

# Assignment 15

## Automata & Theory of Computation

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Student ID: 2023065350

Name: 危立

1. Construct an npda that accepts the language generated by a grammar with productions

$$S \rightarrow aSSSab|\lambda.$$

$$S \rightarrow aSSA|1$$

$$A \rightarrow aB$$

$$B \rightarrow b$$

$$\delta(q_0, \lambda, z) = \{(q, \lambda)\}$$

$$\delta(q_1, a, S) = \{(q_1, SSA)\}$$

$$\delta(q_1, 1, S) = \{(q_1, 1)\}$$

$$\delta(q_1, a, A) = \{(q_1, B)\}$$

$$\delta(q_1, b, B) = \{(q_1, 1)\}$$

$$\delta(q_1, \lambda, z) = \{(q_2, \lambda)\}$$

$$③ (q_0 A q_1) \rightarrow a$$

$$④ (q_1 z q_2) \rightarrow \lambda$$

$$⑤ (q_1 A q_1) \rightarrow \lambda$$

2. Construct a context-free grammar for the language accepted by the npda

$M = (\{q_0, q_1\}, \{a, b\}, \{A, z\}, \delta, q_0, z, \{q_1\})$ , with transitions

$q_2$

$$① \delta(q_1, a, A) = \{(q_1, \lambda)\} \rightarrow \text{③과 ④ } q_1 \text{ 에 } A \text{ 를 } Stack \text{ 에 푸는 } X$$

$$① \delta(q_0, a, z) = \{(q_0, Az)\},$$

$$② \delta(q_0, b, A) = \{(q_0, AA)\},$$

$$③ \delta(q_0, a, A) = \{(q_1, \lambda)\}.$$

$$④ \delta(q_1, \lambda, z) = \{(q_2, \lambda)\} \rightarrow (q_1 z q_2) \rightarrow \lambda$$

$$③ \delta(q_0, a, A) = \lambda \rightarrow (q_0 A q_1) \rightarrow a$$

$$① (q_0 z q_0) \rightarrow a(q_0 A q_0) | a(q_0 A q_1) | a(q_0 A q_2)$$

$$(q_0 z q_1) \rightarrow a(q_0 A q_0) | a(q_0 A q_1) | a(q_0 A q_2)$$

Start

$$(q_0 z q_2) \rightarrow a(q_0 A q_0) | a(q_0 A q_1) | a(q_0 A q_2)$$

loop  $\rightarrow$  시작

$$② (q_0 A q_0) \rightarrow b(q_0 A q_0) | b(q_0 A q_1) | b(q_0 A q_2)$$

$$(q_0 A q_1) \rightarrow b(q_0 A q_0) | b(q_0 A q_1) | b(q_0 A q_2)$$

$$(q_0 A q_2) \rightarrow b(q_0 A q_0) | b(q_0 A q_1) | b(q_0 A q_2)$$

$$(q_0 A q_1) \rightarrow a | b(q_0 A q_1) | \text{cancel}(q_0 A q_1)$$

$$S \rightarrow aA$$

$$\text{cancel}(q_0 A q_0) \rightarrow \lambda$$

$$S \rightarrow a | bA$$

$$\text{cancel}(q_1 z q_2) \rightarrow \lambda$$

$$S (q_0 z q_2) \rightarrow a(q_0 A q_1) | \text{cancel}(q_0 z q_2)$$