

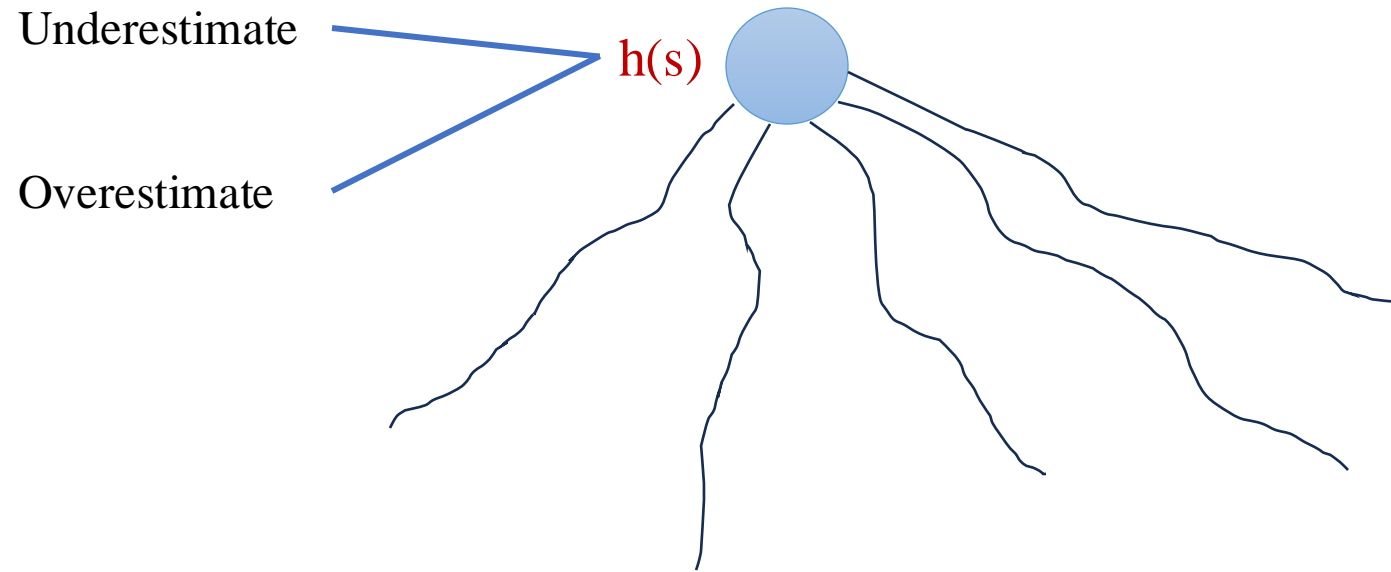
Informed State Space Search

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The notion of heuristics

- Heuristics use domain specific knowledge to estimate the quality or potential of partial solutions



The notion of heuristics

- Examples:
 - Manhattan distance heuristic for 8 puzzle

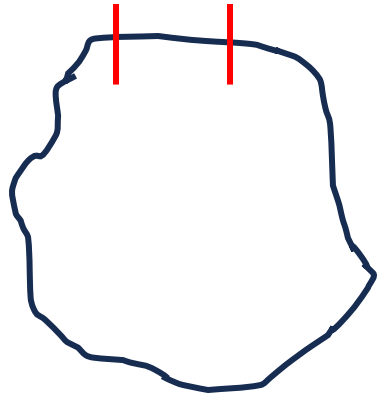
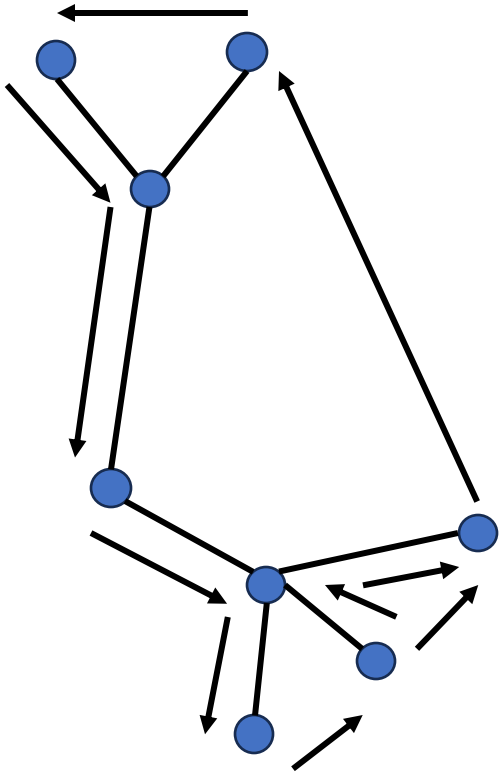
5	6	7
4	1	8
3	9	

1	2	3
4	5	6
7	8	

$$2 + 0 + 4$$

The notion of heuristics

- Examples:
 - Minimum spanning tree heuristic for TSP



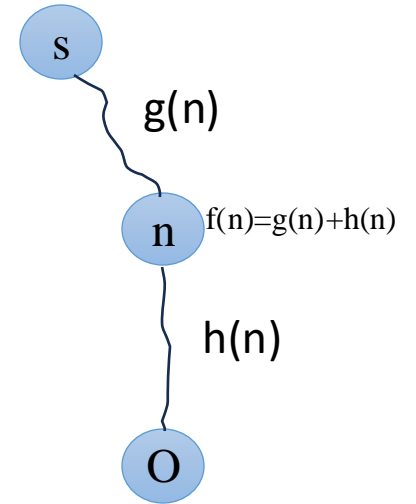
$$C_S < C^* < 2C_S$$

The informed search problem

- **Given:** $[S, s, O, G, h]$ where
 - S is the (implicitly specified) set of states
 - s is the start state
 - O is the set of state transition operators each having some cost
 - G is the set of Goal states
 - $h()$ is a heuristic function estimating the distance to a goal
- **To find:**
 - A minimum cost sequence of transitions to a goal state

Algorithm A*

- **Initialize:** Set $OPEN = \{s\}$, $CLOSED = \{\}$, $g(s)=0$, $f(s) = h(s)$
- **Fail:**
 - If $OPEN = \{\}$, Terminate with failure
- **Select:** Select the minimum cost state, n , from $OPEN$ and save in $CLOSED$
- **Terminate:**
 - If $n \in G$, terminate with success

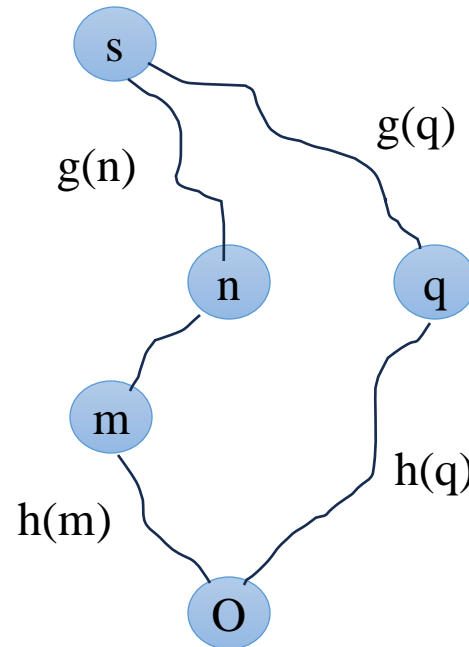


Algorithm A*

- **Expand:**
 - For each successor, m , of n :
 - If $m \notin [\text{OPEN} \cup \text{CLOSED}]$
 - Set $g(m) = g(n) + C(n, m)$
 - Set $f(m) = g(m) + h(m)$
 - Insert m in OPEN
 - If $m \in [\text{OPEN} \cup \text{CLOSED}]$
 - Set $g(m) = \min \begin{cases} g(m) \\ g(n) + C(n, m) \end{cases}$
 - Set $f(m) = g(m) + h(m)$
 - If $f(m)$ has decreased and $m \in \text{CLOSED}$
 - Move m to OPEN
- **Loop:**
 - Go to step 2

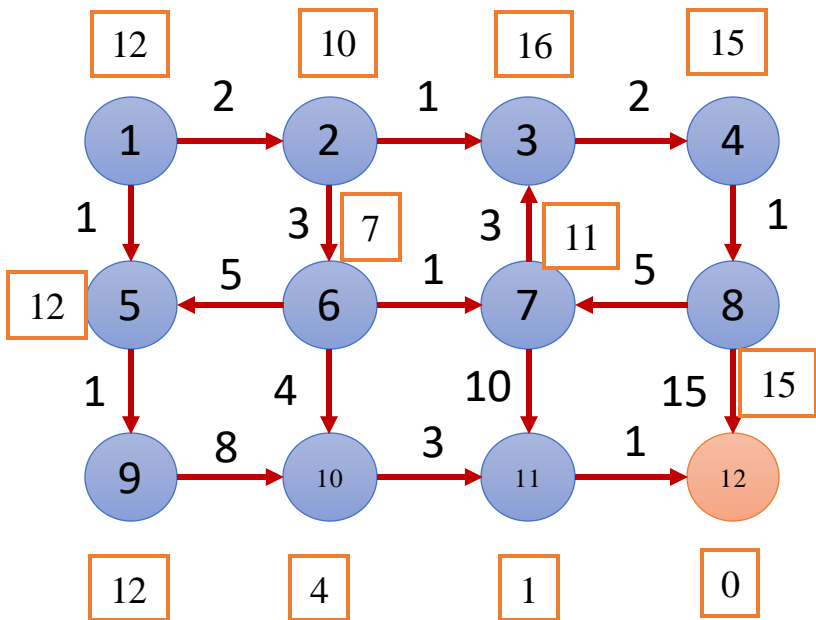
Algorithm A*

- How can a promising path become non-promising?
 - When moving from n to m, we may find a path with less heuristic cost



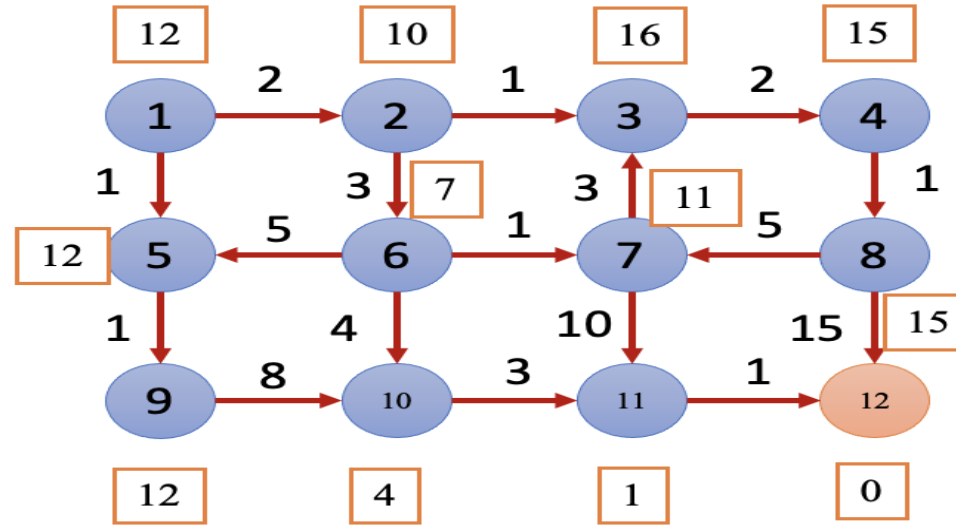
$$h(q) \gg h(m)$$

Algorithm A*



OPEN SET	SELECT	GOAL	EXPANDED	CLOSED
[1(12)]	1(12)	N	[2(12),5(13)]	[1(12)]
[2(12),5(13)]	2(12)	N	[5(13),3(19),6(12)]	[1(12),2(12)]
[5(13),3(19),6(12)]	6(12)	N	[5(13),3(19),7(17),10(13)]	[1(12),2(12),6(12)]
[5(13),3(19),7(17),10(13)]	5(13)	N	[3(19),7(17),10(13),9(14)]	[1(12),2(12),6(12),5(13)]
[3(19),7(17),10(13),9(14)]	10(13)	N	[3(19),7(17),9(14),11(13)]	[1(12),2(12),6(12),5(13),10(13)]

Algorithm A*



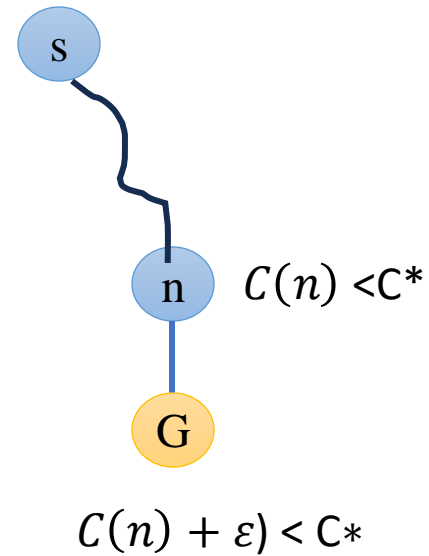
OPEN SET	SELECT	GOAL	EXPANDED	CLOSED
[3(19),7(17),10(13),9(14)]	10(13)	N	[3(19),7(17),9(14),11(13)]	[1(12),2(12),6(12),5(13),10(13)]
[3(19),7(17),9(14),11(13)]	11(13)	N	[3(19),7(17),9(14),12(13)]	[1(12),2(12),6(12),5(13),10(13),11(13)]
[3(19),7(17),9(14),12(13)]	12(13)	Y		

Algorithm A*: Benefit

- Reduces number of expanded nodes
- Performs the lookahead and tells us promising paths
- What about optimality?

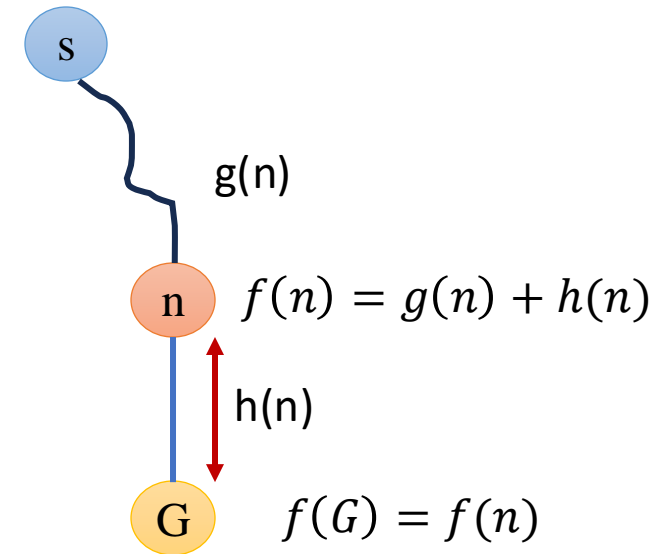
Uniform Cost Search

- **Claim:** If $C(n) < C^*$ (optimal cost) then n must be expanded
- Let algorithm A does not expand n
- For the class of algorithms without any heuristics
 - All states that have cost $< C^*$ will have to be expanded
- Always **expands the minimum cost** node in your frontier
- When we find the goal
 - All the states that we have in the frontier have cost higher than the goal state



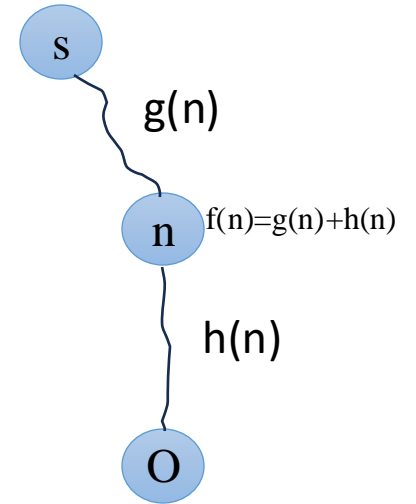
Algorithm A*: Benefit

- **Claim:** $f(n) < C^*$ then n must be expanded
- The heuristic function **underestimates**
 - $h(n) \leq f^*(n)$
 - Cost of reaching goal from n
 - All costs are +ve
- If we do not expand n, we can't find the goal
- If we have a state whose cost is less than C^*
 - Then every algorithm which guarantees finding optimal solution have to expand it



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How to break the tie?

Thank You