1}$

$$a_n = 2a_{n-1} + 1$$

$$= 2[2a_{n-2} + 1] + 1$$

$$= 2^2 a_{n-2} + 2 + 1$$

$$= 2^3 a_{n-3} + 2^2 + 2 + 1$$

$$= 2^{n-1} a_1 + 2^{n-2} + \dots + 2 + 1$$

$$= 2^{n-1} + 2^{n-2} + \dots + 1$$

$$= \frac{2^n - 1}{2 - 1} = 2^n - 1 \quad n \geq 2$$

$\boxed{\forall n \geq 1 \quad a_n = 2^n - 1}$

- Count the # of n -length binary strings which do not have two consecutive 0's.

$$n=1 \quad \overset{\checkmark}{0}, \overset{\checkmark}{1} \quad a_1 =$$

$\underline{a_n} = \# \text{ of}$
such strings

$$n=2 \quad \overset{\checkmark}{0}1, \overset{\checkmark}{1}0, \overset{\checkmark}{1}1, \overset{\times}{0}0 \quad a_2 = 3$$

$$a_3 = \overbrace{001, 000, 010}^{\downarrow}, \overbrace{011, 100, 101}^{\downarrow}, \overbrace{110, 111}^{\downarrow} \quad a_3 = 5$$

$$a_n = [x_n : w]$$

in

$$\omega =$$



\Rightarrow

Assume ω starts with '1'.

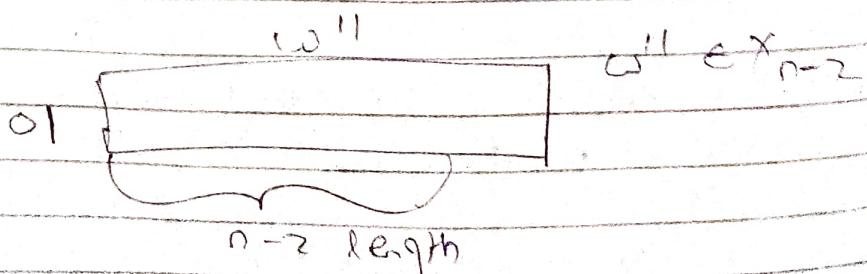
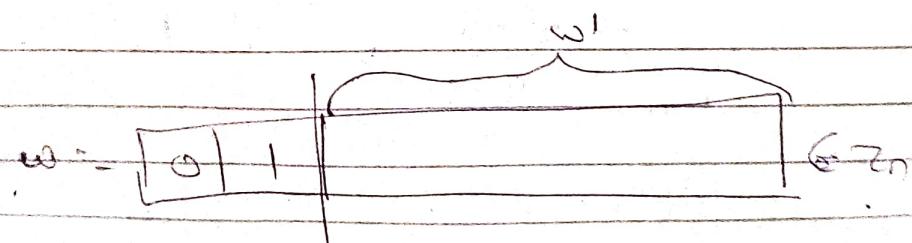
w'' of length $n-1$ and without '00'



$$a_n = (x_n) - (y_n) + (z_n) \quad Y_n \rightarrow x_{n-1}$$

$$1w'' \subset w$$

$$= a_{n-1} +$$



$$z_n \rightarrow x_{n-2}$$

$$n \geq 3$$

$$\left| \begin{array}{l} a_1 = 2, \quad a_2 = 3 \\ n \geq 3 \end{array} \right.$$

$$a_n = a_{n-1} + a_{n-2}$$

$$a_0 = 1$$

$n=0$ (E)



Fibonacci recurrence relation.

- Count the number of regions in which a plane gets divided by n straight lines in "general position" (No two lines are parallel)

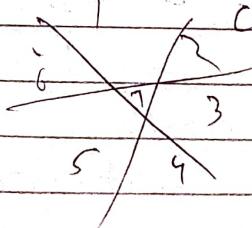
(No two lines are collinear & concurrent.)

$$a_1 = 2$$

$$a_2 = 4$$

$$a_3 = 7$$

For general situation you should get maximum possible count.



n^{th} line

1^{st}
 2^{nd}
 3^{rd}
 4^{th}
 5^{th}

n^{th} line cuts previous $n-1$ lines.

$n-1$ intersections

n^{th} line hits

last 3 lines.
so segments

So c segments get spliced

$$= a_{n-2} + n - \boxed{a_n + n}$$

$$= a_{n-2} + n - 1 + n$$

$$= a_0 + 2 + 3 + \dots + n$$

$$= 2 + 2 + 3 + \dots + n$$

$$= \frac{(1+n)(n+1)}{2}$$

$$= \frac{n^2+n+2}{2}$$

$$\forall n \geq 1 \quad c_n = \frac{n^2+n+2}{2}$$

Count the no. of ways to parenthesize a product of all numbers x_0, x_1, \dots, x_n to decide the order of multiplication.

$$n=1 \quad x_0, x_1,$$

$$n=2 \quad x_0, x_1, x_2$$

$$(m_0, x_1) \cdot x_1 \rightarrow (x_1, x_2)$$

$$n=2 \quad x_0$$

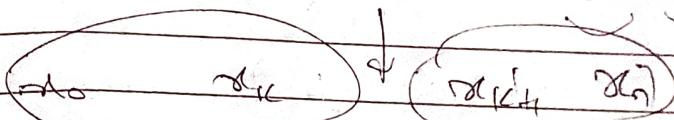
$$a_0 = 1$$

$$a_1 = 1$$

$$a_2 = 2$$

$$a_3 = 6$$

$$n=3 \quad x_0, x_1, x_2, x_3$$



$$(2^k) \quad n-k$$

$$\sum a_{ik} \cdot 2^{n-k-1}$$

R =

Catalan numbers

Also missing notes of last wednesday

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Linear homogeneous recurrence relation of degree
k with "constant" co-effs.

Assume $k=2$:

$$\sum_{n=2}^{\infty} a_n = c_1 c_n + c_2 n + c_3 \neq 0$$

Let $A(x) = \sum_{n=0}^{\infty} a_n x^n$ be the generating

function of (a_n)

$$\sum_{n=2}^{\infty} a_n x^n = \sum_{n=2}^{\infty} (c_1 c_n + c_2 n + c_3) x^n = c_1 \sum_{n=2}^{\infty} c_n x^n + c_2 \sum_{n=2}^{\infty} n x^n + c_3 \sum_{n=2}^{\infty} x^n$$

$$A(x) = a_0 - a_1 x - \dots + c_1 [A(x) - a_0] + \\ (c_2 \text{ or } A(x))$$

$$A(x) [1 - c_1 x - c_2 x^2] = a_0 \quad A_{\text{fact}} \text{ or } A_{\text{fact}}$$

$$A(x) = \frac{P(x)}{1 - c_1 x - c_2 x^2} \quad \deg P(x) \leq 1$$

$$\begin{aligned} & 1 - c_1 x - c_2 x^2 = r_1 x + r_2 x^2 \\ & x^2 - c_1 x - c_2 = (x - r_1)(x - r_2) \\ & (1 - r_1 x)(1 - r_2 x) \quad \therefore 1 - c_1 x - c_2 x^2 \\ & = (1 - r_1 x)(1 - r_2 x) \\ & = \frac{d_1}{1 - r_1 x} + \frac{d_2}{1 - r_2 x} \end{aligned}$$

$$a_n = \text{coeff. of } x^n \text{ of } A(x) = d_1 r_1^n + d_2 r_2^n$$

$$\frac{1}{1 - r_1 x} = 1 + r_1 x + r_1^2 x^2 + r_1^3 x^3 + \dots$$

$$\text{Ex: } x^2 - x - 1 =$$

$$x = \frac{1 \pm \sqrt{1+4}}{2} = \frac{1 \pm \sqrt{5}}{2}$$

$$a_n = d_1 \left(\frac{1+\sqrt{5}}{2} \right)^n + d_2 \left(\frac{1-\sqrt{5}}{2} \right)^n$$

$$1 - c_1 x - c_2 x^2 = (1 - r_1 x)(1 - r_2 x)$$

$$x^2 - c_1 x - c_2 = (x - r_1)(x - r_2)$$

$$A(x) = \frac{P(x)}{(1 - r_1 x)^2} = \frac{d_1}{(1 - r_1 x)^2} + \frac{d_2}{(1 - r_1 x)}$$

Since all constants

$$\frac{1}{(1-y)^2} = \frac{1}{(1-y)(1-y)}$$

$$= \sum_{k=0}^{\infty} (-1)^k (k+1) y^k$$

$$= (1+y) + ky^2$$

$$= (1+y) + ky^2$$

$$= \sum_{k=0}^{\infty} (-1)^k (k+1) y^k$$

$$= \sum_{k=0}^{\infty} (k+1)y^k$$

$$\frac{1}{(1-y)^2} = 1 + y + y^2 + \dots$$

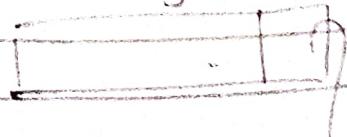
$$\frac{1}{(1-y)^2} = \frac{1}{(1-y)} + \frac{1}{(1-y)}$$

* Linear inhomogeneous recurrence relations of degree k with "constant" coeffs

A valid codeword is a n-digit number in decimal notation with even number of 0's. Let a_n be the number of valid codewords.

$$a_0 = 9 \quad \text{decimal digits } 2, 4, 6, 8$$

$$a_2 = 3 \quad \vdots$$



$$a_n = 9a_{n-2} + 6^{n-1} - a_{n-1}$$

when lost
digit + 0 digit = 0

$$= 8a_{n-1} + 10^{n-1}$$

newton's method

$$\text{if } f'(x_0) \neq 0 \text{ then } x_1 = x_0 - \frac{f(x_0)}{f'(x_0)}$$

and if

if

$$a_1 \leq a_0 \text{ and } f'(a_0) \neq 0$$

$$\left\{ \begin{array}{l} a_2 = a_1 - \frac{f(a_1)}{f'(a_1)} \\ \vdots \end{array} \right.$$

$$A_{\text{new}} = \sum_{n=0}^{\infty} a_n d^n$$

$$\sum_{n=0}^{\infty} a_n d^n = \sum_{n=1}^{\infty} a_n d^{n-1} + \sum_{n=1}^{\infty} d^n$$

$$(A_{\text{old}} - a_0) = \sum_{n=1}^{\infty} a_n d^{n-1}$$

$$A_{\text{new}}(1 - d) = 1 + \frac{d}{1-d} - \frac{d^2}{1-d}$$

$$A_{\text{new}}(1 - d) = 1 + \frac{d}{1-d} - \frac{d^2}{1-d}$$

$$1 - d = \alpha(1 - d) + \beta(1 - d)$$

$$A(x) = \frac{1-q}{(1-qx)(1-xa)}$$

$$\frac{a_1}{1-qx} + \frac{a_2}{1-xa} = \frac{1}{2} \left(\frac{1}{1-qx} + \frac{1}{1-xa} \right)$$

$$a_1 = 1$$

$$a_n = a_{n-1} + a_1 a_{n-2} + a_2 a_{n-3} + \dots + a_{n-4} a_0$$

$$= \sum_{k=1}^n a_k a_{n-k}$$

$$a_1 a_2 a_3 \dots a_n$$

$$a_2 = a_1 a_1 + a_0 = 1+1=2$$

$$A(x) = a_0 + a_1 x + a_2 x^2 + a_3 x^3 + \dots$$

$$a_3 = a_1 a_2 + a_0 = 2+1=3$$

$$a_4 =$$

$$A(x) A(x), A^2(x) = (1+a_1 x + a_2 x^2 + \dots) (1+a_1 x + a_2 x^2 + \dots)$$

$$= a_1 + a_2 x + a_3 x^2 + a_4 x^3 + \dots$$

$$x(A^2(x)) = a_1 x + a_2 x^2 + a_3 x^3 + \dots$$

$$= A(x) - c_0$$

$$x(A^2(x)) - A(x) + 1 = 0$$

$$A(x)$$

$$\sqrt{x^2 - y + 1} = 0$$

$$A(x) = \frac{(1 + \sqrt{1 - 4x})}{2x}$$

$$\sqrt{(1 - 4x)}$$

$$= (1 - 4x)^{\frac{1}{2}}$$

$$= \sum_{k=0}^{\infty} \binom{\frac{1}{2}}{k} (-4)^k x^k$$

$$\sqrt{1 - 4x} = \frac{1 + \sqrt{1 - 4x}}{2x}$$

$$= \sum_{k=0}^{\infty} b_k x^k$$

$$b_k = (-4)^k \binom{\frac{1}{2}}{k}$$

$$= 1 +$$

$$b_1 = (-4)^1 \binom{\frac{1}{2}}{1}$$

$$+ 4x$$

$$= -4 \times \frac{1}{2}$$

$$+ 4x + 4x^2 + 1 + 4x^3 + \dots$$

$$= 1 - 2x - 2x^2$$

$$= -2$$

$$b_2 = (-4)^2 \binom{\frac{1}{2}}{2}$$

$$= 16 \times \frac{1}{2} \times -\frac{1}{2}$$

$$2x^2$$

$$\frac{1 - \sqrt{1 - 4x}}{2x} = 2x + 2x^2 + -2x^3 + \dots$$

$$= -$$

$$a_n = \frac{1}{n+1} \binom{2n}{n}$$

So we can find something like this

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Letting b_k be the number of little squares belonging to $\{k, l\}$ which stay below the diagonal without touching the diagonal.

$$n = \sum_{k=1}^l b_k = \sum_{k=1}^l a_{k+1}$$

$$\text{class } b_k = a_{k+1}$$

$a_0 = 1$	$a_1 = 1$
-----------	-----------

$$(0, 1) (1, 1)$$



$$a_1 = a_0 a_0 \quad a_2 = a_0 a_1 + a_0 a_0$$

$$a_3 = a_0 a_2 + a_1 a_2 + a_2 a_0 = 2 + 1 + 2 = 5$$

$$\text{let } A(n) = a_0 + a_1 n + a_2 n^2 + a_3 n^3 + \dots$$

be the g.f. for $(a_n)_{n=0}^\infty$

$$\begin{aligned} A^2(n) &= a_0 a_1 + (a_0 a_2 + a_1 a_1) n + (a_0 a_3 + a_1 a_2 + a_2 a_1) n^2 \\ &\quad + a_3 n^3 + \dots + a_2 n^2 + a_1 n + a_0 \end{aligned}$$

$$\lambda(A^2(n)) = A(n) - q_0$$

$$\boxed{\lambda(A^2(n)) - A(n) + 1 = 0}$$

~~1000~~ ~~1000~~

~~1000~~ ~~1000~~

~~1000~~

~~1000~~

$3 \times (10) (-3) = -30$

$= 20^2 \cdot (20-3) \cdot 20-5 = 531$

~~300~~ ~~300~~

$= (20-2)(20-2) \cdot 300$

$= (20-2)^2 \cdot 200 = 200$

$= (20-2)$

$\underline{2(20-2) \cdot 200 \cdot 20^2 \cdot 20}$

$\alpha = 2n$

$$\alpha = \frac{2n+1}{2} + \sqrt{\frac{2n+1}{2}^2 - 1}$$

$(2n+1)(n+1)$

$\Delta\alpha = 1 - \frac{1}{2n+1}$

$\approx \frac{1}{2n}$

$\approx \frac{1}{2n}$

$\Delta\alpha = \frac{1}{2n}$

$\approx \frac{1}{2n}$

$\approx \frac{1}{2n}$

$$\alpha = \frac{1}{2} + \frac{1}{2n+1} (2n)$$

$$\alpha = \frac{1}{2} (2n) - (2n) = (2n)$$

we will show

$$\alpha = (2n) - (2n)$$

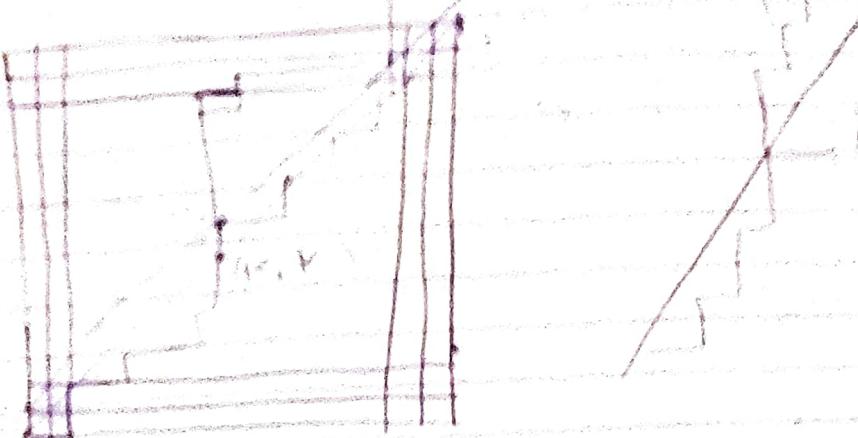
A path from $(2n)$ to $(2n)$ is good if it stays below the diagonal otherwise it is bad.

Case 1:

The number of bad paths from $(2n)$ to $(2n)$ is $\binom{2n}{n}$

5

(part 1) reflection of
path of a particle



60°



60°

A bad path when "reflected" at
the first point when it crosses
the diagonal gives it one extra path
from $(0,2)$ to $(5,1,0,1)$

60°



Catalan numbers part)

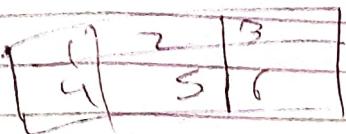
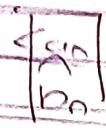
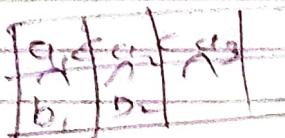
- Number of such parenthesized strings satisfying
open "()" and closed ")" brackets

56 (5)

0000000000

There is bijection between
these problems

Number of $2 \times n$ arrays with distinct entries
elements 1, 2, 3



a_i - the step-number of i th horizontal step

b_i - the step number of i th vertical step



1 2 s
3 4 6

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Lemma: Let (P, \leq) be a finite poset.
Then minimal & maximal elts exist.

def: X_{\min}

$\{x \in P : \forall y \in P, x \leq y\}$

X_{\max}

Let X_{\min} be the set of all minimal elements of P .

claim: Any two distinct elements of X_{\min} are incomparable.

If $x \leq y$ are in X_{\min} and $x \neq y$

If $x \leq y$, then $x \neq y$: y is not minimal. That contradicts the minimality of y .

Similarly,

Def: A partition of P is a subset A of P such that no two elements of A are comparable.

Ex: X_{\min} & X_{\max} are partitions



$P_1(2, 3)$ $\{1, 2\}$

$A = \{ \{1, 2\}, \{3\} \}$

$\{ \{2\}, \{1, 3\} \}$

$P_2(3)$

Let $P = C \times L$ be a poset

C - the course

$L = \{C_1, C_2, \dots, C_n\}$ | either $C_i \in C$, or
 C_i must have completed
in order to register for
 C_j

for each student is a sequence (P_1, P_2, \dots)

the set of subsets of C , such that

1) P_1, P_2, \dots is a partition of C .
2) it respects the " \subseteq "

Poset:

Given $P = C \times L$, P_1, P_2, \dots

Q: What is the minimum number of students
that can be registered?

Definition: A chain C of P is a subset of C
such that any two elements of C are incomparable
by \leq . In other words, if $x, y \in C$, then $x \neq y$ and
 $x \leq y$ and $y \leq x$ are false.



Lemma : If α is an anti-chain in X , then $|\text{Pac } \alpha| \leq c$

$c =$ the maximum size of a chain of P

$$a \geq c$$

$a =$ minimum number of anti-chains required to partition X

$c =$ the maximum size of a chain of P

$$a \geq c$$

Let $\alpha = \{A_1, A_2, \dots, A_l\}$ be a partition of X into l anti-chains and c be a chain of X of size l

Then $a = l = c$

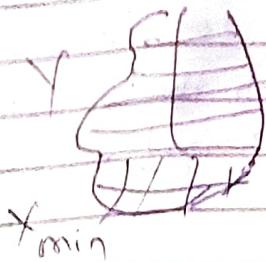
a Best antichain

Smallest number of partitions & largest size stain

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(ΔP)
chain
 $\Delta P'$

By induction Δ has chain say C of size k and an anti-chain decomposition say $C' = \{A_1, A_2, \dots, A_k\}$ using Δ -artition.



let y_0 be the minimum element of chain C . There must exist an element $x \in X_{\min}$ such that $x \leq y_0$.

Now with $C' = C \cup \{x_0\}$ and

$A'_k = \{A_1, A_2, \dots, A_k, x_{\min}\}$ we get a chain C' of P and an anti-chain decomposition of P with $|C'| = k+1 = |\Delta P'| / |\Delta P|$



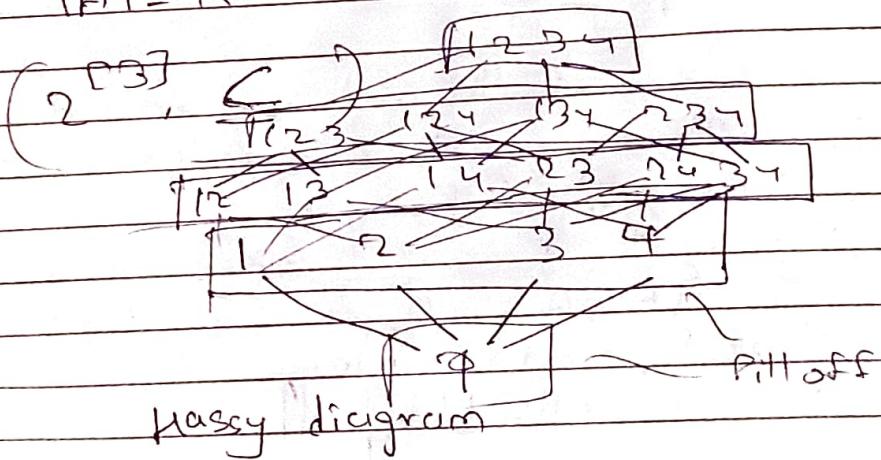
Theorem: Let $P = (\mathcal{X}, \leq)$ be a finite poset. Then \exists a k such that the size of the largest-chain of P is k and the size of the smallest anti-chain decomposition is k .

By the above theorem, $\exists c$ such that $k \leq c$
 and $cA = \{A_1, \dots, A_k\}$
 $|A| = k$

Let D be a chain of P such that $|D| > k$

Let D be a chain of P such that $|D| > k$.

* Dilworth's theorem: let $P = (X, \leq)$. Then
 there exists K such that P has an antichain
 of size K and a chain decomposition of size K .
 $|A| = K \quad C = \{C_1, C_2, \dots, C_K\}$



$$|A| \leq |C|$$

↓
size of

↑ size of chain decomposition

Antichain:

$$\max_A |A| \leq \min_C |C|$$