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Q1 a) There are 4 different paths.

Path 1: 1, 2, 3, 4, 9, 11, 12, 13, 17

Path 2: 1, 2, 3, 4, 9, 11, 15, 16, 17

Path 3: 1, 2, 6, 7, 9, 11, 12, 13, 17

Path 4: 1, 2, 6, 7, 9, 11, 15, 16, 17

b) Path 1:	Edge	Symbolic State (PV)	Path Condition (PC)
	1→2	$x \mapsto X_0, y \mapsto Y_0$	true
	2→3	$x \mapsto X_0, y \mapsto Y_0$	$X_0 + Y_0 > 15$
	3→4	$x \mapsto X_0 + 7, y \mapsto Y_0$	$X_0 + Y_0 > 15$
	4→9	$x \mapsto X_0 + 7, y \mapsto Y_0 - 12$	$X_0 + Y_0 > 15$
	9→11	$x \mapsto X_0 + 9, y \mapsto Y_0 - 12$	$X_0 + Y_0 > 15$
	11→12	$x \mapsto X_0 + 9, y \mapsto Y_0 - 12$	$X_0 + Y_0 > 15 \wedge 2(X_0 + 9 + Y_0 - 12) > 21$
	12→13	$x \mapsto 3(X_0 + 9), y \mapsto Y_0 - 12$	$X_0 + Y_0 > 15 \wedge 2(X_0 + 9 + Y_0 - 12) > 21$
	13→17	$x \mapsto 3(X_0 + 9), y \mapsto 2(Y_0 - 12)$	$X_0 + Y_0 > 15 \wedge 2(X_0 + 9 + Y_0 - 12) > 21$

Path 2

Edge	Symbolic State (PV)	Path Condition (PC)
1→2	$x \mapsto X_0, y \mapsto Y_0$	true
2→3	$x \mapsto X_0, y \mapsto Y_0$	$X_0 + Y_0 > 15$
3→4	$x \mapsto X_0 + 7, y \mapsto Y_0$	$X_0 + Y_0 > 15$
4→9	$x \mapsto X_0 + 7, y \mapsto Y_0 - 12$	$X_0 + Y_0 > 15$
9→11	$x \mapsto X_0 + 9, y \mapsto Y_0 - 12$	$X_0 + Y_0 > 15$
11→15	$x \mapsto X_0 + 9, y \mapsto Y_0 - 12$	$X_0 + Y_0 > 15 \wedge 2(X_0 + 9 + Y_0 - 12) \leq 21$
15→16	$x \mapsto 4(X_0 + 9), y \mapsto Y_0 - 12$	$X_0 + Y_0 > 15 \wedge 2(X_0 + 9 + Y_0 - 12) \leq 21$
16→17	$x \mapsto 4(X_0 + 9), y \mapsto 3*(Y_0 - 12) + 4(X_0 + 9)$	$X_0 + Y_0 > 15 \wedge 2(X_0 + 9 + Y_0 - 12) \leq 21$

Path 3:

Edge	Symbolic State (PV)	Path Condition (PC)
1→2	$x \mapsto X_0, y \mapsto Y_0$	true
2→6	$x \mapsto X_0, y \mapsto Y_0$	$X_0 + Y_0 \leq 15$
6→7	$x \mapsto X_0, y \mapsto Y_0 + 10$	$X_0 + Y_0 \leq 15$
7→9	$x \mapsto X_0 - 2, y \mapsto Y_0 + 10$	$X_0 + Y_0 \leq 15$
9→11	$x \mapsto X_0, y \mapsto Y_0 + 10$	$X_0 + Y_0 \leq 15$
11→12	$x \mapsto X_0, y \mapsto Y_0 + 10$	$X_0 + Y_0 \leq 15 \wedge 2(X_0 + Y_0 + 10) > 21$
12→13	$x \mapsto 3X_0, y \mapsto Y_0 + 10$	$X_0 + Y_0 \leq 15 \wedge 2(X_0 + Y_0 + 10) > 21$
13→17	$x \mapsto 3X_0, y \mapsto 2(Y_0 + 10)$	$X_0 + Y_0 \leq 15 \wedge 2(X_0 + Y_0 + 10) > 21$

Path 4:

Edge	Symbolic State (PV)	Path Condition (PC)
1→2	$x \mapsto X_0, y \mapsto Y_0$	true
2→6	$x \mapsto X_0, y \mapsto Y_0$	$X_0 + Y_0 \leq 15$
6→7	$x \mapsto X_0, y \mapsto Y_0 + 10$	$X_0 + Y_0 \leq 15$
7→9	$x \mapsto X_0 - 2, y \mapsto Y_0 + 10$	$X_0 + Y_0 \leq 15$
9→11	$x \mapsto X_0, y \mapsto Y_0 + 10$	$X_0 + Y_0 \leq 15$
11→15	$x \mapsto X_0, y \mapsto Y_0 + 10$	$X_0 + Y_0 \leq 15 \wedge 2(X_0 + Y_0 + 10) \leq 21$
15→16	$x \mapsto 4X_0, y \mapsto Y_0 + 10$	$X_0 + Y_0 \leq 15 \wedge 2(X_0 + Y_0 + 10) \leq 21$
16→17	$x \mapsto 4X_0, y \mapsto 3(Y_0 + 10) + 4X_0$	$X_0 + Y_0 \leq 15 \wedge 2(X_0 + Y_0 + 10) \leq 21$

C): Path 1 is feasible : $X_0 = 8, Y_0 = 8$.

$$X_0 + Y_0 = 16 > 15 \wedge 2*(8+9+8-12) = 26 > 21$$

Path 2 is not feasible: the path condition is $X_0 + Y_0 > 15 \wedge 2(X_0 + 9 + Y_0 - 12) \leq 21$
 which can be simplified as: $X_0 + Y_0 > 15 \wedge X_0 + Y_0 \leq 13.5$. This condition is always false.

Path 3 is feasible: $X_0 = 8, Y_0 = 7$.

$$X_0 + Y_0 = 15 \leq 15 \wedge 2*(8+7+10) = 50 > 21$$

Path 4 is feasible: $X_0 = 0, Y_0 = 0$

$$X_0 + Y_0 = 0 < 15 \wedge 2*(0+0+10) = 20 < 21$$

- Q2: a) CNF:
- $$\begin{aligned}
 & (a_1 \vee a_2 \vee \neg a_3 \vee \neg a_4) \wedge \\
 & (\neg a_1 \vee \neg a_2 \vee a_3 \vee \neg a_4) \wedge \\
 & (a_1 \vee \neg a_2 \vee \neg a_3 \vee a_4) \wedge \\
 & (\neg a_1 \vee \neg a_2 \vee \neg a_3 \vee \neg a_4) \wedge \\
 & (\neg a_1 \vee a_2 \vee a_3 \vee \neg a_4) \wedge \\
 & (\neg a_1 \vee a_2 \vee \neg a_3 \vee a_4) \wedge \\
 & (\neg a_1 \vee \neg a_2 \vee \neg a_3 \vee a_4) \wedge \\
 & (\neg a_1 \vee \neg a_2 \vee a_3 \vee \neg a_4) \wedge \\
 & (\neg a_1 \vee \neg a_2 \vee \neg a_3 \vee \neg a_4)
 \end{aligned}$$

a1	a2	a3	a4	F
0	0	0	0	1
0	0	0	1	1
0	0	1	0	1
0	0	1	1	0
0	1	0	0	1
0	1	0	1	0
0	1	1	0	0
0	1	1	1	0
1	0	0	0	1
1	0	0	1	0
1	0	1	0	0
1	1	0	0	0
1	1	0	1	0
1	1	1	0	0
1	1	1	1	0

can be simplified as:

$$(\neg a_1 \vee \neg a_2) \wedge (\neg a_1 \vee \neg a_3) \wedge (\neg a_1 \vee \neg a_4) \wedge (\neg a_2 \vee \neg a_3) \wedge (\neg a_2 \vee \neg a_4) \wedge (\neg a_3 \vee \neg a_4)$$

b) It is valid.

① subclaim: $\forall x \exists y (P(x) \vee Q(y)) \Rightarrow (\forall x P(x)) \vee (\exists y Q(y))$.

Prove by contradiction

Assume LHS is true and RHS is false

LHS means for all x there exist some y makes $P(x)$ or $Q(y)$ is true.

RHS is false means for all x $P(x)$ is false and there is no such y that makes $Q(y)$ true.

Suppose $\forall x P(x)$ is false then $\exists y Q(y)$ must be true to make LHS true.

Similarly, suppose $\exists y Q(y)$ is false then $\forall x P(x)$ must be true to make LHS true.

That is a contradiction to our assumption of RHS.

Therefore, we can prove this subclaim is true.

② subclaim: $(\forall x P(x)) \vee (\exists y Q(y)) \Rightarrow \forall x \exists y (P(x) \vee Q(y))$

Suppose LHS is true:

① $\forall x P(x)$ is true and $\exists y Q(y)$ is false, then $P(x)$ is true for every x .

That means RHS is true regardless $Q(y)$

② $\exists y Q(y)$ is true and $\forall x P(x)$ is false, then for any x we can pick a y to make $P(x) \vee Q(y)$ true. RHS is also true.

Therefore, we can prove this FOL is valid.

c) It is invalid.

Proof: $(\forall x \exists y P(x,y) \vee Q(x,y)) \Rightarrow (\forall x \exists y P(x,y)) \vee (\forall x \exists y Q(x,y))$
Suppose $x \in \{0,1\}$, $y \in \{0,1\}$,

$P(1,1)$ is true, $P(1,0)$ is false, $P(0,1)$ is false, $P(0,0)$ is false

$Q(1,1)$ is true, $Q(0,1)$ is true, $Q(0,0)$ is false, $Q(1,0)$ is false

For LHS: when $x=0$, $y=1$, $P(x,y) \vee Q(x,y)$ is true

when $x=1$, $y=1$, $P(x,y) \vee Q(x,y)$ is true

\therefore LHS is true

For RHS $\forall x \exists y P(x,y)$: When $x=0$, there is no y to make $P(x,y)$ true

$\therefore \forall x \exists y P(x,y)$ is false.

For RHS $\forall x \exists y Q(x,y)$: When $x=0$, there is no y to make $Q(x,y)$ true.

$\therefore \forall x \exists y Q(x,y)$ Therefore, RHS is false. We disprove this implication.

d) $\Phi = \exists x \exists y \exists z (P(x,y) \wedge P(z,y) \wedge P(x,z) \wedge \neg P(z,x))$

(a) $M_1 = \langle S_1, P_1 \rangle$, where $S_1 = N$ and $P_1 = \{ (x,y) \mid x, y \in N \wedge x < y \}$.

Suppose $x=1, y=3, z=2$

Then $P(x,y)$ is true because $x < y$.

$P(z,y)$ is true because $z < y$.

$\neg P(z,x), P(x,z)$ are true because $x < z$

$\therefore M_1 \models \Phi$

(b) $M_2 = \langle S_2, P_2 \rangle$ where $S_2 = N$ and $P_2 = \{ (x, x+1) \mid x \in N \}$

Suppose $P(x,y)$ is true then $y = x+1$

$P(z,y)$ is true then $y = z+1$

Then $x = z$

Then $P(x,z)$ is false because $x = z$.

$\therefore M_2 \not\models \Phi$

c) $M_3 = \langle S_3, P_3 \rangle$, where $S_3 = \mathcal{P}(N)$, the powerset of natural number and

$P_3 = \{ (A, B) \mid A, B \in \mathcal{P}(N) \wedge A \subseteq B \}$

Suppose x contains 1 number

y contains 3 numbers

z contains 2 numbers

Then $P(x,y)$ is true because $x \subseteq y$.

$P(z,y)$ is true because $z \subseteq y$

$\neg P(z,x), P(x,z)$ are true because $x \subseteq z$.

$\therefore M_3 \models \Phi$

e) As Boolean variables a_i has disjunction compared to other a_j .

Then we have $\bigwedge_{1 \leq i < j \leq n} (\neg a_i \vee \neg a_j)$.

Q3 a)

Assume n is the size of the square, M is the sum, $r \geq 1$ row number and
Val is the element value, c is column number.

① each cell element is distinct and its value is smaller than n^2 :

$$\bigwedge_{\substack{1 \leq i, j, l, k \leq n \\ (i \neq l \vee j \neq k)}} (\text{Val}_{ij} \neq \text{Val}_{lk})$$

$$\bigwedge_{1 \leq i, j \leq n} (1 \leq \text{Val}_{ij} \leq n^2)$$

② sum of each row, column, diagonal are equal.

$$\bigwedge_{1 \leq r \leq n} \sum_{j=1}^n C_{rj} = M$$

$$\sum_{r=1}^n C_{rr} = M$$

$$\sum_{r=1}^n C_{r(n-r+1)} = M$$

$$\bigwedge_{1 \leq c \leq n} \sum_{j=1}^n C_{jc} = M$$

Q4 e: The program is the last test function in my test_sym.py file. I took 172 seconds to run.

```
#Testing divergence
def test_diverge(self):
    prg1 = "havoc x; c := 1; while c < 5 do {c:=c+1;if x < -100 then {f := x + 1;
x := x/100 + f}; if x < -1000 then {f := 2 * x; x := x/10 * f + x}; if x < 1000 then
{f := 3 / x;}" \
           " x := x * f + 10}; if x > 0 then {f := x / 4; x := x*5/f}; if x >= 10 then {f :=
5+x; x := x-20*f}; if x = 100 then {f := 6-x; x := 100/f}}; print_state "
    ast1 = ast.parse_string(prg1)
    engine = sym.SymExec()
    st = sym.SymState()
    out = [s for s in engine.run(ast1, st)]
    self.assertEqual(len(out), 1)
```

```
pc: [10 < 21, 11 < 21, 12 < 21, 13 < 21, 14 < 21, 15 < 21, 16 < 21, 17 < 21, 18 < 21, 19 < 21]
..x: x!536
pc: [10 > x!536]

x: -1
pc: [Not(10 > x!536)]

.x: 10
y: 10
z: 0
a: 100
b: 1
pc: []

.x: x!543
pc: [10 < x!543, 9 < x!543]

...
-----
Ran 29 tests in 173.496s
OK
```