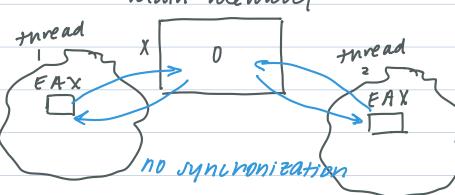
CS 61 Lecture 24	Nov 21, 2019
Concruency Overview	
1 execution unit -> concurr	ency by MULTIPROBRAMMINE
History	
· prolition / 1 program	
· single provisor / many p	WIWI
Many processors / many processes TODAY -> many processors / one process	
· VKr-14WI abstraction of multiple execution resources	
· DM pouls can have multiple	
· NIId More: cache, regist	Cimy Hampusla
pwyss	thread
· individualized view of	
hardware resources	
· isolation: individual views	· shared memory
of numbry, separah addr spairs	running on top of virtual memory
· fix abstraction	/
C++ s+d:: thread	
· Maintain process isolation	_ method defining
· Maintain process isolation · std:: thread (threadfunc, &	n) Argument to turead func
m(i].join() - wait for mr	

threadfunc (unsigned * x) = increment int to 10,000,000

MOVI X, 7. tax add \$ 1, y. cax MOVI Y. Lax, X

main memory

equal to add \$1, X



(++ std:: atomic singh shread modifies

· looking to brild synchronization objects

\$0x ..., 1. Lax

nupl (y. rax)

IDLK add SOXI, (1.rdi)

SUD \$ OXI, Y. Eax

C++ S+d:: MVtx

- · Mutual exclusion, only one thread has across to a pick of code at a time
- · MVHX. (DUKC);

<< code >>

mrtix. Vnlock();

mutix objects

vn10cued

- 1 of 2 states locked
- initialization: starts unlocked
- mvtx:: lock() wait till stan == unlocked

```
· mvtix:: vn10ck() assert 10cked
                    SIT STAM = UNIDURED
  of When 2 threads across same normal variable
                                 undefined benavior
 strult muter
     Std:: atomic < Int7 stan = unlocked;
      void lock () ?
           while Istan == locked) }?
                                           two
                                            TUVEUds
                                           could acress
            stan = 101 ked;
                                             this
       void vnlock () {
           State = unlocked;
            store ()
  STILL WKONЬ!
```

atomically sit stan = locked