12/05/2016

ALGORITHMS

HIW - 12

Kiran Shettar UMIL ID-016058 Show that it HAM-CYCLE & P then the problem of listing the vertices of a hamiltonian cycle, in order, is polynomial time solvable.

Suppose that 'G' is hamiltonian. which means that there is a hamiltonian cycle. Pick any one vertex 'v' in the graph, and consider au the possibilities of deleting all but two of the edges passing through that vertex. For some pair of edges to save, the resulting graph must still be hamiltonian because the hamiltonian graph that existed originally only used two edges. Since the degree of the vertex is bounded by the number of vertices minus one we are only less than squaring that number by looking at all pairs ((n-1) e O(n2)). This means that we are only running the polynomial tester polynomially many independent -t times, so the runtime is polynomial. Once we have all the pair of vertices where deting all the others coming off 'v'.

still results in a hamiltonian graph, we will remember those as special, and once that we will never again try to delete. We repeal the processes with both of the vertices that are now adjacent to 'v', testing hamiltonicity of each way of picking a new vertex to save. We continue in this process until we are left with only |v| eage, and so, we have to just constructed a hamiltonian cycle.

Problem-2: (34.3-2) Pg 1077

Show that the $\leq p$ relation is a transitive relation on languages. That is, show that if:

that is

((stand then),)) was a man to produce the

31 para L. Sp. L3:

Now, det Liép La be (i)

Laép La be (ii)

Let f, be the polynomial time reduction g function such that $x \in L_1$ if g only if