

## ADS Assignment

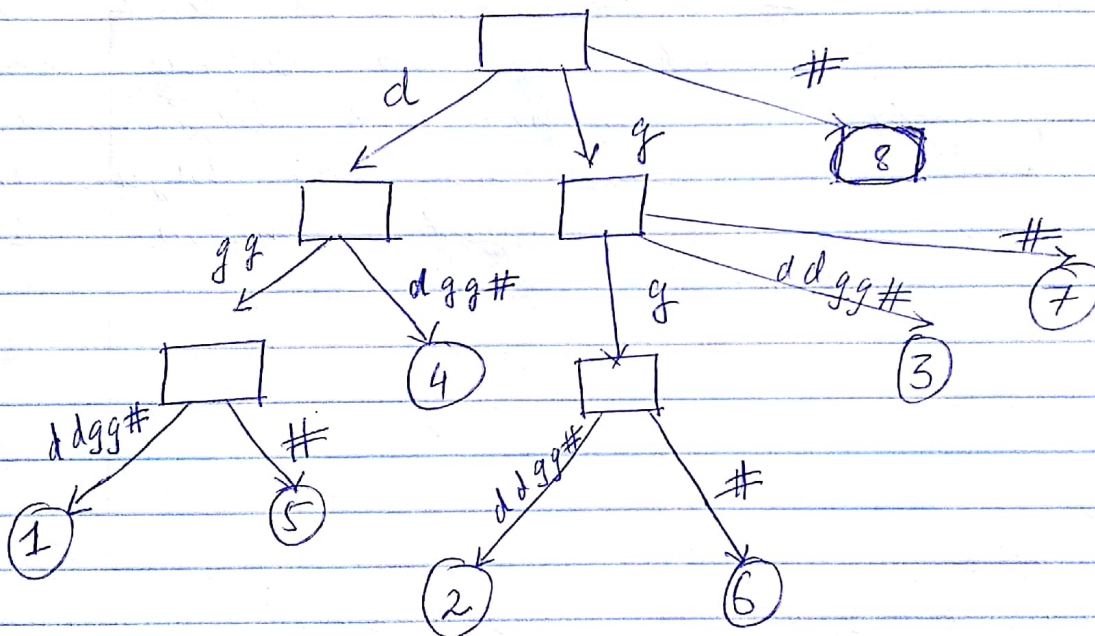
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### Question 1

suffix tree for  $dgddgg\#$



Q-1

(b) Assume 2 strings  $\rightarrow S, Q$

① set  $\$$  as the new character which is not contained in both  $S$  and  $Q$ .

② Put/set  $u = S \$ Q \#$

Now the problem is modified to  $\rightarrow$  finding the longest string that appears later

before and after '\$'

- ③ Construct a suffix tree ~~to~~ for the resulting new string 'u'
- ④ Label all the branch nodes with no of elements present in that subtree
- ⑤ ~~Find~~ Find the branch node with largest character # and has at least one element node in its subtree that represents a suffix that begins in  $S$  and at least one that begins in  $Q$ . This is the longest common substring in  $S \neq Q$ .

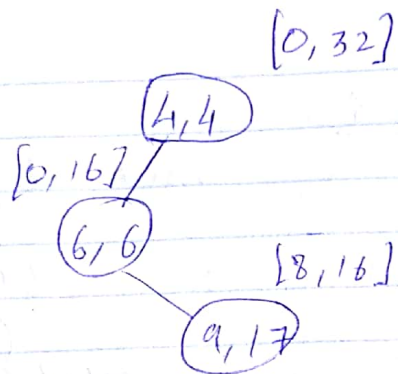
(Q-2)  
(a)

① insert (6, 6)      [ 0, 32 ]  
  (6, 6)

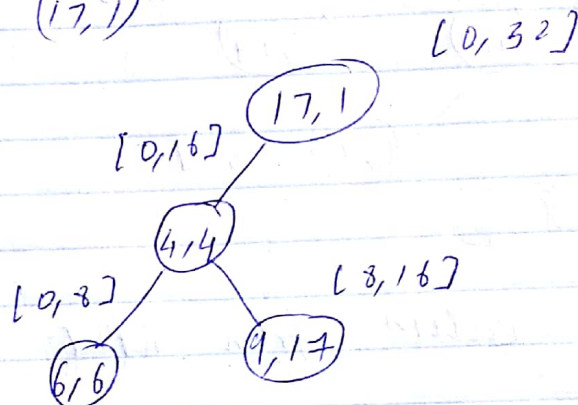
② insert (9, 17)      [ 0, 32 ]  
  (6, 6)  
  [ 0, 16 ]  
  (9, 17)



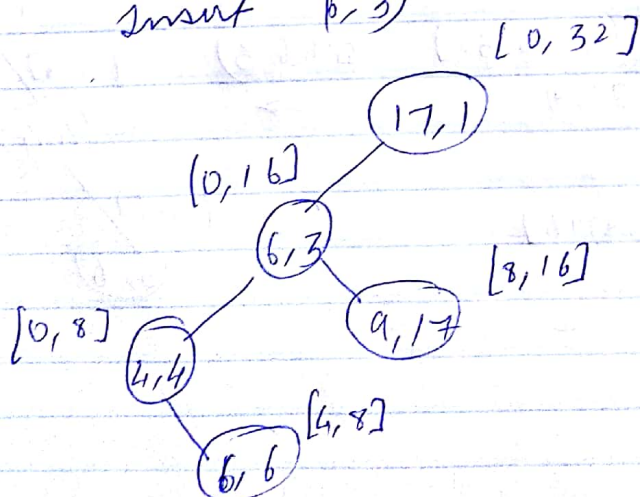
(3) Insert (4, 4)



(4) Insert (17, 1)

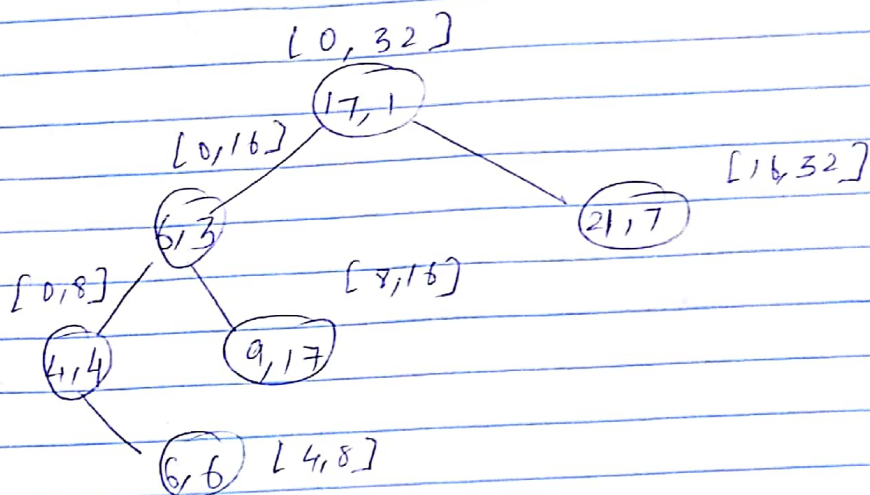


(5) Insert (6, 3)



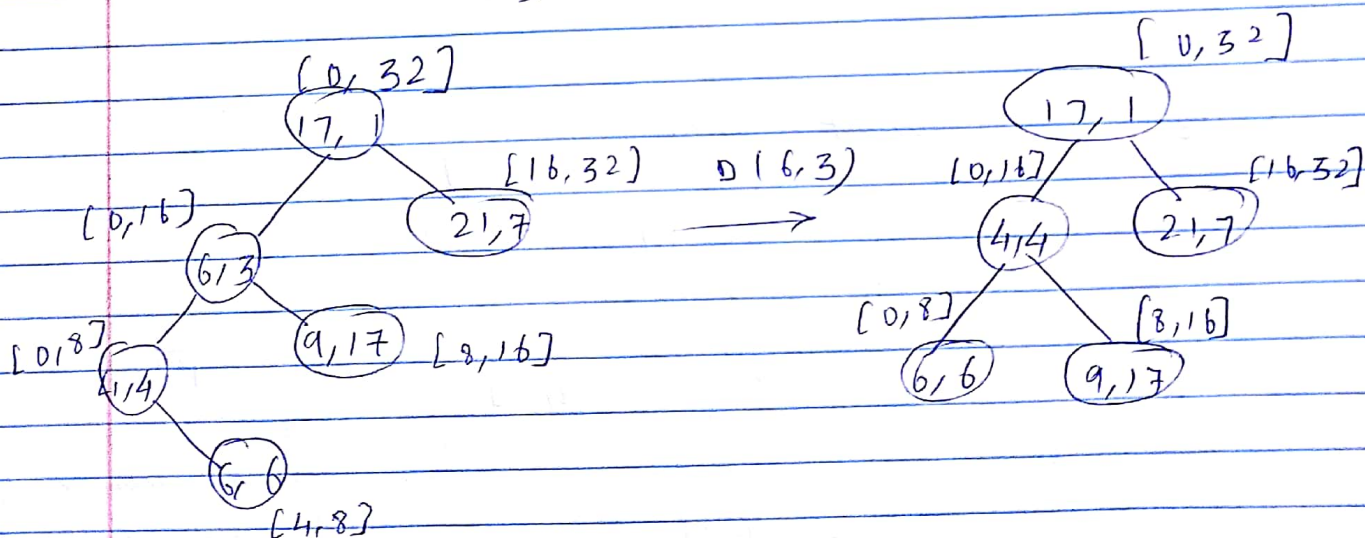
Step 6

Insert (21, 7)



Q-2 (b) Delete from RPS7

Delete (6, 3)



(Q 3)  
(a)

Bloom Filter:

No of hash functions (h) to use =

$$(\ln_2 2) * (m/u)$$

$$\Rightarrow h = 0.693 * (m/u)$$

$$\Rightarrow h = 0.693 * \frac{5000}{1000}$$

$$= 3.465$$

$$h = 3 \text{ or } 4$$

(b) Prob of filter error  $[P(u)]$

$$= e^{-u/n} (1 - e^{-u h/m})$$

For  $h = 3$

$$\frac{u \times h}{m} = \frac{1000 \times 3}{5000} = 3/5$$

$$P(u) = e^{-1000/100000} (1 - e^{-3/5})^3$$

$$= 0.0909$$

For  $h = 4$

$$\frac{u \times h}{m} = \frac{1000 \times 4}{5000} = 4/5$$

$$P(u) = e^{-1/100} (1 - e^{-4/5})^4$$

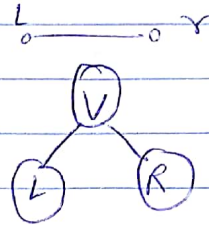
$$= 0.0910$$



(Q-4)

Query interval  $J[l, r]$

3 cases



(1) If  $v.key \in [l, r]$

→ All intervals in root i.e.  $v$  overlaps the query interval

→ Search both  $L$  &  $R$  trees and recursively to search for overlapping intervals.

(2) If  $v.key < l$  and  $l > v.k$

check for all intervals in  $V$ , if for any interval  $(l_i, r_i)$  if  $r_i \geq l$  return it as overlapping interval.

- Search  $R$  for any additional overlapping interval

- No need to search in  $L$  tree

(3) If  $r_i < v.key$  —

check for all intervals in  $V$  if for any interval  $(l_i, r_i)$  if  $l_i \leq r$  return it as overlapping interval.

- Search  $L$  tree for additional overlapping interval

- No need to search in  $R$  tree