

## Section 3.1 — Arithmetic Instructions (Mazidi)

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### Chapter 3 • Section 3.1 — Exercises (Mazidi)

Problems are paraphrased to respect copyright. This section focuses on **ADDS/ADC** and the **C (carry) / Z (zero)** flags.

#### 1) Find C and Z for each case. Also give the result and where it is saved.

(a)

```
MOV    R1, #0x3F
MOV    R2, #0x45
ADDS   R3, R1, R2
```

- **Computation:**  $0x3F + 0x45 = 0x84$
- **Result:**  $R3 = 0x00000084$
- **Flags:**  $C=0, Z=0$  (no carry out; result not zero).

(b)

```
LDR    R0, =0x95999999
LDR    R1, =0x94FFFF58
ADDS   R1, R1, R0
```

- **Computation:**  $0x95999999 + 0x94FFFF58 = 0x12A9998F1 \rightarrow$  low 32 bits  $0x2A9998F1$
- **Result:**  $R1 = 0x2A9998F1$
- **Flags:**  $C=1$  (carry out),  $Z=0$ .

(c)

```
LDR    R0, =0xFFFFFFFF
ADDS   R0, R0, #1
```

- **Computation:**  $0xFFFFFFFF + 1 = 0x1\_0000\_0000 \rightarrow$  low 32 bits  $0x00000000$
- **Result:**  $R0 = 0x00000000$
- **Flags:**  $C=1, Z=1$ .

(d)

```
LDR    R2, =0x00000001
LDR    R1, =0xFFFFFFFF
ADD    R0, R1, R2    ; does NOT set flags
ADCS   R0, R0, #0    ; adds carry-in and sets flags
```

- After **ADD**:  $R0 = 0x00000000$  (flags **unchanged**).
- **ADCS** uses the **previous C** (not set by the **ADD**). Assuming prior  $C=0$  (typical unless set earlier):
  - **Result:**  $R0 = 0x00000000$
  - **Flags set by ADCS:**  $C=0, Z=1$ .  
(If prior  $C=1$ , then  $R0=0x00000001$  and  $Z=0$ .)

(e)

```
LDR    R0, =0xFFFFFFFF
ADDS   R0, R0, #2
ADC    R1, R0, #0    ; uses carry from ADDS; does not set flags
```

- **Computation:**  $0xFFFFFFFF + 2 = 0x1\_0000\_0000 \rightarrow R0 = 0x00000000$
- **Flags after ADDS:**  $C=1, Z=1$
- **Then ADC:**  $R1 = R0 + 0 + C = 0 + 0 + 1 = 0x00000001$  (flags unchanged).

## 2) State the three steps in a subtraction (SUB) and apply them.

### Three steps (A – B):

1. **One's complement** of B  $\rightarrow \sim B$ .
2. **Add 1** to form **two's complement** of B.
3. **Add** to A:  $A + (\sim B + 1)$ . In ARM, the **C flag after subtraction** means:  $C=1 \rightarrow$  no borrow,  $C=0 \rightarrow$  borrow.

Apply to 8-bit examples (showing intermediate two's complement):

- **(a)**  $0x23 - 0x12$ 
  - $\sim 0x12 = 0xED, +1 \rightarrow 0xEE; 0x23 + 0xEE = 0x111 \rightarrow$  result  $0x11, C=1$  (no borrow).
- **(b)**  $0x43 - 0x51$ 
  - $\sim 0x51 = 0xAE, +1 \rightarrow 0xAF; 0x43 + 0xAF = 0xF2 \rightarrow$  result  $0xF2$  (i.e.,  $-0x0E$  in 8-bit),  $C=0$  (borrow occurred).
- **(c)**  $0x99 - 0x39$ 
  - $\sim 0x39 = 0xC6, +1 \rightarrow 0xC7; 0x99 + 0xC7 = 0x160 \rightarrow$  result  $0x60, C=1$  (no borrow).

## Notes for learners

- ADD vs ADDS: only forms with **s** update flags.
- ADC/ADCS add the **carry-in**; ADCS also **updates** flags.
- In ARM subtraction, remember: **C = NOT borrow**.