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Networked Systems Coursework 1

Task 1

- 1. Modify topo1.py to create four hosts (h1-h4) connected to the same switch.
- 2. Run h1 ping h4, iperf h1 h4, dump, and pingall.

Solutions

- 1. topo1.py can be found in the .zip file
- 2. Below are images evidencing the running of the commands in the topo1 network topology

```
vagrant@mnvm:~$ sudo python3 topo1.py
mininet> host1 ping host4
PING 192.168.0.4 (192.168.0.4) 56(84) bytes of data.
64 bytes from 192.168.0.4: icmp seq=1 ttl=64 time=0.000 ms
64 bytes from 192.168.0.4: icmp seq=2 ttl=64 time=0.000 ms
64 bytes from 192.168.0.4: icmp seq=3 ttl=64 time=3.71 ms
64 bytes from 192.168.0.4: icmp seq=4 ttl=64 time=0.000 ms
64 bytes from 192.168.0.4: icmp seq=5 ttl=64 time=0.000 ms
64 bytes from 192.168.0.4: icmp_seq=6 ttl=64 time=7.10 ms
64 bytes from 192.168.0.4: icmp_seq=7 ttl=64 time=4.01 ms
64 bytes from 192.168.0.4: icmp_seq=8 ttl=64 time=3.59 ms
64 bytes from 192.168.0.4: icmp seq=9 ttl=64 time=3.51 ms
64 bytes from 192.168.0.4: icmp seq=10 ttl=64 time=3.48 ms
64 bytes from 192.168.0.4: icmp_seq=11 ttl=64 time=0.310 ms
64 bytes from 192.168.0.4: icmp_seq=12 ttl=64 time=3.89 ms
64 bytes from 192.168.0.4: icmp seq=13 ttl=64 time=0.999 ms
64 bytes from 192.168.0.4: icmp_seq=14 ttl=64 time=0.662 ms
^C
--- 192.168.0.4 ping statistics ---
14 packets transmitted, 14 received, 0% packet loss, time 48592ms
rtt min/avg/max/mdev = 0.000/2.232/7.102/2.145 ms
```

Fig 1.1: Running h1 ping h4

```
vagrant@mnvm:~$ sudo python3 topo1.py
mininet> iperf host1 host4
*** Iperf: testing TCP bandwidth between host1 and host4
*** Results: ['224 Mbits/sec', '248 Mbits/sec']
```

Fig 1.2: Running iperf h1 h4

```
vagrant@mnvm:~$ sudo python3 topo1.py
mininet> iperf host1 host4
*** Iperf: testing TCP bandwidth between host1 and host4
*** Results: ['224 Mbits/sec', '248 Mbits/sec']
mininet> dump
<Host host1: host1-eth0:192.168.0.1 pid=33734>
<Host host2: host2-eth0:192.168.0.2 pid=33736>
<Host host3: host3-eth0:192.168.0.3 pid=33738>
<Host host4: host4-eth0:192.168.0.4 pid=33740>
<OVSSwitch switch1: lo:127.0.0.1,switch1-eth1:None,switch1-eth2:None,switch1-eth3:None,switch1-eth4:None pid=33745>
<OVSController c0: 127.0.0.1:6653 pid=33727>
```

Fig 1.3: Running dump

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```
mininet> pingall

*** Ping: testing ping reachability
host1 -> host2 host3 host4
host2 -> host1 host3 host4
host3 -> host1 host2 host4
host4 -> host1 host2 host3

*** Results: 0% dropped (12/12 received)
```

Fig 1.4: Running pingall

Task 2

- 1. Run h1 ping h4.
- 2. Run dpctl dump-flows and inspect the rules the controller installed. Provide screenshots in your report.
- 3. Briefly explain how 12_learning.py works.
- 4. When you run h1 ping h4 is there any chance you receive an ICMP packet in h2? Explain your answer.

Solutions

```
vagrant@mnvm:~$ sudo python3 topo2.py
mininet> h1 ping h4
PING 10.0.1.3 (10.0.1.3) 56(84) bytes of data.
64 bytes from 10.0.1.3: icmp seq=3 ttl=64 time=0.000 ms
64 bytes from 10.0.1.3: icmp seq=4 ttl=64 time=0.000 ms
64 bytes from 10.0.1.3: icmp seq=5 ttl=64 time=3.46 ms
64 bytes from 10.0.1.3: icmp seq=6 ttl=64 time=0.000 ms
64 bytes from 10.0.1.3: icmp seq=7 ttl=64 time=0.000 ms
64 bytes from 10.0.1.3: icmp_seq=8 ttl=64 time=0.000 ms
64 bytes from 10.0.1.3: icmp_seq=9 ttl=64 time=5.28 ms
64 bytes from 10.0.1.3: icmp seq=10 ttl=64 time=31.6 ms
64 bytes from 10.0.1.3: icmp_seq=11 ttl=64 time=1.60 ms
64 bytes from 10.0.1.3: icmp_seq=12 ttl=64 time=1.95 ms
64 bytes from 10.0.1.3: icmp_seq=13 ttl=64 time=0.801 ms
64 bytes from 10.0.1.3: icmp seq=14 ttl=64 time=1.43 ms
^C
--- 10.0.1.3 ping statistics ---
14 packets transmitted, 12 received, 14.2857% packet loss, time 63315ms
rtt min/avg/max/mdev = 0.000/3.838/31.552/8.502 ms
```

Fig 2.1: Running h1 ping h4

- 1. Above is the console log when h1 ping h4 is run.
- 2. Below is console log when dpctl dump-flows is run.

Fig 2.2: Running dpctl dump-flows

3. The general idea of the learning switch is to dynamically 'learn' and build up a MAC address-to-port mapping table to decide where to forward incoming packets based on a source and destination MAC address.

The process can be split into several phases: * **Learning Phase**: The switch observes which port a source MAC address is received at and populates the mapping table with this relation.

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- **Forwarding Logic**: If the destination MAC address is present in the mapping table, the switch will forward the packet through the associated port. However, if the destination MAC address is unknown, the switch will flood the packet across all channels to ensure the packet has reached its destination.
- **Dropping Rules**: For certain types of traffic (link-local or if the destination address is Bridge Filtered), the switch will drop the packets associated. The switch also conducts a sanity check ensuring that packets with the same source and destination addresses are dropped temporarily as it suggests there to be an error or loop in the network.
- **OpenFlow Rule Set-Up**: Once a packet flow (Source MAC <-> Input Port <-> Destination MAC) is learnt by the switch, the program leverages OpenFlow table by installing a flow rule for that specific flow. This allows for traffic to be automatically forwarded without the involvement of the controller. If the rule is unused for a period of 10s, the rule is discarded. Similarly, after a fixed period of 30s, the installed rule is removed.
- 4. There is no chance for h1 ping h4 to cause ICMP packets to reach h2. The reasoning behind this is that as h1 pings h4, once the initial ARP request is complete and h4's MAC address is known, the controller ensures that subsequent ICMP packets will only be directed to h4's port. Therefore, h2 will never receive any ICMP packets destined for h4 from h1

Task 3

- 1. Run pingall. h1 and h4 should be able to ping each other and h2 and h3 as well. Provide screenshots in your report.
- 2. Run dpctl dump-flows to see the rules you inserted. Provide the equivalent screenshot.

Solutions

Fig 3.1: Running pingall

- 1. Above is the result of runing pingall with the new OpenFlow rules. We can observe that the new rules enable h1 to reach h4 and h2 to reach h3.
- 2. Below is the result of running dpctl_dump-flows which shows the rules inserted in the switch. We can observe the three separate rules for ARP, ICMP and the rest of the traffic. Noting that ARP is given the highest priority to flood the network, followed by ICMP and then the catch-all rule to drop all other packets.

```
mininet> dpctl dump-flows

*** s1

cookie=0x0, duration=115.211s, table=0, n_packets=8, n_bytes=336, priority=100,arp actions=FLOOD cookie=0x0, duration=115.211s, table=0, n_packets=8, n_bytes=784, priority=90,icmp actions=FLOOD cookie=0x0, duration=115.211s, table=0, n_packets=30, n_bytes=2316, priority=10 actions=drop
```

Fig 3.2: Running dpctl_dump-flows

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