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Extracting dependency from code is an almost automatic process. You need to choose a granularity. But once that is chosen, the entire analysis follows.

In the whole activity, you should express the metrics in complexity notation as a function of the parameters of the functions.

1 Coin Collection (from Midterm Spring 2018)

The Coin Collection problem is defined as follows:

Several coins are placed on an $n \times m$ board with at most one coin per cell of the board. A robot is initially located at the upper left cell of the board. The robot can only move to the right or down; it can not move up or left. When the robot visits a cell where a coin is located, it picks it up. At most, how many coins can the robot collect?

This problem can be solved by the following method:

```
void RobotCoin(int n, int m, //size of the board
               int C[n][m] //Is there a coin in (i,j)
               ) {

    int F[n][m]; //How many coins can be collected while on (i,j)

    F[0][0] = C[0][0];

    for (int k=1; k<m; ++k) {
        F[0][k] = F[0][k-1] + C[0][k];
    }

    for (int i=1; i<n; ++i) {
        F[i][0] = F[i-1][0] + C[i][0];
        for (int j=1; j<m; ++j) {
            F[i][j] = max (F[i-1][j], F[i][j-1]) + C[i][j];
        }
    }

    return F[n-1][m-1];
}
```

Question: What is the complexity of this function?

Answer :- The complexity of this function is $O(n*m)$

Question: Extract the dependencies.

Question:-

Answer:-

Steps	Tasks
①	$F[4][4]$ # init $F[n][m]$
②	$F[0][0] = C[0][0]$ # $F[0][0] = C[0][0]$
③	$F[0][1] = F[0][0] + C[0][1]$ for (int k=1; k<m; ++k)
④	$F[0][2] = F[0][1] + C[0][2]$ # $F[0][x] = F[0][x-1] + C[0][x]$
⑤	$F[0][3] = F[0][2] + C[0][3]$

```

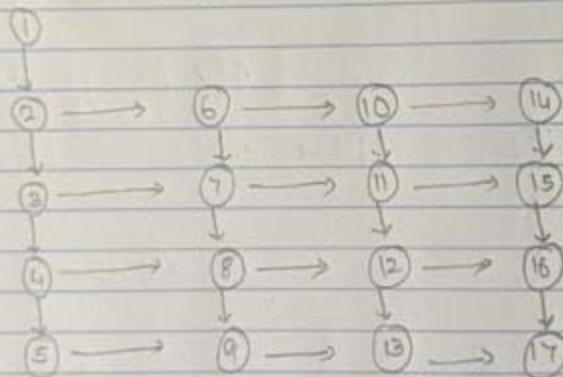
for (int k=1; k<n; ++k) {
    F[k][0] = F[k-1][0] + C[k][0];
    for (int j=1; j<m; ++j) {
        F[k][j] = max(F[k-1][j], F[k][j-1])
        + C[k][j];
    }
}

```

- ⑥ $F[1][0] = F[0][0] + C[1][0]$
- ⑦ $F[1][1] = \max(F[0][1], F[1][0]) + C[1][1]$
- ⑧ $F[1][2] = \max(F[0][2], F[1][1]) + C[1][2]$
- ⑨ $F[1][3] = \max(F[0][3], F[1][2]) + C[1][3]$
- ⑩ $F[2][0] = F[1][0] + C[2][0]$
- ⑪ $F[2][1] = \max(F[1][1], F[2][0]) + C[2][1]$
- ⑫ $F[2][2] = \max(F[1][2], F[2][1]) + C[2][2]$
- ⑬ $F[2][3] = \max(F[1][3], F[2][2]) + C[2][3]$

- (14) $F[3][0] = F[2][0] + C[3][0]$
 (15) $F[3][1] = \max(F[2][1], F[3][0]) + C[3][1]$
 (16) $F[3][2] = \max(F[2][2], F[3][1]) + C[3][2]$
 (17) $F[3][3] = \max(F[2][3], F[3][2]) + C[3][3]$

Graph -



Question:What is the width?

Answer:- Since F[4][4] is only an independent task, the width is 1

Question:What is the work?

Answer:- The work is sum of processing time required to complete $n*m$ tasks.i.e $O(n*m)$

Question:What is the critical path? What is its length?

Answer:- Critical Path is $O(n+m)$ and the length is the sum of processing time of all the task within the critical path.

2 Knapsack

The Knapsack problem aims at finding the best set of objects to pack in a bag. Often the following dynamic programming algorithm is used to solve the problem.

```
void knapsack (int n, int W, int value[], int weight[], int val[][]){
    for (int a = 0; a<=W; ++a) {
        val[0][a] = 0;
    }

    for (int i=1; i<=n; ++i) {
        for (int j=0; j<=W; ++j) {
            val[i][j] = val[i-1][j];
            if (weight[i-1] <= j) {
                val[i][j] = max (val[i-1][j], value[i-1]+val[i-1][j-weight[i-1]]);
            }
        }
    }
}
```

Question:What is the complexity of this function?

Answer: Complexity of algorithm is $O(nW)$

Question:Extract the dependencies.

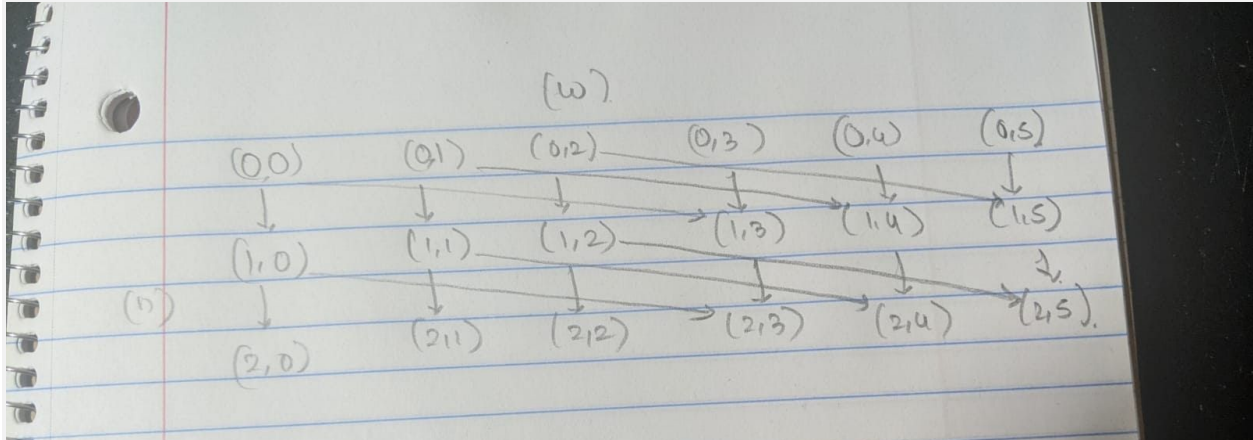
Answer:

Value $\rightarrow [1, 2, 1, 4]$

Weight $\rightarrow [3, 1, 2, 6]$

$N \rightarrow 2$

$W \rightarrow 5$



Question: What is the width?

Answer: As mentioned in first for loop, the value of a depends upon the value on W which is indirectly the reason for independent task, Hence the width is W

Question: What is the work?

Answer: The work will be sum of processing time required to complete $n \times W$ task i.e $O(n \times W)$

Question: What is the critical path? What is its length?

Answer: The critical path would be length of the diagonal as shown in the diagram above ($n \rightarrow W$)

3 Bubble Sort

The bubble sort algorithm can be written like this:

```

void bubblesort(int* A, int n) {
    for (int i=0; i<n; ++i) {
        for (int j=1; j<n; ++j) {
            if (A[j] < A[j-1]) {
                int temp = A[j];
                A[j] = A[j-1];
                A[j-1] = temp;
            }
        }
    }
}

```

Question:What is the complexity of this function?

Answer: The complexity of this function is $O(n^2)$

Question:Extract the dependencies.

Answer:

[3, 6, 1, 4, 7]

N=5



I=0 J=0	I=0 J=1 <u>A</u> [0,1] R <u>A</u> [0,0] R Temp w <u>A</u> [0,1] w <u>A</u> [0,0] w	I=0 J=2 <u>A</u> [0,2] R <u>A</u> [0,1] R Temp w <u>A</u> [0,2] w <u>A</u> [0,1] w	I=0 J=3 <u>A</u> [0,3] R <u>A</u> [0,2] R Temp w <u>A</u> [0,3] w <u>A</u> [0,2] w	I=0 J=4 <u>A</u> [0,4] R <u>A</u> [0,3] R Temp w <u>A</u> [0,4] w <u>A</u> [0,3] w
I=1 J=0	I=1 J=1 <u>A</u> [1,1] R <u>A</u> [1,0] R Temp w <u>A</u> [1,1] w <u>A</u> [1,0] w	I=1 J=2 <u>A</u> [1,2] R <u>A</u> [1,1] R Temp w <u>A</u> [1,2] w <u>A</u> [1,1] w	I=1 J=3 <u>A</u> [1,3] R <u>A</u> [1,2] R Temp w <u>A</u> [1,3] w <u>A</u> [1,2] w	I=1 J=4 <u>A</u> [1,4] R <u>A</u> [1,3] R Temp w <u>A</u> [1,4] w <u>A</u> [1,3] w
I=2 J=0	I=2 J=1 <u>A</u> [2,1] R <u>A</u> [2,0] R Temp w <u>A</u> [2,1] w <u>A</u> [2,0] w	I=2 J=2 <u>A</u> [2,2] R <u>A</u> [2,1] R Temp w <u>A</u> [2,2] w <u>A</u> [2,1] w	I=2 J=3 <u>A</u> [2,3] R <u>A</u> [2,2] R Temp w <u>A</u> [2,3] w <u>A</u> [2,2] w	I=2 J=4 <u>A</u> [2,4] R <u>A</u> [2,3] R Temp w <u>A</u> [2,4] w <u>A</u> [2,3] w
I=3 J=0	I=3 J=1 <u>A</u> [3,1] R <u>A</u> [3,0] R Temp w <u>A</u> [3,1] w <u>A</u> [3,0] w	I=3 J=2 <u>A</u> [3,2] R <u>A</u> [3,1] R Temp w <u>A</u> [3,2] w <u>A</u> [3,1] w	I=3 J=3 <u>A</u> [3,3] R <u>A</u> [3,2] R Temp w <u>A</u> [3,3] w <u>A</u> [3,2] w	I=3 J=4 <u>A</u> [3,4] R <u>A</u> [3,3] R Temp w <u>A</u> [3,4] w <u>A</u> [3,3] w
I=4 J=0	I=4 J=1 <u>A</u> [4,1] R <u>A</u> [4,0] R Temp w <u>A</u> [4,1] w <u>A</u> [4,0] w	I=4 J=2 <u>A</u> [4,2] R <u>A</u> [4,1] R Temp w <u>A</u> [4,2] w <u>A</u> [4,1] w	I=4 J=3 <u>A</u> [4,3] R <u>A</u> [4,2] R Temp w <u>A</u> [4,3] w <u>A</u> [4,2] w	I=4 J=4 <u>A</u> [4,4] R <u>A</u> [4,3] R Temp w <u>A</u> [4,4] w <u>A</u> [4,3] w

Here only row values are dependent. So the dependency graph would be,

00 → 01 → 02 → 03 → 04

10 → 11 → 12 → 13 → 14

20 → 21 → 22 → 23 → 24

.

Question:What is the width?

Answer: So, each row is independent of other rows and which are comparisons of each number with rest in the
So width would be N.

Question:What is the work?

Answer: Total work is $N*(N-1)$

Question: What is the critical path? What is its length?

Answer: Critical path is each row. In the above example $00 \rightarrow 01 \rightarrow 02 \rightarrow 03 \rightarrow 04$ can be the critical path.