

1

The big-Oh space complexity of an adjacency list is $O(V+E)$. V is the number of vertices and E is the number of edges. It needs $O(V)$ memory to store all vertices and $O(E)$ to store all edges. Therefore, the space complexity is $O(V+E)$.

2

The big-Oh space complexity of an adjacency matrix is $O(V^2)$. It needs V rows and V columns to store 1 or 0 for each pair of nodes.

3

The representation of the graph can affect the time complexity. With an adjacency list, the big-Oh time complexity for searching an entire graph using DFS is $O(V+E)$. It needs $O(1)$ time to visit a node and $O(e)$ time to visit all edges starting from this node. If there are V vertices, the sum of the edges of each node e is the number of edges of the graph E . Thus, the time complexity is $O(V+E)$. With an adjacency matrix, the time complexity is $O(V^2)$ because all cells in the V columns and the V rows will be visited once.

4

Similar to DFS, the representation of the graph can affect the time complexity. With an adjacency list, the big-Oh time complexity for searching an entire graph using BFS is $O(V+E)$. With an adjacency matrix, the time complexity is $O(V^2)$.