

# Lokalizacja punktu w przestrzeni dwuwymiarowej – metoda trapezowa

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# Opis zagadnienia

Dany jest obszar z podziałem poligonowym. Zadawany jest punkt  $P$  na płaszczyźnie. Należy zaimplementować algorytm lokalizacji punktu metodą trapezową, który odpowie na pytanie, w którym elemencie znajduje się dany punkt. Program powinien w sposób graficzny prezentować etapy algorytmu dla wybranych przykładów ( w celu objaśnienia działania algorytmu). Program ma służyć jako narzędzie dydaktyczne do objaśnienia działania algorytmu.

# Lokalizowania punktu 2D

Dane wejściowe:

- Punkt 2D
- Obszar poligonowy podziału płaszczyzny

Dane wyjściowe :

- Wielokąt zawierający zadany punkt



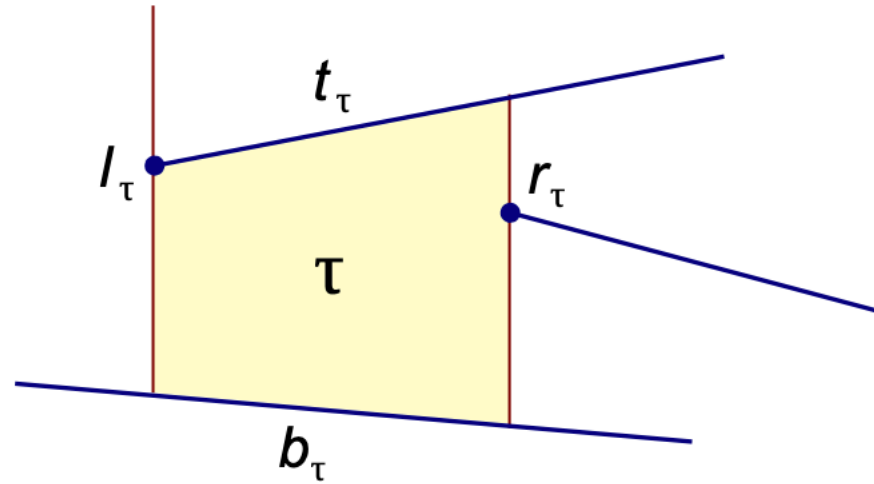
# Tworzenie trapezoid map - algorytm

1. Determine a bounding box, that contains every segment in  $S$ . Initialize the Trapezoidal Map and the corresponding searchstructure for it.
2. Compute a random permutation of the elements in  $S$ .
3. for  $i = 1$  to  $n$  do
4. Find the set of trapezoids that are intersected when the segment  $s_i$  is added.
5. Remove the intersected trapezoids that arise because of the insertion of  $s_i$ .
6. Remove the leaves for the intersected trapezoids from the search structure and create leaves for the new trapezoids. Link the created leaves to the existing nodes in the search structure.

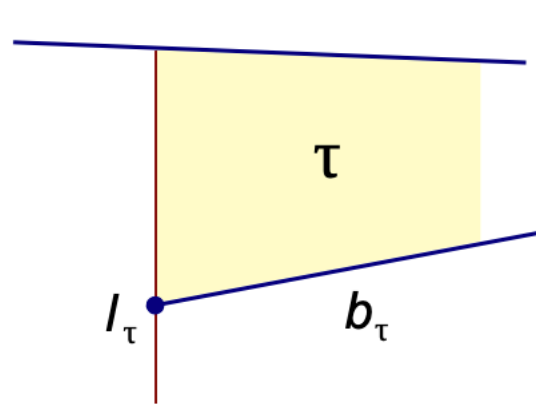
# Definicja trapezu na potrzeby naszego algorytmu

Trapezoid  $\tau$  :

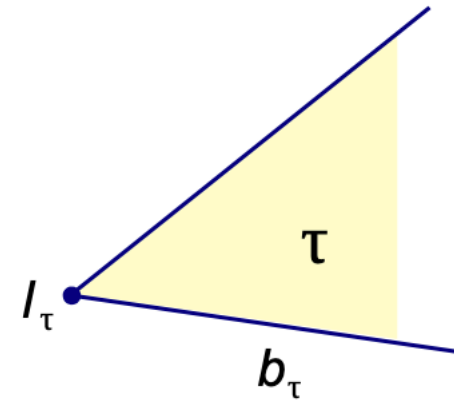
- Top edge:  $t_\tau$
- Bottom edge:  $b_\tau$
- Left vertex:  $l_\tau$
- Right vertex:  $r_\tau$
- (possibly horizontal walls / vertices of the bounding box)



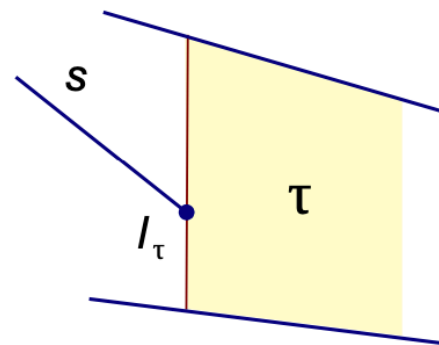
# Przypadki lewego boku



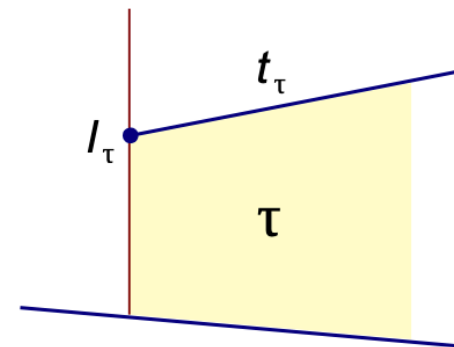
upper vertical extension  
of original endpoint  $l_\tau$



meeting point  $l_\tau$   
of two original segments



whole vertical extension  
of original endpoint  $l_\tau$



lower vertical extension  
of original endpoint  $l_\tau$

# Algorithm steps

- Initially the map contains only the *bounding box*
- $\rightarrow$  one-node DAG
- For each edge  $e \in S$  in randomized order...
  - remove the trapezoids  $\mathcal{T}_1, \mathcal{T}_2, \dots, \mathcal{T}_k$  in *conflict* with  $e$
  - replace them with the new trapezoids determined by  $e$
  - remove the DAG's leaves linked to  $\mathcal{T}_1, \mathcal{T}_2, \dots, \mathcal{T}_k$
  - replace these leaves with  $x$ -/ $y$ -nodes as appropriate
  - create and link leaves for the new trapezoids

## Finding trapezoids in conflict with a new edge

- Point location of  $e$ 's left endpoint (current DAG)
- $\rightarrow$  leftmost trapezoid  $\tau_1$  in conflict with  $e$
- Follow right-neighbor links from  $\tau_1$  to the trapezoid  $\tau_k$  which contains  $e$ 's right endpoint (edges do not cross)
- The correct neighbor  $\tau_{i+1}$  of  $\tau_i$  is identified by testing where  $r_{\tau_i}$  lies relative to  $e$



## Updating the map

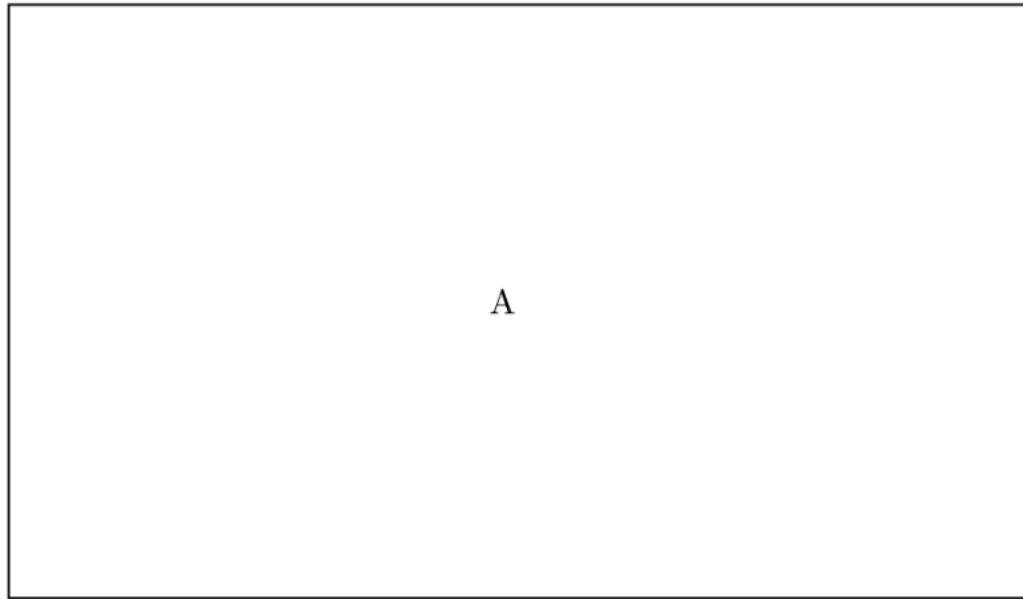
- $\mathcal{T}_1$  and  $\mathcal{T}_k$  are partitioned in three parts (four if  $\mathcal{T}_1 = \mathcal{T}_k$ )
- $\mathcal{T}_2, \mathcal{T}_3, \dots, \mathcal{T}_{k-1}$  are split
- Whenever possible, the resulting trapezoids bounded by  $e$  are merged
- All operations can be done in  $O(k)$   
(in constant time for each involved trapezoid)

## Updating the DAG

- Cross links between leaf nodes and trapezoids
- At most three new  $x$ -/ $y$ -nodes for each removed trapezoid
- Several nodes are linked to a new “merged” trapezoid
- All arrangements can be done in  $O(k)$

## Step 1: Initialization of Trapezoidal Map and History

TrapezoidalMap

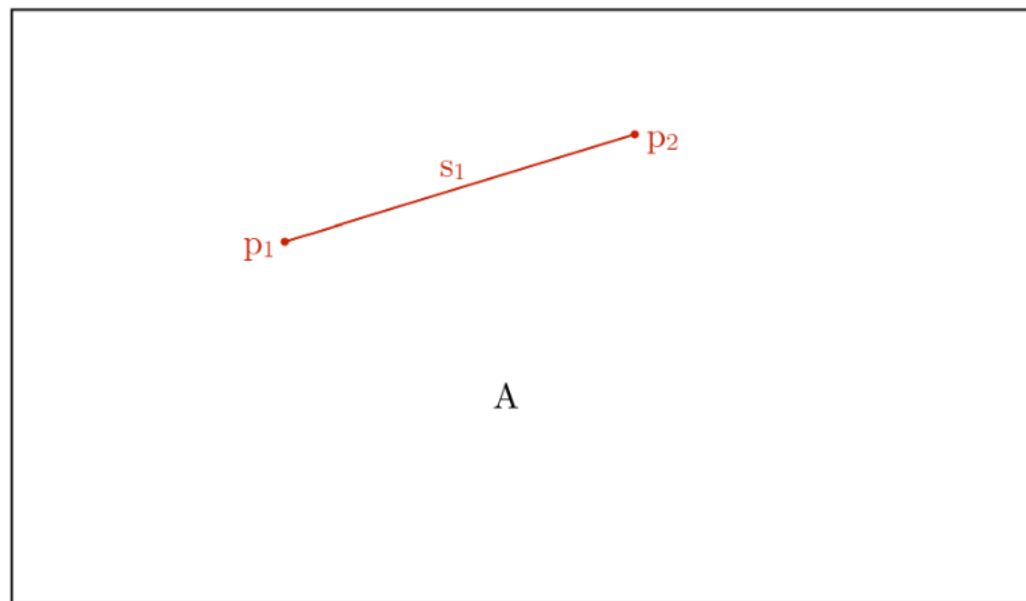


History



### Step 3: First segment to be added

TrapezoidalMap

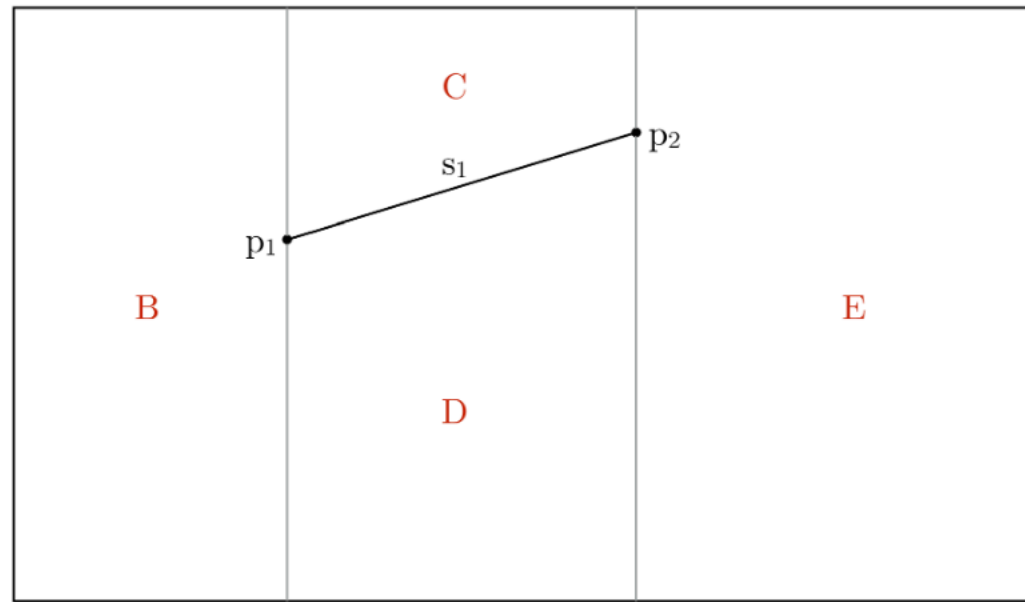


History

A

## Step 5: Vertical extension of the segment's endpoints

TrapezoidalMap

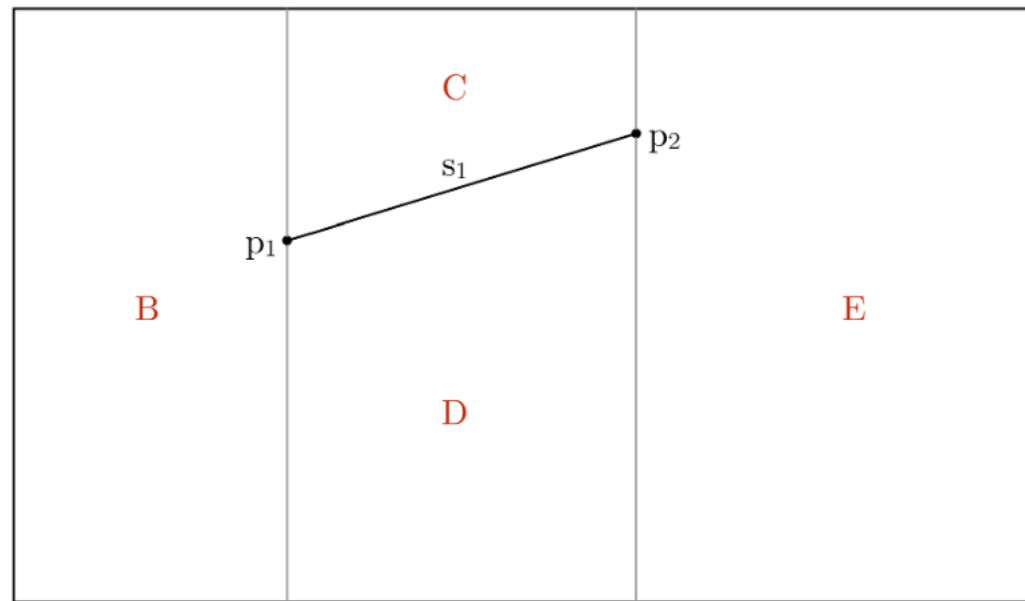


History

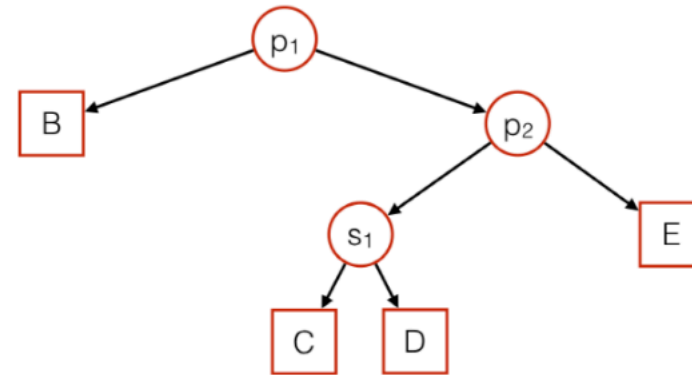
A

Step 6: The History is updated with new nodes for the new trapezoids

TrapezoidalMap

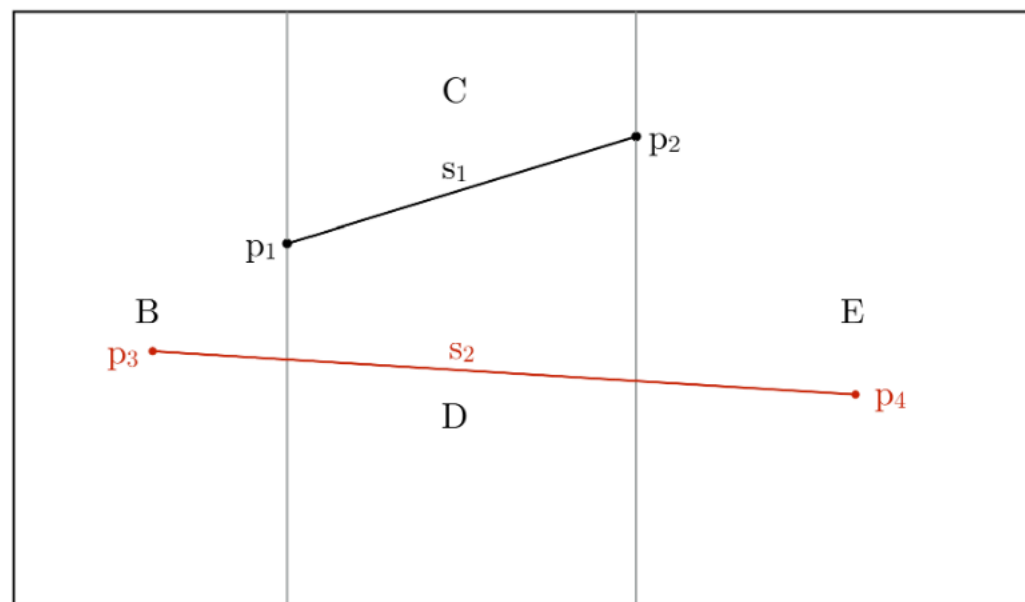


History

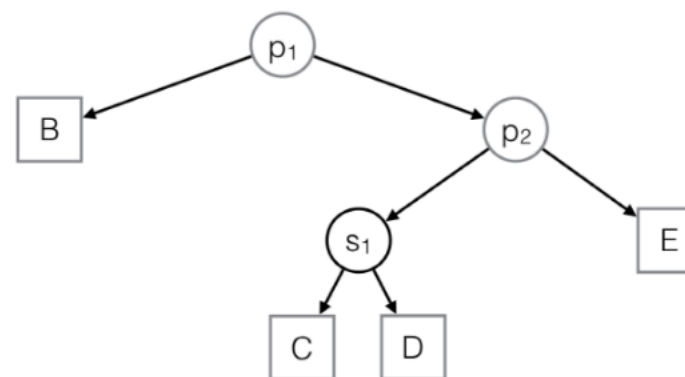


### Step 3: Second segment to be added

#### TrapezoidalMap

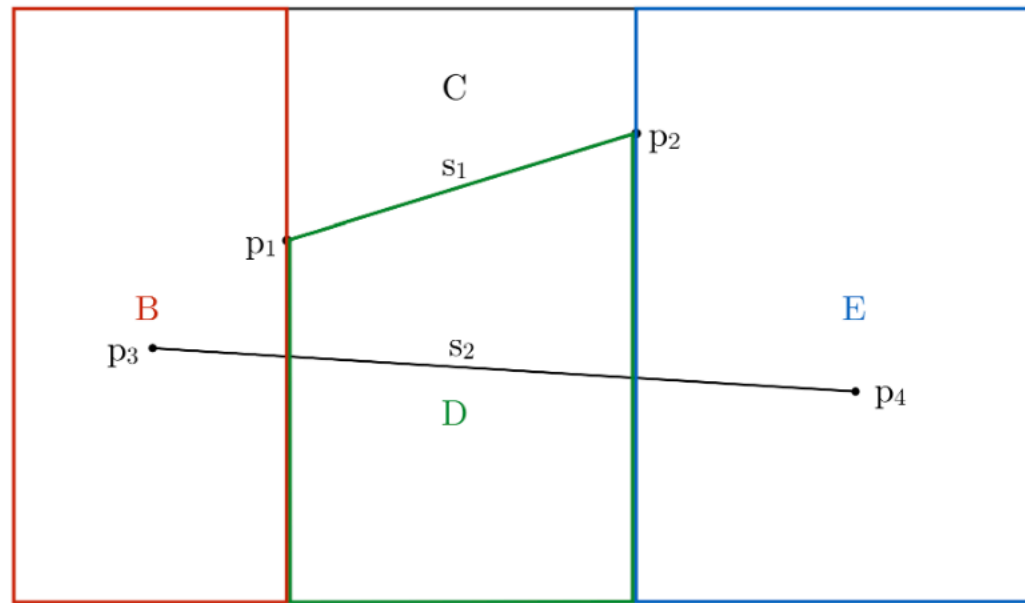


#### History

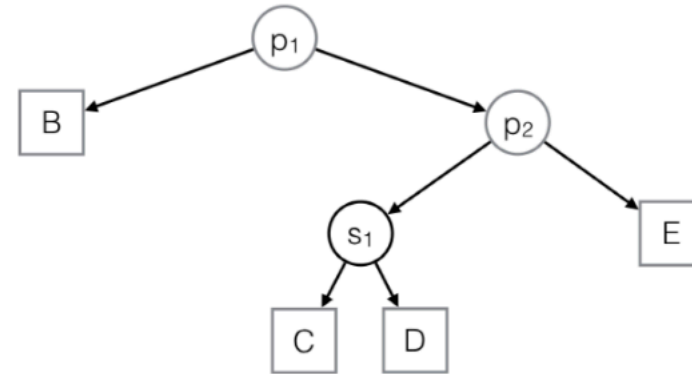


## Step 4: Intersected trapezoids are highlighted

### TrapezoidalMap



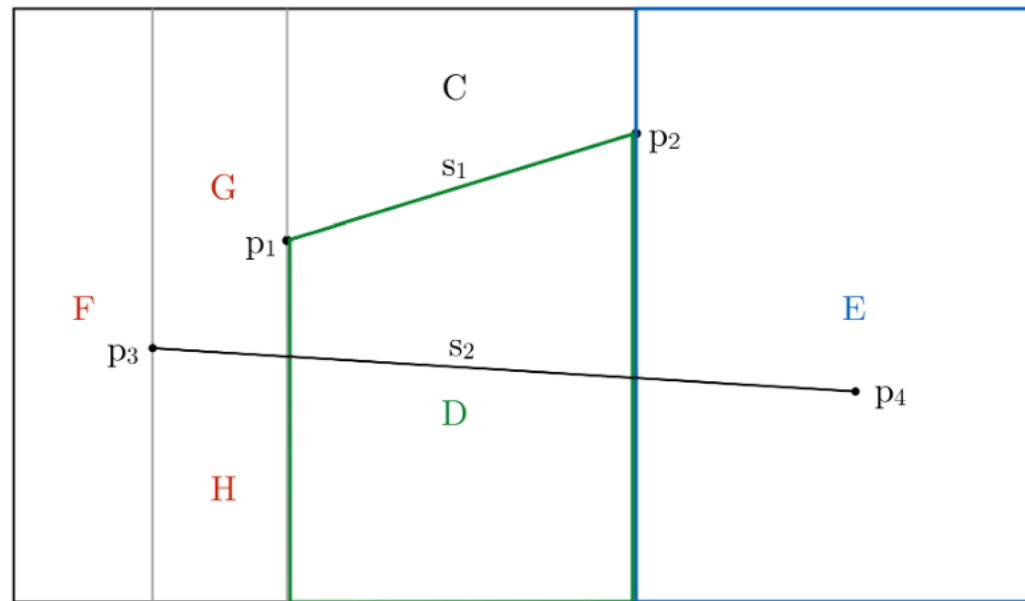
### History



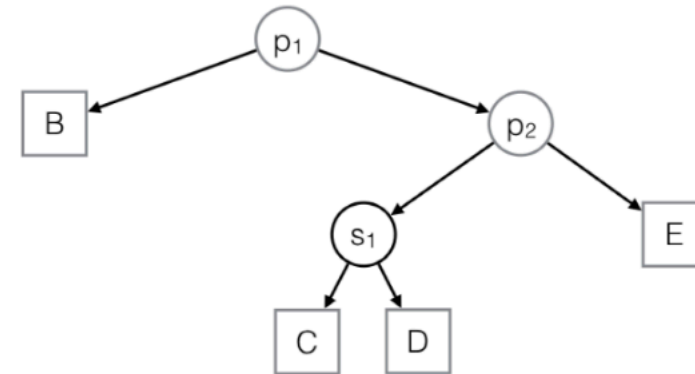


Step 5: The first intersected trapezoid is replaced by three new trapezoids

## TrapezoidalMap

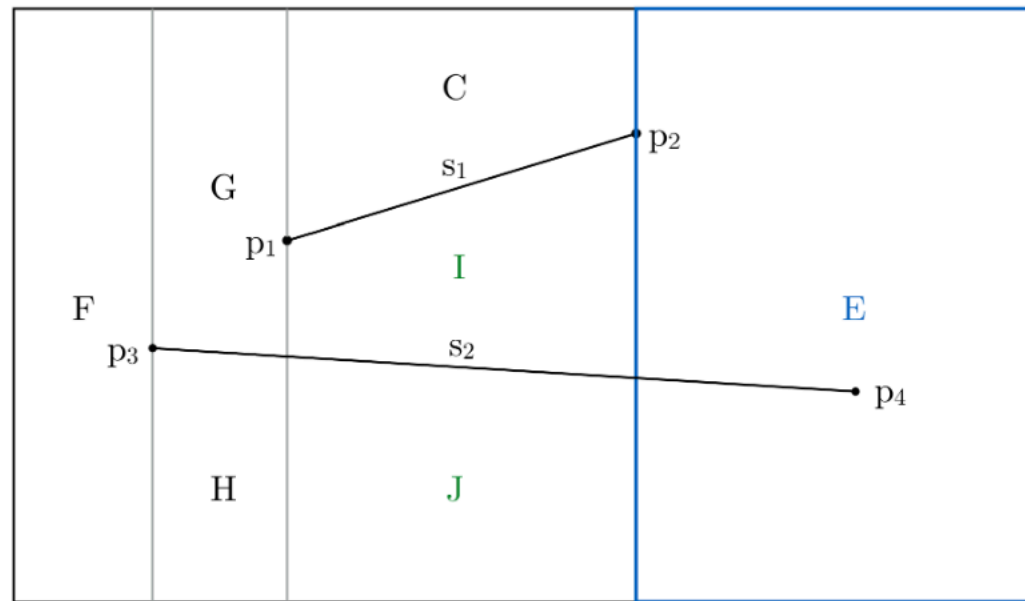


## History

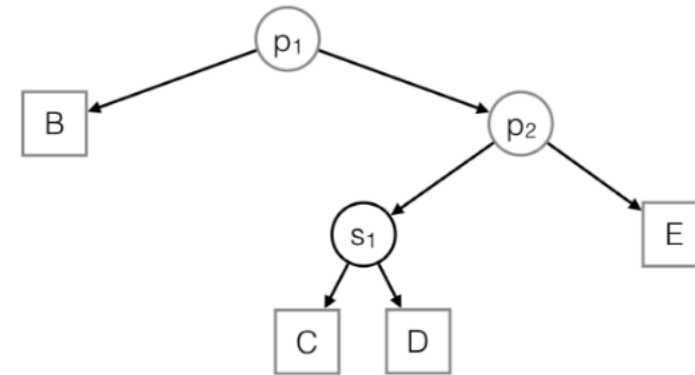


Step 5: The second intersected trapezoid is replaced by two new trapezoids

## TrapezoidalMap

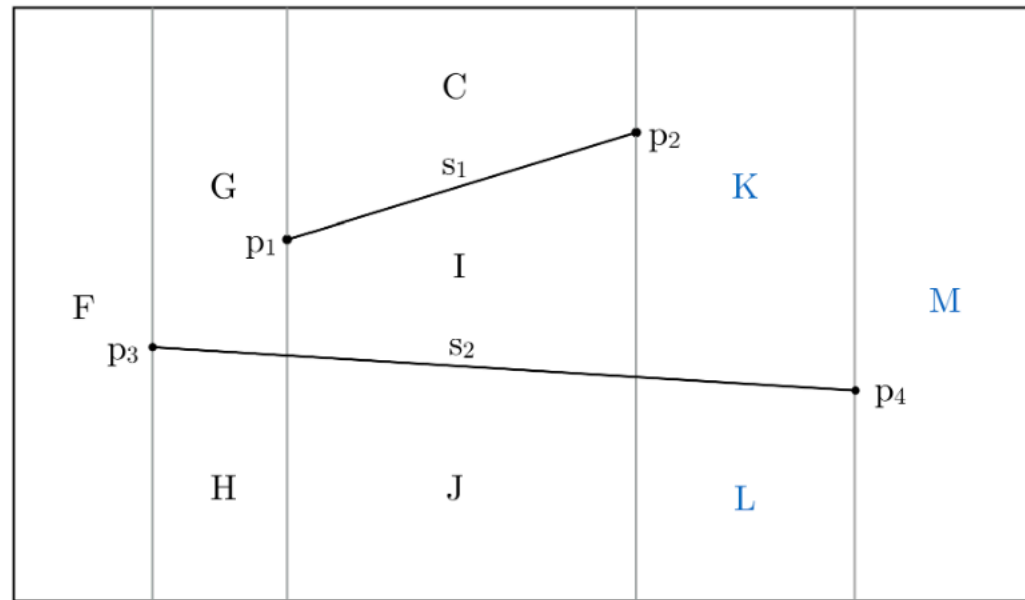


## History

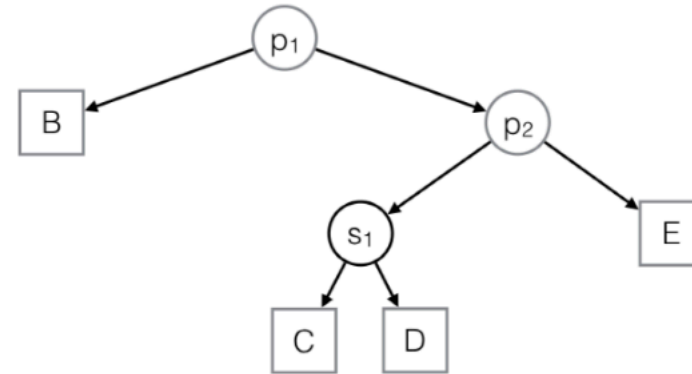


Step 5: The third intersected trapezoid is replaced by three new trapezoids

## TrapezoidalMap

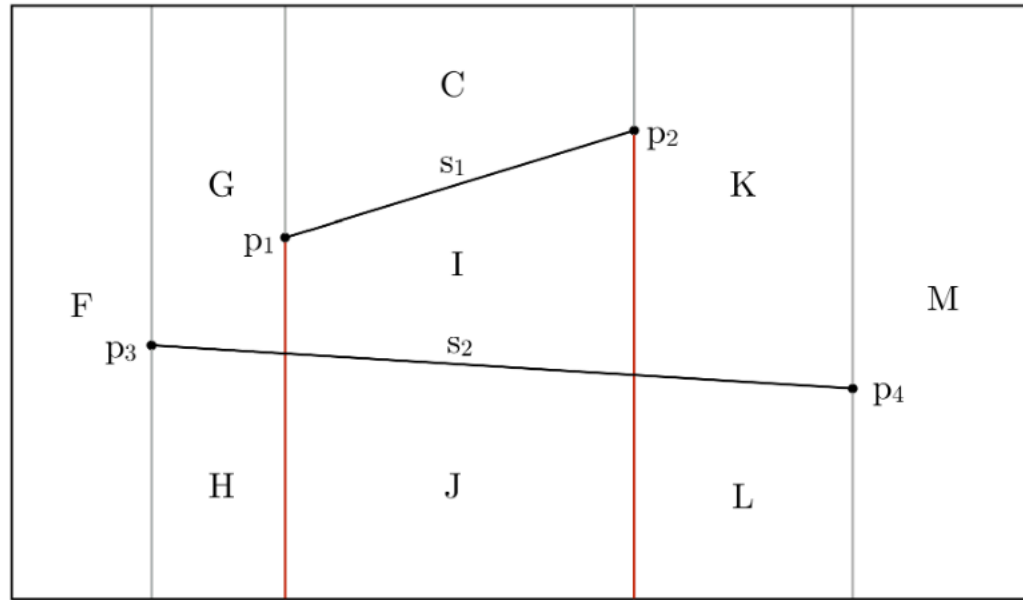


## History

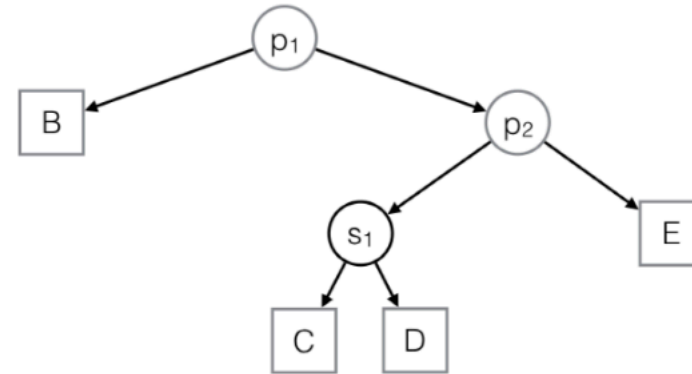


Step 5: Highlighted vertical extensions of the first segment are too long

## TrapezoidalMap

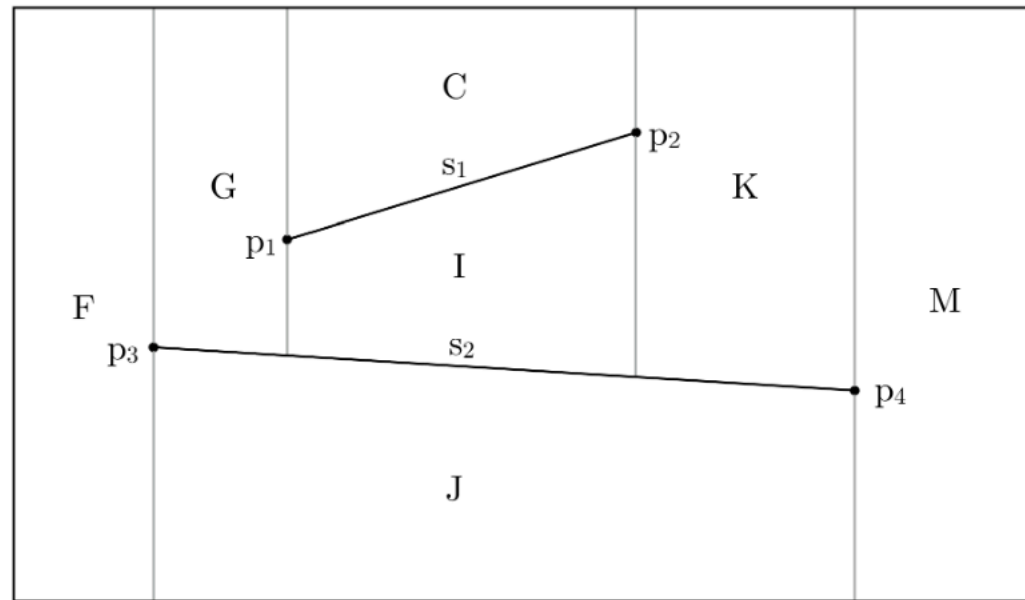


## History

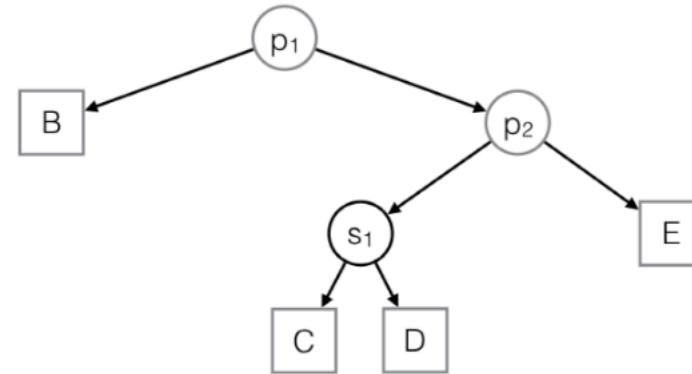


## Step 5: Shortening of the vertical extensions

### TrapezoidalMap

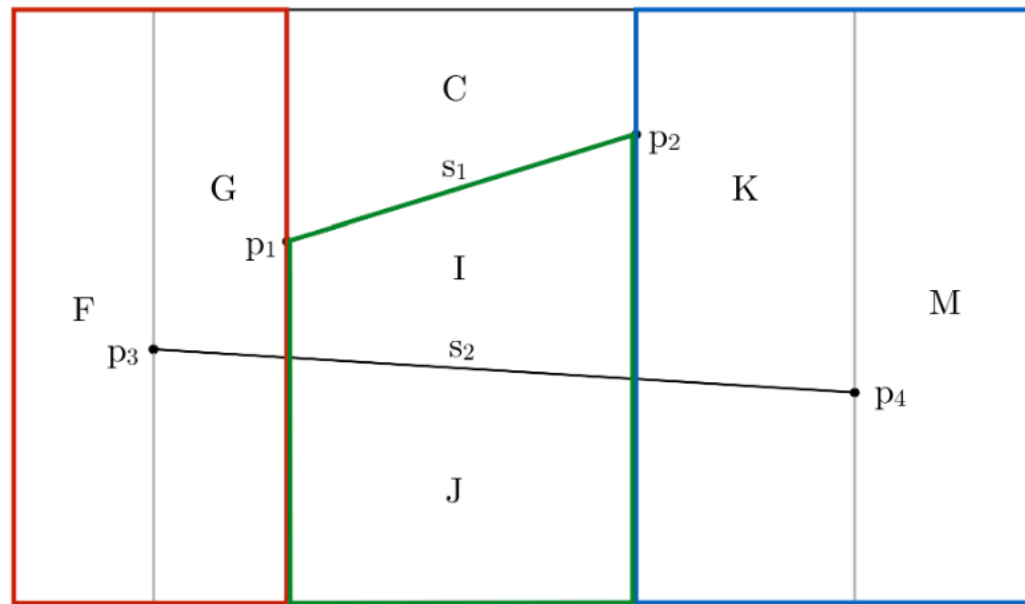


### History

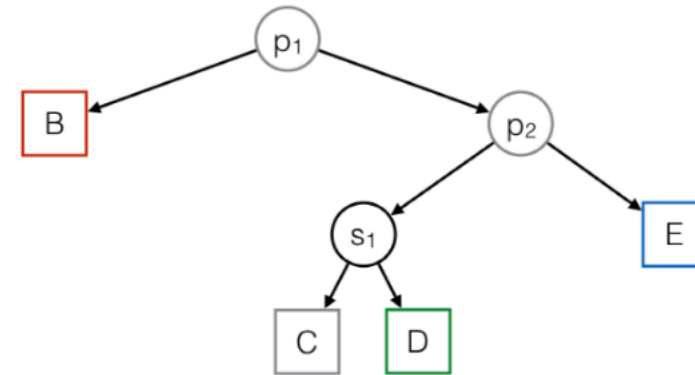


Intersected trapezoids are highlighted in the Trapezoidal Map and the History

TrapezoidalMap

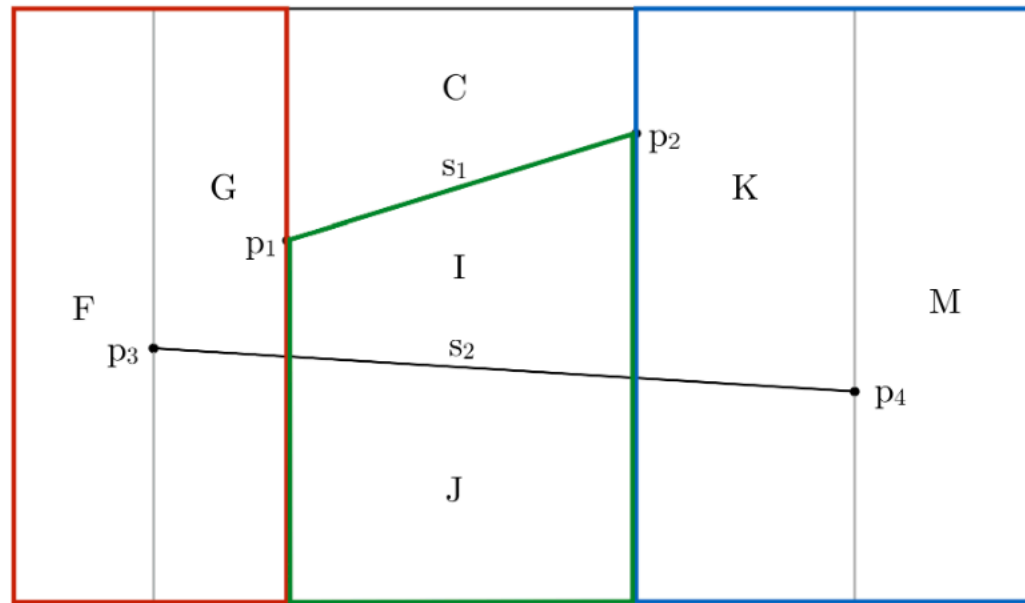


History

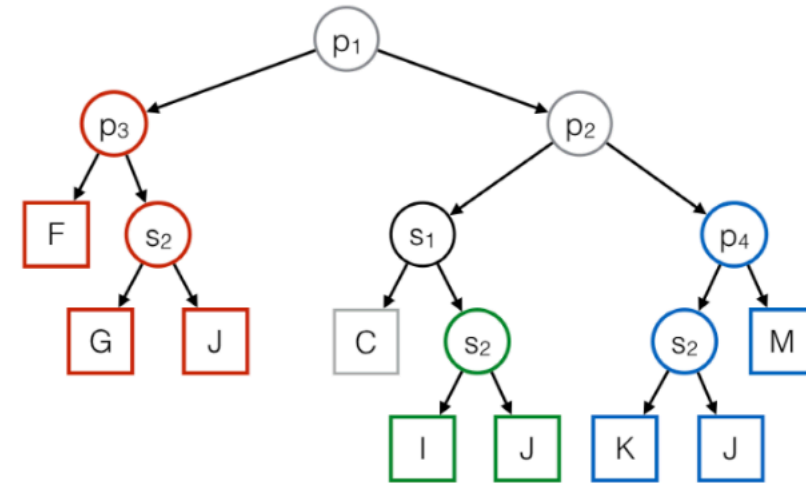


Step 6: Nodes of intersected trapezoids are replaced by new ones

TrapezoidalMap

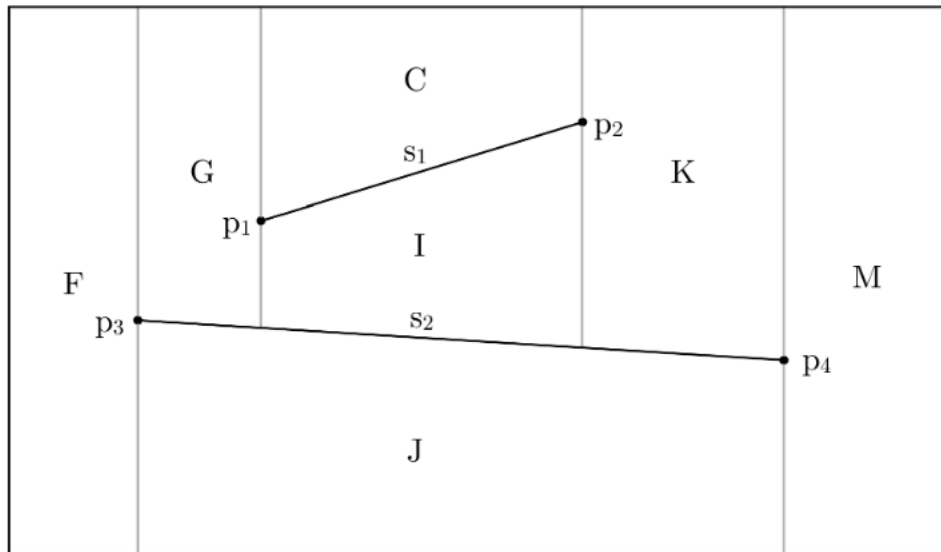


History

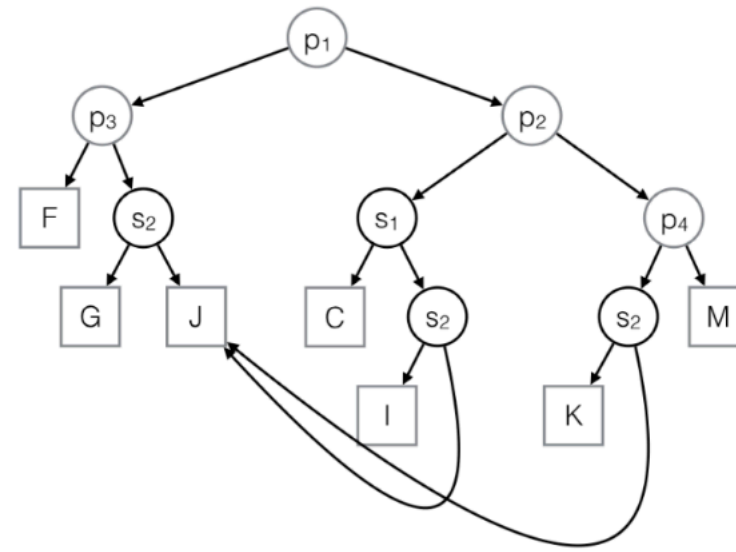


Same trapezoid nodes are linked together

## TrapezoidalMap

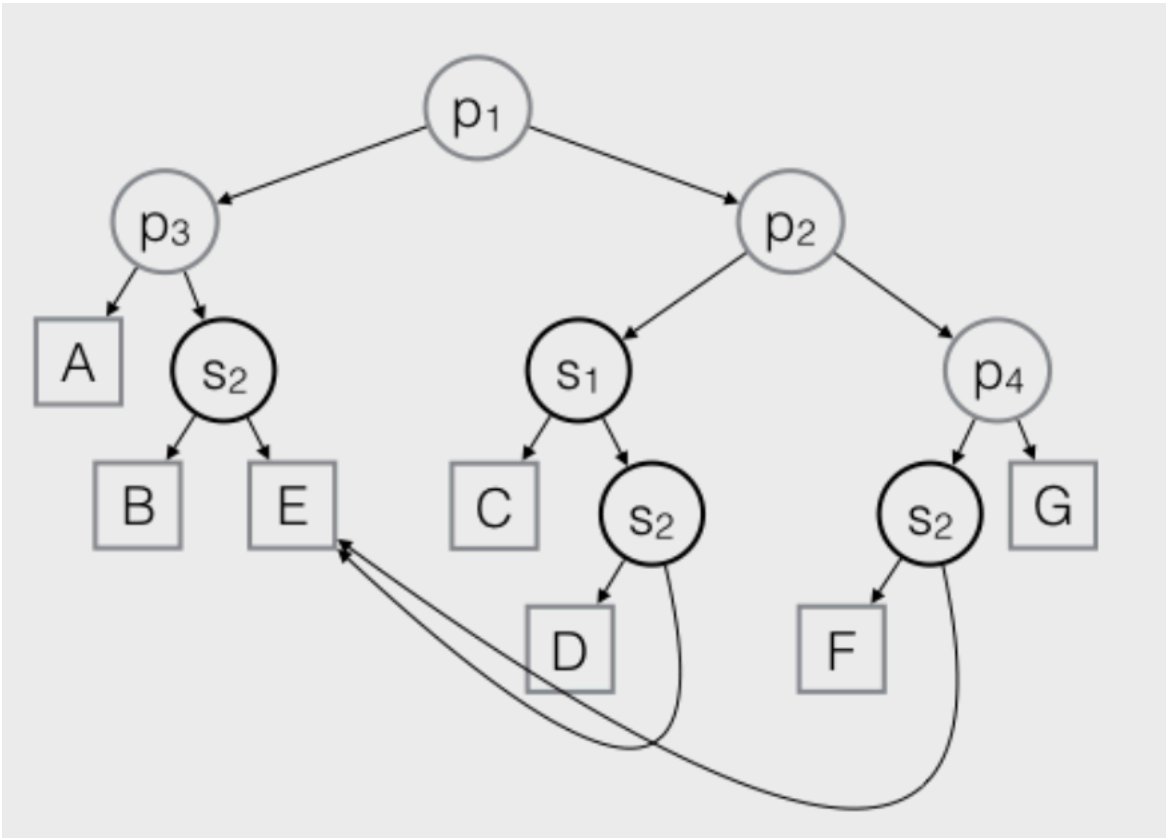


## History

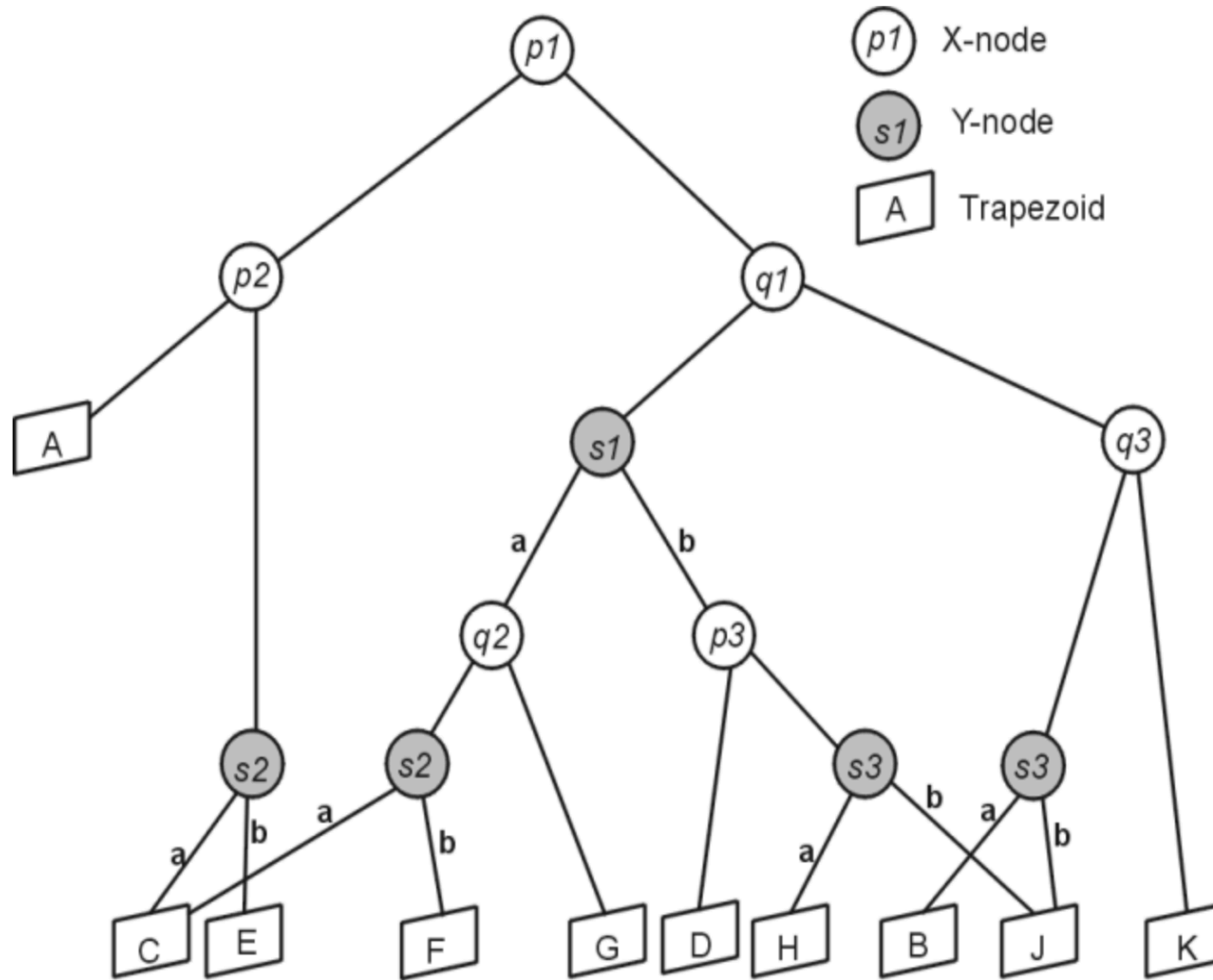




# Wyszukiwanie wielokąta, w którym znajduje się punkt wg struktury history



- X-nodes: Inner node storing an endpoint of a segment
- Y-nodes: Inner node storing a segment
- Leaf-nodes: Representing trapezoids



**Algorithm QueryTrapezoidMap( $D, n, p$ )**

*Input:*  $T$  is the trapezoid map search structure,  $n$  is a node in the search structure and  $p$  is a query point.

*Output:* A pointer to the node in  $D$  for the trapezoid containing the point  $p$ .

1. **if** ( $n$  is a Trapezoid Node)
  - return**  $n$ ;
2. **if** ( $n$  is an X-node)
  - a. **if** the x-coordinate of  $p$  is less than the x-coordinate of the point stored at this node then
    - return**  $\text{QueryTrapezoidMap}(T, \text{leftChild}(n), p)$ .
  - else**
    - return**  $\text{QueryTrapezoidMap}(T, \text{rightChild}(n), p)$
3. **if** ( $n$  is a Y-node)
  - a. **if**  $p$  is above the segment stored at  $n$  then
    - return**  $\text{QueryTrapezoidMap}(T, \text{aboveChild}(n), p)$ .
  - else**
    - return**  $\text{QueryTrapezoidMap}(T, \text{belowChild}(n), p)$

# Zalety podejścia

- Otrzymujemy strukturę trapezoid map. Złożoność czasowa operacji jej budowy wynosi  $O(n \log n)$
- Lokacja punktu dzięki strukturze history ma złożoność czasową  $O(\log n)$

# Bibliografia

<https://www.ti.inf.ethz.ch/ew/lehre/CG12/lecture/Chapter%209.pdf>

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[https://users.dimi.uniud.it/~claudio.mirolo/teaching/geom\\_comput/presentations/trapezoidal\\_map.pdf](https://users.dimi.uniud.it/~claudio.mirolo/teaching/geom_comput/presentations/trapezoidal_map.pdf)

<https://janrollmann.de/projects/thesis/>

<https://isotropic.org/papers/point-location.pdf>