

## Chapter 4

# Forward Differencing

Horner's algorithm is the fastest method for evaluating a polynomial at a single point. For a degree  $n$  polynomial, it requires  $n$  multiplies and  $n$  adds.

If a polynomial is to be evaluated at several evenly spaced values  $t, t + \delta, t + 2\delta, \dots, t + k\delta$ , the fastest method is to use *forward differences*.

Consider a degree 1 polynomial

$$f(t) = a_0 + a_1 t.$$

The difference between two adjacent function values is

$$\Delta_1(t) = f(t + \delta) - f(t) = [a_0 + a_1(t + \delta)] - [a_0 + a_1 t] = a_1 \delta.$$

Thus,  $f(t)$  can be evaluated at several evenly spaced points using the algorithm:

```
 $\Delta_1 = a_1 \delta$   
 $t_0 = 0$   
 $f(0) = a_0$   
for  $i = 1$  to  $k$  do  
   $t_i = t_{i-1} + \delta$   
   $f(t_i) = f(t_{i-1}) + \Delta_1$   
endfor
```

Thus, each successive evaluation requires only one add, as opposed to one add and one multiply for Horner's algorithm.

This idea extends to polynomials of any degree. For the quadratic case,

$$f(t) = a_0 + a_1 t + a_2 t^2.$$

The difference between two adjacent function values is

$$\Delta_1(t) = f(t + \delta) - f(t) = [a_0 + a_1(t + \delta) + a_2(t + \delta)^2] - [a_0 + a_1 t + a_2 t^2]$$

$$\Delta_1(t) = a_1 \delta + a_2 \delta^2 + 2a_2 t \delta.$$

We can now write

```

 $t_0 = 0$ 
 $f(0) = a_0$ 
for  $i = 1$  to  $k$  do
   $t_i = t_{i-1} + \delta$ 
   $\Delta_1(t_i) = a_1\delta + a_2\delta^2 + 2a_2t_{i-1}\delta$ 
   $f(t_i) = f(t_{i-1}) + \Delta_1(t_{i-1})$ 
endfor

```

In this case,  $\Delta(t)$  is a linear polynomial, so we can evaluate it as above, by defining

$$\Delta_2(t) = \Delta_1(t + \delta) - \Delta_1(t) = 2a_2\delta^2$$

and our algorithm now becomes

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 $t_0 = 0$ 
 $f(0) = a_0$ 
 $\Delta_1 = a_1\delta + a_2\delta^2$ 
 $\Delta_2 = 2a_2\delta^2$ 
for  $i = 1$  to  $k$  do
   $t_i = t_{i-1} + \delta$ 
   $f(t_i) = f(t_{i-1}) + \Delta_1$ 
   $\Delta_1 = \Delta_1 + \Delta_2$ 
endfor

```

It should be clear that for a degree  $n$  polynomial, each successive evaluation requires  $n$  adds and no multiplies! For a cubic polynomial

$$f(t) = a_0 + a_1t + a_2t^2 + a_3t^3,$$

the forward difference algorithm becomes

```

 $t_0 = 0$ 
 $f(0) = a_0$ 
 $\Delta_1 = a_1\delta + a_2\delta^2 + a_3\delta^3$ 
 $\Delta_2 = 2a_2\delta^2 + 6a_3\delta^3$ 
 $\Delta_3 = 6a_3\delta^3$ 
for  $i = 1$  to  $k$  do
   $t_i = t_{i-1} + \delta$ 
   $f(t_i) = f(t_{i-1}) + \Delta_1$ 
   $\Delta_1 = \Delta_1 + \Delta_2$ 
   $\Delta_2 = \Delta_2 + \Delta_3$ 
endfor

```

Several questions remain. First, what are the initial values for the  $\Delta_i$  if we want to start at some value other than  $t = 0$ . Second, what is a general equation for the  $\Delta_i$  for a general degree  $n$  polynomial  $f(t)$ . Also, what if our polynomial is not in power basis.

These questions can be answered almost trivially by observing the following. Since  $t_{i+1} = t_i + \delta$ , we have

$$\begin{aligned}\Delta_1(t_i) &= f(t_{i+1}) - f(t_i); \\ \Delta_j(t_i) &= \Delta_{j-1}(t_{i+1}) - \Delta_{j-1}(t_i), \quad j = 2, \dots, n; \\ \Delta_n(t_i) &= \Delta_n(t_{i+1}) = \Delta_n(t_{i+k}) = \text{a constant} \\ \Delta_{n+1} &= 0\end{aligned}$$

Thus, our initial values for  $\Delta_j(t_i)$  can be found by simply computing  $f(t_i), f(t_{i+1}), \dots, f(t_{i+n})$  and from them computing the initial differences. This lends itself nicely to a table. Here is the table for a degree four case:

$f(t_i)$	$f(t_{i+1})$	$f(t_{i+2})$	$f(t_{i+3})$	$f(t_{i+4})$
$\Delta_1(t_i)$	$\Delta_1(t_{i+1})$	$\Delta_1(t_{i+2})$	$\Delta_1(t_{i+3})$	
$\Delta_2(t_i)$	$\Delta_2(t_{i+1})$	$\Delta_2(t_{i+2})$		
$\Delta_3(t_i)$	$\Delta_3(t_{i+1})$			
$\Delta_4(t_i)$				
0	0	0	0	0

To compute  $f(t_{i+5})$ , we simply note that every number  $R$  in this table, along with its right hand neighbor  $R_{right}$  and the number directly beneath it  $R_{down}$  obey the rule

$$R_{right} = R + R_{down}.$$

Thus, we can simply fill in the values

$$\begin{aligned}\Delta_4(t_{i+1}) &= \Delta_4(t_i) + 0 \\ \Delta_3(t_{i+2}) &= \Delta_3(t_{i+1}) + \Delta_4(t_{i+1}) \\ \Delta_2(t_{i+3}) &= \Delta_2(t_{i+2}) + \Delta_3(t_{i+2}) \\ \Delta_1(t_{i+4}) &= \Delta_1(t_{i+3}) + \Delta_2(t_{i+3}) \\ f(t_{i+5}) &= f(t_{i+4}) + \Delta_1(t_{i+4})\end{aligned}$$

Note that this technique is independent of the basis in which  $f(t)$  is defined. Thus, even if it is defined in Bernstein basis, all we need to do is to evaluate it  $n + 1$  times to initiate the forward differencing.

For example, consider the degree 4 polynomial for which  $f(t_i) = 1$ ,  $f(t_{i+1}) = 3$ ,  $f(t_{i+2}) = 2$ ,  $f(t_{i+3}) = 5$ ,  $f(t_{i+4}) = 4$ . We can compute  $f(t_{i+5}) = -24$ ,  $f(t_{i+6}) = -117$ , and  $f(t_{i+7}) = -328$  from the following difference table:

$t :$	$t_i$	$t_{i+1}$	$t_{i+2}$	$t_{i+3}$	$t_{i+4}$	$t_{i+5}$	$t_{i+6}$	$t_{i+7}$
$f(t) :$	1	3	2	5	4	-24	-117	-328
$\Delta_1(t) :$	2	-1	3	-1	-28	-93	-211	
$\Delta_2(t) :$	-3	4	-4	-27	-65	-118		
$\Delta_3(t) :$	7	-8	-23	-38	-53			
$\Delta_4(t) :$	-15	-15	-15	-15	-15			
$\Delta_5(t) :$	0	0	0	0	0	0	0	0

**Example**

For a certain cubic polynomial  $f(t)$ , we have:

$$f(1) = 1; \quad f(2) = 2; \quad f(3) = 4; \quad f(5) = 15.$$

Solve for  $f(4)$  using forward differencing.

**Solution**

$$\begin{array}{rcccccc} t : & 1 & 2 & 3 & 4 & 5 \\ f(t) : & 1 & 2 & 4 & x & 15 \\ \Delta_1(t) : & 1 & 2 & x-4 & 15-x & \\ \Delta_2(t) : & 1 & x-6 & 19-2x & & \\ \Delta_3(t) : & x-7 & 25-3x & & & \\ \Delta_4(t) : & 32-4x & & & & \end{array}$$

For a cubic polynomial,  $\Delta_4(t) = 0$ , so  $32 - 4x = 0$  and  $x = 8$ .

## 4.1 Choosing $\delta$

This raises the important question of how to select an appropriate value for  $\delta$  when using forward differencing to plot a curve. One way to determine  $\delta$  is so that distance from the the curve to its polygonal approximation lies within a tolerance. A discussion of how to chose  $\delta$  that will satisfy such a requirement is found in Section 10.6.

A second criterion that might be used to choose  $\delta$  arises in the problem of rasterizing a curve. This means to “turn on” a contiguous series of pixels that the curve passes through, without any gaps. If the control points of the degree  $n$  Bézier curve are  $\mathbf{P}_i = (x_i, y_i)$  in *pixel coordinates*, then let

$$d = \max\{x_i - x_{i-1}, y_i - y_{i-1}\} \quad i = 1, \dots, n$$

If you now evaluate the curve at  $n * d + 1$  evenly spaced values of  $t$  and paint each resulting pixel, there will be no gaps in the drawing of the curve. Another way to say this, compute

$$\delta = \frac{1}{n * d}.$$

Then, compute the points  $\mathbf{P}(i * \delta)$ ,  $i = 0 \dots, n * d$ . Note that  $n * d$  is an upper bound on the magnitude of the  $x$  or  $y$  component of the first derivative of the curve.