

# Matrix Berlekamp-Massey

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## 1 Definitions

First thing: what are linearly recurrent sequences and minimal generators, when talking about matrix sequences? The definition is essentially the same as in the scalar case. There is a similar notion of generator of a sequence (which is a matrix): it is related to the denominator in some minimal fraction description of the generating series of the sequence.

We consider the following definition which extends the scalar case. It can be found in [2, Sec.3], and in [6, Def.4.2].

**Definition 1.1.** Let  $\mathcal{S} = (S_k)_{k \in \mathbb{Z}_{\geq 0}} \subset \mathbb{K}^{m \times n}$ . A vector  $\mathbf{p} = \sum_{0 \leq k \leq d} p_k X^k \in \mathbb{K}[X]^{1 \times m}$  of degree at most  $d$  is said to be a (linear recurrence) relation for  $\mathcal{S}$  if  $\sum_{k=0}^d p_k S_{\delta+k} = 0$  for all  $\delta \geq 0$ . Then,  $\mathcal{S}$  is said to be linearly recurrent if there exists a nontrivial relation for  $\mathcal{S}$ .

The set of relations for  $\mathcal{S}$  is a  $\mathbb{K}[X]$ -submodule of  $\mathbb{K}[X]^{1 \times m}$ , which **has rank  $m$  if  $\mathcal{S}$  is linearly recurrent** (I cannot find this result in the literature..!). This is showed in [3, Fact 1] for sequences having a *scalar* recurrence relation. Let us now link this with the property of being linearly recurrent as in the above definition.

**Lemma 1.2.** The sequence  $\mathcal{S}$  is linearly recurrent if and only if there exists a polynomial  $P(X) = \sum_{0 \leq k \leq d} p_k X^k \in \mathbb{K}[X]$  such that  $\sum_{k=0}^d p_k S_{\delta+k} = 0$  for all  $\delta \geq 0$ .

*Proof.* ? □

Then, the fact that the relation module has rank  $m$  is straightforward: since the sequence is linearly recurrent, it has a scalar recurrence polynomial  $P(X)$ , hence for each  $i$  the vector  $[0 \cdots 0 P(X) 0 \cdots 0]$  with  $P(X)$  at position  $i$  is a relation for  $\mathcal{S}$ .

A generating matrix for the sequence is a matrix whose rows form a generating set for the module of relations; it is said to be

- *minimal* if the matrix is reduced [7, 1];
- *ordered weak Popov* if the matrix is in weak Popov form [4] with pivots on the diagonal;
- *canonical* if the matrix is in Popov form [5, 1].

**Lemma 1.3.** Consider  $\mathcal{S} = (S_k)_{k \in \mathbb{Z}_{\geq 0}} \subset \mathbb{K}^{m \times n}$  and its generating series  $\mathbf{S} = \sum_{k \geq 0} S_k X^{-k} \in \mathbb{K}[[X]]^{m \times n}$ . Let  $\mathbf{P} \in \mathbb{K}[X]^{m \times m}$ , and consider  $\mathbf{Q} = \mathbf{P}\mathbf{S}$ .

*is a generating matrix for  $\mathcal{S}$  if and only if*

## 2 Computing a minimal matrix generators

We consider the following problem.

**Problem 1** – *Minimal generator*

*Input:*

- sequence  $\mathcal{S} = (S_k)_k \subset \mathbb{K}^{m \times n}$ ,
- degree bound  $\delta \in \mathbb{Z}_{\geq 0}$ .

*Assumptions:*

- the sequence  $\mathcal{S}$  is linearly recurrent,
- the canonical generating matrix of  $\mathcal{S}$  has degree at most  $\delta \in \mathbb{Z}_{\geq 0}$ .

*Output:* a minimal generating matrix for  $\mathcal{S}$ .

## References

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