Analysis of Algorithms 2020/2021

Practice 3

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| Code | Plots | Documentation | Total |
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**1. Introduction.**

The main goal of this assignment is to implement searching algorithms such as *linear search* and *binary search* using the programming language C. We also have to implement a *linear search* “variation”, which whenever the key is found, it is swapped with the previous element.

Once the algorithms are implemented, we have to create functions which are going to test the performance of these algorithms and store the data related with this performance.

The keys are going to be searched in dictionaries, this is going to be the data structure used in order to store our tables. We will implement some primitives to manage the content and the memory usage of the dictionary.

We are also given 2 generators which generates keys. This is going to be useful for having specific keys to search, but these 2 generators doesn’t work in the same way, the first one *uniform key generator* generates all the keys given a range, for example if we give to the generator a number N, the generator will fill an array with the numbers from 1 to N, n\_times times. The second generator *potential key generator*, creates keys in way which the smaller ones are more likely to appear than the bigger ones.

Once we execute make in the terminal we would have the executables for exercise 1 and 2. In order to execute the first exercise we have to introduce this in the terminal:

./exercise1 -size <size> -key <key>

We just have to introduce the size and the key we want to look for, it will output something like this:

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Search done with linear search:

Key 12 found in position 7278 in 7279 basic op.

Search done with auto-linear search:

Key 12 found in position 7278 in 7279 basic op.

Search done with binary search:

Key 12 found in position 11 in 13 basic op.

\*We modified the exercise1.c, so now it uses all algorithms.

For the exercise 2 we have to introduce this:

./exercise2 -num\_min <tamaño minimo> -num\_max <tamaño máximo> -incr <incremento> -n\_times <número de veces> -outputFile <nombre del archivo>

After executing the previous command we will have of the information of the performance stored in the *outputFile* and the following output:

Running exercise2

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Correct output

**2. Objectives**

2.1 Implementation of the dictionary and the searching algorithms (search.c).

2.1.1 Dictionary.

For this section we have to build the primitives for a given data structure:

**typedef** **struct** dictionary {

**int** size; /\* table size \*/

**int** n\_data; /\* number of entries in the table \*/

**char** order; /\* sorted or unsorted table \*/

**int** \*table; /\* data table \*/

} DICT, \*PDICT;

This is the data structure, it contains a table, the size of the table, the number of elements in the table and the order of the table (sorted or not sorted).

We also have to implement the following primitives:

PDICT **init\_dictionary** (**int** size, **char** order);

**void** **free\_dictionary**(PDICT pdict);

**int** **insert\_dictionary**(PDICT pdict, **int** key);

**int** **massive\_insertion\_dictionary** (PDICT pdict,**int** \*keys, **int** n\_keys);

With these primitives we will manage the content of each dictionary.

2.1.2 Searching algorithms.

Now that we have the structure implemented, we need to implement the different searching algorithms, but first we need to create a generic function to search:

**int** **search\_dictionary**(PDICT pdict, **int** key, **int** \*ppos, pfunc\_search method);

With this function we can search in a dictionary with the searching function we want.

**int** **bin\_search**(**int** \*table,**int** F,**int** L,**int** key, **int** \*ppos);

**int** **lin\_search**(**int** \*table,**int** F,**int** L,**int** key, **int** \*ppos);

**int** **lin\_auto\_search**(**int** \*table,**int** F,**int** L,**int** key, **int** \*ppos);

**bin\_search** refers to the binary search algorithm, **lin\_search** refers to the linear search algorithm, and **lin\_auto\_search** refers to a variation of the linear search algorithm which whenever the key is found it is being swapped with the previous element.

2.2 Comparison of search efficiency (times.c).

In this section we have to implement 2 new functions which will measure the efficiency of the previously mentioned searching algorithms. These are the 2 functions:

**short** **generate\_search\_times**(pfunc\_search method, pfunc\_key\_generator generator,

**int** order, **char**\* file,

**int** num\_min, **int** num\_max,

**int** incr, **int** n\_times);

**short** **average\_search\_time**(pfunc\_search metodo, pfunc\_key\_generator generator,

**int** order,

**int** N,

**int** n\_times,

PTIME\_AA ptime);

**generate\_search\_times** calls **average\_search\_time** with the sizes starting from *num\_min* to *num\_max* with an increment of *incr*. In order to store all the information, we will call the function *save\_time\_table* which was created in the assignment 1.

These 2 functions are very similar to the ones implemented in the assignment 1.

**3. Tools and methodology**

Here, you include the development environment (Windows, Linux, MacOS) and the tolos you have used (Netbeans, Eclipse, gcc, Valgrind, Gnuplot, Sort, uniq, etc) and what methodologies and solutions for the problemas you have used for each exercise, as well as the tests you have applied to the developed programs.

3.1 Section 1

Methodology and solution for section 1

3.2 Section 2

Methodology and solution for section 2

**4. Source code**

Here, include the source code only **for the routines you have developed** for each exercise.

4.1 Section 1

4.2 Section 2

**5. Results, plots**

Results obtained for the different exercises, and their plots (if any).

5.1 Section 1

Results for section 1.

5.2 Section 2

Section 2 results.

Plot comparing the average number of BOs of linear and binary search approaches, comments to the plot.

Plot comparing the average clock time for the linear and binary search approaches, comments to the plot.

Plot comparing the average number of BOs of the binary and auto-organized linear search (for n\_times=1, 100 y 10000), comments to the plot.

Plot comparing the maximum number of BOs of the binary and auto-organized linear search (for n\_times=1, 100 y 10000), comments to the plot.

Plot comparing the minimum number of BOs of the binary and auto-organized linear search (for n\_times=1, 100 y 10000), comments to the plot.

Plot comparing the average clock time for the binary and auto-organized linear search (for n\_times=1, 100 y 10000), comments to the plot.

**5. Response to the theoretical questions.**

5.1 *Which is the basic operation of lin search, bin search and lin auto search?*

**The comparison between the key and an element of the table**.

In linear search and linear auto search, the basic operation is a key comparison between the key and the element in the table traversed at that moment. **if** (table[i] == key)

For binary search the basic operation is a key comparison between the key and the middle element of the table in that iteration of the loop (the middle element changes in every iteration of the loop). **if**(key == table[middle])

5.2 *Give the execution times, in terms of the input size n for the worst WSS(n) and best BSS(n) cases of bin search and lin search. Use the asymptotic notation (O,Θ,o,Ω,etc) as long as you can.*

In each iteration of the algorithm we divide the table and we check the middle element of that division, so given a table with size n, the maximum times we can divide the table is , so in the last division we would have in the middle of the table (the table length is 1) the element we are looking for.

In the case that the key is in the middle of the table.

Linear search goes from the beginning of the table to the end, so in the case that the element is in the last position of the table, we would have made comparisons.

In the case that the key is in the first element.

5.3 *When lin auto search and the given not-uniform key distribution are used, how does it vary the position of the elements in the list of keys as long as the number of searches increases?*

The potential key generator generates more the smaller keys, so, having a table with numbers going from 1 to N (not sorted), using *lin auto search* and using the potential key generator for generating keys in a range of 1 to N several times, we would see that the after some searching the smaller keys are going to be at the beginning of the table. If we continue searching the table would have the smaller values at the beginning and the bigger ones at the end, this doesn’t mean that is going to be sorted.

1 will be the first, it has a probability of 50%. It will be 1st most of the times not always, because 2 will be the second because it is the second most probable number that can be generated (17%), but if 2 is generated then it is swapped with one, and now the 2 is the first element, then 1 could be generated so it will be swapped with 2 and so on …

To sum up, **the smaller numbers are going to be at the beginning**.

5.4 *Which is the average execution time of lin auto search as a function of the number of elements in the dictionary n for the given not uniform key distribution? Consider that a large number of searches have been conducted and the list is in a quite stable state.*

The average execution time function looks like a 1/n where n ≥1. At the beginning the algorithm will struggle to find keys, but once several searches are done, the most likely to search keys will be at the beginning, so the searches will be faster.

5.5 *Justify as formally as you can the correction (in other words, why it searches well) of the bin search algorithm.*

It works well because it is sorted, and it divides de table.

This means that we look for a key which ,

we would have to look in the first half of the table ().

In case that , we would look in the second half of the table.

then we would have found the key.

The process of diving the table by 2 in each iteration is very powerful, because we can divide a table by 2 a maximum number of times equal to .

We can check this with an example:

Let us imagine that we have a table of 2.000.000 elements. It is a big table but, how many times do we have to divide it until the result of the division is 1?

2.000.000/2🡪1.000.000/2🡪500.000/2🡪250.000/2🡪125.000/2🡪62.500/2🡪31.250/2🡪15.625/2🡪7812/2🡪3906/2🡪1953/2🡪976/2🡪488/2🡪244/2🡪122/2🡪61/2🡪 30/2🡪15/2🡪7/2🡪3/2🡪1.

20 = log2 2.000.000 divisions we performed until we had one, this will be as if we have the key in the last or first position. The algorithm is very fast looking for keys.

**6. Conclusions.**

Final discussion about the practice and the obtained results