

Determine RGBW LED PWM from CIE Chromaticity

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June 23, 2020

Abstract

This article is a private note to develop `kch-rgbw-lib`, available in [github](#) and [npm](#). `kch-rgbw-lib` is written in typescript, its main functions include color space conversions between HSV, RGB, XYZ and xyY, and calculation of color mixing. This article gives a general solution of multi-number (more than RGB) LEDs to represent different colors used in `kch-rgbw-lib`. It is not intended to carry new, accurate, or most efficient ideas. This document is granted under MIT License.

1 Define the Problem

1.1 Past works

Obtaining accurate color by mixing RGB color sources has been utilized as color displays since 1950s. As various colors are available by LED, recent topics are solution to expand additional colors or light source typically for OLED applications [1, 2]. Usually while light sources are used as an additional light source [3, 1, 2]. For display purposes colors other than RGB have been also used to expand the possible color ranges[4]. Sharp once added yellow in Aquos flat-panel TV, but they researched 5-primary color display [5]. Amber [6], turquoise, and violet can be other colors to expand the gamut.

1.2 Given parameters and assumptions

We have $n \geq 3$ color sources (LEDs) with chromaticity (x_i, y_i) and maximum luminosity Y_i , where $i = 1 \dots n$. Our problem is to find the optimum PWM output $\alpha = [\alpha_1 \dots \alpha_n]^T$ where $0 \leq \alpha_i \leq 1$ to represent a given color input with chromaticity (x, y) and luminosity Y .

Here, we set our goal of optimization as the following:

1. Under the physical constraint $0 \leq Y \leq \alpha_1 Y_1 + \dots + \alpha_n Y_n$,
2. Minimize the error of color to represent,

3. If (x, y) is outside the gamut of color source, nearest color in it is used,
 4. When possible, minimize energy consumption,
 5. When possible, set α_i to null when it's very small.
- Condition 3 can be achieved by projecting the input to the gamut contour.
 - Condition 4, energy consumption is obtained as

$$E = \sum_1^n \alpha_i W_i \quad (1.2.1)$$

where W_i is the power (W) of each LED at the maximum luminosity.

- Condition 5 is preferable to avoid jitter at low PWM output.

It is a typical linear programming (LP) problem. When $n = 3$, e.g. R-G-B color sources only, it's a deterministic and not an optimization problem. And when $n = 4$, e.g. R-G-B-W LEDs, there is only one parameter to optimize, which makes the problem as simple as we don't need to use sophisticated LP solver. We first derive a general description of the problem and solve it for $n = 3$, $n = 4$ and $n > 4$ cases.

1.3 Description of problem

We use XYZ color space because a composite of color source in XYZ color space can be obtained as a simple sum of each term. In XYZ color space, our problem is to determine $\boldsymbol{\alpha} = [\alpha_1 \dots \alpha_n]$ which gives an equation between input color $[X, Y, Z]^T$ and color source $[X_i, Y_i, Z_i]^T$, $i = 1 \dots n$;

$$\begin{bmatrix} X \\ Y \\ Z \end{bmatrix} = \alpha_1 \begin{bmatrix} X_1 \\ Y_1 \\ Z_1 \end{bmatrix} + \dots + \alpha_n \begin{bmatrix} X_n \\ Y_n \\ Z_n \end{bmatrix} \quad (1.3.2)$$

Color space $[X, Y, Z]^T$ is expressed by using (x_i, y_i, Y_i) :

$$X_i = \frac{x_i}{y_i} Y_i \quad (1.3.3)$$

$$Y_i = Y_i \quad (1.3.4)$$

$$Z_i = \frac{1 - x_i - y_i}{y_i} Y_i \quad (1.3.5)$$

Using matrix representation, Eq. 1.3.2 is written as

$$[\mathbf{X}] = [\mathbf{A}] [\boldsymbol{\alpha}] \quad (1.3.6)$$

where

$$[\mathbf{X}] = [X, Y, Z]^T \quad (1.3.7)$$

$$[\mathbf{A}] = \begin{bmatrix} \frac{x_1}{y_1} Y_1 & \dots & \frac{x_n}{y_n} Y_n \\ Y_1 & \dots & Y_n \\ \frac{1-x_1-y_1}{y_1} Y_1 & \dots & \frac{1-x_n-y_n}{y_n} Y_n \end{bmatrix} \quad (1.3.8)$$

$$[\boldsymbol{\alpha}] = [\alpha_1, \dots, \alpha_n]^T \quad (1.3.9)$$

Our goal is to solve Eq. 1.3.6 for $\boldsymbol{\alpha} = [\alpha_1 \dots \alpha_n]^T$. To solve it, \mathbf{A}^{-1} , the pseudo-inverse matrix of \mathbf{A} is obtained by the singular value decomposition (SVD) (Eq. 1.3.10) [7].

$$[\mathbf{A}] = [\mathbf{U}] \begin{bmatrix} \omega_1 & & & \\ & \omega_2 & & \\ & & \omega_3 & \\ & & & 0 \dots 0 \end{bmatrix} [\mathbf{V}^T] \quad (1.3.10)$$

where \mathbf{A} is $3 \times n$, \mathbf{U} is 3×3 , $[\omega_1 \dots \omega_3, 0 \dots 0]$ is $3 \times n$, \mathbf{V}^T is $n \times n$ matrixes. In this specific case, since \mathbf{A} is a $3 \times n$ matrix, there are upto 3 ω 's. When $n \geq 4$, $[\omega_1 \dots \omega_3]$ is null-padded in $3 \times (n - 3)$. \mathbf{A}^{-1} is obtained by

$$[\mathbf{A}]^{-1} = [\mathbf{V}_{1-3}] \begin{bmatrix} 1/\omega_1 & & \\ & 1/\omega_2 & \\ & & 1/\omega_3 \end{bmatrix} [\mathbf{U}^T] \quad (1.3.11)$$

where \mathbf{V}_{1-3} is the first 3 columns of \mathbf{V} , those correspond to $\omega_1 \dots \omega_3$. Using \mathbf{A}^{-1} , one obtains

$$[\boldsymbol{\alpha}] = [\mathbf{A}]^{-1} [\mathbf{X}] \quad (1.3.12)$$

By the way, what about the rest of columns in \mathbf{V} ? They are null vectors of \mathbf{A} . A null vector \mathbf{n} of \mathbf{A} is such vector that satisfies $\mathbf{A}\mathbf{n} = [0]$. By denoting these columns as $\mathbf{n}_4 \dots \mathbf{n}_n$, Eq. 1.3.12 can be extended as

$$[\boldsymbol{\alpha}] = [\mathbf{A}]^{-1} [\mathbf{X}] + \beta_4 \mathbf{n}_4 + \dots + \beta_n \mathbf{n}_n \quad (1.3.13)$$

where $\beta_4 \dots \beta_n$ are arbitrary numbers. By choosing these using other constraints, one can obtain the optimum solution.

2 Solution of $n = 3$ case

When $n = 3$, Eq. 1.3.12 gives a deterministic solution. No optimization. However, the obtained α_i should be physically meaningful, i.e., $0 \leq \alpha_i \leq 1$. It can be out of range when

- Input color $[X]$ is out of the gamut defined by the color sources,
- Y of $[X]$ is greater (brighter) than the color sources.

When $\alpha_i < 0$, it should be truncated to 0. In this case, the color has certain error.

When $\alpha_i > 1$, All α s should be normalized by the largest α . This way the color will be correctly obtained, but it will be darker than expected.

3 Solution of $n = 4$ case

When $n = 4$, Eq. 1.3.13 is modified as

$$[\alpha] = [\mathbf{A}]^{-1} [\mathbf{X}] + \beta_4 \mathbf{n}_4 \quad (3.0.1)$$

Parameter β_4 will be determined by the assumptions in section 1.2. Since it is always $0 \leq \alpha_i \leq 1$, by solving it for β (hereafter omitting '4'), we obtain the following conditions.

$$\beta \geq -\frac{b_i}{n_i} \quad i = 1 \dots 4, \text{ if } n \neq 0 \quad (3.0.2)$$

$$\beta \leq \frac{1 - b_i}{n_i} \quad i = 1 \dots 4, \text{ if } n \neq 0 \quad (3.0.3)$$

$$E = \sum_{i=1}^4 (\beta n_i + b_i) W_i \rightarrow \min \quad (3.0.4)$$

where $[b_1, \dots, b_4]^T = [\mathbf{A}]^{-1} [\mathbf{X}]$, n_i are the elements of \mathbf{n} . Eq. 3.0.4 is from Eq. 1.2.1. Finding the largest and smallest values of the right hand side of Eqs. 3.0.2 and 3.0.3, denoted as β_{min} and β_{max} , Eqs. 3.0.2 and 3.0.3 are rewritten as

$$\beta_{min} \leq \beta \leq \beta_{max} \quad (3.0.5)$$

Since Eq. 3.0.4 is rewritten as $E = s_1 \beta + s_2$, here s_1 and s_2 are constants determined by calculating the sums in Eq. 3.0.4, optimized β is determined as

$$\beta = \begin{cases} \beta_{min} & \text{if } s_1 = \sum_1^4 n_i W_i > 0 \\ \beta_{max} & \text{else} \end{cases} \quad (3.0.6)$$

3.1 When β is not determined

When $\beta_{min} > \beta_{max}$, there is no possible answer. This can happen when Y in Eq. 1.3.7 is greater (brighter) than the color source's capacity. When it happens, since Eq. 3.0.2 cannot compromise, set β to β_{min} and normalize the largest α_i to be 1.

3.2 When α_i is small

To implement the assumption 5 in section 1.2, you introduce allowance of small α , α_ε , and exchange Eq. 3.0.2 as

$$\beta \geq \frac{\alpha_\varepsilon - b_i}{n_i} \quad (3.2.7)$$

But this may also result in $\beta_{min} > \beta_{max}$. When it is the case, you need to relax the constraint by α_ε . The easiest is to go back to Eq. 3.0.2.

4 Solution of $n > 4$ case

You need to optimize Eq. 1.2.1 under constraints of $0 \leq \alpha_i \leq 1$ and Eq. 1.3.13 using a linear programming solution. We will implement it in future.

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