

RSA

If we can factor \rightarrow we can break RSA.

Extended Euclid's Algorithm

$\gcd(a, b)$ but also ints x, y s.t. $ax + by = d$

\hookrightarrow multiplicative
inverses
 $ax + by = 1$

$EE(a, b)$:

(gcd, x, y)

if $b = 0$: return $(a, 1, 0)$

$a = b * k + (a \bmod b)$

$-bkx + kky$
 $b x - (a \bmod b) y = d$

$(d, x, y) = EE(b, a \bmod b)$

$ay + (x - ky)b = d$

return $(d, y, x - k * y)$

$y * (a \bmod b + bk)$

a	b	k	x	y
313	37	8	-13	110
			6	-13
37	17	2	-1	6
17	3	5	1	-1
3	2	1	0	1
2	1	2	1	0
<div style="border: 1px solid black; padding: 2px;">1</div>	0			

$17 = 313 \bmod 37$ $8 \times 37 = 296$

\downarrow \uparrow

$\boxed{1}$ gcd

$y \rightarrow x$
 $x - ky \rightarrow y$

claim:

$$313(-13) + 37(110) = 1$$

Duplicates

Alta Vista - problem w/ duplicate documents on the web

↳ people don't want the same page in search results

↳ storage problem: no use storing duplicates / near duplicates

64-bit hashes → look for collisions?

doesn't work for near-duplicate documents

hash value should be same or nearly the same

document = set of numbers

set A, set B

Jaccard Coefficient

$$R(A, B) = \frac{|A \cap B|}{|A \cup B|}$$

$$0 \leq R(A, B) \leq 1$$

↑
disjoint

↑
exactly same set

$O(n^2)$ → dumbest method

$O(n \log n)$ → sorting and comparing

$O(n)$ → hash table and look for collisions

Method to Estimate Resemblance

calling card set (sketch)

$$TT: \{0, 1\}^{64} \rightarrow \{0, 1\}^{64}$$

$$\text{Black Box } (1, x) = \pi_1(x)$$

$$\text{Black Box } (50, y) = \pi_{50}(y)$$

On 4-bit integers: $\pi: [0, 15] \rightarrow [0, 15]$

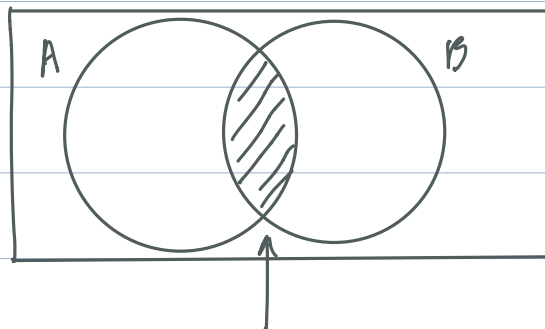
x	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
$\pi(x)$	9	2	14	11	6	3	7	8	15	10	4	13	12	0	1	5

$$\text{Set } A = \{3, 5, 11\} \longrightarrow \pi_1(A) = \{11, \boxed{3}, 13\}$$

calling card: $\min(\pi_1(A)), \min(\pi_2(A)), \min(\pi_3(A)), \dots$
 $\min(\pi_1(B)), \min(\pi_2(B)), \dots$

constant-size

$$\text{Prob}[\min(\pi_1(A)) = \min(\pi_1(B))] = \frac{|A \cap B|}{|A \cup B|} = R$$



min(union) for them to be =

$$\Pr[\min \pi_1(A \cup B) = \min \pi_1(A \cap B)]$$

Estimate resemblance: $\frac{\text{count \# of matches}}{\text{\# permutations}} \approx R$

$$E[\# \text{ matches}] = R \times \# \text{ permutations}$$

Documents \rightarrow Sets

4-shingling

Four score and seven years ago

↓ hash ↓ hash ↓ hash

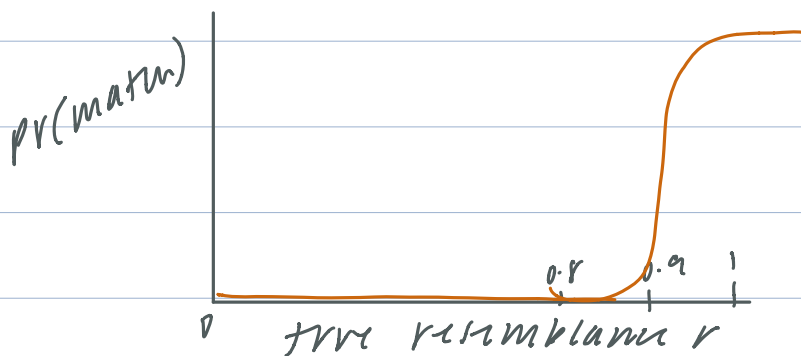
Subsample: only take hash values $= 0 \bmod 8$

Working in Practice

similar if estimated resemblance $\geq \theta$

true resemblance $= r$

$$\Pr[\text{match} \geq 90 \text{ matches in } 100\text{-len calling card}] \\ = \sum_{k=90}^{100} \binom{100}{k} r^k (1-r)^{100-k}$$



$$r = 0.05 \rightarrow 10^{-18}$$

$$r = 0.95$$

78.8% time similar

2-SAT

$$(x_1 \vee \bar{x}_2) \wedge (\bar{x}_1 \vee \bar{x}_3) \wedge (x_1 \vee x_2) \wedge (x_4 \vee \bar{x}_3) \wedge (x_4 \vee \bar{x}_1)$$

1SAT: normally NP-complete

2SAT: poly-time

Randomized Algo (slower than deterministic)

Start w/ randomized assignment

while \exists unsatisfied clause until timed

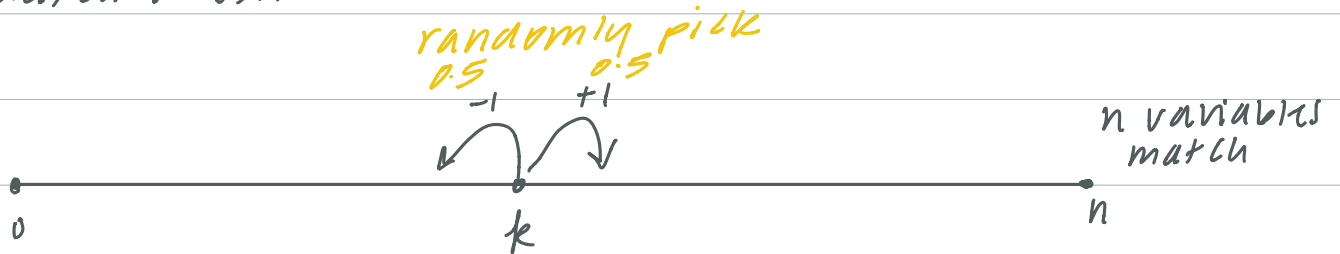
Flip a variable at random in an unsatisfied clause

return "unsatisfiable" or truth assignment

A = assignment

S = "start" - truth assignment

matches S, A



Random Walk on a Line

- Gambler's Ruin Problem