

Q7(d) (i) Suppose a computer spends 60% of its time handling a particular type of computation when running a given program and its manufacturers make a change that improve its performance on that type of computation by a factor of 10. If the program takes 100 sec to execute. what will its execution time be after the change?

→ The execution time of program before change = 100 sec  
 it spends 60% of its time, handling a particular type of computation = 60% of 100 sec = 60 sec.

∴, Remaining time for execution =  $(100 - 60) = 40$  sec.

Manufacturers improves its computation by a factor of 10 by some change ⇒ which increases its efficiency by a factor of 10 and hence decreases the computation time by factor of 10, i.e.,  $\frac{60}{10} = 6$  sec.

Now, Execution time after the change =  $40 + 6 = 46$  sec.

(ii) If  $A \oplus B = AB' + A'B$ , find the value of  $x \oplus y \oplus z$ .

$$\begin{aligned} \Rightarrow x \oplus y \oplus z &= (xy' + x'y) \oplus z = xy' \oplus z + x'y \oplus z \\ &= xy'z' + (xy')'z + x'yz' + (x'y)'z \\ &= xy'z' + (y'' + x')z + x'yz' + (x'' + y')z \\ &= z'(xy' + x'y) + z(y + x' + x + y') \quad [\because A'' = A] \\ &= z'(xy' + x'y) + z\{(x + x') + (y + y')\} \\ &= z'(xy' + x'y) + z(1 + 1) \quad [\because A + A' = 1] \\ &= z'(xy' + x'y) + z \quad [\because 1 + 1 = 1] \end{aligned}$$

$$\therefore x \oplus y \oplus z = z'(xy' + x'y) + z$$

Q7(c) (i) Find the Hexadecimal equivalent of the decimal number  $(587632)_{10}$ .

$\Rightarrow$ 

16	587632	
16	36727	0
16	2295	7
16	143	7
	8	15

 $\uparrow$ 
 $\therefore 8 \ 15 \ 7 \ 7 \ 0$   
 $= (8F770)_{16}$

So,  
 $(587632)_{10} = (8F770)_{16}$

7) (c) (ii) For the given set of data points  $(x_1, f(x_1)), (x_2, f(x_2)), \dots, (x_n, f(x_n))$  write an algorithm to find the value of Lagrange's interpolation formula.

$\Rightarrow$  Step 1. Start the program.

Step 2. Input number of terms  $n$ .

Step 3. Input the array  $ax$ .

Step 4. Input the array  $ay$ .

Step 5. for  $i=0; i < n; i++$

Step 6.  $nx = 1$

Step 7.  $dx = 1$

Step 8. for  $j=0; j < n; j++$

Step 9. if  $j \neq i$

a.  $nx = nx * (x - ax[j])$

b.  $dx = dx * (ax[i] - ax[j])$

Step 10. End Loop  $j$

Step 11.  $y+ = (nx/dx) * ay[i]$

Step 12. End Loop  $i$

Step 13. print output  $x, y$ .

Step 14. End of the program.



Q(c) (iii) use Boolean algebra, simplify the following expressions

⇒ i)  $a + a'b + a'b'c + a'b'c'd + \dots$

ii)  $x'y'z + yz + xz$ , where  $x'$  is complement of  $x$

⇒ i)  $y = a + a'b + a'b'c + a'b'c'd + \dots$

$$= (a + a'b) + a'b'(c + c'd) + \dots$$

$$= (a + b) + a'b'(c + d) + \dots$$

$$[\because A + A'B = A + B]$$

$$= (a + a'b'c) + (b + a'b'd) + \dots$$

$$= (a + b'c) + (b + a'd) + \dots$$

$$= (a + a'd) + (b + b'c) + \dots$$

$$= a + d + b + c + \dots$$

$$= a + b + c + d + \dots$$

ii) Given,

$$A = x'y'z + yz + xz$$

$$= (x'y' + y)z + xz$$

$$= (x' + y)z + xz$$

$$[\because A'B + A = A + B]$$

$$= x'z + yz + xz$$

$$= yz + (x + x')z$$

$$= yz + z$$

$$[\because A + A' = 1]$$

$$= z(1 + y)$$

$$= z$$

$$[\because 1 + A = 1]$$