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Affidavit



CS 480 OPERATING SYSTEMS

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Assignment 03

Part I Due: 11:59pm, Mar 22nd

You must work on this part I on you own

Part I 6 Questions – Each has 10 points, question 5 has 15 points

You must work on your own for this part, it would be a red flag if we find submissions from two students are with exact same incorrect answers in multiple parts.

SINGLE EXAMINEE AFFIDAVIT

"I, the undersigned, promise that this assignment submission is my own work. I recognize that should this not be the case; I will be subject to plagiarism penalties as outlined in the course syllabus."

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Date: 10 Mar 2022

Question 1

1. Three processes A, B, C start at the same time in that order, each has a series of alternating CPU bursts and I/O bursts as follows:

A: 14ms (CPU), 3ms (I/O), 6ms (CPU)
 B: 4ms (CPU), 8ms (I/O), 7ms (CPU)
 C: 2ms (CPU), 1ms(I/O), 3ms (CPU), 1ms (I/O), 1ms (CPU)

What would be the **turnaround (completion) and wait times** of all processes using each of the following scheduling algorithms?

Round Robin (RR) with **5ms** time quantum,
 Shortest Job First (SJF).

For simplicity, assume context switch time and other overheads are comparatively negligible. **NOTE:**

- **Wait time** is the wait time each process spends in the process READY queue.
- **Turnaround time = CPU bursts + I/O bursts + Wait time**

Fill in the following table with your answers, and you **must show** the steps of your calculation (follow the hints below).

	Turnaround Time		Wait Time	
	RR	SJF	RR	SJF
A	50 ms	50 ms	22 ms	27 ms
B	47 ms	27 ms	23 ms	8 ms
C	28 ms	8 ms	10 ms	0 ms
Average	41.66 ms	28.33 ms	21.66 ms	11.66 ms

Hint: Use notation p (e, s, r) to track CPU execution of a process over the time:

e - total CPU burst time that has been executed for a particular process,
 s - total system time passed from the beginning of executing first process,
 r - the reason why the process is stopped along the execution (either suspended or completed) due to:

- t - time quantum expired during RR (only applicable to RR),
- i - process initiated an I/O,
- c - process completed execution
- For example: A(10, 16, t) means total CPU burst time that has been executed for process A is 10ms, total system time passed is 16ms since the beginning of executing the first CPU burst, and process A is suspended due to expiration of time quantum.

RR

A(5, 5, T)

B(4, 9, I)

C(2, 12, I)

A(10, 20, T)

B(4, 10, T)

C(4, 15, T)

A(14, 32, I)

B(4, 25, T)

C(5, 26, I)

A(14, 33, T)

B(7, 38, T)

C(6, 28, C)

A(17, 43, T)

B(11, 47, C)

A(20, 50, C)

STF

C(6, 8, C)

B(11, 27, C)

A(20, 50, C)

Question 2

2. You are designing a hard real-time autonomous driving system. It has following sensors in the inertial navigation subsystem for tracking and controlling vehicle attitude, velocity, and position:

Sensors	Sampling Frequency (Hz)	CPU time required (ms)
Accelerator	40	5
Gyroscope	25	4
Magnetometer	20	8

Suppose you also want to add wheel speed sensors to help the traffic collision analysis, each wheel speed sensor sampling frequency is 10 Hz, and the required CPU time for processing the wheel speed sensor data is 10ms.

Determine how many wheel speed sensors you could add to system to still allow system to be schedulable. (Assume the overall overhead of the inertial navigation subsystem is .25) **Show your work.**

$$\begin{array}{l}
 40 \text{ Hz} = 40 \frac{\text{cycle}}{\text{s}} \rightarrow \frac{1 \text{ s}}{40 \text{ cycles}} = .025 \text{ s/cycle} = 25 \text{ ms} \\
 25 \text{ Hz} = 25 \frac{\text{cycle}}{\text{s}} \rightarrow \frac{1 \text{ s}}{25 \text{ cycles}} = .04 \text{ s/cycle} = 40 \text{ ms} \\
 20 \text{ Hz} = 20 \frac{\text{cycle}}{\text{s}} \rightarrow \frac{1 \text{ s}}{20 \text{ cycles}} = .05 \text{ s/cycle} = 50 \text{ ms}
 \end{array}
 \left\{
 \begin{array}{l}
 \text{Acel} \quad 25 \text{ ms} \quad 5 \\
 \text{Gyro} \quad 40 \text{ ms} \quad 4 \\
 \text{Magnet} \quad 50 \text{ ms} \quad 8
 \end{array}
 \right.$$

$$\frac{5}{25} + \frac{4}{40} + \frac{8}{50} + .25 = .71 \leq 1 \quad \text{schedulable w/o new sensors} \checkmark$$

$$\text{Wheel speed sensors: } 10 \text{ Hz} = 10 \frac{\text{cycle}}{\text{s}} \rightarrow \frac{1 \text{ s}}{10 \text{ cycles}} = .10 \text{ s/cycle} = 100 \text{ ms}$$

$$\frac{C_i}{P_i} = \frac{10}{100} = .1$$

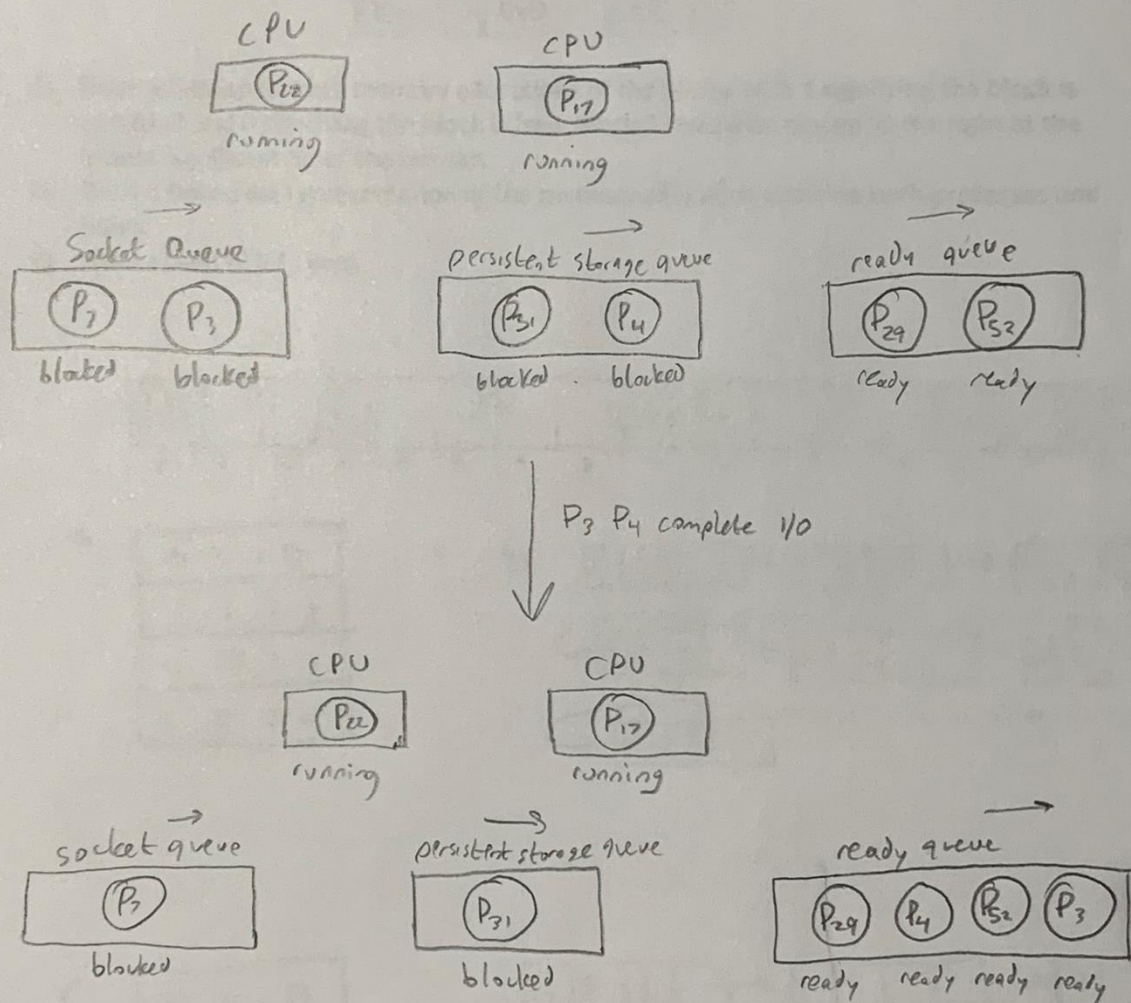
$$\frac{5}{25} + \frac{4}{40} + \frac{8}{50} + \frac{10}{100} + \frac{10}{100} + .25 = .91 \leq 1$$

2 wheel speed sensors max

Question 3

3. Recall question 4 prompt from assignment 2, we now apply priority scheduling to the system. The higher the priority number is, the higher the process's priority. Processes P22 (with priority 6) and P17 (with priority 7) are executing. P7 (with priority 6) and P3 (with priority 7, first in the queue) are waiting for a socket write to complete. P4 (with priority 5, first in the queue) and P31 (with priority 4) are waiting on persistent storage access. P52 (with priority 6, first in the queue) and P29 (with priority 4) are awaiting to be scheduled on the CPU.

Suppose P3 and P4 both complete I/O operations at the same time, what would the process queueing diagram for these processes become? (Note each type of I/O has its own waiting queue)



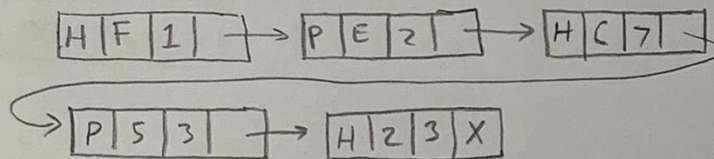
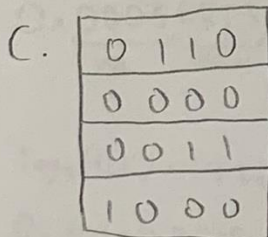
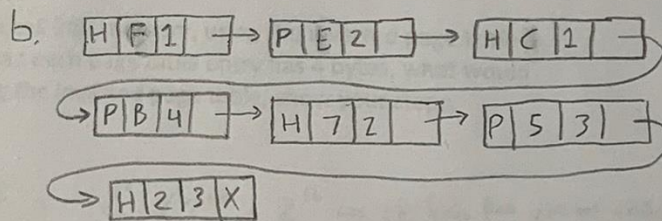
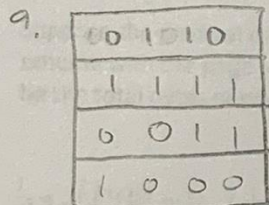
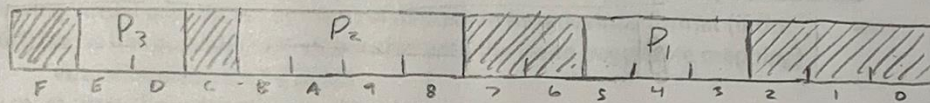
Question 4

4. Suppose a small system has a physical memory size of 16KB, which is divided into 16 blocks (each block has 1KB) for applying a contiguous memory allocation scheme. The blocks are numbered 0x0 through 0xF.

Three processes are loaded in memory for execution:

Process	Start block	End block
P ₁	0x3	0x5
P ₂	0x8	0xB
P ₃	0xD	0xE

- Draw a bitmap to track memory allocations of the blocks with 1 signifying the block is occupied and 0 signifying the block is free. Block 0 should be placed to the right at the lowest significant bit of the bitmap.
- Draw a linked list representation of the memory allocation showing both processes and holes.
- Redo a) and b) if P₂ exits.



	Page Table	
Virtual Page Number	Valid	Physical Frame Number
0x0	1	0xE
0x1	0	0x2
0x2	1	0x20B
0x3	1	0xA2
0x4	1	0x6
0x5	0	0x30
0x6	1	0x725

- $0x0020B \rightarrow 0000\ 0010\ 0000\ 0000\ 1011 \rightarrow 1011 = 1000$
 $0000\ 0010\ 0000\ 0000\ 1000\ 0111\ 0001\ 0101$
 $= 0x020008715$

3. physical space $\rightarrow 2GB = 2^{31}$
Page size $\rightarrow 8KB = 2^{13}$
of entries $\rightarrow 2^{31} / 2^{13} = 2^{18} = 256K \text{ entries}$
 $256K \times 4 \text{ bytes} = 2^{20} \text{ bytes}$

Question 6

6. On a specific architecture, each physical memory access time costs 20 ns. Suppose a TLB access requires 3 ns and has a miss rate of 0.03. If the MMU is capable of accessing a 3-level page table with negligible overhead other than the memory access times, what is the effective access time? Show your steps.

access time: 20 ns

page table access time: 20 ns

hit rate: .97

TLB access time: 3 ns

$$.97(3 + 20) + .03(3 + 3(20) + 20)$$

$$= 24.8 \text{ ns}$$