

Surreal # ($\mathbb{S} \Rightarrow \{ \text{all the surreal } \#\}'$)

Part 1 Definitions.

Def 1: A surreal number x is a pair of previously created surreal numbers, known as the left set \bar{X}_L and the ~~right~~ set \bar{X}_R with no member of the right set may be less than or equal to any member of \bar{X}_L .

Def 2: $x, y \in \mathbb{S}$ $x \leq y$ iff the followings:

$$\bullet \exists x_L \in \bar{X}_L : y \leq x_L \quad (*1)$$

$$\bullet \exists y_R \in \bar{X}_R : y_R \leq x \quad (*2)$$

Def (zero). We define zero as: $0 := \{1\} := \{\emptyset | \emptyset\}$

Thm: $0 \leq 0$

Proof: Since O_L and O_R are empty, then the conditions in def 2 holds. $\Rightarrow 0 \leq 0$ \square

Def $-1 := \{0\}$, $1 := \{0\}'$

Rmk: Check the well-definedness of -1 and 1 . Since $(-1)_L$ and $(1)_R$ are empty, then the condition for well-definedness in def 1 holds naturally.

Thm $0 \leq 1$ and $-1 \leq 0$

Proof: For " $0 \leq 1$ ", $O_L \& I_R$ are empty. Then " $0 \leq 1$ " holds. For " $-1 \leq 0$ ", $(-1)_L$ and O_R are empty. Then " $-1 \leq 0$ " holds. Hence, we're done. \square

Notational Conventions

$x \neq y$ means $\neg(x \leq y)$

$x < y$ means $x \leq y$ and $\neg(y \leq x)$

$x \neq y$ means $\neg(x < y) \Leftrightarrow x \neq y \vee y \leq x$

$x = y$ means $x \leq y$ and $y \leq x$

$x \neq y$ means $\neg(x = y)$

Theorem, $0 = 0$

Proof: $0 \leq 0$ & $0 \leq 0 \Rightarrow 0 = 0$

Thm: $\stackrel{\textcircled{1}}{1=1}$, $\stackrel{\textcircled{2}}{-1=-1}$, $\stackrel{\textcircled{3}}{0<1}$, $\stackrel{\textcircled{4}}{-1<0}$, $\stackrel{\textcircled{5}}{-1<1}$

Proof:

① W.T.S: $1 \leq 1 \Leftrightarrow \exists x \in 1_L : 1 \leq x \wedge \exists y \in 1_R : y \leq 1$

Since $1_R = \emptyset$, then (ii) holds. For (i), $1_L = \{0\}$, so (i) holds because $0 \leq 1$ by the previous thm.

② $\neg \exists x \in (-1)_L : -1 \leq x \wedge \neg \exists y \in (-1)_R : y \leq -1$

Since $(-1)_L = \emptyset$, then (iii) holds. For (iv), $(-1)_R = \{0\}$, so (iv) holds because $-1 \leq 0$ by the previous thm.

③ W.T.S: $1 \neq 0 \Leftrightarrow \neg(1 \leq 0) \Leftrightarrow \exists x \in 1_L : 0 \leq x \vee \exists y \in 0_R : y \leq 1$

Since $1_L = \{0\}$ and $0 \leq 0$, then " $\exists x \in 1_L : 0 \leq x$ " holds.

④ W.T.S: $0 \neq -1 \Leftrightarrow \exists x \in 0_L : -1 \leq 0_L \vee \exists y \in (-1)_R : y \leq 0$

Since $(-1)_R = \{0\}$ and $0 \leq 0$, then " $\exists y \in (-1)_R : y \leq 0$ " holds.

⑤ W.T.S: $-1 \leq 1 \wedge 1 \neq -1$

(a) $\neg \exists x \in (-1)_L = \emptyset \wedge \neg \exists y \in 1_R = \emptyset \Rightarrow \neg \exists x \in (-1)_L : x \leq -1 \wedge \neg \exists y \in 1_R : y \leq 1 \Rightarrow -1 \leq 1$

(b) $-1 \leq 0 \Rightarrow \exists x \in 1_L = \{0\} (x=0) : -1 \leq x \Rightarrow 1 \neq -1$

Def $z := \{1\}$, $-z := \{-1\}$, $\frac{1}{2} := \{01\}$, $-\frac{1}{2} := \{-10\}$

Rmk: The well-definedness of $z, -z, \frac{1}{2}, -\frac{1}{2}$ can be easily check by the thms on the left side of this page.

Thm $0 = \{-11\}$

Proof:

① $0 \leq \{-11\} : \neg \exists x \in 0_L : \{11\} \leq x \Rightarrow \neg \exists x \in 0_L : \{11\} \leq x$

$0 < 1 \Rightarrow 1 \neq 0 \Rightarrow \neg \exists y \in \{1\} : y \leq 0$

$\Rightarrow \neg \exists y \in \{-11\}_R : y \leq 0 \quad (\#2)$

Then by (#1) (#2), $0 \leq \{-11\}$

② $\{-11\} \leq 0$:

$\neg \exists y \in 0_R = \emptyset \Rightarrow \neg \exists y \in 0_R : y \leq \{-11\}$

$-1 < 0 \Rightarrow 0 \neq -1 \Rightarrow \neg \exists x \in \{-1\} : 0 \leq x$

$\Rightarrow \neg \exists x \in \{-11\}_L : 0 \leq x \quad (\#1)$

By (#1) (#2), $\{-11\} \leq 0$

then by ① ②, we get that $0 = \{-11\}$. \square

Def (Representative).

The different forms that represent the same surreal number are called representatives of that surreal number

Example: $\{1\}$ and $\{-1\}\{1\}$ are two representatives of the surreal number 0 .

Def (Birthday)

- $0 = \{1\}$ is born on day 0
- $-1 = \{1\}$, $1 = \{01\}$ are born on day 1 .
- ⋮

Part 2 Basic Properties of Surreal Numbers

Def (Set comparisons). Let A be a set of surreal numbers, c be a surreal number.

Then we define:

$$A \leq c \Leftrightarrow \forall a \in A : a \leq c$$

$$c \leq A \Leftrightarrow \forall a \in A : c \leq a$$

$$A \leq B \Leftrightarrow \forall a \in A, \forall b \in B : a \leq b.$$

Similarly for $A < c$, $A \neq B$.

Rank: $\rightarrow (A \leq b)$ is NOT equivalent to $A \leq b$

Example: " $\{3, 5\} \leq 4$ " is false $\Rightarrow \neg(\{3, 5\} \leq 4)$ is true

however, $\{3, 5\} \neq 4$ is also false \Rightarrow

$\Rightarrow (\{3, 5\} \neq 4)$ is NOT equivalent to $\neg(\{3, 5\} \leq 4)$

Def (Set Equality) $A, B \subseteq S$.

$$A = B \Leftrightarrow (\forall a \in A \exists b \in B : a = b) \wedge (\forall b \in B \exists a \in A : b = a)$$

Rank: A and B are not necessarily identical forms

$$\underline{\text{Example}}: \{\{1\}, \{11\}\} = \{\{-1\}, \{1\}, \{-1, 0, 11\}\},$$

$$\text{since } \{1\} = \{-1\}, \{1\} = \{1\}, \{11\} = \{-1, 0, 11\}.$$

Thm (Alternative Form of Comparison)

$$\text{Let } x = \{X_L | X_R\}, y = \{Y_L | Y_R\} \text{ Then}$$

$$x \leq y \Leftrightarrow (Y_R \notin x) \wedge (y \leq X_L)$$

$$\text{where } \begin{cases} Y_R \notin x \text{ means } \neg \exists y_R \in Y_R : y_R \leq x \\ y \notin X_L \text{ means } \neg \exists x_L \in X_L : y \leq x_L. \end{cases}$$

Rank: It's just a notational rewriting of Def 2.

Def (Parents (or options))

The members of the left set \bar{X}_L and the right set \bar{X}_R of a surreal number $x = \{\bar{X}_L \mid \bar{X}_R\}$ are called the parents (or options) of x .

Thm (Reflexivity of \leq) If $x \in S$, then $x \leq x$.

Proof: By induction.

Basic Case = $0 \leq 0$, which we've shown

Inductive Step. Suppose that for all parents of x , the theorem holds, that is, $\forall x_L \in X_L \wedge \forall x_R \in X_R$ we have $x_L \leq x_L \wedge x_R \leq x_R$. W.T.S.: $x \leq x$

$$\Leftrightarrow \underbrace{\neg \exists x_L \in X_L : x_B \leq x_L}_{(i)} \quad \wedge \quad \underbrace{\neg \exists x_R \in X_R : x_R \leq x}_{(ii)}$$

For (i), suppose not, i.e. $\exists x_L \in X_L : x \leq x_L$

$$\Rightarrow \neg \exists x'_L \in \bar{X}_L : x_L \leq x'_L \quad \wedge \quad \neg \exists x_{LR} \in \bar{X}_{LR} : x_{LR} \leq x$$

Since $x_L \leq x_L'$, then let $x_L' = x_L$. Then contradict!

Hence, (i) is true. Similarly, (ii) is true.

Coro (Reflexivity of $=$). If $x \in S$, then $x = x$.

Proof: $x \leq x \Leftrightarrow x = x$, by def of " $=$ ". \square

Thm Let A, A', B, B' be sets of surreal numbers, and let a_1, a_2, \dots be the members of A , let a'_1, a'_2, \dots be the members of A' .

Let b_1, b_2, \dots be the members of B , let

b'_1, b'_2, \dots be the members of B' . If

$\forall a_i \in A, \exists a'_j \in A', \text{ s.t. } a_i \leq a'_j \text{ and } \forall b_k \in B$

$\exists b_p^i \in B^i$, s.t. $b_k \leq b_p^i$, then $\{A|B\} \leq \{A'|B'\}$

Proof: W.T.S.: $\neg \exists a \in A : a \in \{A' \mid B'\}^{\leq a} \wedge \neg \exists b' \in B' : b' \in \{A \mid B\}$

For (i), suppose that $\exists a \in A : \{A' \mid B\} \leq a$. Then

$\neg \exists a' \in A' : a \leq a'$, Contradict ! \Rightarrow (i) holds.

For (ii), suppose that $\exists b' \in B' : b' \leq \{A|B\}^y$. Then
 $\neg \exists b \in B : b' \leq b$, Contradict! \Rightarrow (ii) holds. \square

Cor. If $A = A'$ and $B = B'$, then $\{A|B\} = \{A'|B'\}$

Proof: $A = A' \Rightarrow \forall a \in A, \exists a' \in A', \text{ s.t. } a = a' \Rightarrow a \leq a'$. Similarly for B and B' $\Rightarrow \{A|B\} \subseteq \{A'|B'\}$ & $\{A'|B'\} \subseteq \{A|B\}$ \square

Thm Let $x = \{A \mid B\}$ be a surreal number. Then $\forall a \in A : a < x$ and $\forall b \in B : x < b$, that is,

$A < x$ and $x < B$.

Proof: W.T.S.: ~~\exists~~ (a) $\forall a \in A : a < x \Leftrightarrow a \leq x \wedge a \neq x$
 ~~\exists~~ (b) $\forall b \in B : x < b \Leftrightarrow x \leq b \wedge b \neq x$

(a1): $a \leq x \Leftrightarrow \neg \exists a_L \in A_L : a \leq a_L \wedge \neg \exists b \in B : b \leq a$
 holds by def 1.

Suppose $\exists a_L \in A_L : a \leq a_L$. Then $\neg \exists a \in A : a_L \leq a$, which contradicts to the inductive step. \Rightarrow (a1) holds.

(a2): W.T.S.: $x \neq a \Leftrightarrow \neg(x \leq a) \Leftrightarrow \exists a' \in A : a \leq a' \vee \exists b \in B : a \leq b$
 The right-hand-side holds by def 1. \Rightarrow (a2) holds

(b1): $x \leq b \Leftrightarrow \neg \exists a \in A : a \leq b \wedge \neg \exists b_R \in B_R : b_R \leq x$
 holds by def 1

Suppose $\exists b_R \in B_R : b_R \leq x$. Then $\neg \exists b \in B : b \leq b_R$, which contradicts to the inductive step. \Rightarrow (b1) holds.

(b2): W.T.S.: $b \neq x \Leftrightarrow \neg(b \leq x) \Leftrightarrow \exists b_L \in B_L : b \leq b_L \vee \exists b \in B : b \leq b$
 By inductive step, left-hand-side holds
 \Rightarrow (b2) holds.

Then we're done.

Coro. If $x = \{a\}$, then $a < x$.

Thm (Transitive law for \leq). If $x, y, z \in S$, then $x \leq y \wedge y \leq z$ can imply $x \leq z$.

Proof: Suppose not, i.e.: $x \neq z$ when $x \leq y \wedge y \leq z$.

Take $p(x, y, z) \Leftrightarrow x \leq y \wedge y \leq z \wedge x \neq z$. Then

$\neg \exists x_L \in X_L : y \leq x_L$ (1); $\neg \exists y_R \in Y_R : y_R \leq z$ (2),

$\neg \exists y_L \in Y_L : z \leq y_L$ (3); $\neg \exists z_R \in Z_R : z_R \leq y$ (4)

$\exists x_L \in X_L : z \leq x_L \vee \exists z_R \in Z_R : z_R \leq x$.

Case 1: $\exists x_L \in X_L : z \leq x$. By (1), we have

$p(y, z, x_L) \Leftrightarrow y \leq z \wedge z \leq x_L \wedge y \neq x_L$

Case 2: $\exists z_R \in Z_R : z_R \leq x$. By (4), we have

$p(z_R, x, y) \Leftrightarrow z_R \leq x \wedge x \leq y \wedge z_R \neq y$

Hence, we have $p(x, y, z) \Leftrightarrow \exists x_L \in X_L : p(y, z, x_L) \vee$

$\exists z_R \in Z_R : p(z_R, x, y)$

Note that x_L, z_R are the parents of x and z respectively. Then by induction, we can prove it.

Basic case: $0 \leq 0 \wedge 0 \leq 0 \Rightarrow 0 = 0$.

Inductive step: Suppose that the transitive law holds for all parents of x, y, z . Then

$p(x, y, z) \Leftrightarrow \exists x_L \in X_L : p(y, z, x_L) \vee \exists z_R \in Z_R : p(z_R, x, y)$

Then repeat that ~~process~~ procedure and we
~~will~~ reach 0. Then we have contradiction,
since $p(0, 0, 0)$ ~~is~~ ^{is} false. Hence, the assumption
never holds. Then we're done. \square